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## Chain development of metric compacts <sup>☆</sup>



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### ABSTRACT

Chain distance between points in a metric space is defined as the infimum of  $\varepsilon$  such that there is an  $\varepsilon$ -chain connecting these points. We call a mapping of a metric compact into the real line a chain development if it preserves chain distances. We give a criterion of existence of the chain development for metric compacts. We prove the diameter of any chain development of a given compact to be the same iff the compact is countable.

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**Notions and basic facts.** Let  $(X, d)$  be a metric space. We call a sequence of points  $x = x_0, x_1, x_2, \dots, x_n = y$  an  $\varepsilon$ -chain if  $d(x_i, x_{i+1}) \leq \varepsilon$  for all  $i$ . Define *chain distance*  $c(x, y)$  as the infimum of  $\varepsilon$  such that there exists an  $\varepsilon$ -chain from  $x$  to  $y$ .

Chain distance satisfies strong triangle inequality:  $c(x, z) \leq \max(c(x, y), c(y, z))$ ; hence it is ultrametric if it does not degenerate. Obviously,  $c = d$  if  $d$  is already ultrametric.

**Definition.** A function  $f: X \rightarrow \mathbb{R}$  is called *chain development* if  $f$  preserves chain distance:

$$c(x, y) = \tilde{c}(f(x), f(y)) \quad \text{for } x, y \in X,$$

where  $c$  is the chain distance on  $(X, d)$  and  $\tilde{c}$  is the chain distance on the set  $f(X)$  with usual distance  $\tilde{d}(s, t) = |s - t|$ .

Chain development was firstly introduced by E.V. Schepin for finite sets as a tool for fast hierarchical cluster analysis. Note that chain development always exists for finite spaces and can be effectively constructed using minimum weight spanning tree of the corresponding graph; see [1] and [2, Section 4] for more details.

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An equivalent construction appeared in the paper [3] by A.F. Timan and I.A. Vestfid: they proved that points of any finite ultrametric space can be enumerated in a sequence  $x_1, \dots, x_n$  such that  $c(x_i, x_j) = \max(c(x_i, x_j), c(x_j, x_k))$  for  $i < j < k$ .

The goal of this paper is to discuss some properties of chain development for infinite spaces. So, there are compacts with no chain developments, e.g. the square  $C \times C$  of a Cantor set. Necessary and sufficient condition of existence of chain developments is given below in Theorem 2.

By *diameter of a chain development*  $f: X \rightarrow \mathbb{R}$  we mean  $\text{diam } f(X) = \sup f(X) - \inf f(X)$ . It is proven in [1] that for finite spaces  $X$  the diameter of chain developments is determined uniquely. It turns out that this is not true in general case.

**Theorem 1.** *Let  $(X, d)$  be a compact metric space. Then the diameter of chain developments (if there are any) is determined uniquely if and only if  $X$  is countable.*

Throughout this paper by  $(Z, d)$  we denote a zero-dimensional compact metric space. We focus on such spaces because study of chain developments for arbitrary compacts essentially reduces to the zero-dimensional case.<sup>1</sup> We have the following property:

- (i)  $(Z, c)$  is an ultrametric space, i.e. chain distance does not degenerate.

Indeed, take  $x, y \in Z$ . The set  $\{x\}$  is a connected component, hence  $x \in U \not\rightarrow y$  for some closed open set  $U$ , so

$$c(x, y) \geq \min_{\substack{u \in U \\ v \in X \setminus U}} d(u, v) > 0.$$

The transition from metric  $d$  to ultrametric  $c$  (which can be seen as a functor) preserves topology:

- (ii) The identity map  $\text{id}: Z \rightarrow Z$  is a homeomorphism between  $(Z, d)$  and  $(Z, c)$ .

Indeed,  $\text{id}$  is 1-Lipshitz ( $c(x, y) \leq d(x, y)$ ), hence it is a continuous bijection from compact to Hausdorff space, hence a homeomorphism.

