Simulating a dark matter halo using a parallel grid method

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1 INTRODUCTION

Dark matter appears to be the dominant form of matter in the universe. While no direct observations of dark matter exist, it has been observed indirectly through its gravitational interaction with visible matter and radiation (?). The effect of dark matter on structure formation in the universe is the subject of an increasing scientific interest and a useful tool in the search for possible candidates of dark matter. Current progress in the investigation of structure formation is mainly driven by advances in computational methods and capabilities (?).

In this work we simulate a dark matter halo, neglecting baryonic matter, by using a particle based method and solving the Poisson's equation of gravity. The software used is *Uintah* available from http://uintah.utah.edu/, a parallel grid framework with a support for particle interactions and adaptive mesh refinement and which is aimed at solving partial differential equations. A major part of the project was to learn how to use the tools provided by Uintah.

2 METHODS

2.1 Theory

A distribution of mass density ρ gives rise to a gravitational potential ϕ according to the Poisson's equation of gravity

$$\nabla^2 \phi = 4\pi G \rho, \tag{1}$$

where $G\approx 6.674\times 10^{-11}~{\rm Nm^2/kg^2}$ is the gravitational constant. In our simulation we normalize G=1. The corresponding force field can be solved from the gradient

$$\mathbf{F} = -\nabla \phi. \tag{2}$$

In our simulation a distribution of massive particles create the mass density. A new position and velocity for the particles after a time step dt can be solved from the Newton's equation of motion.

$$\mathbf{v}(t+dt) = \int_{t}^{t+dt} \frac{\mathbf{F}(t')}{m} dt' + \mathbf{v}(t)$$
(3)

$$\mathbf{x}(t+dt) = \int_{t}^{t+dt} \mathbf{v}(t')dt' + \mathbf{x}(t)$$
(4)

The use of this theory assumes that we do not have relativistic speeds or masses and that the maximum grid size is small enough that expansion of the universe can be neglected.

One approach would be to try to calculate the particle-particle interactions but this is not computationally feasible since we would have 2^N interactions where N is the particle number and we are using thousands of particles.

Numerical methods

The simulation is set up with respect to a three dimensional grid that supports particles (Fig. 1).

The overall algorithm is expressed in Algorithm 1.

Algortihm 1. Main program.

set Dirichlet or periodic boundary conditions set initial \mathbf{x}_p , \mathbf{v}_p and m_p for all particles pinterpolate initial ρ from particle mass m_p set initial guess for ϕ **loop** solve ϕ at the nodes using the SOR algorithm calculate $\mathbf{F} = -\nabla \phi$ with respect to each particle p **for** every particle p **do** $\mathbf{v}_p(t + \Delta t) = (\mathbf{F}_p/m)\Delta t + \mathbf{v}_p(t)$ $\mathbf{x}_p(t + \Delta t) = \mathbf{v}_p(t)\Delta t + \mathbf{x}_p(t)$ **end for** interpolate new ρ from particle mass m_p $t \leftarrow t + \Delta t$

The Poisson's equation of gravity is discretized to

end loop

$$4\pi G \rho_{i,j,k} = \frac{\phi_{i+1,j,k} - 2\phi_{i,j,k} + \phi_{i-1,j,k}}{(\Delta x)^2} + \frac{\phi_{i,j+1,k} - 2\phi_{i,j,k} + \phi_{i,j-1,k}}{(\Delta y)^2} + \frac{\phi_{i,j,k+1} - 2\phi_{i,j,k} + \phi_{i,j,k-1}}{(\Delta z)^2}.$$
(5)

The gravitational potential ϕ is then solved at each node using using the successive over-relaxation (SOR) algorithm (Algorithm 2). In the calculation of the potential gradient we obtain the potential values near the particles by interpolation and then calculate the gradient by numerical differentiation. For exmaple at the pth particle the gradient is calculated in the x direction as

$$(\nabla \phi)(x_p) = \frac{\phi(x_p + dx) - \phi(x_p - dx)}{2dx},\tag{6}$$

where x_p is the particle's x coordinate and dx is a small distance.

The particle velocity and position are evolved over a small time step dt, assuming that the force remains constant. Essentially this means calculating

for each particle

$$= \mathbf{v}_p(t + \Delta t) = (\mathbf{F}_p/m)\Delta t + \mathbf{v}_p(t) \tag{7}$$

$$= \mathbf{x}_p(t + \Delta t) = \mathbf{v}_p(t)\Delta t + \mathbf{x}_p(t). \tag{8}$$

After this using the new positions the particle masses are interpolated back to the nodes to obtain a new mass density ρ .

