# Lecture #23 State Space Search (1)

Algorithm
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### In This Lecture

#### ☐ State Space Search

Concept of state space tree

- Techniques for state space search
  - Backtracking
  - Bounded branch

### Outline

Overview

☐ State Space Tree

■ Backtracking

☐ Bounded Branch

### Algorithmic Strategies

- ☐ Algorithmic strategies commonly used for designing algorithms to solve a problem
  - ✓ Brute-force method
    - Enumerate all possible cases of the problem
  - ✓ Divide and conquer
    - Divide the problem into sub-problems & solve them recursively
  - V Dynamic programing
    - Memo the optimal solutions of sub-problems and re-use them
  - ✓ Greedy algorithm
    - Always make the choice that seems to be the best at that moment
  - State space (combinatorial) search 

    Today's topic

#### Motivation

- ☐ Suppose we need to obtain an optimal solution of a problem
  - What if it is hard to find optimal sub-structure?
  - What if there is no overlapping sub-problems?
  - What if it does not have a greedy choice property?
- ☐ If there are no other options, how can do solve it?
  - A brute-force method solves it definitely
- ☐ Q. But it's too slow! Is there an efficient way?
  - A. State space search with bounded branch!

### Outline

- Overview
- ☐ State Space Tree

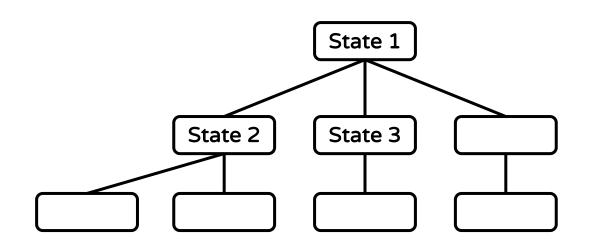
■ Backtracking

☐ Bounded Branch

### State Space Tree

#### ☐ State space tree of an algorithm for a problem

- State: an intermediate state in a process solving a problem
- State space tree: a tree constructed from all of the possible states as nodes
  - Connected via state transitions from initial state (root) to terminal state (leaf)

