

Required Problems

1. Using truth tables, prove both of DeMorgan's Laws for logical connectives.

(a) $\neg(P \wedge Q)$ is logically equivalent to $\neg P \vee \neg Q$

P	Q	$\neg P$	$\neg Q$	$(P \wedge Q)$	$\neg(P \wedge Q)$	$\neg P \vee \neg Q$
T	T	F	F	T	F	F
T	F	F	T	F	T	T
F	T	T	F	F	T	T
F	F	T	T	F	T	T

The last two columns, for $\neg(P \wedge Q)$ and $\neg P \vee \neg Q$ respectively, have the same truth values for all truth assignments of P and Q ; thus, they are logically equivalent.

(b) $\neg(P \vee Q)$ is logically equivalent to $\neg P \wedge \neg Q$

P	Q	$\neg P$	$\neg Q$	$(P \vee Q)$	$\neg(P \vee Q)$	$\neg P \wedge \neg Q$
T	T	F	F	T	F	F
T	F	F	T	T	F	F
F	T	T	F	T	F	F
F	F	T	T	F	T	T

Once again, the last two columns, for $\neg(P \vee Q)$ and $\neg P \wedge \neg Q$ respectively, have the same truth values for all truth assignments of P and Q ; thus, they are logically equivalent.

2. Find the contrapositive and converse of each of the following statements:

(a) "If squares have four sides, then triangles have four sides."

This has the form: (squares have four sides) \implies (triangles have four sides).

- Contrapositive: "If triangles do not have four sides, then squares do not have four sides."
- Converse: "If triangles have four sides, then squares have four sides."

(b) "A sequence a is bounded whenever a is convergent."

This has the form: (a is convergent) \implies (a is bounded).

- Contrapositive: "If a is not bounded, then a is not convergent."
- Converse: "If a is bounded, then a is convergent."

(c) "The differentiability of a function f is sufficient for f to be continuous."

This has the form: (f is differentiable) \implies (f is continuous)

- Contrapositive: "If f is not continuous, then f is not differentiable."
- Converse: "If f is continuous, then f is differentiable."

3. Let x and y be integers. Prove that if x and y are even, then $x + y$ is even.

- \mathbb{Z} is the set of integers
- Def. of even: $x \in \mathbb{Z}$ is even iff $\exists k \in \mathbb{Z} \ni x = 2k$

To show: There exists an integer h such that $x + y = 2h$

Proof:

Let $x, y \in \mathbb{Z}$ be even (by hypothesis)
 $\implies (\exists k \in \mathbb{Z} \ni x = 2k) \wedge (\exists j \in \mathbb{Z} \ni y = 2j)$ (by def. of even)
 $\implies x + y = 2k + 2j$ (summing)
 $\implies x + y = 2(k + j)$ (by distributivity)
 $k + j \in \mathbb{Z}$ (by closure)
Let $h = k + j$ (defining an integer h)
 $\implies \exists h \in \mathbb{Z} \ni x + y = 2h$ (substituting for $j + k$)
 $\implies x + y$ is even (by def. of even) ■

4. Let A and B be sets. Prove that $A \subset B$ if and only if $A - B = \emptyset$.

- Statement using biconditional: $(A \subset B) \iff (A - B = \emptyset)$
- Def. of subset: $A \subset B \iff (x \in A \implies x \in B)$
- Def. of set difference: $A - B = \{x | x \in A \wedge x \notin B\}$
- Theorem (T1): $(P \implies Q)$ is logically equivalent to $(\neg P) \vee Q$
- Theorem (T2): For all sets A , $\emptyset \subset A$
- Contrapositive of (\implies) : $A - B \neq \emptyset \implies A \not\subset B$

Proof by contraposition, to show: $A \not\subset B$.

Proof (\implies) :

Let $A - B \neq \emptyset$ (by hypothesis)
 $\implies \exists x \in A - B$ (by def. of non-empty)
 $\implies x \in A \wedge x \notin B$ (by def. of set difference)
 $\implies \neg(x \notin A \vee x \in B)$ (by negation)
 $\implies \neg(x \in A \implies x \in B)$ (by T1)
 $\implies \neg(A \subset B)$ (by def. of subset)
 $\implies A \not\subset B$ (by negation)

To show: $A \subset B$.