Algorithm 2. Calculating potential ϕ using the SOR algorithm.

```
function SOR(\phi, tolerance, max_iterations)

for n=0,1,\ldots max_iterations do

error \sigma\leftarrow 0

for every node \phi_{i,j,k}

\phi_{i,j,k}^{(n+1)}\leftarrow (1-\omega)\phi_{i,j,k}^{(n)}+\frac{\omega}{6}(\phi_{i+1,j,k}^{(n)}+\phi_{i-1,j,k}^{(n)}+\phi_{i,j+1,k}^{(n)}+\phi_{i,j-1,k}^{(n)}+\phi_{i,j,k+1}^{(n)}+\phi_{i,j,k-1}^{(n)}+h^3\rho_{i,j,k})

update \sigma

end for

if \sigma\leq tolerance, break

end for

return \phi

end function
```

2.2 Initialization of the simultaion

The specific case we want to study with our simulation is a dark matter halo. Dark matter halos are structures composed of dark matter, believed to envelope galaxy disks. We initialize the halo by letting a sphere of randomly distributed dark matter particles to collapse by the effect of gravity.

2.3 Verification of results

To verify the simulation results we test for energy conservation, compliance with the virial theorem

$$\langle T \rangle_t = -\frac{1}{2} \langle V \rangle_t. \tag{9}$$

Previous N-body simulations suggest that dark matter halos follow the

Navarro-Frenk-White radial mass distribution

$$\rho(r) = \frac{\rho_0}{\frac{r}{R_s} \left(1 + \frac{r}{R_s}\right)^2},\tag{10}$$

where r is distance from the center of the halo and (ρ_0, R) are parameters. As part of our verification procedure we fit this distribution to the simulated dark matter halo.

IMPLEMENTATION

In this section, we will often refer to Section 2.1 and describe each numerical method's implementation in detail. During the whole project, we made a lot of effort to make sure the algorithms fit the Uintah framework. Because of this, we will have extensive description of the Uintah framework in the Section 2.3, describing how the Uintah patches which are used in the Poisson solver work, how particle masses are interpolated to the grid and how the gradient of potential is calculated.

In addition to describing our implementation and the Uintah framework, we will discuss some of the limitations and bugs of Uintah software which drained a painstaking amount of time to solve. For more information on Uintah framework, please refer to the *uintah_presentation.pdf* we made on the software.

Uintah

As mentioned in the title, we made a *fully parallel* poisson solver utilizing particle-in-cell approach.

Before we start with the actual algorithms, let's go through the basics of Uintah. There are a few concepts to understand; cells, patches, ghost cells, cell-centered variables, node-centered variables, face-centered variables and particles.

We present an explanation for each of these accompanied by a Figure (Figure 1) demonstrating a visual description.

1. Cells are finite boxes in the grid. These are the basic elements of our simulation. In the case of our poisson solver, we discretize our poisson equation based on cells in such a way that our discretization length h in Equation 5 is the length between cells.

- 2. Patches are collections of cells in our grid. The reason to have a concept of patches is that Uintah uses patches to distribute workloads for different processes. In Uintah, whenever we iterate through cells, we iterate through them patch-by-patch. This introduces some complications, such as how to handle the boundary between patches. The boundary issue will be further elaborated on.
- 3. Ghost cells are cells which are one or more step beyond the border of a patch. The whole concept of ghost cells is meaningless if we do not work in parallel. In a parallel simulation (employing more than one process), having ghost cells is of utmost importance because it is used to handle MPI communication between different processes' patches.
- 4. Cell-centered variables are variables which are defined at the center of each cell.
- 5. Node-centered variables are variables which are defined at the corners of each cell.
- 6. Face-centered variables are variables which are defined at the faces of each cell.
- 7. Particles are point variables defined at any arbitrary point in the simulation grid. In our simulation, cell/node/face-centered variables interact with each other by interpolating particle variables to cell/node/face-centered positions and by interpolating cell/node/face-centered variables to particle positions.

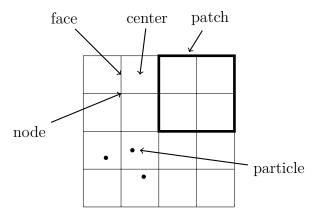


Figure 1. A two dimensional view of the grid with its various components and properties shown.

