



Antonius RC

SISTEM OPERASI 12 VIRTUAL MACHINE

OVERVIEW

- ▶ **Fundamental idea** – abstract hardware of a single computer into several different execution environments
 - **layered approach model**
- ▶ Several components:
 - **Host** – underlying hardware system
 - **Virtual machine manager (VMM)** or **hypervisor** – creates and runs virtual machines by providing interface that is **identical** to the host
 - **Guest** – process provided with virtual copy of the host
 - Usually an operating system
- ▶ Single physical machine can run **multiple** operating systems concurrently, each in its **own virtual machine**



VIRTUAL MACHINE (1)

- Mesin virtual memperlakukan hardware dan sistem operasi seolah-olah berada pada level yang **sama** sebagai hardware.
- Pendekatan mesin virtual menyediakan sebuah interface yang **identik** dengan seluruh hardware yang ada.
- Sistem Operasi **host** membuat **ilusi** dari banyak proses, masing-masing dieksekusi pada prosesoranya sendiri dengan virtual memorinya sendiri.
- Setiap **guest** menyediakan sebuah (virtual) copy dari semua hal yang ada pada komputer
- VM ada 2: **system VM dan application VM**



VIRTUAL MACHINE (2)

- First appeared commercially in IBM mainframes in 1972
- Fundamentally, **multiple execution environments** (different operating systems) can share the same hardware
- **Protect** from each other VM
- **Communate** with each other, other physical systems via networking
- It using “**Open Virtual Machine Format**”, standard format of virtual machines, allows a VM to run within many different virtual machine (host) platforms



KEUNTUNGAN VM

- Keamanan
 - Dari virus, serangan hacker
- Kemudahan instalasi
 - Tinggal copy paste image, cocok untuk semua hardware
- Cocok digunakan dalam penelitian dan percobaan
- Mendukung green computing
 - Run multiple, different OSes on a single machine
- Murah



KEUNTUNGAN VM (2)

- **Templating** – create an OS + application VM, provide it to customers, use it to create multiple instances of that combination
- **Live migration** – move a running VM from one host to another!
 - No interruption of user access
- When all those features taken together -> **cloud computing**

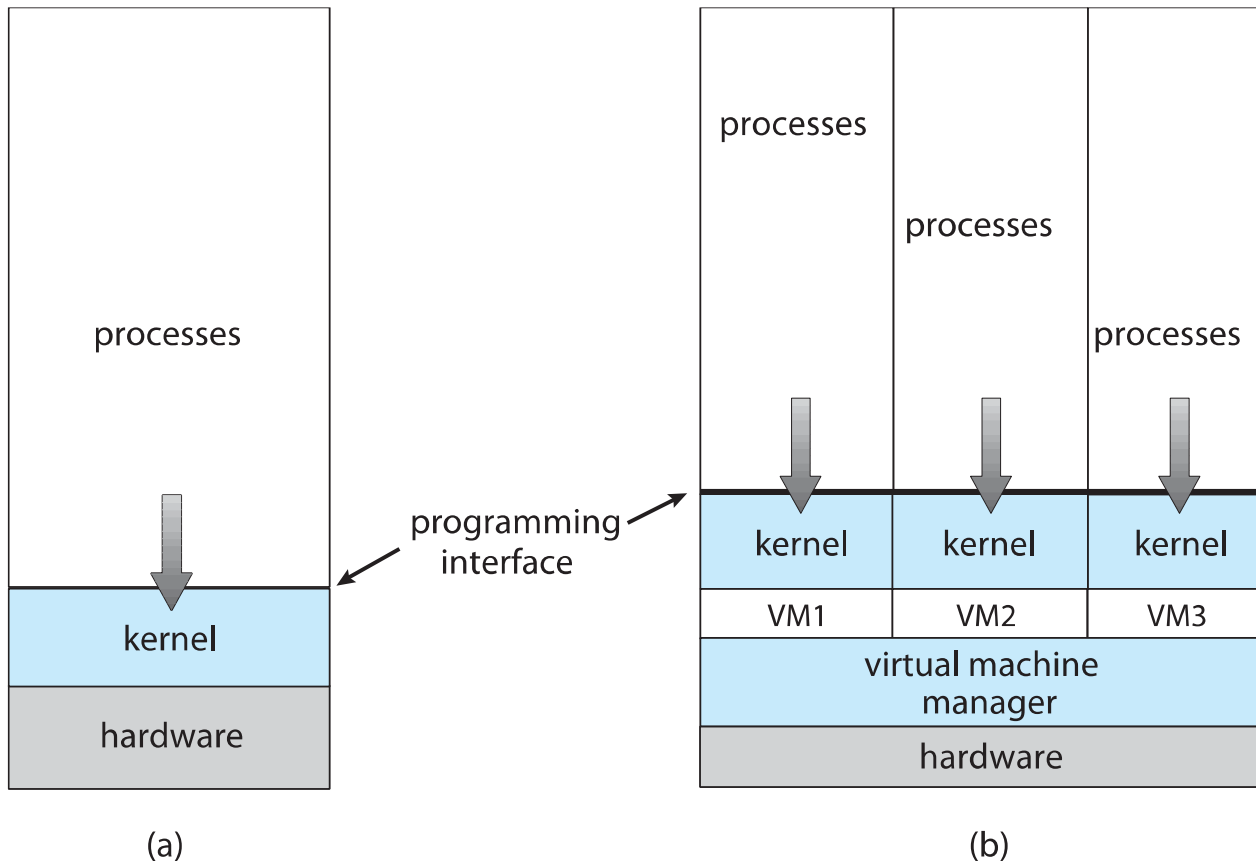


KERUGIAN VM

- Waktu yang dibutuhkan I/O bisa lebih cepat (karena ada spooling), tapi bisa lebih lambat (karena diinterpreted)
- Tidak semua aplikasi kompatibel
- Membutuhkan hardware dengan spesifikasi yang cukup tinggi