- (iii) Any chain development  $f: Z \rightarrow \mathbb{R}$  is continuous (with usual topology on  $\mathbb{R}$ ). Hence,  $f(Z)$  is compact and  $f$  is a homeomorphism between  $Z$  and  $f(Z)$ .

Let  $x_n \rightarrow x^*$  in  $Z$ ; prove that  $t_n := f(x_n) \rightarrow t^* := f(x^*)$ . Suppose that  $t_n \not\rightarrow t^*$ , say,  $t_n > t^* + \varepsilon$  for some  $\varepsilon > 0$ . If there are no points of  $f(Z)$  in  $(t^*, t^* + \varepsilon)$ , then  $\tilde{c}(t_n, t^*) \geq \varepsilon$  (where  $\tilde{c}$  is the chain distance on  $f(Z)$ ). And if there is some  $t = f(x) \in (t^*, t^* + \varepsilon)$ , then  $\tilde{c}(t_n, t^*) \geq \tilde{c}(t, t^*) = c(x, x^*) > 0$ . In both cases  $\tilde{c}(t_n, t^*) \not\rightarrow 0$ , which contradicts that  $\tilde{c}(t_n, t^*) = c(x_n, x^*) \leq d(x_n, x^*) \rightarrow 0$ . So,  $f$  is continuous.

The chain distance on a compact  $K \subset \mathbb{R}$  is determined by the lengths of the intervals of the open set  $U_K := [\min K, \max K] \setminus K$ .

- (iv) Chain distance between points  $s, t$  of  $K$  is equal to the maximal length of the intervals of  $U_K$ , lying between  $s$  and  $t$ .

**Existence of chain development.** There is a well-known correspondence between ultrametric spaces and labeled trees; here we describe it for our purposes. Let  $(X, d)$  be a compact metric space; we will construct a

<sup>1</sup> One can identify points of  $(X, c)$  with  $c(x, y) = 0$  to obtain zero-dimensional ultrametric compact  $(Z_X, c)$ ; a chain development of  $(X, d)$  exists if and only if there is a chain development of  $(Z_X, c)$ .

labeled tree  $T(X, d)$  with a vertex set  $V$  and a labeling function  $r: V \rightarrow \mathbb{R}$ . We take an arbitrary point  $v_0$  as a root of our tree and assign to it the  $c$ -diameter of  $X$ , i.e.  $r(v_0) = \max_{x, y \in X} c(x, y)$ . The relation  $c(x, y) < r(v_0)$  is an equivalence relation; hence,  $X$  breaks into finite number of “clusters”  $Q_1, \dots, Q_n$  of points with pairwise chain distance less than  $r(v_0)$ . Next, we connect the root with  $n$  children, say  $v_1, \dots, v_n$ , with  $v_j$  corresponding to  $Q_j$ . Then we repeat the construction for each of  $Q_j$ : we assign  $r(v_j) = \max_{x, y \in Q_j} c(x, y)$ , and connect  $v_j$  with children corresponding to the clusters  $Q_{j,k} \subset Q_j$  with  $c(x, y) < r(v_j)$ ,  $x, y \in Q_{j,k}$ . And so on. The process stops if  $c$ -diameter of a cluster becomes zero.

So, with each vertex  $v$  of  $T(X, d)$  we associate:

- $n(v)$  — the number of children of  $v$ ;
- $C(v)$  — the set of children of  $v$ ;
- $Q(v)$  — the cluster of points, corresponding to  $v$ ; e.g.  $Q(v_0) = X$ ;
- $r(v)$  — the  $c$ -diameter of  $Q(v)$ .

**Definition.** The width of the space  $(X, d)$  is defined as

$$w(X, d) := \sum_v r(v)(n(v) - 1),$$

where the sum is over all vertices of the tree  $T(X, d)$ .

**Theorem 2.** Let  $(X, d)$  be a compact metric space. Then there exists a chain development  $f: X \rightarrow \mathbb{R}$  if and only if  $w(X, d) < \infty$ . Moreover,  $w(X, d)$  is the minimal possible diameter of a chain development of  $X$ .

