

ON THE ORDERED SET OF REFLECTIVE SUBCATEGORIES

G.M. KELLY

Given a category A , we consider the (often large) set $\text{Ref } A$ of its reflective (full, replete) subcategories, ordered by inclusion. It is known that, even when A is complete and cocomplete, wellpowered and cowellpowered, the intersection of two reflective subcategories need not be reflective. Supposing that A admits (i) small limits and (ii) arbitrary (even large) intersections of strong subobjects, we prove that an infimum $\bigwedge_i C_i$ in $\text{Ref } A$ must necessarily be the intersection $\bigcap_i C_i$. Accordingly $\text{Ref } A$ is not in general, even for good A , a complete lattice. We show, however, under the same conditions on A , that $\text{Ref } A$ does admit *small* suprema $\bigvee_i C_i$, given by the closure in A of the union $\bigcup_i C_i$ under the limits of types (i) and (ii) above.

1. Strongly complete categories

We suppose an inaccessible cardinal ∞ chosen once for all, and call a set *small* if its cardinal is less than ∞ . The morphisms of any category A form a set, and A is *small* if this set is small; while A

Received 28 August 1986. The author gratefully acknowledges the assistance of the Australian Research Grants Scheme.

Copyright Clearance Centre, Inc. Serial-fee code: 0004-9727/87 \$A2.00 + 0.00.

is *locally small* if each hom-set $A(A, B)$ is small. We use *large* to mean *not small*. Although a category is said to be *complete* when it admits all *small* limits, an ordered set is called a *complete lattice* only when it admits *all* suprema and infima - even large ones if it is a large set.

We refer to [7] for the definition and properties of *strong monomorphisms*. We call a category *strongly complete* if, besides admitting small limits, it admits the fibred product (or *intersection*) of any family $(f_\alpha: A_\alpha \rightarrow B)$ of strong monomorphisms; in other words, it admits *all intersections of strong subobjects*. Of course a category is strongly complete if it is complete and wellpowered, or even *weakly wellpowered* - in the sense that each object has but a small set of *strong* subobjects.

In a strongly-complete category, the epimorphisms and the strong monomorphisms form a factorization system: given $f: A \rightarrow B$ we take for $i: C \rightarrow B$ the intersection of those strong subobjects of B through which f factorizes then $f = ip$, and p is epimorphic by Lemma 4.4 of [7].

We use *subcategory* to mean *full replete* subcategory. A subcategory B of the strongly-complete A is said to be *closed in A under strong limits* if the limit in A of a diagram in B lies in B , whenever the diagram is either small or else of the form $(f_\alpha: A_\alpha \rightarrow B)$ where the maps f_α in B are strong monomorphisms in A . Clearly the intersection of subcategories closed under strong limits is itself so closed. Given any subcategory D of A , there is accordingly a smallest subcategory D^* containing D which is closed in A under strong limits; we call D^* the *strong-limit closure of D in A* .

PROPOSITION 1. *If B is closed in the strongly-complete A under strong limits, every strong monomorphism $f: C \rightarrow B$ in B is a strong monomorphism in A . Accordingly B is itself strongly complete, and the inclusion $B \rightarrow A$ preserves small limits and arbitrary intersections of strong monomorphisms.*

Proof. Consider the totality of those maps $j: E_j \rightarrow B$ through which f factorizes and which are strong monomorphisms in A with $E_j \in B$. Then their intersection $k: E \rightarrow B$ in A is a strong

monomorphism in A , and $E \in B$ since B is closed in A under strong limits. Let the factorization of f through k be $f = kp$; we show that p is invertible, so that f like k is a strong monomorphism in A . Since f is by hypothesis a strong monomorphism in B , so is p ; thus it suffices to show that p is epimorphic in B . Let $x, y: E \rightarrow F$ in B satisfy $xp = yp$, and let $z: G \rightarrow E$ be the equalizer in A of x and y . Since B is closed in A under small limits, we have $G \in B$; and f factorizes through kz , which is a strong monomorphism in A . By the definition of k , therefore, z is invertible and $x = y$. \square

2. Strong cogenerators

The definition in [7] of strong monomorphism admits an evident generalization to a definition of a *strongly monomorphic family* $(f_\alpha: A \rightarrow B_\alpha)$, in such a way that, if the product $\prod B_\alpha$ exists, the family is strongly monomorphic if and only if the corresponding map $A \rightarrow \prod B_\alpha$ is so.

