# A Summary of Methods of Solution for Linear Systems

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MTH 499-02

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## Outline

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• Types of methods

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- Direct methods and examples

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- Indirect methods and examples

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Note: methods can be general or exploit certain matrix characteristics



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- $\bullet$  Recall: GE is equivalent to computing LU factorization

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    - Unscaled: compare absolute value of entries directly
    - Scaled: compare ratios of pivot entries to maximum entries in magnitude in each row

$$A_1 = \begin{bmatrix} 1 & 3 & 6 \\ 2 & 1 & 1 \\ 1 & 3 & 3 \end{bmatrix}$$

$$p = \begin{bmatrix} 1, & 2, & 3 \end{bmatrix}$$

1: function GE\_PP(
$$n, (a_{ij}), (p_i)$$
)
2: for  $k = 1$  to  $n - 1$  do
3: for  $i = k + 1$  to  $n$  do
4:  $z \leftarrow a_{p_i k}/a_{p_k k}$ 
5:  $a_{p_i k} \leftarrow 0$ 

6: **for** 
$$j = k + 1$$
 **to**  $n$  **do**

7: 
$$a_{p_ij} \leftarrow a_{p_ij} - z a_{p_kj}$$

8: **return** 
$$(a_{ij})$$

5:

$$\max(|1|, |2|, |1|) = 2$$

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$$R_3 \leftarrow R_3 - \frac{1}{2}R_2$$

$$R_2 \leftarrow \frac{1}{2}R_2$$



$$A_2 = \begin{bmatrix} 0 & 2.5 & 5.5 \\ 1 & 0.5 & 0.5 \\ 0 & 2.5 & 2.5 \end{bmatrix}$$

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$$A_3 = \begin{bmatrix} 0 & 1 & 2.2 \\ 1 & 0.5 & -0.6 \\ 0 & 0 & -3 \end{bmatrix}$$

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## Direct Methods: Pivoting Examples

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## Direct Methods: Pivoting Examples

$$A_4 = \begin{bmatrix} 0 & 1 & 2.2 \\ 1 & 0.5 & -0.6 \\ 0 & 0 & 1 \end{bmatrix}$$
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# Iterative Methods

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### Theorem

For the linear system Ax = b with A invertible, define the iteration formula

$$x^{(k)} = Gx^{(k-1)} + c.$$

The sequence  $\{x^{(k)}\}$  will converge to  $(I-G)^{-1}c$  provided that  $\rho(G) < 1$ .

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Method	Iteration Formula: $x^{(k)} = (I - Q^{-1}A)x^{(k-1)} + Q^{-1}b$
Richardson	$x^{(k)} = (I - A)x^{(k-1)} + b = x^{(k-1)} + r^{(k-1)}$
Jacobi	$x^{(k)} = (I - D^{-1}A)x^{(k-1)} + D^{-1}b$
Gauss-Seidel	$x^{(k)} = (I - L^{-1}A)x^{(k-1)} + L^{-1}b$
SOR	$x^{(k)} = (I - (\alpha D - C)^{-1} A) x^{(k-1)} + (\alpha D - C)^{-1} b$

#### where

- D: diagonal matrix where  $d_{ii} = a_{ii}$
- L: lower triangular matrix where  $l_{ij} = a_{ij}, i \geq j$
- $C: C + C^* = D A, C^*$  Hermitian,  $\alpha \in \mathbb{R} \ni \alpha > \frac{1}{2}$

Richardson

System Stipulations

• A invertible

Iteration Formula

• 
$$x^{(k)} = (I - A)x^{(k-1)} + b = x^{(k-1)} + r^{(k-1)}$$

Convergence Criteria

• 
$$\rho(I - A) < 1$$

• Consider matrix 
$$A = \begin{bmatrix} 6 & 1 & 1 \\ 2 & 4 & 0 \\ 1 & 2 & 6 \end{bmatrix}$$
 with solution  $b = \begin{bmatrix} 12 \\ 0 \\ 6 \end{bmatrix}$ .

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0

$$x^{(1)} = \begin{bmatrix} 0 & -1/6 & -1/6 \\ -1/3 & 1/3 & 0 \\ -1/6 & -1/3 & 0 \end{bmatrix} \begin{bmatrix} 2 \\ 2 \\ 2 \end{bmatrix} + \frac{1}{6} \begin{bmatrix} 12 \\ 0 \\ 6 \end{bmatrix}$$
$$= \begin{bmatrix} 4/3 & 0 & 0 \end{bmatrix}^{T}$$

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$$= \begin{bmatrix} 4/3 & 0 & 0 \end{bmatrix}^{T}$$

• After 12 iterations,  $x^{(12)} \approx \begin{bmatrix} 2 & -1 & 1 \end{bmatrix}^T$ .

• Jacobi

System Stipulations

- A invertible
- ullet A diagonally dominant

Iteration Formula

• 
$$x^{(k)} = (I - D^{-1}A)x^{(k-1)} + D^{-1}b$$

• 
$$D = a_{ii}$$

Convergence Criteria

• 
$$\rho(I - D^{-1}A) < 1$$

• Gauss-Seidel

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$$L = a_{ij}, i \geq j$$

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• SOR

System Stipulations

- Complex system
- A: positive definite, Hermitian

Iteration Formula

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$$x^{(k+1)} = (I - (\alpha D - C)^{-1}A)x^{(k)} + (\alpha D - C)^{(-1)}b$$

• D: positive definite, Hermitian

• 
$$C: C + C^* = D - A$$

• 
$$\alpha \in \mathbb{R} \ni \alpha > \frac{1}{2}$$

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Typical values

• 
$$D = a_{ii}, C = -a_{ij} \ni i \ge j$$