The construction of the tree uses only the chain distance, so  $T(Z, d) = T(Z, c)$  and  $w(Z, d) = w(Z, c)$ . On the other hand, the ultrametric structure is fully captured by the tree  $T(Z, d)$ . Each point  $x \in Z$  lies in some sequence of clusters; hence, it corresponds to a path in the tree.

**Lemma 1.** Let  $x, y \in Z$ . If  $x \neq y$ , then they lie in different path of the tree, and  $c(x, y)$  is equal to  $r(v)$ , where  $v$  is the lowest common ancestor of  $x, y$ , i.e. the farthest from root vertex lying on both paths.

**Proof.** Assume  $x, y$  lie in the same path  $\{v_0, v_1, \dots\}$  of the tree. The compactness of  $Z$  implies that diameters of the clusters  $Q(v_j)$  tend to zero. Then  $c(x, y)$  is less than any diameter of the corresponding clusters, hence,  $c(x, y) = 0$ , and  $x = y$ .

Let  $v$  be the lowest common ancestor of  $x$  and  $y$ . Then  $c(x, y) \leq r(v)$  by the definition of  $r(v)$  and  $c(x, y) = r(v)$  because  $x, y$  lie in different sub-clusters of  $Q(v)$ .  $\square$

Let us prove [Theorem 2](#).

**Proof.** Consider the case of zero-dimensional ultrametric compact space  $(Z, c)$ . The construction of the set  $f(Z)$  is equivalent to the construction of the tree  $T(Z, c)$ . Pick an interval  $[a, b]$  of length  $w(Z, c)$ ; we know that

$$w(Z, c) = \sum_{v \in C(v_0)} w(Q(v), c) + (n(v_0) - 1)r(v_0).$$

One can remove  $n(v_0) - 1$  disjoint open intervals of length  $r(v_0)$  from  $[a, b]$  so that the remaining  $n(v_0)$  closed intervals will have lengths  $\{w(Q(v), c)\}_{v \in C(v_0)}$ . Those closed intervals correspond to each of  $Q(v)$  and we proceed with them as with  $[a, b]$ .

After removal all of the open intervals we arrive at some closed set  $K \subset [a, b]$ . Every point  $x \in Z$  corresponds to a path in  $T(Z, c)$  and to a nested sequence of closed intervals with non-empty intersection  $t \in K$ ; we put  $f(x) = t$ . Intersection is always a point because  $\mu(K) = 0$  (here and after  $\mu$  stands for the standard Lebesgue measure). The proof that  $f$  is chain development is straight-forward using Lemma 1 and property (iv). Note that  $\text{diam } f(Z) = w(Z, c)$ .

Now, let  $f: Z \rightarrow \mathbb{R}$  be a chain development. Define

$$U_{f(Z)} := [\min f(Z), \max f(Z)] \setminus f(Z).$$

We prove that

$$w(Z, c) = \mu(U_{f(Z)}) = \text{diam } f(Z) - \mu(f(Z)). \quad (1)$$

Remind that  $r(v_0)$  is the  $c$ -diameter of  $Z$  and the  $\tilde{c}$ -diameter of  $f(Z)$ . It is obvious from (iv) that there are exactly  $n(v_0) - 1$  intervals of  $U$  of length  $r(v_0)$ . Repeating this argument with sets  $f(Q(v))$ ,  $v \in C(v_0)$ , we will count all of the intervals of  $U$  and find that each vertex  $v$  corresponds to  $n(v) - 1$  intervals of  $U$  of length  $r(v)$ . That implies (1). Hence,  $w(Z, c) < \infty$  and  $\text{diam } f(Z) \geq w(Z, c)$ .

The general case follows easily.  $\square$

We will make use of the following standard construction.