VM SYSTEM MODELS

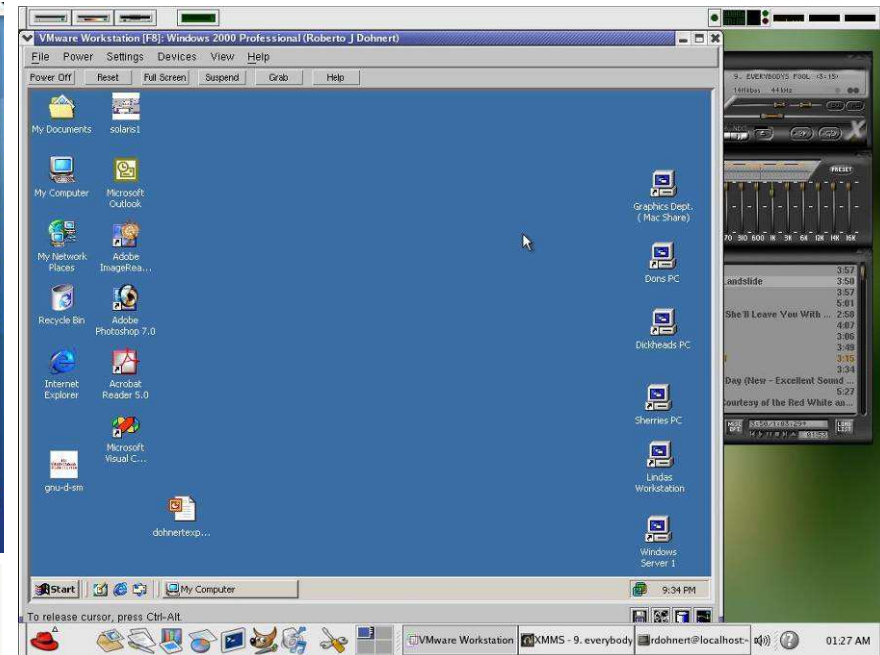
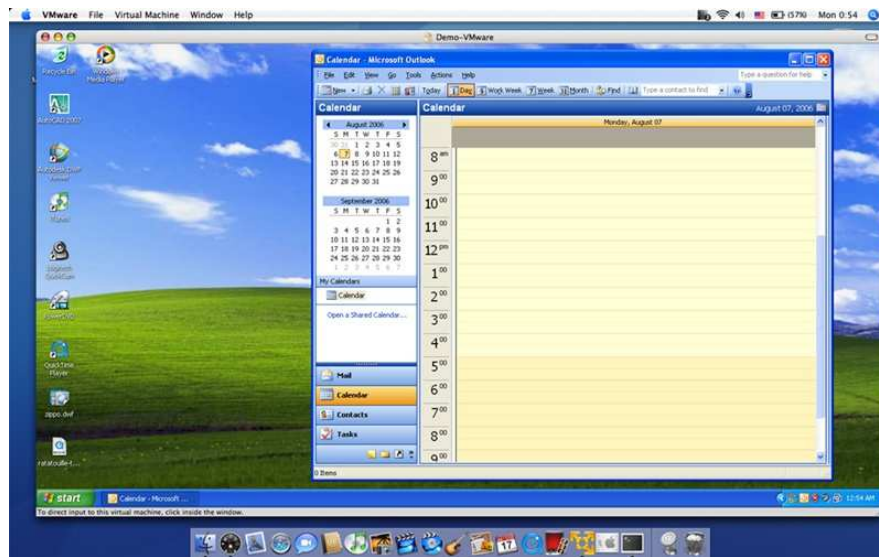


(a) Nonvirtual machine

(b) Virtual machine



CONTOH: VMWARE



IMPLEMENTATION OF VMMS

- **Type 0 hypervisors - Hardware-based solutions** that provide support for virtual machine creation and management via **firmware**
 - Hypervisor is unhosted
 - IBM LPARs and Oracle LDOMs are examples
- **Type 1 hypervisors (OS)** – includes **general-purpose operating systems** that provide standard functions as well as VMM functions
 - Hypervisor is controlled and assisted by a root or parent operating system
 - Including Microsoft Windows Server with HyperV and RedHat Linux with KVM

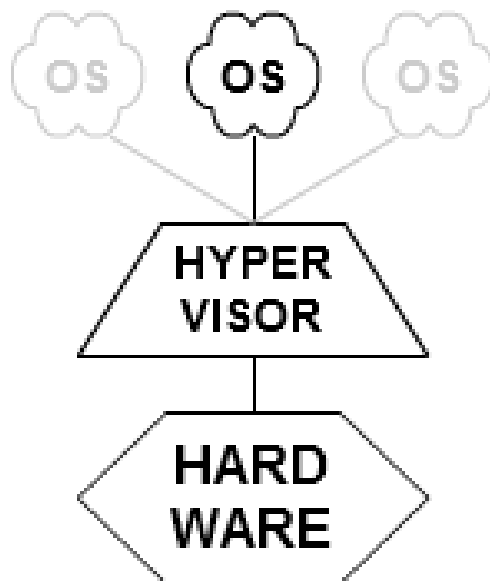


IMPLEMENTATION OF VMMS

- **Type 2 hypervisors (Application) - Applications** that run on standard operating systems but provide VMM features to guest operating systems
 - hypervisor is integrated into a host OS
 - Including VMware Workstation and Fusion, Parallels Desktop, and Oracle VirtualBox