Proof (\Leftarrow) :

Let $A - B = \emptyset$ (by hypothesis)
Case 1: $A = \emptyset$ (by hypothesis)
 $\implies A \subset B$ (by T2)
Case 2: $A \neq \emptyset$ (by hypothesis)
 $\implies \exists x \in A$ (by def. of non-empty)
 $\implies x \notin A - B$ (by def. of empty)
 $\implies \neg(x \in A \wedge x \notin B)$ (by def. of set difference)
 $\implies x \notin A \vee x \in B$ (by negation)
 $\implies (x \in A \implies x \in B)$ (by T1)
 $\implies A \subset B$ (by def. of subset) ■

Practice Problems

5. If P , Q , and R are true while S and T are false, which of the following are true?

(a) $Q \wedge (R \wedge S)$

Q	R	S	$(R \wedge S)$	$Q \wedge (R \wedge S)$
T	T	F	F	F

Thus, $Q \wedge (R \wedge S)$ is false.

(b) $(P \vee Q) \wedge (R \vee S)$

P	Q	R	S	$(P \vee Q)$	$(R \vee S)$	$(P \vee Q) \wedge (R \vee S)$
T	T	T	F	T	T	T

Thus, $(P \vee Q) \wedge (R \vee S)$ is true.

(c) $(P \vee S) \wedge (P \vee T)$

P	S	T	$(P \vee S)$	$(P \vee T)$	$(P \vee S) \wedge (P \vee T)$
T	F	F	T	T	T

Thus, $(P \vee S) \wedge (P \vee T)$ is true.

6. Make truth tables for these propositional forms:

(a) $P \implies (Q \wedge P)$

P	Q	$Q \wedge P$	$P \implies (Q \wedge P)$
T	T	T	T
T	F	F	F
F	T	F	T
F	F	F	T

(b) $(\neg P \implies Q) \vee (Q \implies P)$

P	Q	$\neg P$	$\neg P \implies Q$	$Q \implies P$	$(\neg P \implies Q) \vee (Q \implies P)$
T	T	F	T	T	T
T	F	F	T	T	T
F	T	T	T	F	T
F	F	T	F	T	T

(c) $\neg Q \implies (Q \iff P)$

P	Q	$\neg Q$	$Q \iff P$	$\neg Q \implies (Q \iff P)$
T	T	F	T	T
T	F	T	F	F
F	T	F	F	T
F	F	T	T	T

7. Rewrite each of the following sentences to be symbolic sentences using logical connectives.

(a) If $x = 1$ or $x = -1$, then $|x| = 1$.

$$(x = 1 \vee x = -1) \implies |x| = 1$$

- (b) B is invertible is a necessary and sufficient condition for $|B| \neq 0$.

$$\exists B^{-1} \iff |B| \neq 0$$

- (c) $6 \geq n - 3$ only if $n > 4$ or $n > 10$.

$$\neg(n > 4 \vee n > 10) \implies \neg(6 \geq n - 3)$$

- (d) S is compact iff S is closed and bounded.

$$(S \text{ is compact}) \iff (S \text{ is closed and bounded})$$

8. Rewrite each of the following sentences to be symbolic sentences using quantifiers. The universe of discourse for each is given in parentheses.

- (a) Every nonzero real number is positive or negative. (Real Numbers)

$$\forall x \in \mathbb{R}, (x \geq 0 \wedge x \leq 0)$$

- (b) Every integer is greater than some integer. (Integers)

$$\forall x \in \mathbb{Z} \exists y \in \mathbb{Z} \ni x > y$$

- (c) There is a smallest positive real number. (Real Numbers)

$$\exists x \in \mathbb{R} \ni x \geq 0 \wedge \forall y \in \mathbb{R} \ni y \geq 0, x \leq y$$

9. The qualifier $\exists!$ is defined as follows:

$$\exists! x \ni A(x) \iff (\exists x \ni A(x)) \wedge (\forall y \wedge \forall z, A(y) \wedge A(z) \implies y = z)$$

Describe in plain english the qualifer $\exists!$.