**Lemma 2.** *Let  $K$  be an uncountable compact in  $[a, b]$ . Then for any  $c > 0$  there is a continuous increasing function  $\theta: [a, b] \rightarrow \mathbb{R}$  such that  $\mu(\theta(K)) = \mu(K) + c$  and  $\mu(\theta(I)) = \mu(I)$  for any interval  $I \subset [a, b] \setminus K$ .*

**Proof.** Write  $K$  as  $N \cup P$ , where  $N$  is countable and  $P$  is perfect. Let  $\varkappa: [a, b] \rightarrow [0, 1]$  be an analog of the Cantor's ladder for the set  $P$ ; we need that  $\varkappa$  is continuous and non-decreasing,  $\varkappa([a, b]) = [0, 1]$  and  $\varkappa|_I \equiv \text{const}$  for any interval  $I \subset [a, b] \setminus P$ . It remains to take  $\theta(t) = t + c\varkappa(t)$ .  $\square$

Now we are ready to prove Theorem 1.

**Proof.** We consider only the zero-dimensional case. If  $Z$  is countable, then  $\mu(f(Z)) = 0$  and from (1) we get  $\text{diam } f(Z) = w(Z, c)$ . Suppose  $Z$  is uncountable. Take any chain development  $f: Z \rightarrow \mathbb{R}$  and apply Lemma 2 to  $K = f(Z)$  with some  $c > 0$ . Then  $\theta \circ f$  gives us a chain development with another diameter.  $\square$

It appears that the diameter of a chain development of an uncountable compact may be any number greater or equal than  $w(X, d)$ .

**Example.** Consider the set  $C \times C$ , where  $C \subset [0, 1]$  is the usual Cantor set. Let  $d((x_1, y_1), (x_2, y_2)) = \max(|x_1 - x_2|, |y_1 - y_2|)$  for  $(x_i, y_i) \in C \times C$ . Then there is no chain development for the space  $(C \times C, d)$ .

**Proof.** Let us compute  $w(C \times C, d)$ . In the tree  $T(C \times C, d)$  each node has four children; for example, the children of the root correspond to the clusters

$$\left(C \cap \left[\frac{2i}{3}, \frac{2i+1}{3}\right]\right) \times \left(C \cap \left[\frac{2j}{3}, \frac{2j+1}{3}\right]\right), \quad i, j = 0, 1. \quad (2)$$

We have  $r(v_0) = 1/3$  for the root  $v_0$  and  $r(u) = \frac{1}{3}r(v)$  for each child  $u$  of  $v$ , by self-similarity of  $C$ . Hence,  $w(C \times C, d) = \sum_{k=0}^{\infty} 4^k 3^{-k} = \infty$  and the claim follows from Theorem 2.  $\square$

### Measure of disconnectivity.

**Definition.** Let  $(X, d)$  be a metric space. Define *measure of disconnectivity* of  $(X, d)$  as

$$\text{dis}(X, d) = \inf_{x_i \sim y_i} \sum_i d(x_i, y_i),$$

where the infimum is taken over sequences (finite or infinite) of pairs  $(x_i, y_i) \in X \times X$ , such that the space  $(X, d)$  with identified points  $x_i \sim y_i$  is a connected topological space.

This notion is closely related to the minimum spanning trees of graphs. Indeed, if  $X$  is finite, then  $\text{dis}(X, d)$  is equal to the weight of a minimum spanning tree for  $X$  (we regard points of  $X$  as vertices and take weights of edges equal to the corresponding distances).

**Theorem 3.** Let  $(X, d)$  be a compact metric space. Then  $\text{dis}(X, d) = w(X, d)$ .

We need one more notation for vertices of a tree  $T(X, d)$ : by  $\text{level}(v)$  we denote the length of the path from the root to  $v$ .

**Proof.** Note that for finite sets  $X$  the theorem follows from [1]. We prove there that  $w(X, d)$  is the diameter of any chain development of  $X$ , and it is clear from the proof that it is equal to the weight of a minimum spanning tree of  $X$ .