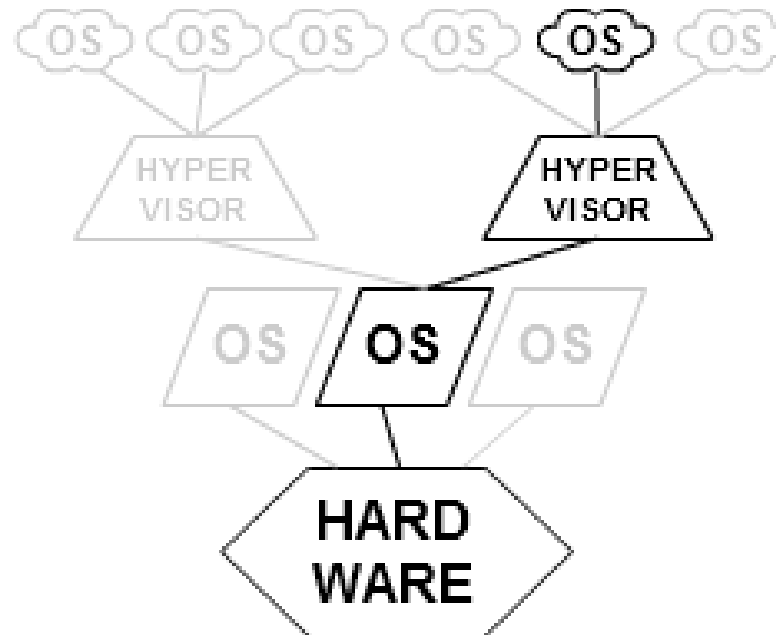


TYPE 1 VS TYPE 2



TYPE 1

native
(bare metal)



TYPE 2

hosted

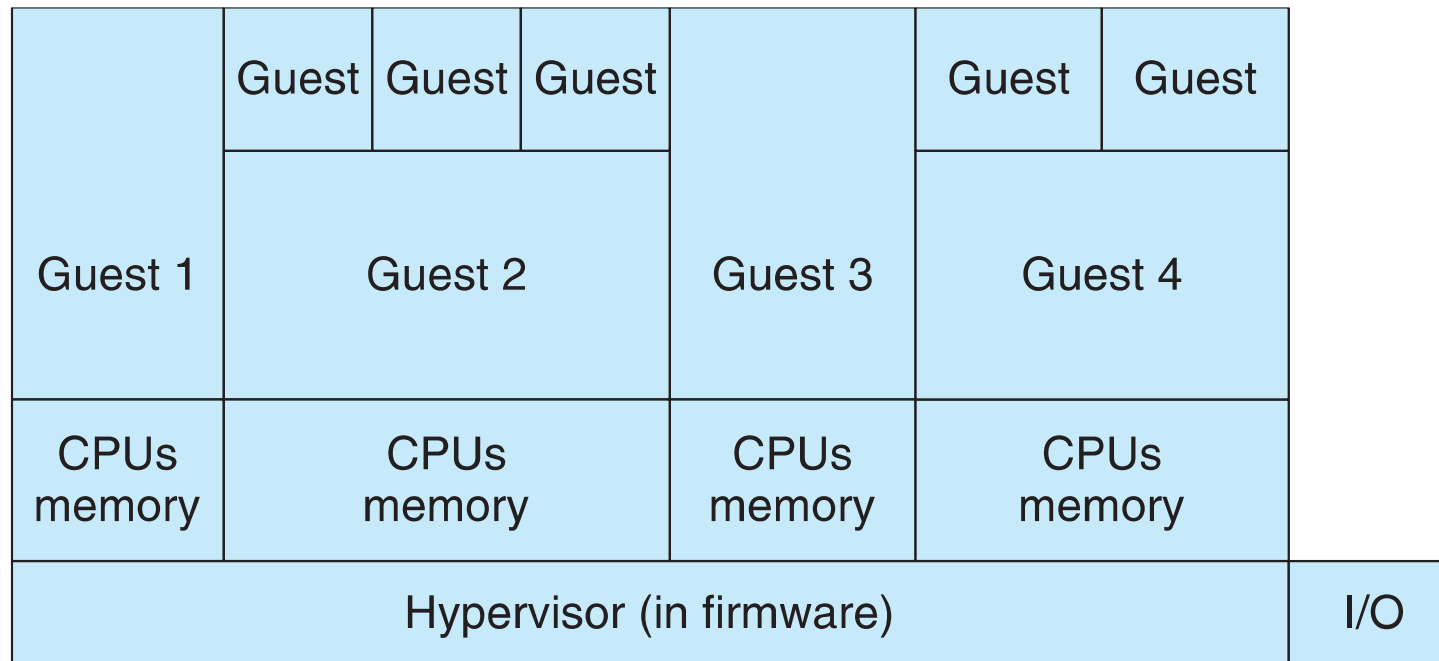


TYPES OF VMS – TYPE 0 HYPERVISOR

- Old idea, under many names by Hardware manufacturers
 - A HW feature implemented by **firmware**
 - OS need to nothing special, because VMM is in firmware
 - **Smaller** feature set than other types
 - Each guest has dedicated HW
 - Can be guest in guest VMM
- I/O a challenge as: **difficult to have enough** devices, controllers to dedicate to each guest



TYPE 0 HYPERVISOR



TYPES OF VMS – TYPE 1 HYPERVISOR

- Commonly found in **company data centers**
- **Special purpose operating systems** that run natively on HW
 - Run in **kernel mode**
 - Guests generally don't know they are running in a VM
 - Implement **device drivers** for host HW
 - In many ways, treat **guests OSes as just another process**
- Ex: RedHat Enterprise Linux with KVM, Windows with Hyper-V, Oracle Solaris



TYPES OF VMS – TYPE 2 HYPERVISOR

- **Very little** OS involvement in virtualization
- VMM is **simply another process**, run and managed by host
 - Even the host doesn't know they are a VMM running guests
- Tend to have poorer overall performance because **can't take advantage of some HW features**
- But also a benefit because **require no changes to host OS**



IMPLEMENTATION OF VMMS (CONT.)