A *small* set G of objects of A is said to be a *strong cogenerator* if, for each $A \in A$, the family $(h: A \rightarrow G_h)$ of all maps with codomain in G is strongly monomorphic. When A is complete and locally small, this is to say that the canonical map $\eta_A: A \rightarrow \prod_{G \in G} \{A(A, G) \nrightarrow G\}$ (where $X \nrightarrow G$ denotes the product of X copies of G) is a strong monomorphism. From the form of the Special Adjoint Functor Theorem given by Börgers *et al.* in Corollary 1.10 of [3], it follows that:

PROPOSITION 2. *Let the strongly-complete and locally-small A have a strong cogenerator, and let B be locally small. Then a functor $U: A \rightarrow B$ has a left adjoint if and only if it preserves small limits (and hence monomorphisms), and sends an intersection of strong monomorphisms to an intersection of monomorphisms.* \square

If G is a strong cogenerator, the set-valued functors $A(-, G)$ for $G \in G$ jointly reflect isomorphisms. Taking for simplicity the case of a locally-small complete A , suppose $f: A \rightarrow B$ to be such that each $A(f, G)$ is invertible. The naturality of η gives $u\eta_A = \eta_B f$ where u

is the invertible map $\prod_G \{A(f, G) \nrightarrow G\}$; since $u\eta_A$ is a strong monomorphism, so is f . On the other hand, G being *a fortiori* a cogenerator since the η_A are *a fortiori* monomorphic, the $A(-, G)$ are jointly faithful and hence jointly reflect epimorphisms, so that f is epimorphic and hence invertible.

Suppose conversely that the small set G is such that the $A(-, G)$ jointly reflect isomorphisms. Then if η_A factorizes through an epimorphism $f: A \rightarrow B$, this f is invertible. For to say that η_A factorizes through f is to say that every $h: A \rightarrow G$ with $G \in G$ factorizes through f , or that each $A(f, G)$ is surjective; but $A(f, G)$ is injective since f is epimorphic.

If A is such that $g: A \rightarrow C$ is a strong monomorphism whenever g factorizes through no non-trivial epimorphism $f: A \rightarrow B$ (which is the case whenever A admits finite colimits, or whenever epimorphisms and strong monomorphisms constitute a factorization system on A), a small set G is a strong cogenerator exactly when the $A(-, G)$ for $G \in G$ jointly reflect isomorphisms; this is really a special case of Proposition 4.3 of Im and Kelly [6]. Since we have seen that epimorphisms and strong monomorphisms do form a factorization system when A is strongly complete, we have:

PROPOSITION 3. *A small set G of objects of a locally-small and strongly-complete A form a strong cogenerator if and only if a map f is invertible whenever $A(f, G)$ is so for each $G \in G$.* □

PROPOSITION 4. *Let \mathcal{D} be a small subcategory of the locally-small and strongly-complete A , and let \mathcal{D}^* be the strong-limit closure of \mathcal{D} in A . Then \mathcal{D} is a strong cogenerator for \mathcal{D}^* (which is itself strongly complete by Proposition 1.)*

Proof. By Proposition 3, we have to show that a map f in \mathcal{D} is invertible if $A(f, D)$ is so for each $D \in \mathcal{D}$. Consider the subcategory \mathcal{B} of A given by all those objects B of A for which $A(f, B)$ is invertible. If any diagram in \mathcal{B} admits a limit in A , this limit lies

in B ; for $A(f, \lim B_i) \cong \lim A(f, B_i)$. So B is closed in A under strong limits; therefore, since it contains \mathcal{D} , it contains \mathcal{D}^* . Thus f is invertible since it is a map in \mathcal{D}^* with $A(f, B)$ invertible for each $B \in \mathcal{D}^*$. \square

Propositions 1, 2, and 4 now give:

PROPOSITION 5. *If \mathcal{D} is a small subcategory of the locally-small and strongly-complete A and \mathcal{D}^* is the strong-limit closure of \mathcal{D} in A , then \mathcal{D}^* is reflective in A .* \square

3. Infima of reflective subcategories

We write $\text{Ref } A$ for the set of reflective subcategories of A , ordered by inclusion. We use \bigvee and \bigwedge to denote suprema and infima in $\text{Ref } A$, so far as they exist, retaining \cup and \cap for union and intersection.

