# CS 61A Fall 2016

# Structure and Interpretation of Computer Programs

Orders of Growth Review

For Exercises 1-6 (below), identify the order of growth of the runtime as a function of n. As an example, your answer might be  $O(\sqrt{n})$ ... or even  $O(n^5 \log n)$ . Note that this code will also be on the slides, along with the solutions to all of these problems!

```
1. (2 points) Exercise 1
  def mystery1(n):
      n, result = str(n), ''
      num_digits = len(n)
      for i in range(num_digits):
          result += n[num_digits - i - 1]
      return result
  def mystery2(n):
      n, result = 5, 0
      while n <= 3000:
          result += mystery1(n // 2)
          n += 1
      return result
  _____
2. (3 points) Exercise 2
  def mystery3(n):
      if n < 0 or n \le sqrt(n):
          return n
      return n + mystery3(n // 3)
  def mystery4(n):
      if sqrt(n) \le 50:
          return 1
      return n * mystery4(n // 2)
  def mystery5(n):
      for _ in range(int(sqrt(n))):
          n = 1 + 1
      return n
```

```
3. (2 points) Exercise 3
  def mystery6(n):
      while n > 1:
          x = n
          while x > 1:
              print(n, x)
              x = x // 2
          n -= 1
  -----
  def mystery7(n):
      result = 0
      for i in range(n // 10):
          result += 1
          for j in range(10):
              result += 1
              for k in range(10 // n):
                  result += 1
      return result
  _____
4. (2 points) Exercise 4
  def mystery8(n):
      if n == 0:
          return ''
      result, stringified = '', str(n)
      for digit in stringified:
          for _ in range(n):
             result += digit
      result += mystery8(n - 1)
      return result
  _____
  def mystery9(n):
      total = 0
      for i in range(1, n):
          total *= 2
          if i % n == 0:
              total *= mystery9(n - 1)
              total *= mystery9(n - 2)
          elif i == n // 2:
              for j in range(1, n):
                  total *= j
      return total
```

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```
5. (2 points) Exercise 5
  def mystery10(n):
      if n > 0:
          r1 = mystery10(-n)
          r2 = mystery10(n - 1)
          return r1 + r2
      return 1
  _____
  def mystery11(n):
      if n < 1:
          return n
      def mystery12(n):
          i = 1
          while i < n:
              i *= 2
          return i
      return mystery11(n / 2) + mystery11(n / 2) + mystery12(n - 2)
  _____
6. (2 points) Exercise 6
  The orders of growth should now be functions of m and n.
  def mystery13(m, n):
      if n <= 1:
          return 0
      result = 0
      for i in range(3 ** m):
          result += i // n
      return result + mystery13(m - 5, n // 3)
  -----
  def mystery14(m, n):
      result = 0
      for i in range(1, m):
          j = i * i
          while j \le n:
              result += j
              j += 1
          return result
```

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#### 7. (1 points) Exercise 7

Define n to be the length of the input list. How much memory does the following program use as a function of n?

```
def weighted_random_choice(lst):
    temp = []
    for i in range(len(lst)):
        temp.extend([lst[i]] * (i + 1))
    return random.choice(temp)
```

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#### 8. (7 points) Exercise 8

Provide an algorithm that, given a sorted list A of distinct integers, determines whether there is an index i for which A[i] = i. Your algorithm should run in time  $O(\log n)$ , where n is the length of the list.

## 9. (3 points) Summer 2013 MT2 | Q2

(a) (1 pt) What is the order of growth for a call to fizzle(n)?

```
def fizzle(n):
    if n <= 0:
        return n
    elif n % 23 == 0:
        return n
    return fizzle(n - 1)</pre>
```

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```
(b) (1 pt) What is the order of growth for a call to explode(n)?
       def boom(n):
            if n == 0:
                 return 'BOOM!'
            return boom(n - 1)
       def explode(n):
            if n == 0:
                return boom(n)
            i = 0
            while i < n:
                boom(n)
                 i += 1
            return boom(n)
    (c) (1 pt) What is the order of growth for a call to dreams(n)?
       def dreams(n):
            if n \le 0:
                 return n
            if n > 0:
                 return n + dreams (n // 2)
       _____
10. (4 points) Summer 2014 MT2 | Q6
   Consider the following function (assume that parameter S is a list):
   def umatches(S):
        result = set()
        for item in S:
            if item in result:
                 result.remove(item)
             else:
                 result.add(item)
```

(a) (1 pt) Fill in the blank: The function umatches returns the set of all

(b) (1 pt) Let's assume that the operations of adding to, removing from, or checking containment in a set each take roughly constant time. Give an asymptotic bound (the tightest you can) on the worst-case time

for umatches as a function of N = len(S).

(c) (1 pt) Suppose that instead of having result be a set, we make it a list (so that it is initialized to [] and we use .append to add an item). What now is the worst-case time bound? You can assume that .append is a constant-time operation, and .remove and the in operator require time that is  $\Theta(L)$  in the worst case, where L is the length of the list operated on. Since we never add an item that is already in the list, each value appears at most once, just as for a Python set.

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return result

(d) (1 pt) Now suppose that we consider only cases where the number of different values in list S is at most 100, and we again use a list for result. What is the worst-case time now?

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11. (2 points) Summer 2015 MT2 | Q5(d)
```

```
def append(link, value):
    """Mutates LINK by adding VALUE to the end of LINK."""
    if link.rest is Link.empty:
        link.rest = Link(value)
    else:
        append(link.rest, value)

def extend(link1, link2):
    """Mutates LINK_1 so that all elements of LINK_2
    are added to the end of LINK_1.
    """
    while link2 is not Link.empty:
        append(link1, link2.first)
        link2 = link2.rest
```

(a) (1 pt) What order of growth describes the runtime of calling append? Give your function in terms of n, where n is the number of elements in the input LINK.

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(b) (1 pt) Assuming the two input linked lists both contain n elements, what order of growth best describes the runtime of calling extend?

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#### 12. (2 points) Summer 2012 Final | Q2

(a) (1 pt) What is the order of growth in n of the runtime of collide, where n is its input?

```
def collide(n):
    lst = []
    for i in range(n):
        lst.append(i)
    if n <= 1:
        return 1
    if n <= 50:
        return collide(n - 1) + collide(n - 2)
    elif n > 50:
        return collide(50) + collide(49)
```

(b) (1 pt) What is the order of growth in n of the runtime of into\_me, where n is its input?

```
def crash(n):
    if n < 1:
        return n
    return crash(n - 1) * n</pre>
```

```
def into_me(n):
    lst = []
    for i in range(n):
        lst.append(i)
    sum = 0
    for elem in lst:
        sum = sum + crash(n) + crash(n)
    return sum
```

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## 13. (4 points) Spring 2014 Final | Q5(c)

Give worst-case asymptotic  $\Theta$  bounds – you guys can write them as Big-O bounds – for the running time of the following code snippets. As a reminder, it is meaningful to write things with multiple arguments like  $\Theta(a+b)$ , which you can think of as " $\Theta(N)$  where N=a+b."

```
(a) (1 pt)
   def a(m, n):
       for i in range(m):
           for j in range(n // 100):
               print('hi')
   -----
(b) (1 pt)
   def b(m, n):
       for i in range(m // 3):
           print('hi')
       for j in range(n * 5):
           print('bye')
   _____
(c) (1 pt)
   def d(m, n):
       for i in range(m):
           j = 0
           while j < i:
               j = j + 100
(d) (1 pt)
   def f(m):
       i = 1
       while i < m:
           i = i * 2
       return i
```

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