- Other variations of Virtual Machine:
 - **Paravirtualization** - Technique in which the **guest operating system is modified** to work in cooperation with the vmm to optimize performance
 - Ex: VMWare dan Linux
 - **Programming-environment virtualization** - VMMS do not virtualize real hardware but instead create an **optimized virtual system framework**
 - Used by Oracle Java and Microsoft.Net by its Framework

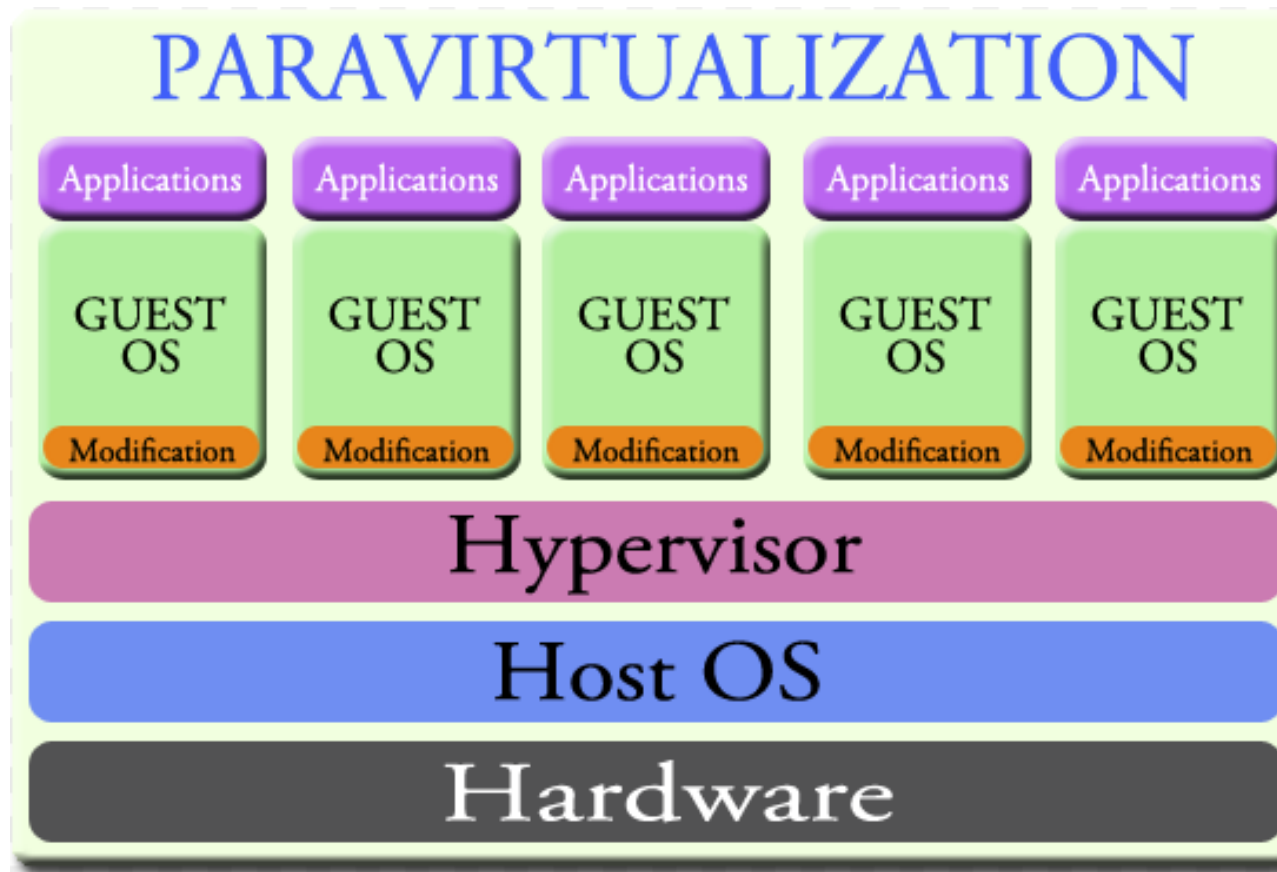


TYPES OF VMS – PARAVIRTUALIZATION

- ▶ Does not fit the definition of virtualization – VMM not presenting an **exact duplication** of underlying hardware
 - Guest had to be **modified** to use run on paravirtualized VMM
 - Less needed as hardware support for VMs grows
- ▶ Ex: Xen, Microsoft Hyper-V and VMware ESX Server



PARAVIRTUALIZATION ARCHITECTURE



TYPES OF VMS – PROGRAMMING ENVIRONMENT VIRTUALIZATION

- n Programming language is designed to run within **custom-built virtualized environment**
 - l For example Oracle Java has many features that depend on running in **Java Virtual Machine (JVM)**
- n Its using API
- n Its create **new layer** between OS and programming language
- n Example: JVM
 - n JVM compiled to run on many systems (Windows, Linux, etc)
 - n Programs written in Java run in the JVM no matter the underlying system
 - n It's interpreted



IMPLEMENTATION OF VMMS (CONT.)