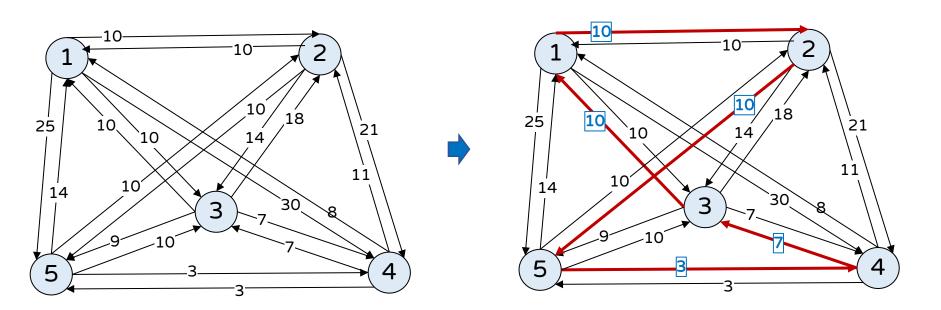


Imagine you enumerate all possible cases

### **Example: SST of TSP**

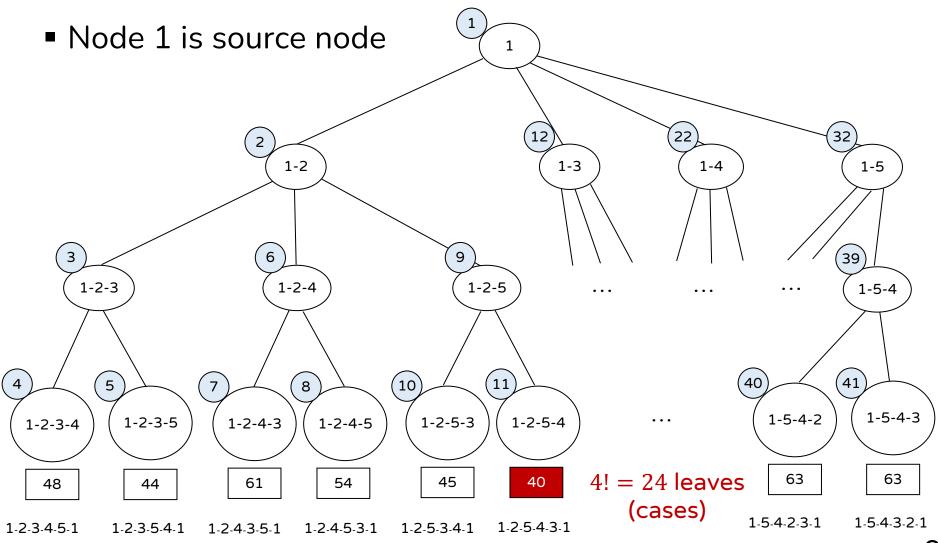
#### □ Asymmetric TSP

- Find the shortest Hamiltonian cycle in a directed graph?
  - Hamiltonian cycle is a directed cycle visiting all nodes exactly once
- Also, it's NP-Hard, but we use this problem to describe SST and bounded branch.



### Example: SST of TSP

#### ☐ SST of Brute-force Search for TSP



### Outline

- Overview
- ☐ State Space Tree

■ Backtracking

☐ Bounded Branch

### Backtracking

#### ☐ Solving a problem is equal to

Searching for an answer in its state space tree ⇒ state
 space search

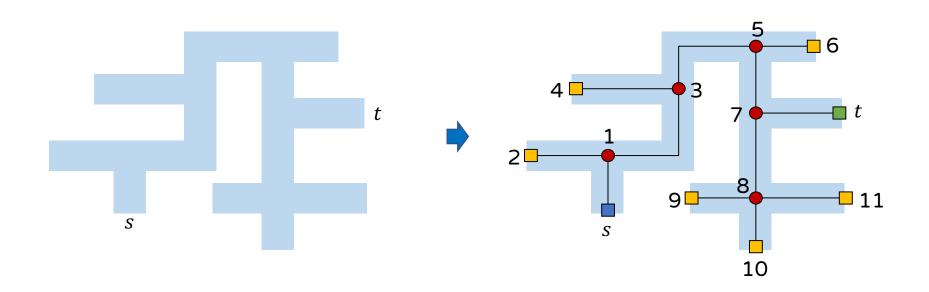
#### ☐ Backtracking is a technique for state space search

- A backtracking algorithm performs DFS (depth first search)
   in the state space tree (SST)
  - DFS goes as deeply as possible, and backtracks when it reaches a dead-end
  - Don't need to explicitly construct an SST because DFS's (or backtracking's) process corresponds to the state space search

### Example Of Backtracking (1)

#### ■ Maze problem

- Find a path from a starting point s to a target point t in a maze ⇒ Do DFS from s
  - Branch (O) and dead-end (□)



### Example Of Backtracking (2)

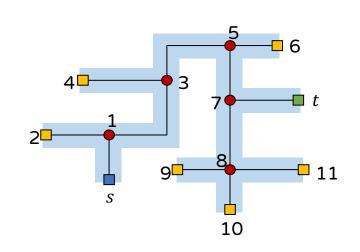
#### ☐ Backtracking algorithm for Maze problem

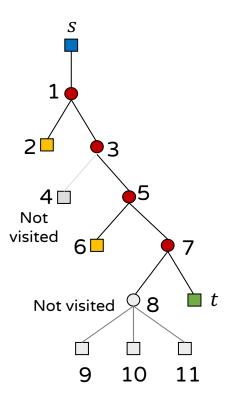
■ Do DFS from s in the maze

```
def maze(u):
    visited[u] ← true

if u == t:
    return true

for each v in N(u):
    if visited[v] == false:
        prev[v] ← u
        maze(v)
```





Backtracking algorithm for Maze problem

State space tree incurred by DFS

### Backtracking v.s. Brute-force

#### ☐ Backtracking can be more efficient than brute-force

- Brute-force method enumerates all possible cases
  - $\circ$  e.g., There are  $k^n$  cases for graph coloring problem
  - However, many of them are not valid ⇒ don't need to see all of them

- Backtracking tries to search for all possible cases
  - But unpromising states could be pruned during DFS!
    - If the backtracking algorithm has pruning strategies
  - The searching cost from pruned states is saved
- ☐ Let's check a pruning technique called bounded branch