One of the biggest reasons we chose to use the Uintah framework was that it has a very advanced $parallel\ task\ system$. As already mentioned previously, in Uintah parallelization goes over patches. Every process has N patches which they are in charge of, and those patches are distributed according to load balancing settings. In our case, we chose to use a load balancing option that takes the number of particles into account.

The parallel task system takes care of most data communication and working with data dependencies. For each *function*, we must define a task which takes in the function, the function's *dependencies* and the function's *output*.

For instance, as we can see in the Equation 5, for one poisson iteration, we need the *last iteration's* $\phi = \phi_{previous}$ value as well as the distribution of particle density ρ , and one poisson iteration calculates the next $\phi = \phi_{next}$ value. Additionally, because according to Equation 5 we need neighbor values of every ϕ in our iterations, we need to tell the task that it needs to fetch 1 layer of *ghost cells*. The concept of *ghost cells* was demonstrated in Figure 1.

In pseudocode, this would be:

Algortihm 3. Demonstrating Uintah taskgraph system

 $Task \leftarrow new task(poissonIterationStep)$ $Task \leftarrow \phi_{previous}$ with 1 layer of ghost cells $Task \leftarrow \rho_{current}$ $Task \ computes \ \phi_{next}$ $Scheduler \leftarrow Task$

In addition to tasks and different patch hierarchies, we may also explain the way Uintah parallelizes things, and why it is currently regarded as perhaps the most scalable grid library. Normally in most grid libraries, things advance in a *semi-serial* way due to normal MPI programming. Whenever a function is done calculating, there is some communication between MPI processes (e.g. process #1 needs to tell #2 some result of a calculation). Until the communication is done, the process is often left hanging. This phenomena is often referred to as *load imbalance*.

In the case of Uintah, the problem of load imbalance is overcome by introducing task dependencies, which we already went through shortly in this section. Task dependencies refer to tasks having certain dependencies before they can be executed. For instance, say process #1 can move to interpolating particle values to a grid or to calculating poisson iterations and poisson iteration still requires a variable which has not been calculated. In this case, the process #1 would simply move on to interpolating particle values, instead of hanging around doing nothing, as would be the case in most normal parallel simulations.

For demonstration, we present a task graph taken from Uintah website in Figure 2.

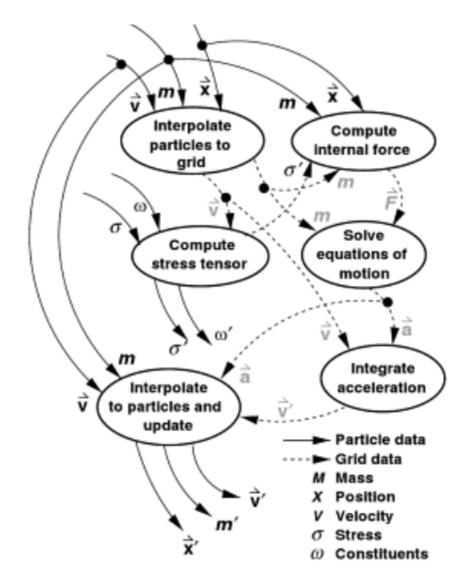


Figure 2. Task graph example from Uintah's website

In Uintah, there are essentially 2 different ways to parallelize the simulation. Either one distributes the workload among processes, or one distributes the workload among both processes and threads. We did not use the thread-based parallelization because unfortunately neither of our MPICH versions had support for the required level 3 parallelization, and tackling that problem simply requires too much of our valuable time. In fact, even FMI's supercomputer Voima did not have support for level 3 parallelization. We were

still intrigued by the concept of making a *hybrid-parallelization*, especially since our simulation should have it working out of the box, so we read more about it.

The way hybrid-parallelization in Uintah works is that every process is assigned N threads, from which 1 is used for MPI communication, and N-1 are used for execution of tasks. This should guarantee fast execution on supercomputers, as long as $N\approx 15-30$. Using hybrid-parallelization would reduce MPI communication, as threads within a process would have shared memory, and would not require for instance communication of ghost cells through MPI.

In addition to having a nice task-graph system, Uintah has put some effort into having general-use interfaces for implementing things. Patches can be sorted with *Hilbert spatial filling curves*¹, Uintah has tools for interpolating particles to nodes and nodes to particles and it has tools for calculating interpolated derivatives. Additionally, Uintah has support for implementing adaptive mesh refinement². Despite the fact that these tools have a steep learning curve, they will become invaluable when implementing parallel and optimized algorithms for mesh-based methods.