- ▶ **Emulators** – Allow applications written for one hardware environment **to run on a very different hardware** environment, such as a different type of CPU
 - ▶ NeoGeo Emulator, Android Emulator, Windows Phone Emulator, Cygwin, MAME
- ▶ **Application containment** - Not virtualization at all but rather provides virtualization-like features by **segregating applications** from the operating system, making them **more secure, manageable**
 - ▶ **Sandboxie**: runs your programs in an isolated space which prevents them from making permanent changes to other programs and data in your computer.
 - Including Oracle Solaris Zones, Sanboxie (<http://www.sandboxie.com/>)

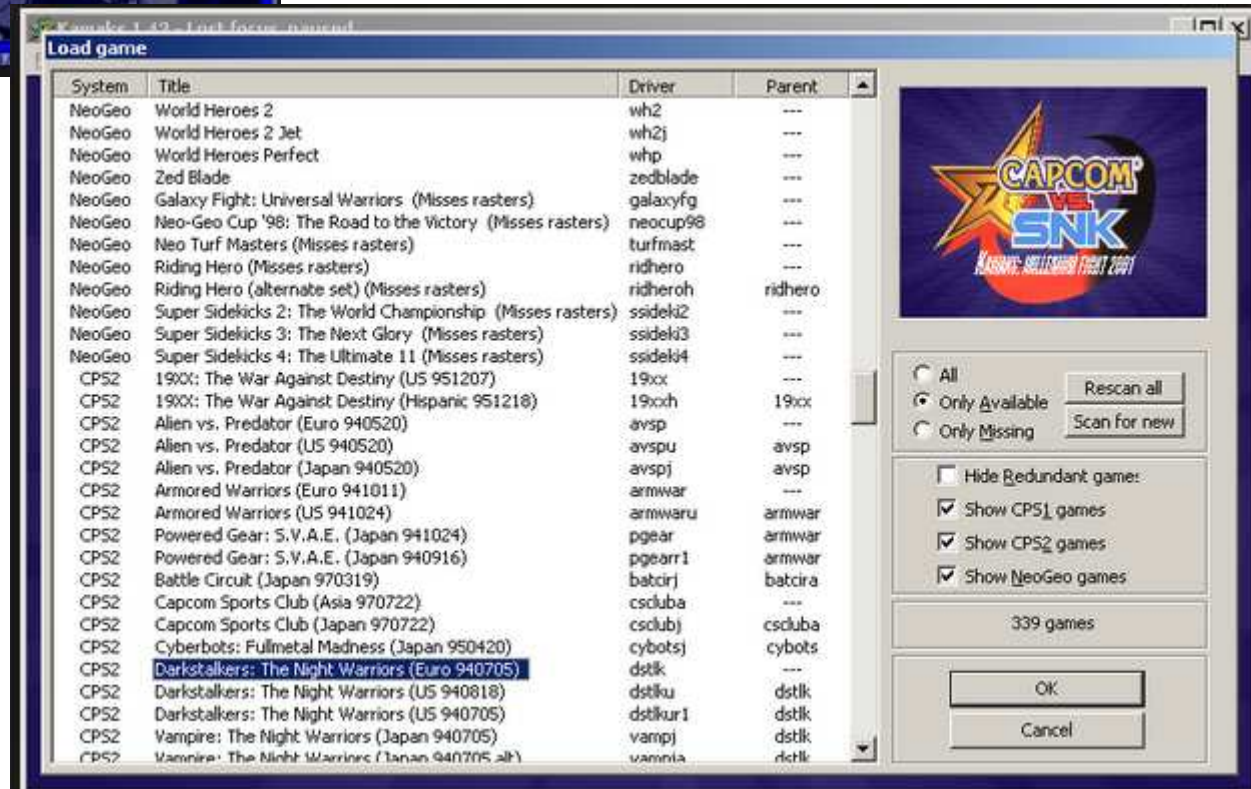


TYPES OF VMS – EMULATION / EMULATOR

- ▶ Another (older) way for running one operating system on a different operating system
 - Virtualization requires underlying **CPU to be same as guest was compiled for**
 - Emulation allows **guest to run on different CPU**
- ▶ Necessary to **translate** all guest instructions from guest CPU to native CPU
- ▶ Useful when host system has **one architecture**, guest compiled for **other architecture**
 - Ex: CPU is new but still want to run old applications
- ▶ **Very popular** – especially in gaming where old consoles emulated on new



NeoGEO



BUILDING BLOCK OF VM

- Most VMMs implement **virtual CPU (VCPU)** to represent **state** of CPU per guest
- Dual mode CPU means: **guest executes in user mode**
 - Kernel runs in kernel mode, but it's not safe to let guest kernel run in kernel mode too
 - So VM needs **two modes** – **virtual user mode** and **virtual kernel mode**
 - Actions in guest that usually cause switch to kernel mode must cause switch to virtual kernel mode