### Outline

- Overview
- ☐ State Space Tree

■ Backtracking

■ Bounded Branch

#### **Bounded Branch**

#### ☐ Let's consider TSP problem again

- There are (n-1)! cases in a graph of node n
- However, we don't need to see all of them
  - If the length of a partial path is greater than the best solution we've found so far, the partial path is very unlikely to be a better solution

#### ☐ Main idea of bounded branch (branch & bound)

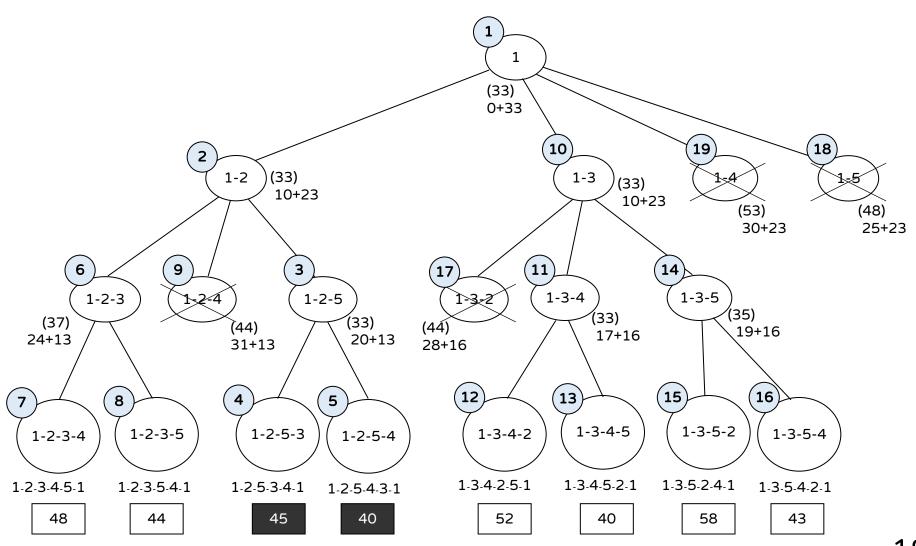
- Let's prune unpromising states during backtracking to make the branches of SST bounded
  - Need to search for all possible cases (backtracking)
  - Need to measure the unpromising-ness of a state (pruning)

### Pruning Technique

#### ☐ How to judge whether a state is pruned?

- In general, we judge it based on the best solution found so for (:= base solution)
- The better the quality of the base solution, the more branches are pruned
  - If a good solution is initially found, the bounded branch is likely to quickly end
  - Otherwise, it could check all possible cases for a worst case
- A practical method for setting base solution
  - Use a good approximate algorithm & set its output to the base solution
  - Then, start the bounded branch algorithm with the base solution

#### ☐ SST of a bounded branch algorithm for TSP



#### ■ Notations

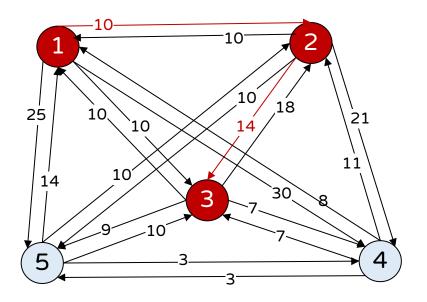
Visited order

6

1-2-3 ← Nodes of a partial HAM-CYCLE

Lower bound of the cost when  $\rightarrow$  (37) this cycle is complete

 $24+13 \leftarrow$  Reason how the lower bound is obtained



HAM-CYCLE: 1-2-3-(4,5)-1

- 24: 10 + 14 from 1-2-3
- 13: 7 + 3 + 3
  - Node 3: 7
  - Node 4: 3
  - Node 5: 3