Remark 6. Even when A is locally small, complete and cocomplete, wellpowered and cowellpowered, and has a generator and a strong cogenerator, the intersection of two reflective subcategories need not be reflective. The following simple counter-example is contained in a forthcoming article [1] by Adámek and Rosický. An object of A is a set with two topologies, a first and a second; and a map $f: A \rightarrow B$ is a function continuous both for the first topology and for the second. The subcategory C_1 [respectively C_2] consists of those objects for which the first [respectively second] topology is compact Hausdorff. A reflexion of $A \in A$ into C_1 is given by taking its Stone-Čech compactification $r: A \rightarrow B$ with respect to the first topology, and giving to B as its second topology the final one with respect to r ; similarly for the reflexion into C_2 . Yet $C_1 \cap C_2$ is not reflective; the formal proof is in [1], but the intuition is clear enough - if we take the reflexion of A into C_1 , and then the reflexion of this into C_2 , and then the reflexion of this last into C_1 , and so on alternately, the cardinal of the successive reflexions increases unboundedly. {The author knows of no

such counter-example for a well-behaved A with a *strong* generator.}

This counter-example does not of itself show that $\text{Ref } A$ for such an A may lack binary infima; but it does so when combined with the following:

THEOREM 7. *Let A be locally small and strongly complete. If a subset $\{C_i\}$ of $\text{Ref } A$ admits an infimum $\bigwedge C_i$, this must be the intersection $\bigcap C_i$.*

Proof. If *any* diagram in a reflective C_i admits a limit in A , this limit lies in C_i ; hence each C_i is closed in A under strong limits. Thus if $D \in \bigcap C_i$ and if $\{D\}$ denotes the subcategory of A consisting of D alone, its strong-limit closure $\{D\}^*$ is contained in $\bigcap C_i$. Since $\{D\}^*$ is reflective in A by Proposition 5, we have $\{D\}^* \subset \bigwedge C_i$. Since this is true for every $D \in \bigcap C_i$ and since $D \in \{D\}^*$, we have $\bigcap C_i \subset \bigwedge C_i$; whence $\bigcap C_i = \bigwedge C_i$. \square

Remark 8. Since $\text{Ref } A$ for a strongly-complete A , or even for an A so good as that of Remark 6, may lack binary infima, it is not in general a complete lattice. Yet we show in Theorem 14 below that $\text{Ref } A$ admits *small* suprema when A is strongly complete. It follows that $\text{Ref } A$ is a *large* set for the A of Remark 6. Since this A has a strong cogenerator $G = \{G_1, G_2, G_3\}$, where each G_i is a two-element set and the respective pairs of topologies are (chaotic, chaotic), (chaotic, Sierpinski), and (Sierpinski, chaotic), $\text{Ref } A$ can be large even when the locally-small strongly-complete A has a strong cogenerator. The following result, therefore, has no converse:

PROPOSITION 9. *Let A be locally small and strongly complete. If $\text{Ref } A$ is a small set, A has a strong cogenerator.*

Proof. Since $\{D\}^*$ is reflective in A for each $D \in A$ by Proposition 5, the set of *distinct* $\{D\}^*$ is small; let it consist of the $\{G\}^*$ where G runs through the small set G . Given $A \in A$, let

$\{A\}^* = \{H\}^*$ where $H \in G$. By Proposition 4, H is a strong cogenerator for $\{A\}^*$, so that the canonical $\zeta: A \rightarrow A(A, H) \not\downarrow H$ is a strong monomorphism in $\{A\}^*$. Here $A(A, H) \not\downarrow H$ is in the first instance the product in $\{A\}^*$; but it is equally the product in A . Moreover ζ is a strong monomorphism in A by Proposition 1. Since ζ factorizes through $\eta_A: A \rightarrow \prod_{G \in G} \{A(A, G) \not\downarrow G\}$, the latter too is a strong monomorphism in A ; so that G is a strong cogenerator for A . \square

4. Reflective factorization systems

We call a set M of maps in a category A a *skein* if it contains the isomorphisms and is closed under composition. A skein M is *stable under pullbacks* if the pullback of a map in M along any map in A is again in M . The skein M is *stable under fibred products* if, whenever a family $(f_\alpha: A_\alpha \rightarrow B)$ of maps in M admits a fibred product $h: C \rightarrow B$, we have $h \in M$; similarly for stability under *small* or *finite* fibred products. (Note the distinction between stability under pullbacks and stability under binary fibred products.) *Stability under the intersections of strong monomorphisms* means stability under those fibred products, possibly large, in which each f_α is a strong monomorphism.