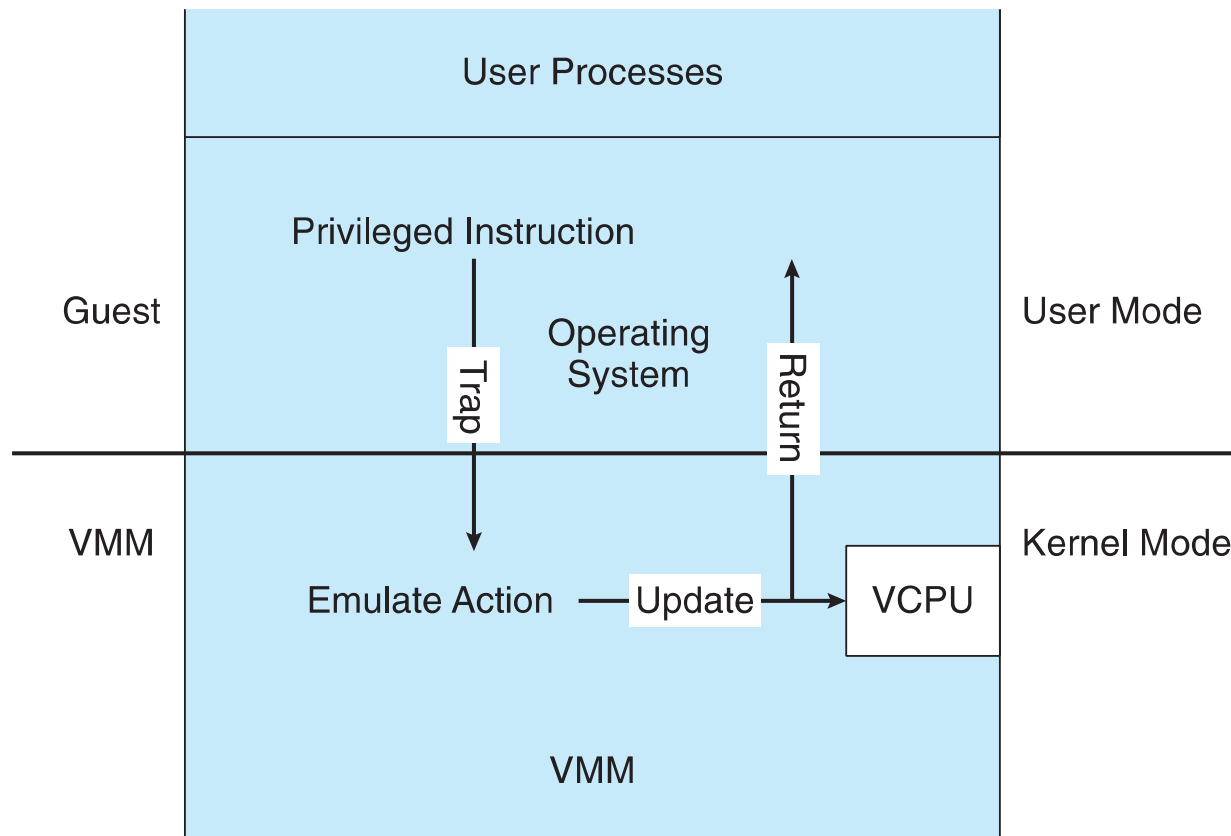


TEKNIK 1: TRAP-AND-EMULATE

- How does switch from virtual user mode to virtual kernel mode occur?
 - Attempting a privileged instruction in user mode causes an error -> trap
 - VMM gains control, analyzes error, executes operation as attempted by guest
 - Known as **trap-and-emulate**
- User mode code in guest runs at same speed
- But kernel mode privilege mode code runs **slower** due to trap-and-emulate



TRAP-AND-EMULATE VIRTUALIZATION IMPLEMENTATION

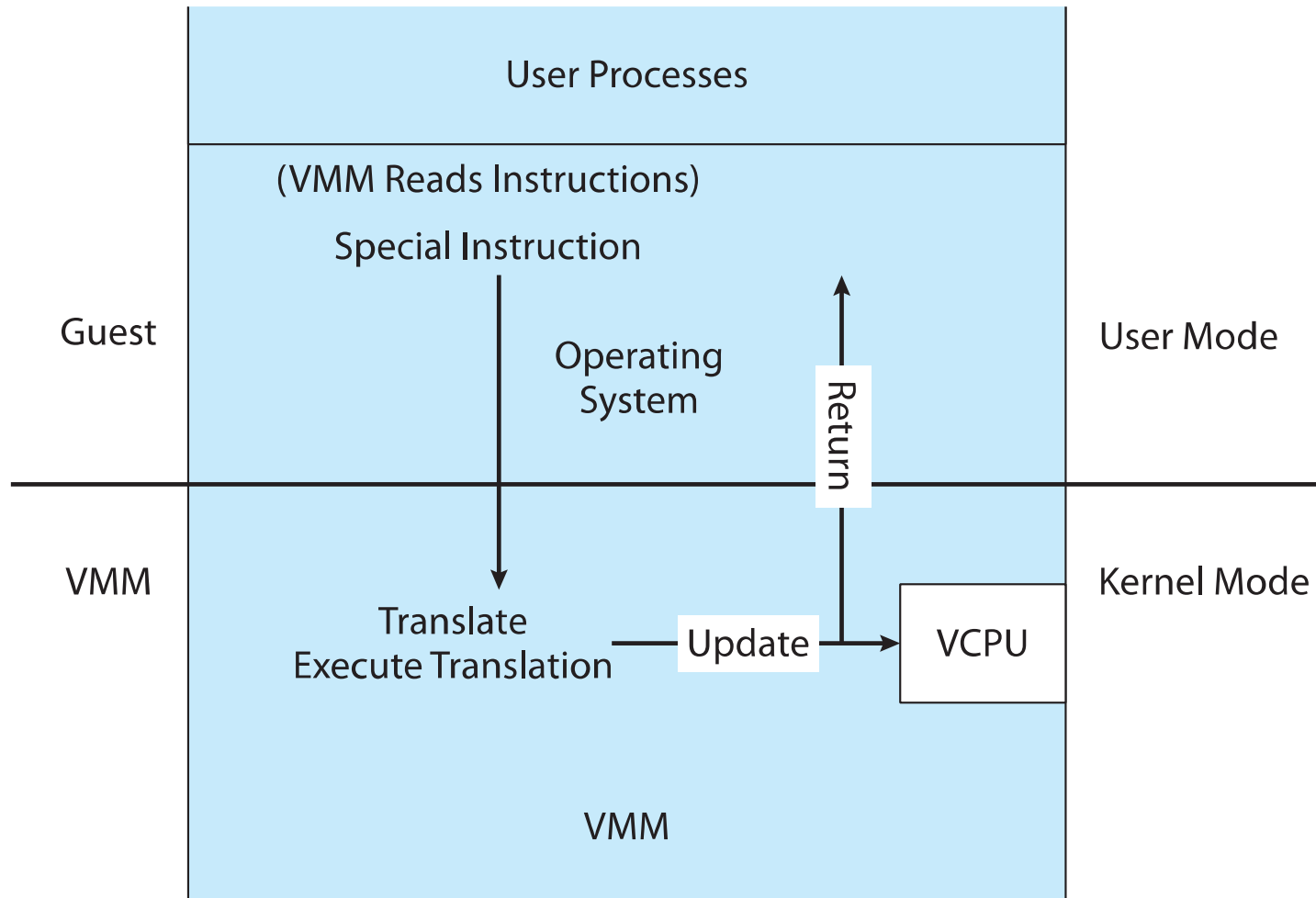


TEKNIK 2: BINARY TRANSLATION

- Implemented by **translation of code within VMM**
- Code reads native instructions dynamically from guest, generates **native binary code** that executes in place of original code
- Performance of this method would be **poor without optimizations**
 - Products like VMware use **caching**



