

STCLang: State Thread Composition as a Foundation for Monadic Dataflow Parallelism

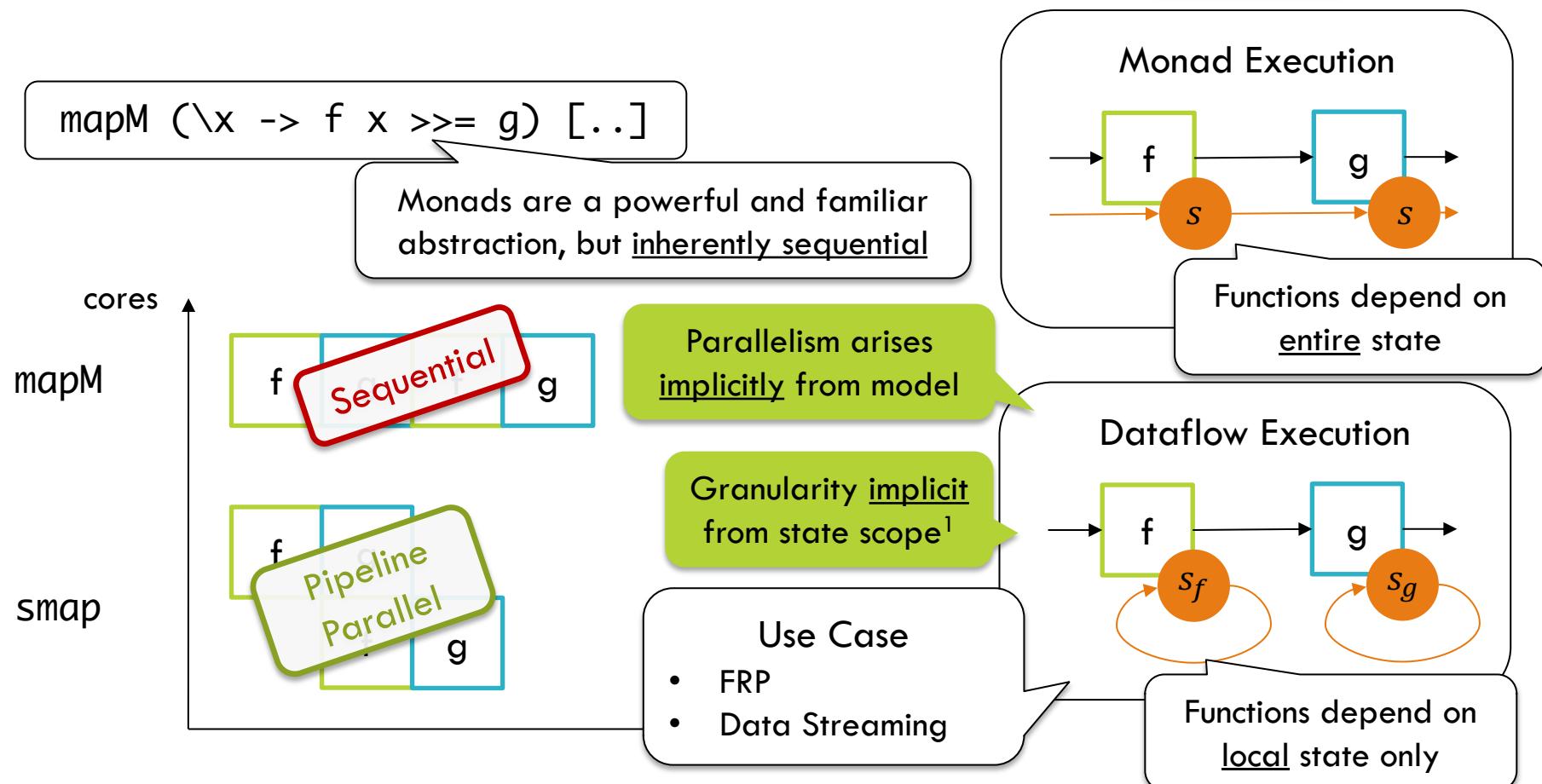
By Sebastian Ertel, Justus Adam, Norman A. Rink,
Andrés Goens and Jerónimo Castrillón-Mazo

We want ...