For maps e and m in A we write $e \downarrow m$ if, for every commutative square $ve = mu$, there is a unique "diagonal" w with $we = u$ and $mw = v$. If N is any set of maps in A we write $N^\downarrow = \{m \mid n \downarrow m \text{ for all } n \in N\}$ and $N^\uparrow = \{e \mid e \downarrow n \text{ for all } n \in N\}$.

By Proposition 2.1.1 of Freyd and Kelly [5], N^\downarrow is a skein, stable under pullbacks and fibred products, such that

(1) if fg and f are in N^\downarrow , so is g .

We recall from [5] that a *factorization system* (E, M) on A consists of two skeins E and M such that every map f in A admits a factorization $f = me$ with $m \in M$ and $e \in E$, and such that $e \downarrow m$ for each $e \in E$ and $m \in M$. This last requirement may be expressed equivalently as $M \subset E^\downarrow$, or as $E \subset M^\uparrow$; in fact, by Proposition 2.2.1

of [5], every factorization system satisfies $M = E^\perp$ and $E = M^\perp$, and hence is fully determined by the knowledge of either E or M . In any factorization system $E \cap M$ consists of the isomorphisms, since $e \downarrow e$ only for invertible e . We order factorization systems by setting $(E, M) \leq (E', M')$ when $M \subset M'$; or equivalently when $E' \subset E$.

By the dual of (1), any factorization system (E, M) satisfies

(2) if fg and g are in E , so is f ;

we call (E, M) a *reflective* factorization system if it also satisfies

(3) if fg and f are in E , so is g .

It was proved by Cassidy, Hébert, and Kelly in Section 2 of [4] that, when A is *strongly finitely complete* (that is, when A admits finite limits and arbitrary intersections of strong monomorphisms), there is an order-preserving bijection between reflective factorization systems on A and reflective subcategories of A . We now give a modification of their proof of that result; the new proof provides extra information on M needed in our applications below.

PROPOSITION 10. *Let C be a reflective subcategory of the strongly-finitely-complete A , the reflexion of A into C being $\rho_A: A \rightarrow rA$.*

Let $\Phi C = (E, M)$ where E is the set of maps inverted by $r: A \rightarrow A$ and where M is the smallest skein in A containing $\text{mor } C$ and stable under pullbacks and arbitrary intersections of strong monomorphisms. Then

(4) $C = E^\perp =_{\text{df}} \{A \in A \mid A(e, A) \text{ is invertible for all } e \in E\}$,

and (E, M) is a reflective factorization system with factorizations constructed as in (5) and (6) below.

Proof. If $C \in C$ and $e \in E$, we have $A(e, C)$ invertible since it is conjugate to the invertible $C(r(e), C)$: thus $C \in E^\perp$. We may always so choose the reflexion that $\rho_C: C \rightarrow rC$ for $C \in C$ is $1: C \rightarrow C$; then $r(\rho_A) = 1$ for any $A \in A$. If $A \in E^\perp$ we have $A(\rho_A, A)$ invertible,

Reflective subcategories

since $\rho_A \in E$; so there is an $h: rA \rightarrow A$ with $h\rho_A = 1$, and now $\rho_A h\rho_A = \rho_A$ gives $\rho_A h = 1: rA \rightarrow rA$; thus ρ_A is invertible and $A \in \mathcal{C}$. This proves (4).

E is clearly a skein, and M is a skein by definition. For $C \in \mathcal{C}$, we have $C \in E^\perp$ by (4); this is equally the assertion that $C \rightarrow 1$, where 1 is the terminal object of A , lies in E^\perp . It follows from (1) that $\text{mor } C \subset E^\perp$. Since the skein E^\perp , by the remarks preceding (1), is stable under pullbacks and fibred products, it follows from the definition of M that $M \subset E^\perp$.