BINARY TRANSLATION VIRTUALIZATION IMPLEMENTATION



BUILDING BLOCKS – HARDWARE ASSISTANCE

- ▶ All virtualization needs some **HW support**
- ▶ More support -> more feature rich, stable, better performance of guests
- ▶ Intel added new **VT-x** instructions in 2005 and AMD the **AMD-V** instructions in 2006
 - CPUs with these instructions **remove need for binary translation**
 - In guest mode, guest OS thinks it is running natively



OS COMPONENT – CPU SCHEDULING

- Single-CPU systems act like **multiprocessor** ones when virtualized
 - One or more virtual CPUs per guest
- Generally VMM has **one or more physical CPUs and number of threads to run on them**
 - Guests configured with certain number of VCPUs
 - Can be adjusted throughout life of VM
 - When enough CPUs for all guests -> VMM can allocate dedicated CPUs, each guest much like native operating system managing its CPUs
 - Usually not enough CPUs -> CPU **overcommitment**



OS COMPONENT – NETWORKING

- Networking also complex as VMM and guests all need network access
 - VMM can **bridge** guest to network (allowing direct access)
 - And / or provide **network address translation (NAT)**
 - NAT address local to machine on which guest is running, VMM provides address translation to guest to hide its address

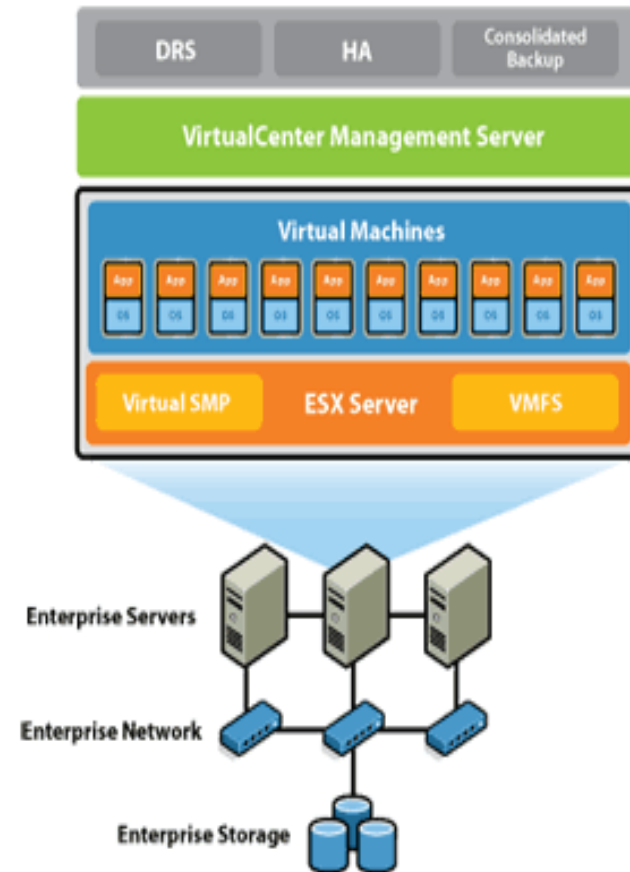
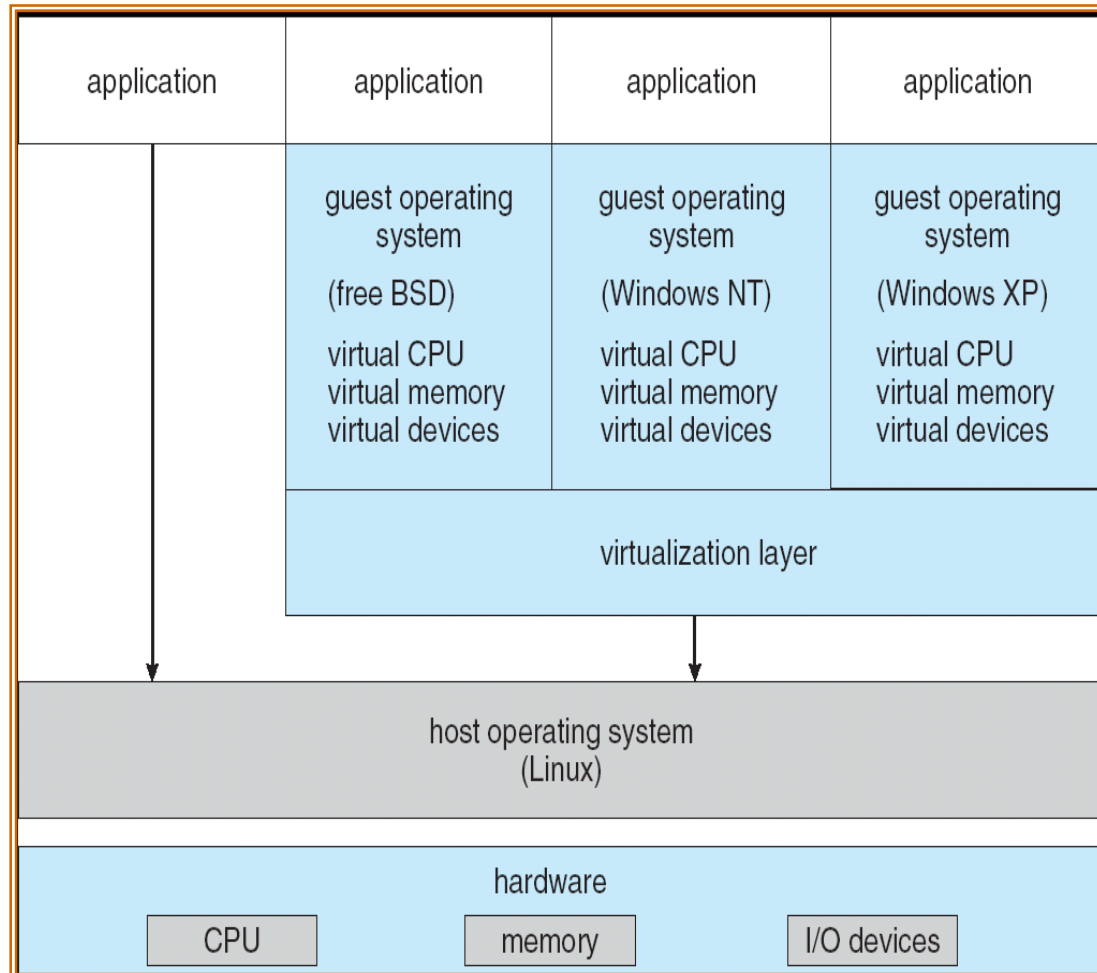