Select the minimum

weight of out-going edges

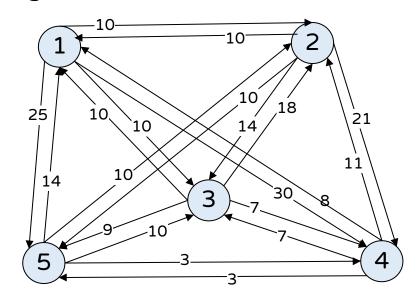
for each node

At least 13 cost is needed to compte the cycle

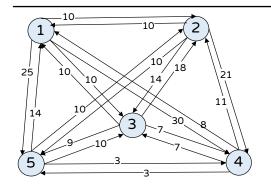
#### ☐ Initial phase

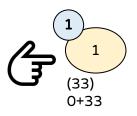
- Find the lower bound of HAM-CYCLE selecting the minimum weight of out-going edges for each node
  - Lower bound: 33
    - Node 1: 10
    - Node 2: 10
    - Node 3: 7
    - Node 4: 3
    - Node 5: 3

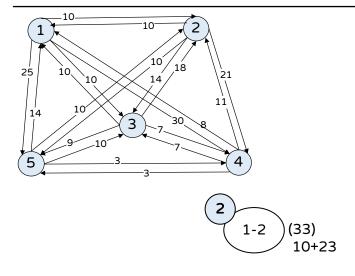
#### Why?



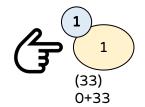
- $\circ$  A Hamiltonian cycle has n edges building a path of n nodes
  - In the path, a node has an out-going edge
- Thus, an optimal solution cannot be less than the lower bound



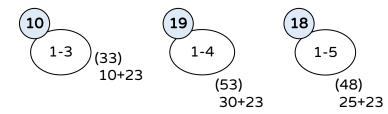


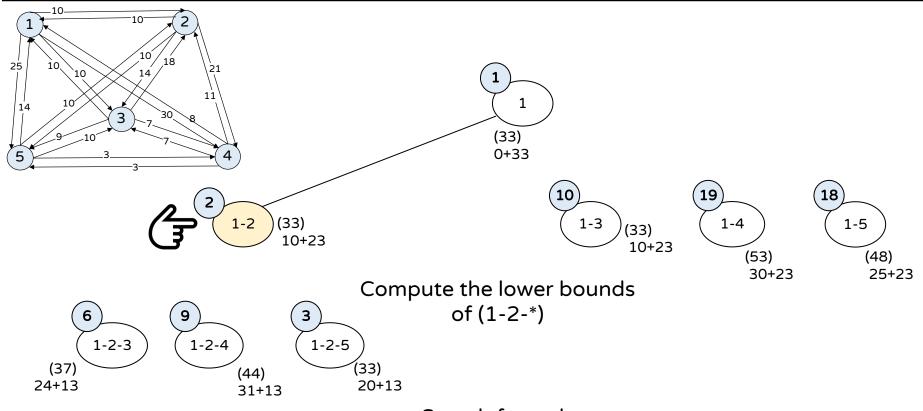


Search from the lowest lower bound!