Since E clearly satisfies (3), it remains to show that any $f: A \rightarrow B$ in A has an (E, M) factorization. In the diagram

$$(5) \quad \begin{array}{ccccc} A & & & & \\ & \searrow w & & \nearrow \rho_A & \\ & D & \xrightarrow{v} & rA & \\ & \downarrow u & & \downarrow r(f) & \\ & B & \xrightarrow{\rho_B} & rB & \end{array} ,$$

let the square be a pullback. Since $r(f) \in \text{mor } \mathcal{C} \subset M$, we have $u \in M$ by the definition of M . Now let $k: E \rightarrow D$ be the intersection of all those strong subobjects of D , lying in M , through which w factorizes; then we have a factorization

$$(6) \quad A \xrightarrow{w} D = A \xrightarrow{e} E \xrightarrow{k} D +$$

where the strong monomorphism k lies in M by the definition of M . We now show that $e \in E$, so that f has $(uk)e$ as its (E, M) factorization.

Since $r(\rho_A) = 1$, applying r to the top triangle of (5) shows that $r(w)$ is a coretraction; whence $r(e)$ is a coretraction by (6). Let the analogue of (5) with e in place of f be

(7)

$$\begin{array}{ccccc}
 & & A & & \\
 & \nearrow \rho_A & & \searrow & \\
 & D' & \xrightarrow{v'} & rA & \\
 & \downarrow u' & & \downarrow r(e) & \\
 & E & \xrightarrow{\rho_E} & rE &
 \end{array}$$

Diagram (7) is a commutative diagram with five nodes: A at the top left, D' below it, E below D' , rA to the right of D' , and rE to the right of E . Arrows are: $\rho_A: A \rightarrow rA$, $w': A \rightarrow D'$, $e: A \rightarrow E$, $u': D' \rightarrow E$, $v': D' \rightarrow rA$, $\rho_E: E \rightarrow rE$, and $r(e): rA \rightarrow rE$.

Since $r(e)$, being a coretraction, is a strong monomorphism, so is its pullback u' . Being a strong monomorphism in M through which e factorizes, u' is invertible by the definition of k ; so we may as well take $D' = E$, $u' = 1$, and $w' = e$. Applying r to the top triangle of (7) gives $r(v')r(e) = 1$, while applying r to the square in (7) gives $r(e)r(v') = 1$. So $r(e)$ is invertible and $e \in E$. \square

PROPOSITION 11. For any factorization system (E, M) on an A with a terminal object 1 , define a subcategory $C = \Psi(E, M)$ by

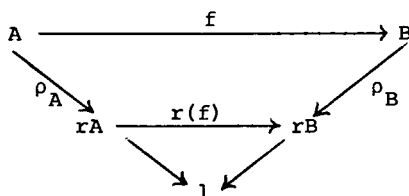
(8) $C = M/1 =_{Df} \{A \in A \mid A \rightarrow 1 \text{ lies in } M\}.$

Then C is reflective in A , the reflexion $\rho_A: A \rightarrow rA$ being the E -part of the (E, M) factorization of $A \rightarrow 1$. \square

THEOREM 12. For a strongly-finitely-complete A , the (clearly order-preserving) functions ϕ and ψ constitute a bijection between reflective factorization systems on A and reflective subcategories of A .

Proof. $\Psi\phi C = M/1$, where M is as in Proposition 10. Since M is E^\perp , we have $M/1 = E^\perp$, which is C by (4); hence $\Psi\phi = 1$.

Let $\Psi(E, M)$, where (E, M) is now reflective, be C as in Proposition 11, and let ΦC be (E', M') . Then $r(f)$ for $f: A \rightarrow B$ in A is the unique map rendering commutative



Since $rA \rightarrow 1$ and $rB \rightarrow 1$ lie in M , we have $r(f) \in M$ by (1). Hence $r(f)$ is invertible, so that $f \in E'$, if and only if $r(f) \in E$. Since $\rho_A \in E$, this is equivalent by (2) to $r(f)\rho_A \in E$, and hence to $\rho_B f \in E$. Since $\rho_B \in E$, this last is equivalent by (3) to $f \in E$. So $E' = E$, and $\Phi\Psi = 1$. \square

The following is contained in Borceux and Kelly [2] in the special case in which C is a *localization*, in the sense that the reflexion r is left exact:

PROPOSITION 13. *Let the reflective subcategory C of the strongly-finitely-complete A correspond as in Theorem 12 to the reflective factorization system (E, M) . Then if $gf \in E$ and $f \in M$, the map f is a strong monomorphism.*

Proof. Since $r(gf) = r(g)r(f)$ is invertible, $r(f)$ is a coretraction and hence a strong monomorphism, whence its pullback u in (5) is a strong monomorphism. Since k is a strong monomorphism in (6) and e in (6) is invertible because $f \in M$, the composite $f = uke$ is a strong monomorphism. \square

5. Small suprema of reflective subcategories

Consider reflective subcategories C_i of a strongly-finitely-complete A and the reflective factorization systems (E_i, M_i) that correspond to them by Theorem 12. If there is a factorization system (E, M) where $E = \cap E_i$, it is clearly reflective by (3), and it is

obviously the supremum of the (E_i, M_i) in the set of all factorization systems, and *a fortiori* in the set of reflective factorization systems. Accordingly the corresponding $C = M/1$ is the supremum of the C_i in *Ref A*. The following is an analogue of Theorem 3.1 of Borceux and Kelly [2], who dealt with the simpler case of small suprema in the ordered set *Loc A* of localizations of *A*:

THEOREM 14. *Let A be strongly complete [respectively strongly finitely complete] and let $\{C_i\}_{i \in I}$ be a set of reflective subcategories with I small [respectively finite]. Set $E = \cap E_i$, and let M be the smallest skein containing $\cup M_i$ which is stable under small [respectively finite] fibred products and all intersections of strong monomorphisms. Then (E, M) is a reflective factorization system, so that the corresponding $C = M/1$ is the supremum of the C_i in *Ref A*.*

Proof. E is a skein since each E_i is so, and M is a skein by definition. Since $M_i = E_i^\perp \subset E^\perp$, and since E^\perp is a skein stable under fibred products by the remarks preceding (1), we have $M \subset E^\perp$. It remains to show that any $f: A \rightarrow B$ has an (E, M) factorization. Let $f = m_i e_i$ be its (E_i, M_i) factorization, and form the fibred product

$$\begin{array}{ccc} & D_i & \\ p_i \nearrow & & \searrow m_i \\ B' & \xrightarrow{n} & B \end{array};$$

then $f = n f'$, where $f': A \rightarrow B'$ is the unique map satisfying $p_i f' = e_i$ for all i . Factorize f' as

$$A \xrightarrow{f'} B' = A \xrightarrow{e} E \xrightarrow{k} B',$$

where k is the intersection of all those strong subobjects of B' , lying in M , through which f' factorizes. By the definition of M , we have $n \in M$ and $k \in M$; it remains to show that $e \in E$, or

equivalently that $e \in E_i$ for each i .

Let $e = m'_i e'_i$ be its (E_i, M_i) factorization. Since

$$e_i = p_i f' = p_i k e = p_i k m'_i e'_i$$

and since e_i and e'_i belong to E_i , we have $p_i k m'_i \in E_i$ by (2).

Because $m'_i \in M_i$, it follows from Proposition 13 that m'_i is a strong monomorphism. Since e factorizes through the strong monomorphism m'_i lying in M , it follows from the definition of k that m'_i is invertible. Thus $e \in E_i$ for each i , and hence $e \in E$. \square

The following, in its statement and in its proof, is a modification of the result given for localizations in Theorem 3.3 of Borceux and Kelly [2]:

THEOREM 15. *In the situation of Theorem 14, C is the strong-limit closure [respectively strong-finite-limit closure] in A of $\cup C_i$.*

Proof. Write \mathcal{D} for the closure in question of $\cup C_i$; clearly $\mathcal{D} \subset C$ since the reflective C is closed in A under all limits that exist; and it remains to show the converse.

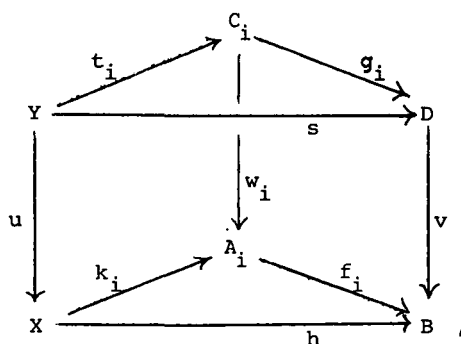
Define as follows a set N of maps in A : the map $f: A \rightarrow B$ lies in N if, for every pullback

$$(9) \quad \begin{array}{ccc} C & \xrightarrow{g} & D \\ u \downarrow & & \downarrow v \\ A & \xrightarrow{f} & B \end{array}$$

with $D \in \mathcal{D}$, we have $C \in \mathcal{D}$.