Let us prove that  $\text{dis}(X, d) \geq w(X, d)$ . Pick some  $N \in \mathbb{N}$  and consider all clusters  $Q(v)$  with either  $\text{level}(v) = N$  or  $\text{level}(v) < N$  and  $r(v) = 0$ . We denote by  $(X_N, c_N)$  the ultrametric space, which comes from  $(X, c)$  when we identify points in each cluster. To make  $X$  connected, we should connect all of the mentioned clusters, so  $\text{dis}(X, d) \geq \text{dis}(X_N, c_N)$ . For finite sets,  $\text{dis} = w$ , so  $\text{dis}(X_N, c_N) = w(X_N, c_N)$ . Obviously,  $T(X_N, c_N)$  is obtained from  $T(X, d)$  by deleting vertices of level  $> N$ , and assigning  $r(v) = 0$  for the new leaves. So

$$w(X_N, c) = \sum_{\text{level}(v) < N} r(v)(n(v) - 1) \rightarrow w(X, c) \quad \text{as } N \rightarrow \infty,$$

hence  $\text{dis}(X, d) \geq w(X, d)$ .

Let us prove that  $\text{dis}(X, d) \leq w(X, d)$ . For each vertex  $v$  we connect the clusters  $\{Q(u)\}_{u \in C(v)}$  to each other by picking appropriate pairs  $(x_i, y_i) \in C(u') \times C(u'')$ . It is easy to show that one can make the set of that clusters connected using pairs with  $\sum d(x_i, y_i) = r(v)(n(v) - 1)$ . In total, the sum is  $w(X, d)$ . Let us prove that the image  $\tilde{X}$  of  $X$  after projection  $\pi: X \rightarrow \tilde{X}$  of identification  $x_i \sim y_i$ , is connected. If  $\tilde{U} \subset \tilde{X}$  is non-empty, open and closed, then  $U = \pi^{-1}\tilde{U} \subset X$  is also non-empty, open and closed; besides that, if  $x_i \sim y_i$  and  $x_i \in U$ , then  $y_i \in U$ . It remains to prove that  $U = X$ .

If  $x \in U$ , then  $x \in Q(v) \subset U$  for some  $v$ . Indeed,  $\delta := \min_{u \in U, v \in X \setminus U} d(u, v) > 0$ , so if we take  $Q(v) \ni x$  with sufficiently small diameter,  $r(v) < \delta$ , then  $Q(v) \subset U$ . So,  $U$  is a union of clusters; since  $U$  is compact, it is a finite union. Now one can prove via induction on  $N$  that for all  $v$  of level  $\geq N$  either  $Q(v) \subset U$  or  $Q(v) \cap U = \emptyset$ . Indeed,  $U$  is a union of finite number of clusters, so this is true for large  $N$ . Let us make an induction step from  $N$  to  $N - 1$ . Suppose there is  $Q(v)$ ,  $\text{level}(v) = N - 1$ , with  $Q(v) \cap U \neq \emptyset$ . We have  $Q(v) = \sqcup_{u \in C(v)} Q(u)$  so  $Q(u') \cap U \neq \emptyset$  for some  $u' \in C(v)$ . As  $\text{level}(u') = N$ ,  $Q(u') \subset U$ . There is some  $u'' \in C(v)$  and a pair  $x_i \sim y_i$ ,  $(x_i, y_i) \in Q(u') \times Q(u'')$ . As  $x_i \in U$ , we have  $y_i \in U$  and  $Q(u'') \subset U$ . As all the clusters  $\{Q(u)\}_{u \in C(v)}$  are connected, we will prove that  $Q(u) \subset U$  for all  $u \in C(v)$ , i.e.  $Q(v) \subset U$ . The claim follows.

Finally,  $Q(v_0) \subset U$  so  $U = X$  and  $\tilde{X}$  is connected.  $\square$

**Corollary.** *For any metric compact  $(X, d)$  three quantities are equal:*

- *the minimal diameter of a chain development of  $X$ ;*
- *the width  $w(X, d)$ ;*
- *the measure of disconnectivity  $\text{dis}(X, d)$ .*

Note that first two quantities definitely have ultrametric nature, but this is not obvious for the third quantity.

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