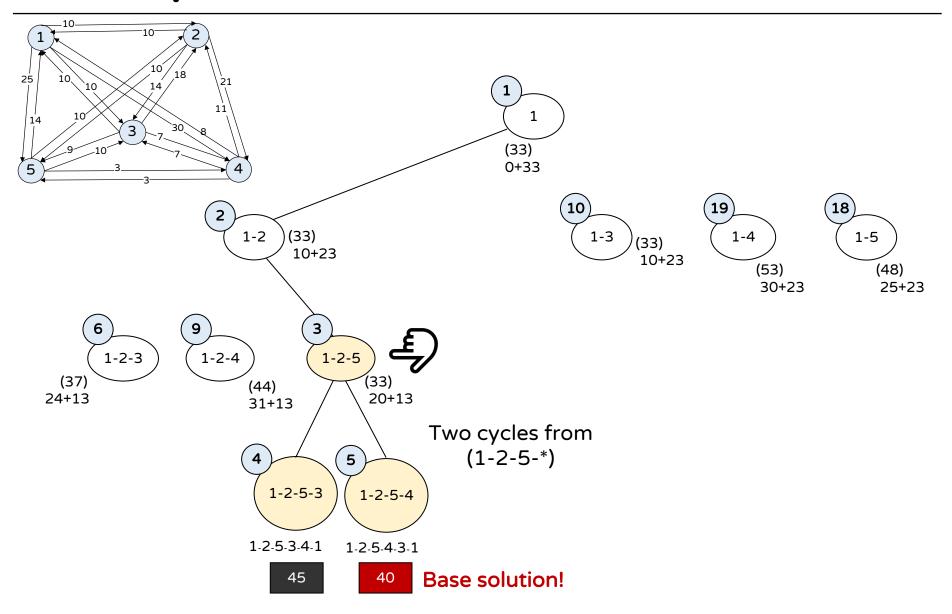


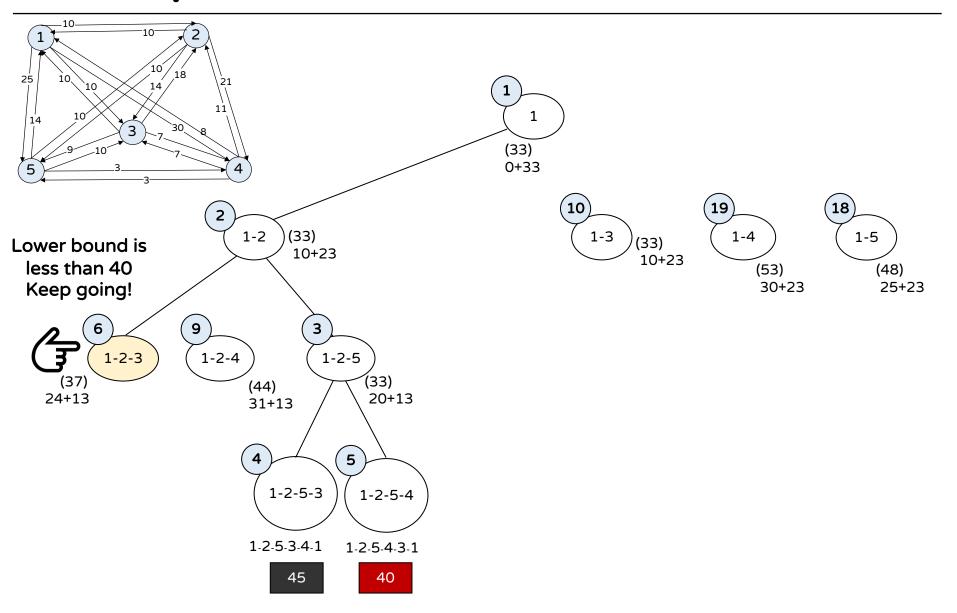
Compute the lower bounds of (1-\*)