Finally, we will discuss some of the bugs and limitations in Uintah. For one, interpolating particle and cell variables back and forth are viable only for node-based variables out-of-the-box. This is not unusual, as the visualization software VisIt and most other mesh visualization softwares have a similar take on this. Making sparse data is also not possible in Uintah, meaning that solving e.g. Vlasov equation on a sparse grid is not possible, and neither is solving many-dimensional partial differential equations with methods such as adaptive $radial\ basis\ method^3$. Additionally, there was one great limitation we encountered; the particle interpolation does not work when the particle variable has values around 1.0e40. It gives spiked results for the interpolated variable, and it took us a day to find out this was why our simulation was producing incorrect results.

¹Hilbert spatial filling curves are used to optimize memory cache. In the case of our simulation, cells are sorted in the way that neighbors of each cell are closer in memory, reducing cache misses in fetching neighbor cell data. For more information on space filling curves, see http://www.diss.fu-berlin.de/diss/servlets/MCRFileNodeServlet/FUDISS_derivate_000000003494/2_kap2.pdf?hosts=

²See Section 3

 $^{^3}$ Solving partial differential equations by assuming a basis of radial functions and often neglecting the functions which contribute less

For some additional information, please refer to the Uintah presentation we made which can be found in:

./doc/Uintah/uintah_presentation.pdf

The program

Now that we know the basics on Uintah, we can move on to looking at our program. Our main algorithm is described in the Algorithm 1, and here we attempt to break this down into smaller details. All of this can be found also in our project files ParticleTest1.cc and ParticleTest1.h.

set Dirichlet or periodic boundary conditions. Setting either dirichlet or periodic boundary conditions is fairly straightforward. In fact, this is something that we found Uintah can do for us. We need to simply input our boundary conditions in a run file and take it into account in our poisson solver. In our poisson solver, we simply skip iterating the cells that are on our grid border.

set initial \mathbf{x}_p , \mathbf{v}_p and m_p for all particles p. For this purpose, Uintah specifies a function for initializing our data. We set particle position $\mathbf{x}_p = rand(grid)$, velocity $\mathbf{v}_p = 0$ and mass $m_p = 1.0e4$.

interpolate initial ρ from particle mass m_p . For this, we use linear interpolation method. Our ρ is specified at each cell corner $(node^4)$ and we interpolate the particle mass variable to nearest 8 nodes. Since ρ is density, we also divide the interpolated mass variable by one cell volume $\Delta x \Delta y \Delta z$. The reason we chose to interpolate only to 8 nearest nodes (e.g. use linear interpolation) is because it has been shown that this is produces a sufficient accuracy for the Poisson equation(?). By sufficient, we of course mean that it does not introduce a global error of higher order than what is already in our simulation due to discretizing poisson equation. For this purpose, we found that Uintah has a class called ParticleInterpolator that produces interpolation. This required us to change our ρ variable to node-centered form from cell-centered form, however. This method is widely used for communication between grids and particles, and commonly methods using this are referred to as a cloud-in-cell or particle-in-cell methods.

set initial guess for ϕ . We set this to zero everywhere.

solve ϕ at the nodes. For this, we created an iterative solver that goes through each patch⁵ and iterates through each node in that patch. Basically,

 $^{^4\}mathrm{See}$ section 2.3

⁵See Section 2.3

we solve the Equation 5. We require that each patch fetches 1 layer of ghost nodes⁶ to deal with patch boundaries, as our algorithm needed the neighboring nodes⁷ for each node. Additionally, we skip the cells at the patch boundaries due to boundary conditions. Our algorithm for solving the poisson equation is demonstrated in Algorithm 2. For every iteration, we also created a task to utilize the full functionality of *Uintah*. For task dependencies in the Poisson solver, please refer to Algorithm 3.

calculate $\mathbf{F} = -\nabla \phi$ for each particle p. For this, we use linear interpolation. For every particle p, we pick the 8 nearest nodes and interpolate to values $x_i + dx_i$ and $x_i - dx_i$ for every i. This means that our gradient calculation is also based on linear interpolation. This introduces some error, but again it has been shown to be sufficient for solving poisson equation(?).

propagating every particle p. This is already documented in Algorithm 1, so we do not elaborate it further here.

interpolate new ρ from particle mass m_p . See the step interpolate initial ρ from particle mass m_p .

 $t \leftarrow t + \Delta t$. This is done in the Uintah framework and is determined by the variable *Delt* in the runtime *XML* file. We use constant timestepping, so this step needs nothing complicated.

RESULTS

3 CONCLUSIONS

⁶See Section 2.3

⁷neighboring node = node which is one node to the left,right,down,up,forward or backward