The qualifier $\exists!$ states that there exists an element making $A(x)$ true and if $A(x)$ is true for two values, those values must be the same. in other words, the qualifier $\exists!$ reads “there exists a unique,” so it is the unique existential qualifier.

10. Let x and y be integers. Prove the following propositions:

- (a) If x and y are even, then xy is even.

To show: there exists an integer h such that $xy = 2h$.

Proof:

$$\begin{array}{ll} \text{Let } x \text{ and } y \text{ be even integers} & \text{(by hypothesis)} \\ \implies (\exists k \in \mathbb{Z} \ni x = 2k) \wedge (\exists j \in \mathbb{Z} \ni y = 2j) & \text{(by def. of even)} \\ \implies xy = (2j)(2k) & \text{(multiplying)} \\ \implies xy = 2(2jk) & \text{(by associativity)} \\ 2jk \in \mathbb{Z} & \text{(by closure)} \\ \text{Let } h = 2jk & \text{(defining an integer } h) \\ \implies xy = 2h & \text{(substituting for } 2jk) \\ \implies xy \text{ is even} & \text{(by def. of even)} \end{array}$$

■

(b) **If x and y are odd, then $x + y$ is even.**

To show: there exists an integer h such that $x + y = 2h$

Proof:

Let x and y be odd integers (by hypothesis)
 $\implies (\exists k \in \mathbb{Z} \ni x = 2k + 1) \wedge (\exists j \in \mathbb{Z} \ni y = 2j + 1)$ (by def. of odd)
 $\implies x + y = (2k + 1) + (2j + 1)$ (summing)
 $\implies x + y = 2k + 2j + 2$ (by associativity/commutativity)
 $\implies x + y = 2(k + j + 1)$ (by distributivity)
 $k + j + 1 \in \mathbb{Z}$ (by closure)
Let $h = k + j + 1$ (defining an integer h)
 $\implies x + y = 2h$ (substituting for $k + j + 1$)
 $\implies x + y$ is even (by def. of even) ■

(c) **If x is even and y is odd, then $x + y$ is odd.**

To show: there exists an integer h such that $x + y = 2h + 1$

Proof:

Let x be an even integer and let y be an odd integer (by hypothesis)
 $\implies (\exists k \in \mathbb{Z} \ni x = 2k) \wedge (\exists j \in \mathbb{Z} \ni y = 2j + 1)$ (by def. of even and odd)
 $\implies x + y = 2k + 2j + 1$ (summing)
 $\implies x + y = 2(k + j) + 1$ (by distributivity)
 $k + j \in \mathbb{Z}$ (by closure)
Let $h = k + j$ (defining an integer h)
 $\implies x + y = 2h + 1$ (substituting for $k + j$)
 $\implies x + y$ is odd (by def. of odd) ■

11. **Let a and b be real numbers. Prove that $|a + b| \leq |a| + |b|$.**

- Lemma (L1): if a is a real number, then $-|a| \leq a \leq |a|$
- Lemma (L2): $|b| \leq c \iff -c \leq b \leq c$

To show: $|a + b| \leq |a| + |b|$

Proof:

Let a and b be real numbers (by hypothesis)
 $\implies (-|a| \leq a \leq |a|) \wedge (-|b| \leq b \leq |b|)$ (by L1)
 $\implies -|a| - |b| \leq a + b \leq |a| + |b|$ (summing)
 $\implies |a + b| \leq |a| + |b|$ (by L2) ■

12. **Let x be an integer. Write a proof by contraposition to show that if x is even, then $x + 1$ is odd.**

- Contrapositive: if $x + 1$ is not odd, then x is not even

- Theorem (T1): x is not even if and only if x is odd.

