

# Math Camp: PS 1 Logic and Sets

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## 1. Using truth tables, prove both of DeMorgan's Laws for logical connectives.

(a)  $\neg(P \wedge Q)$  is logically equivalent to  $\neg P \vee \neg Q$

$P$	$Q$	$\neg P$	$\neg Q$	$\neg(P \wedge Q)$	$\neg P \vee \neg Q$
T	T	F	F	<b>F</b>	<b>F</b>
T	F	F	T	<b>T</b>	<b>T</b>
F	T	F	F	<b>T</b>	<b>T</b>
F	F	F	T	<b>T</b>	<b>T</b>

The bolded columns are identical, thus  $\neg(P \wedge Q)$  is logically equivalent to  $\neg P \vee \neg Q$ .

(b)  $\neg(P \vee Q)$  is logically equivalent to  $\neg P \wedge \neg Q$

$P$	$Q$	$\neg P$	$\neg Q$	$\neg(P \vee Q)$	$\neg P \wedge \neg Q$
T	T	F	F	<b>F</b>	<b>F</b>
T	F	F	T	<b>F</b>	<b>F</b>
F	T	F	F	<b>F</b>	<b>F</b>
F	F	F	T	<b>T</b>	<b>T</b>

The bolded columns are identical, thus  $\neg(P \vee Q)$  is logically equivalent to  $\neg P \wedge \neg Q$ .

## 2. Find the contrapositive and converse of each of the following statements:

(a) "If squares have four sides, then triangles have four sides."

- Contrapositive: "If triangles do not have four sides, then squares do not have four sides."
- Converse: "If triangles have four sides, then squares have four sides."

(b) "A sequence  $a$  is bounded whenever  $a$  is convergent."

- Rewritten in  $p \Rightarrow q$  form: "If a sequence  $a$  is convergent, then  $a$  is bounded."
- Contrapositive: "If sequence  $a$  is not bounded, then  $a$  is not convergent."
- Converse: "If sequence  $a$  is bounded, then  $a$  is convergent."

(c) "The differentiability of a function  $f$  is sufficient for  $f$  to be continuous."

- Rewritten in  $p \Rightarrow q$  form: "If function  $f$  is differentiable, then  $f$  is continuous."
- Contrapositive: "If function  $f$  is not continuous, then  $f$  is not differentiable."
- Converse: "If function  $f$  is continuous, then  $f$  is differentiable."

## 3. Let $x$ and $y$ be integers. Prove that if $x$ and $y$ are even, then $x + y$ is even.

To show:  $x + y$  is even.

**Proof:**

Let $x$ and $y$ be even integers.	(by hypothesis)
$\Rightarrow \exists i, j \in \mathbb{Z} \ni x = 2i \wedge y = 2j$	(by definition of even)
$\Rightarrow x + y = 2i + 2j$	(substitution)
$\Rightarrow x + y = 2(i + j)$	(distributivity)
$\Rightarrow \exists k \in \mathbb{Z} \ni k = i + j$	(closure of integer addition)
$\Rightarrow x + y = 2k$	(substitution)
$\Rightarrow x + y$ is even ■	(by definition of even)

4. **Let  $A$  and  $B$  be sets. Prove that  $A \subset B$  if and only if  $A - B = \emptyset$ .** By contrapositive,  $A - B \neq \emptyset \iff A \not\subset B$

Proof by contraposition, to show:  $A \not\subset B$ , and then back to  $A - B \neq \emptyset$ .

**Proof:**

Let $A$ and $B$ be sets such that $A - B \neq \emptyset$	(toward contrapositive)
$\iff \exists C \neq \emptyset \ni A - B = C$	(definition of not empty set)
$\iff \exists x \in C \ni x \in A \wedge x \notin B$	(definition of set subtraction)
$\iff A \not\subset B$ ■	(by definition of subset)

While I believe each of these steps should be biconditional, it seems like by creating a set  $C$  to contain an element  $x$  (on the way toward  $A \not\subset B$ ) may not be so reversible, since moving upward I would invoke  $\exists x \in C$  before  $\exists C$ . Maybe I can skip invoking set  $C$  entirely? Looking forward to the answer key!