

A Summary of Methods of Solution for Linear Systems

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MTH 499-02

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- Types of methods

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- Direct methods and examples

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- Indirect methods and examples

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Note: methods can be general or exploit certain matrix characteristics

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- Recall: GE is equivalent to computing LU factorization

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 - Complete: analyze entire submatrix
 - Unscaled/scaled: pertains to how the potential pivot locations are compared
 - Unscaled: compare absolute value of entries directly
 - Scaled: compare ratios of pivot entries to maximum entries in magnitude in each row

$$A_1 = \begin{bmatrix} 1 & 3 & 6 \\ 2 & 1 & 1 \\ 1 & 3 & 3 \end{bmatrix}$$

$$p = [1, \quad 2, \quad 3]$$

```

1: function GE_PP( $n, (a_{ij}), (p_i)$ )
2:   for  $k = 1$  to  $n - 1$  do
3:     for  $i = k + 1$  to  $n$  do
4:        $z \leftarrow a_{p_i k} / a_{p_k k}$ 
5:        $a_{p_i k} \leftarrow 0$ 
6:       for  $j = k + 1$  to  $n$  do
7:          $a_{p_i j} \leftarrow a_{p_i j} - z a_{p_k j}$ 
8:   return ( $a_{ij}$ )

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Theorem

For the linear system $Ax = b$ with A invertible, define the iteration formula

$$x^{(k)} = Gx^{(k-1)} + c.$$

The sequence $\{x^{(k)}\}$ will converge to $(I - G)^{-1}c$ provided that $\rho(G) < 1$.

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Methods

- Richardson
- Jacobi
- Gauss-Seidel
- SOR
- Steepest Descent
- Conjugate Gradient

Method

- Richardson

System Stipulations

- A invertible

Iteration Formula

- $x^{(k)} = (I - A)x^{(k-1)} + b = x^{(k-1)} + r^{(k-1)}$

Convergence Criteria

- $\rho(I - A) < 1$

- Consider matrix $A = \begin{bmatrix} 6 & 1 & 1 \\ 2 & 4 & 0 \\ 1 & 2 & 6 \end{bmatrix}$ with solution $b = \begin{bmatrix} 12 \\ 0 \\ 6 \end{bmatrix}$.

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$$\begin{aligned} x^{(1)} &= \begin{bmatrix} 0 & -1/6 & -1/6 \\ -1/3 & 1/3 & 0 \\ -1/6 & -1/3 & 0 \end{bmatrix} \begin{bmatrix} 2 \\ 2 \\ 2 \end{bmatrix} + \frac{1}{6} \begin{bmatrix} 12 \\ 0 \\ 6 \end{bmatrix} \\ &= [4/3 \quad 0 \quad 0]^T \end{aligned}$$

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- After 12 iterations, $x^{(12)} \approx [2 \quad -1 \quad 1]^T$.

Method

- Jacobi

System Stipulations

- A invertible
- A diagonally dominant

Iteration Formula

- $x^{(k)} = (I - D^{-1}A)x^{(k-1)} + D^{-1}b$
- $D = a_{ii}$

Convergence Criteria

- $\rho(I - D^{-1}A) < 1$

Method

- Gauss-Seidel

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- A diagonally dominant

Iteration Formula

- $x^{(k)} = (I - L^{-1}A)x^{(k-1)} + L^{-1}b$
- $L = a_{ij}, i \geq j$

Convergence Criteria

- $\rho(I - L^{-1}A) < 1$

Method

- SOR

System Stipulations

- Complex system
- A : positive definite, Hermitian

Iteration Formula

- $x^{(k+1)} = (I - (\alpha D - C)^{-1}A)x^{(k)} + (\alpha D - C)^{-1}b$
- D : positive definite, Hermitian
- $C : C + C^* = D - A$
- $\alpha \in \mathbb{R} \ni \alpha > \frac{1}{2}$

Convergence Criteria

- $\rho(I - (\alpha D - C)^{-1}A) < 1$

Typical values

- $D = a_{ii}, C = -a_{ij} \ni i \geq j$