

# Are We There Yet?

## Iterative Methods for Solving Linear Systems

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- Description and examples of Richardson, Jacobi, and Gauss-Seidel methods



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  - Example: finding solutions to linear systems (today's topic!)





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- Note: often the initial vector  $x^{(0)}$  is an estimate of the solution or arbitrary ( $x = 0$ )



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$$\|x^{(k)} - x\| \leq \|I - Q^{-1}A\|^k \|x^{(0)} - x\|$$

- Thus, if  $\|I - Q^{-1}A\| < 1$ , then  $\lim_{k \rightarrow \infty} \|x^{(k)} - x\| = 0$  [[How does this follow from the last step?]]

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### Corollary

*The iteration formula*

$$Qx^{(k)} = (Q - A)x^{(k-1)} + b$$

*will produce a convergent sequence provided that  $\rho(I - Q^{-1}A) < 1$ .*



Method	$Q$	Iteration Formula: $x^{(k)} = (I - Q^{-1}A)x^{(k-1)} + Q^{-1}b$
Richardson	$I$	$x^{(k)} = (I - A)x^{(k-1)} + b = x^{(k-1)} + r^{(k-1)}$
Jacobi	$D$	$x^{(k)} = (I - D^{-1}A)x^{(k-1)} + D^{-1}b$
Gauss-Seidel	$L$	$x^{(k)} = (I - L^{-1}A)x^{(k-1)} + L^{-1}b$

where

- $D$ : diagonal matrix where  $d_{ii} = a_{ii}$
- $L$ : lower triangular matrix where  $l_{ij} = a_{ij}, i \geq j$