Since the pullback along $f h$ is the pullback along h of the pullback along f , it is clear that N is a skein. Since the pullback along $w v$ is the pullback along v of the pullback along w , it is clear

that N is stable under pullbacks. Moreover N is stable under such fibred products as exist, and hence under small [respectively finite] ones and all intersections of strong monomorphisms. To see this, consider a diagram



where the base of the prism is a fibred product, and where u, s constitute the pullback of v and h . If we define w_i and g_i as the pullback of v and f_i , there are unique t_i that complete the commutative diagram. Since pullback along v is a right-adjoint functor $A/B \rightarrow A/D$, the top of the prism is again a fibred product. Now if each $f_i \in N$ and $D \in \mathcal{D}$, we have each $C_i \in \mathcal{D}$, whence $Y \in \mathcal{D}$ because \mathcal{D} is closed under small [respectively finite] limits, so that $h \in N$.

Since $C_i \in \mathcal{D}$ and \mathcal{D} is closed under finite limits, it is clear that $\text{mor } C_i \in N$. Because N is stable under pullbacks and fibred products, it follows from Proposition 10 that each $M_i \in N$; and now Theorem 14 gives $M \in N$.

If $A \in \mathcal{C}$, the unique map $f: A \rightarrow 1$ is in M by Proposition 10 (since of course $1 \in \mathcal{C}$) and is hence in N . Applying (9) with v the identity map of 1 , and recalling that $1 \in \mathcal{D}$, we have $A \in \mathcal{D}$. Thus $\mathcal{C} \subset \mathcal{D}$, as required. \square

Remark 16. We have seen in Remark 8 above that large suprema need not exist in $\text{Ref } A$ for a strongly-complete A , even under extra hypotheses on A . We now observe that, even when a large supremum $C = \bigvee C_i$ does exist, it need not be the strong-limit closure of $\bigcup C_i$. The counter-example is contained in Example 3.7 of [2]. We take for A the ordinal sum $1 + \omega^{\text{op}}$, the dual of the ordered set $\omega + 1$ consisting of all ordinals $\leq \omega$. This A is locally small and admits all limits and colimits, even large ones; it is trivially weakly wellpowered and weakly cowellpowered, although neither wellpowered nor cowellpowered; and it trivially admits a generator and a cogenerator. The reflective subcategories $C_\beta = [\beta^{\text{op}}, 0^{\text{op}}]$ for $\beta \in \omega$, where β^{op} denotes the ordinal $\beta < \omega$ as an element of ω^{op} , have as their supremum A itself; but $\bigcup C_\beta = \omega^{\text{op}} \subset A$ is already closed in A under small limits and all intersections of strong monomorphisms.

References

- [1] J. Adámek and J. Rosický, "Intersections of reflective subcategories", (in preparation).
- [2] F. Borceux and G.M. Kelly, "On locales of localizations", *J. Pure Appl. Algebra*, 46 (1987), (to appear).
- [3] R. Börger, W. Tholen, M.B. Wischnewsky and H. Wolff, "Compact and hypercomplete categories", *J. Pure Appl. Algebra* 21 (1981), 129-144.
- [4] C. Cassidy, M. Hébert and G.M. Kelly, "Reflective subcategories, localizations and factorization systems", *J. Austral. Math. Soc. Ser. A* 38 (1985), 287-329; Corrigenda *Ibid* 41 (1986), 286.
- [5] P.J. Freyd and G.M. Kelly, "Categories of continuous functors, I", *J. Pure Appl. Algebra* 2 (1972), 169-191.
- [6] G.M. Im and G.M. Kelly, "Some remarks on conservative functors with left adjoints", *J. Korean Math. Soc.* 23 (1986), 19-33.
- [7] G.M. Kelly, "Monomorphisms, epimorphisms, and pull-backs", *J. Austral. Math. Soc.* 9 (1969), 124-142.

Pure Mathematics Department,
University of Sydney,
New South Wales, 2006
Australia.