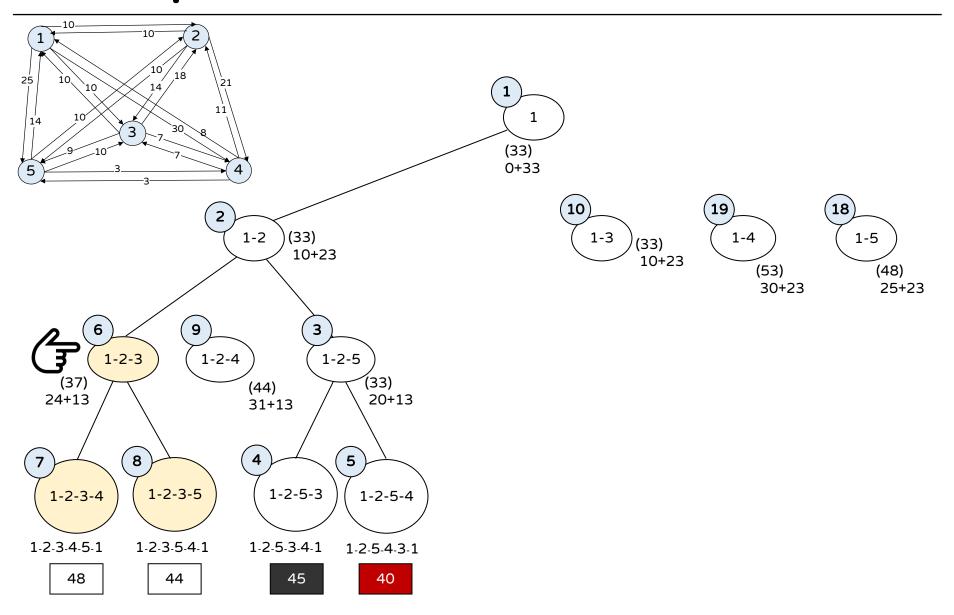




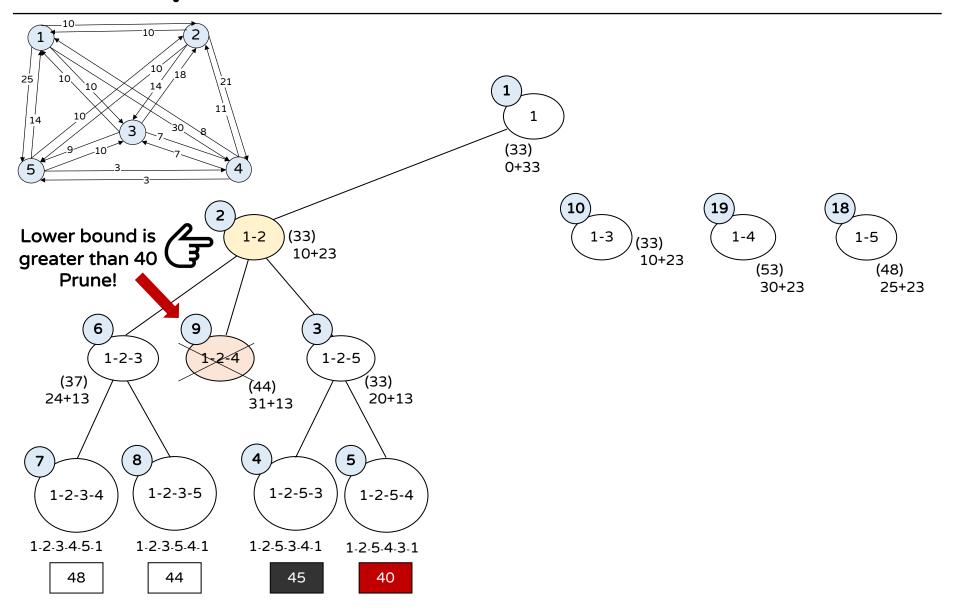
Search from the lowest lower bound!

