CMPE 321 Summer 2019 Project1-Designing Storage Manager System

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1 Introduction

In this assignment I want to design a storage manager system which consists of a system catalogue and some data files. Data will be stored in records which will be stored in pages and pages will be stored in files. I will make some assumptions and determine some limitations about inputs. Then I will give pseudo codes for DDL and DML operations.

First of all I will explain the assumptions then give the detailed pseudo codes. Following are the relevant DDL and DML operations that are available to user.

DDL Operations

- 1) Create a type
- 2) Delete a type
- 3) List all types

DML Operations

Create a record
Delete a record
Search for a record (by primary key)
Update a record (by primary key)
List all records of a type

1.1 Assumptions

Page size is 1KB

File size is 6KB

In each page there is a page header which has an unique page id, information about the remaining size in the page and a pointer to the next page

In each record there is a record header which has also an unique record id(primary key), info about whether record empty or not is and a pointer to the next record

Max number of fields a type can have is 10

Max length of a type name is 20

Max length of a field name is 16

User always enters valid input.

All fields shall be integers. However, type and field names shall be alphanumeric.

A disk manager already exists that is able to fetch the necessary pages when addressed.

2 Structure Design

We begin by creating the system catalogue.

number of types					
first type	number of files	number of pages	number of records	first field	•••
second type	number of files	number of pages	number of records	first field	
••••	••••	••••			

Above table shows the design of system catalogue. It stores the metadata about the system

2.1 File

A file has multiple pages in it. When all the pages in a file are full we can create a new a file and a page of the desired type and add, update or delete records from that new file until its all pages are also full or deleted.

page1
page2
page3
••••
••••

Above table shows the structure of a file which can be created for any type

2.2 Page

A page consists of two different things. One of them is page header which includes an unique page id and a pointer to the next page and the other is several records which is another structure I will explain in the next section.

Page ID (4byte)
A pointer to the next page(8 byte)
Remaining size in the page(4byte)

Above table shows the structure of the page header

Page Header		
Record 1		
Record 2		
•••		
•••		

Above table shows the structure of a page which can be created for any type

2.3 Record

A record consists of two different things. First one is the record header which has the record id, info about whether a record is empty or not and also a pointer to the next record. Second one is that a record has at most 10 fields store the relevant data for its specific type

Record ID (4byte)
isEmpty (1byte)
A pointer to the next record (8byte)

Above table shows the structure of the record header

Record Header		
First Field		
Second Field		

Above table shows the structure of a record

3 Algorithms for Operations

```
First DDL operations:

Input: typeName,number of fields,names of fields systemCatalog.open(); construct dataType;

while i \leq numberofFields do

| dataType.fieldName[i]=namesofFields[i]

end

numberofTypes++
systemCatalogue.push(dataType);
systemCatalog.close();
file \leftarrow createFile(typeName)

Algorithm 1: Create a Type
```

```
Input: typeName
systemCatalog.open();
numberofTypes - -;
k \leftarrow number of Files Intype Name;
while i \leq k do
   delete file named typeName[i];
end
delete line associated with typeName in system catalog;
systemCatalog.close();
                               Algorithm 2: Delete a Type
systemCatalog.open();
k \leftarrow number of Types \ in \ system catalog \ ;
while i \leq k do
   j \leftarrow reads\ line[i]\ in\ system catalog;
   print(j.typeName);
end
systemCatalog.close();
                               Algorithm 3: List all Types
```

```
Input: typeName,recordFields
systemCatalog.open();
k \leftarrow typeName.number of Files in system catalog;
while i \leq k \operatorname{do}
   if typeName.file[i].remainingSize <math>\geq record.size then
       for all pages p in typeName.file[i] do
           p.remainingSize=p.header.readByte[12];
           if p.remainingSize \ge record.size then
              construct new record r with recordFields;
              p.push(r);
              typeName.numberofRecords ++;
              systemCatalog.close();
              return r
           else
             next page;
           end
        end
    else
     next file;
    end
   end
   construct a new file typeName[k+1];
   construct a new page p;
   construct a new record r with recordFields;
   p.push(r);
   typeName[k+1].push(p);
   typeName.numberOfFiles++;
   typeName.numberofRecords ++;
   systemCatalog.close();
   return r
```

Algorithm 4: Create a record

```
Input: typeName,recordID
systemCatalog.open();
k \leftarrow typeName.number of Files \ in \ system catalog \ ;
while i \leq k \operatorname{do}
   for all pages p in typeName.file[i] do
       for all records r in p do
           r.recordID=r.header.readByte[0](read first 4 bytes of header);
           if r.recordID==recordID then
               delete r;
               if typeName.file[i].isEmpty==TRUE then
                   delete typeName.file[i];
                   rename other files accordingly;
               else
               | print("File is not empty")
               end
               typeName.numberofRecords - -;
               print("Record deleted successfully");
               systemCatalog.close();
               return
           else
               next record;
           end
       end
   end
print("Record does not exist");
systemCatalog.close();
return
                              Algorithm 5: Delete a Record
```

```
Input: typeName,recordID
systemCatalog.open();
k \leftarrow typeName.number of Files in system catalog;
while i \leq k \operatorname{do}
   for all pages p in typeName.file[i] do
       for all records r in p do
           r.recordID=r.header.readByte[0](read first 4 bytes of header);
           if r.recordID==recordID then
               print("Record is found");
               return r;
           else
              next record;
           end
       end
   end
end
print("Record does not exist");
systemCatalog.close();
return
                    Algorithm 6: Search a Record (by primary key)
Input: typeName,recordID,recordFields
systemCatalog.open();
k \leftarrow number of Files Intype Name;
while i \leq k \text{ do}
   for all pages p in typeName.file[i] do
       for all records r in p do
           r.recordID=r.header.readByte[0](read first 4 bytes of header);
           if r.recordID==recordID then
               r.updateFields(recordFields) updating the given fields;
               print("Fields are updated");
               systemCatalog.close();
               return;
           else
               next record;
           end
       end
   end
end
print("Record does not exist");
systemCatalog.close();
return:
```

Algorithm 7: Update a Record(by primary key)

```
Input: typeName
systemCatalog.open();
k \leftarrow number of Files Intype Name;
while i \leq k \text{ do}
   for all pages p in typeName.file[i] do
       for all records r in p do
           r.isEmpty=r.header.readByte[4](read byte 4 to 5 in header);
           if r.isEmpty = = false then
               print(r);
           else
               next record;
           end
       end
   end
end
systemCatalog.close();
```

Algorithm 8: List all Records

4 Conclusions and Assessment

In this project I had an intuition about how to design a database management storage. My designs advantages are since I insert records linearly without any specific order insertion is fast however, as a disadvantage searching is costly since it looks page by page and record by record.

Another advantage is that I use page ids and record ids they make easier to find a page in a file or find a record in a page but as a disadvantage it allocates more memory. This ids make easier to implement my design. Hopefully in the next project I will improve my design.