

Are We There Yet?

Iterative Methods for Solving Linear Systems

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- Description and examples of Richardson, Jacobi, and Gauss-Seidel methods

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 - Example: finding solutions to linear systems (today's topic!)

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 - $[x^{(k)}]$ is easy to compute
- Note: often the initial vector $x^{(0)}$ is an estimate of the solution or arbitrary ($x = 0$)

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$$\|x^{(k)} - x\| \leq \|I - Q^{-1}A\|^k \|x^{(0)} - x\|$$

- Thus, if $\|I - Q^{-1}A\| < 1$, then $\lim_{k \rightarrow \infty} \|x^{(k)} - x\| = 0$ [[How does this follow from the last step?]]

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Corollary

The iteration formula

$$Qx^{(k)} = (Q - A)x^{(k-1)} + b$$

will produce a convergent sequence provided that $\rho(I - Q^{-1}A) < 1$.