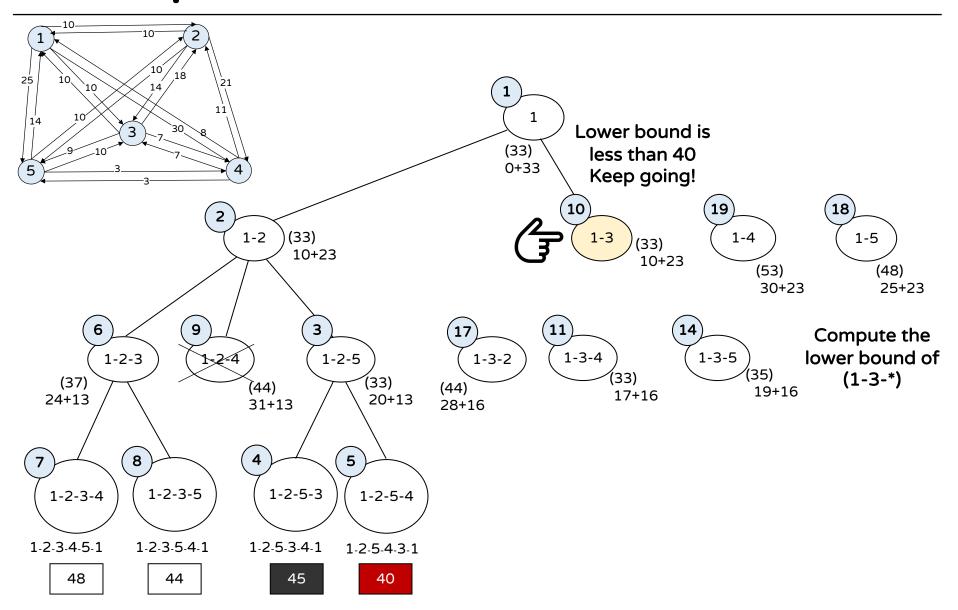


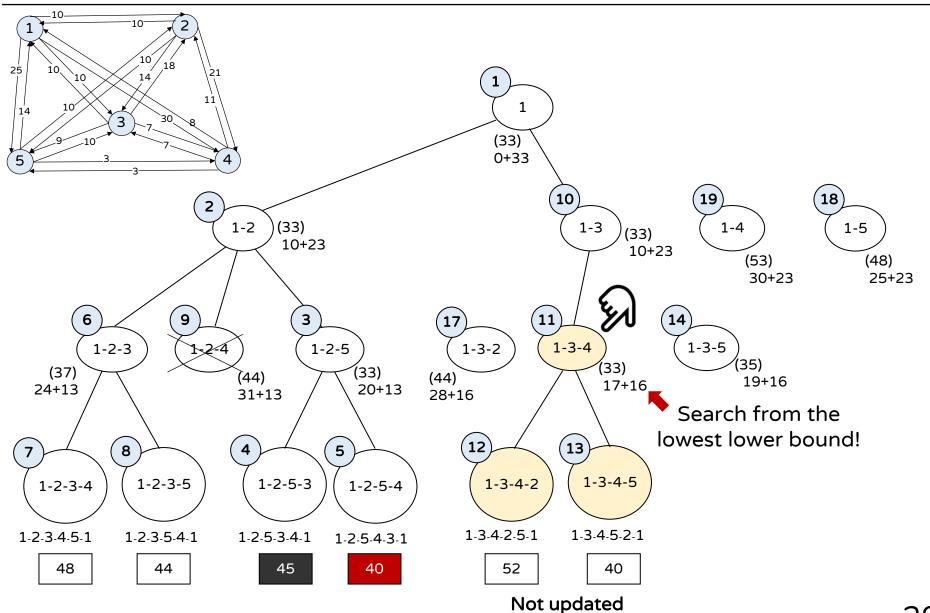


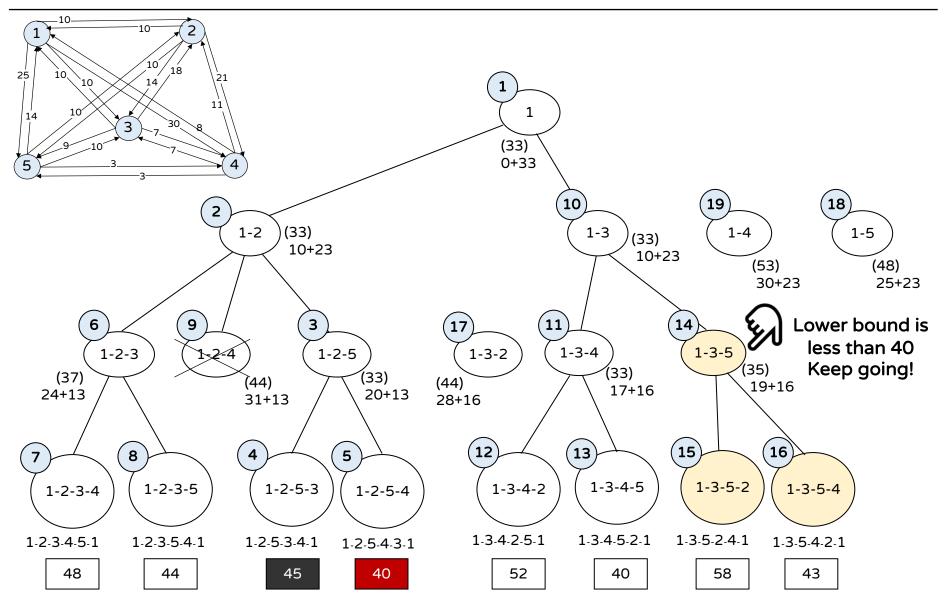


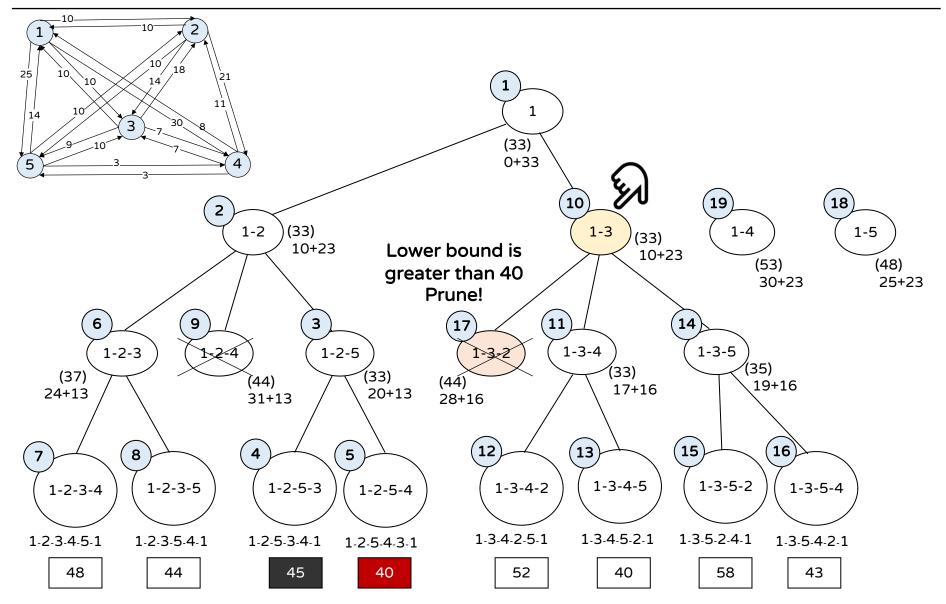
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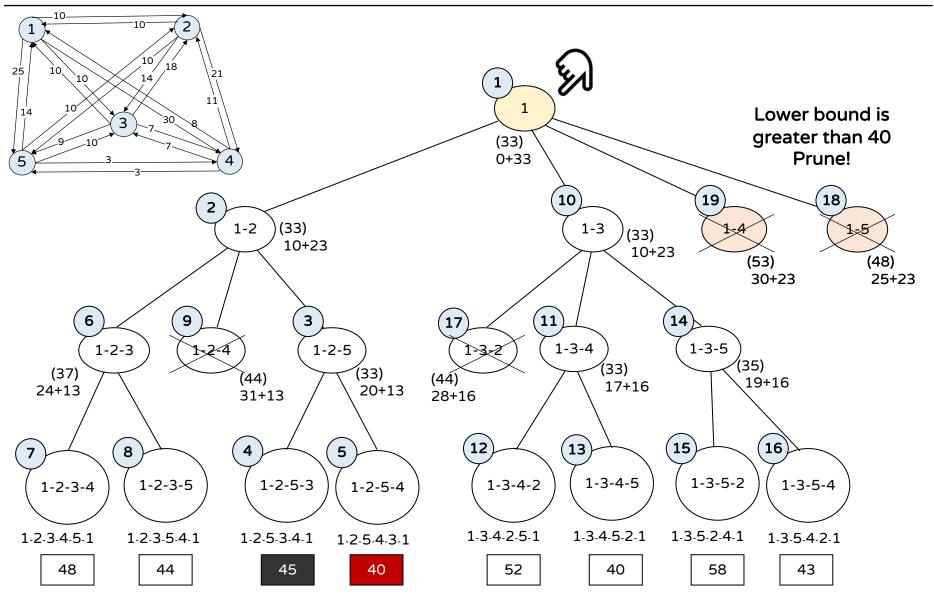












**Final Solution!** 

#### Discussion

#### ☐ Bounded branch v.s. brute-force

- Brute-force algorithm for TSP considers 24 cases
- Bounded branch algorithm considers 8 cases

#### ☐ On base solutions

- In general, the quality of an initial base solution is poor
- The base solution monotonically decreases (or improves) during backtracking
  - Pruning is highly likely to occur at the latter part of backtracking
- Approximate solution is a good option for base solution

#### ☐ Better way for measuring the lower bound

 Increase the probability of pruning by quickly finding a good base solution (it's specific to problems – research)

#### What You Need To Know

- ☐ State space search of an algorithm
  - Searching for an answer in SST depicted by the algorithm
- Backtracking
  - Performs DFS in the state space tree
- ☐ Bounded branch (backtracking + pruning)
  - Prune unpromising states during backtracking to make the branches of SST bounded
- ☐ Pruning technique
  - Prune a state based on a base solution found so far
  - Problem specific

### In Next Lecture

#### □ A\* Search Algorithm

Searching technique on SST with bounded branch

## Thank You