EXAMPLES - VMWARE

- VMware Workstation runs **on x86 or x64**, provides VMM for guests
- Runs as application on other native, installed host operating system -> **Type 2**
- Lots of guests possible, including Windows, Linux, etc all runnable concurrently (as resources allow)



VMWARE ARCHITECTURE



EXAMPLES – JAVA VIRTUAL MACHINE

- Example of **programming-environment virtualization**
- Very popular language / application environment invented by Sun Microsystems in 1995
- Write once, run anywhere (WORE)
- Includes language specification (Java), API library, Java virtual machine (JVM)
- Each Java object compiled into architecture-neutral **bytecode** output (`.class`) which JVM **class loader** loads
- JVM compiled per architecture, reads bytecode and executes
- Includes **garbage collection** to reclaim memory no longer in use

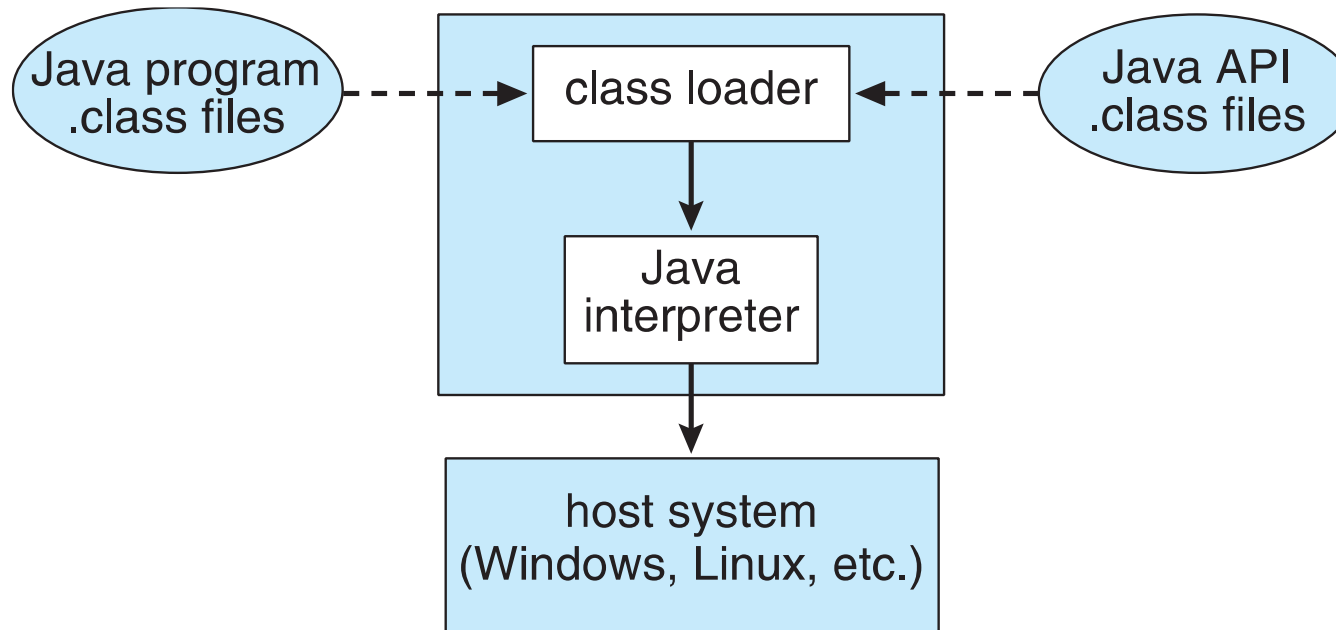


JAVA VIRTUAL MACHINE

- Program Java yang telah dicompile adalah **platform-neutral bytecodes** yang dieksekusi oleh Java Virtual Machine(JVM)
- JVM terdiri dari:
 - Class loader
 - Class verificatier
 - runtime interpreter
- Mendukung **Just In-Time (JIT)** compilers yang meningkatkan performance



THE JAVA VIRTUAL MACHINE



NEXT

- **Presentasi final**
- TAS