Proof by contraposition to show: x is not even

Proof:

Let x be an integer such that $x + 1$ is not odd (by hypothesis)
 $\implies x + 1$ is even (by T1)
 $\implies \exists k \in \mathbb{Z} \ni x + 1 = 2k$ (by def. of even)
 $\implies x = 2k - 1$ (subtracting 1 from both sides)
 $\implies x = 2k - 2 + 1$ (rearranging the r.h.s)
 $\implies x = 2(k - 1) + 1$ (by distributivity)
 $k - 1$ is an integer (by closure)
let $j = k - 1$ (defining an integer j)
 $\implies x = 2j + 1$ (substituting for $k - 1$)
 $\implies x$ is odd (by def. of odd)
 $\implies x$ is not even (by T1) ■

13. Suppose a and b are positive integers. Write a proof by contradiction to show that if ab is odd, then both a and b are odd.

- Theorem (T1): x is not even if and only if x is odd.

Proof by contradiction to show: if a or b is not odd, then ab is odd and not odd.

Proof:

Let $ab \in \mathbb{Z}_+ \ni ab$ is odd and a or b is not odd (towards a contradiction)
 $\implies a$ is even or b is even (by T1)
Case 1: a is even
 $\implies \exists k \in \mathbb{Z} \ni a = 2k$ (by def. of even)
 $\implies ab = (2k)b$ (substituting for a)
 $\implies ab = 2(kb)$ (by associativity)
 kb is an integer (by closure)
 $\implies ab$ is even (by def. of even)
 $\implies ab$ is not odd
Thus, a contradiction
Case 2: b is even
(similar to Case 1)
Case 3: both a and b are even
 $\implies ab$ is even (by proof in problem 10(a))
 $\implies ab$ is not odd (by T1)
Thus, a contradiction

Because every case produces the contradiction ab is odd and not odd, it must be the case that a and b are both odd. ■

14. **Prove that if $x \notin B$ and $A \subset B$, then $x \notin A$.**

To show: if $x \notin B$, then $x \notin A$:

Proof:

$$\begin{aligned}
 &\text{Let } A \subset B && \text{(by hypothesis)} \\
 &\implies (x \in A \implies x \in B) && \text{(by def. of subset)} \\
 &\implies [(\neg x \in B) \implies (\neg x \in A)] && \text{(the contrapositive)} \\
 &\implies [x \notin B \implies x \notin A] && \text{(by negation)}
 \end{aligned}$$

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15. **Let $A = \{1, 3, 5, 7, 9\}$, $B = \{0, 2, 4, 6, 8\}$, and $C = \{1, 2, 4, 5, 7, 8\}$ and $D = \{1, 2, 3, 5, 6, 7, 8, 9, 10\}$. Find the following:**

(a) $A \cup B$

$$A \cup B = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$$

(b) $A - B$

$$A - B = \{1, 3, 5, 7, 9\}$$

(c) $(A \cap C) \cap D$

$$(A \cap C) \cap D = \{1, 5, 7\}$$

(d) $A \cup (C \cap D)$

$$A \cup (C \cap D) = \{1, 2, 3, 5, 7, 8, 9\}$$

16. **Let A , B , C , and D be sets prove that if $C \subset A$ and $D \subset B$ and A and B are disjoint, then C and D are disjoint.**

- Def. of disjoint: A and B are disjoint $\iff A \cap B = \emptyset$
- Contrapositive: C and D are not disjoint implies A and B are not disjoint

Proof by contraposition to show: A and B are not disjoint

Proof:

$$\begin{aligned}
 &\text{Let } C \subset A \text{ and } D \subset B \text{ such that } C \text{ and } D \text{ are not disjoint} && \text{(by hypothesis)} \\
 &\implies C \cap D \neq \emptyset && \text{(by def. of disjoint)} \\
 &\implies \exists x \in C \cap D && \text{(by def. of non-empty)} \\
 &\implies x \in C \wedge x \in D && \text{(by def. of } \cap \text{)} \\
 &\implies x \in A \wedge x \in B && \text{(by def. of subset)} \\
 &\implies x \in A \cap B && \text{(by def. of } \cap \text{)} \\
 &\implies A \cap B \neq \emptyset && \text{(by def. of non-empty)} \\
 &\implies A \text{ and } B \text{ are not disjoint} && \text{(by def. of disjoint)}
 \end{aligned}$$

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