- Performance
- Responsiveness

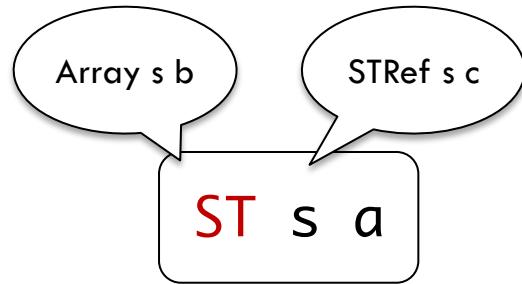
Parallelism in the presence of State

- Efficient
- Intuitive

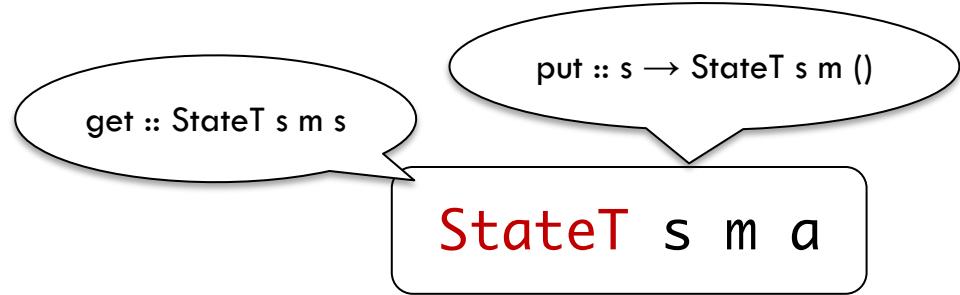


1. Tim Harris and Satnam Singh. 2007. Feedback directed implicit parallelism. In *Proceedings of the 12th ACM SIGPLAN international conference on Functional programming (ICFP '07)*. ACM, New York, NY, USA, 251-264.

Lots of references strewn about

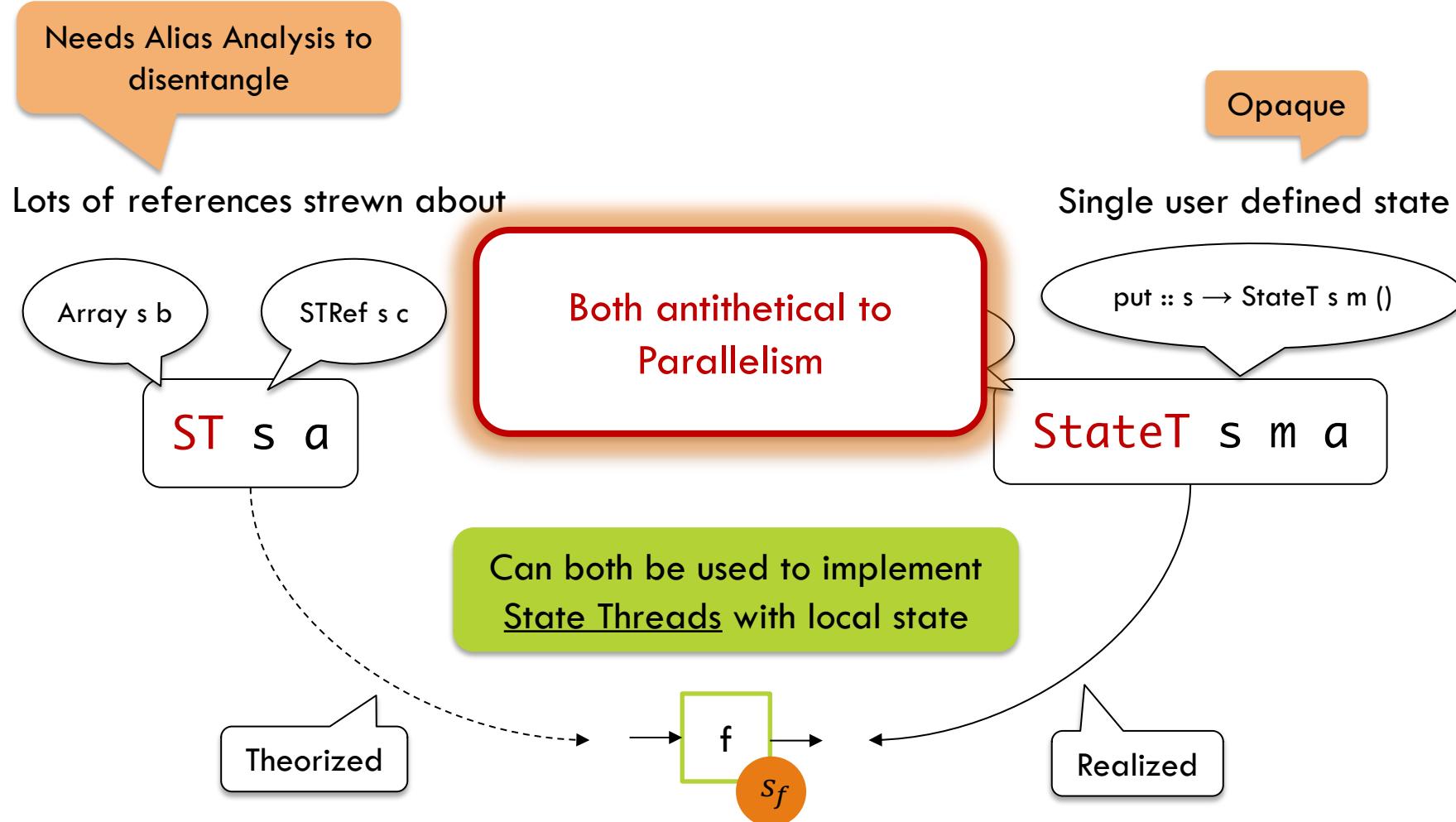


Single user defined state



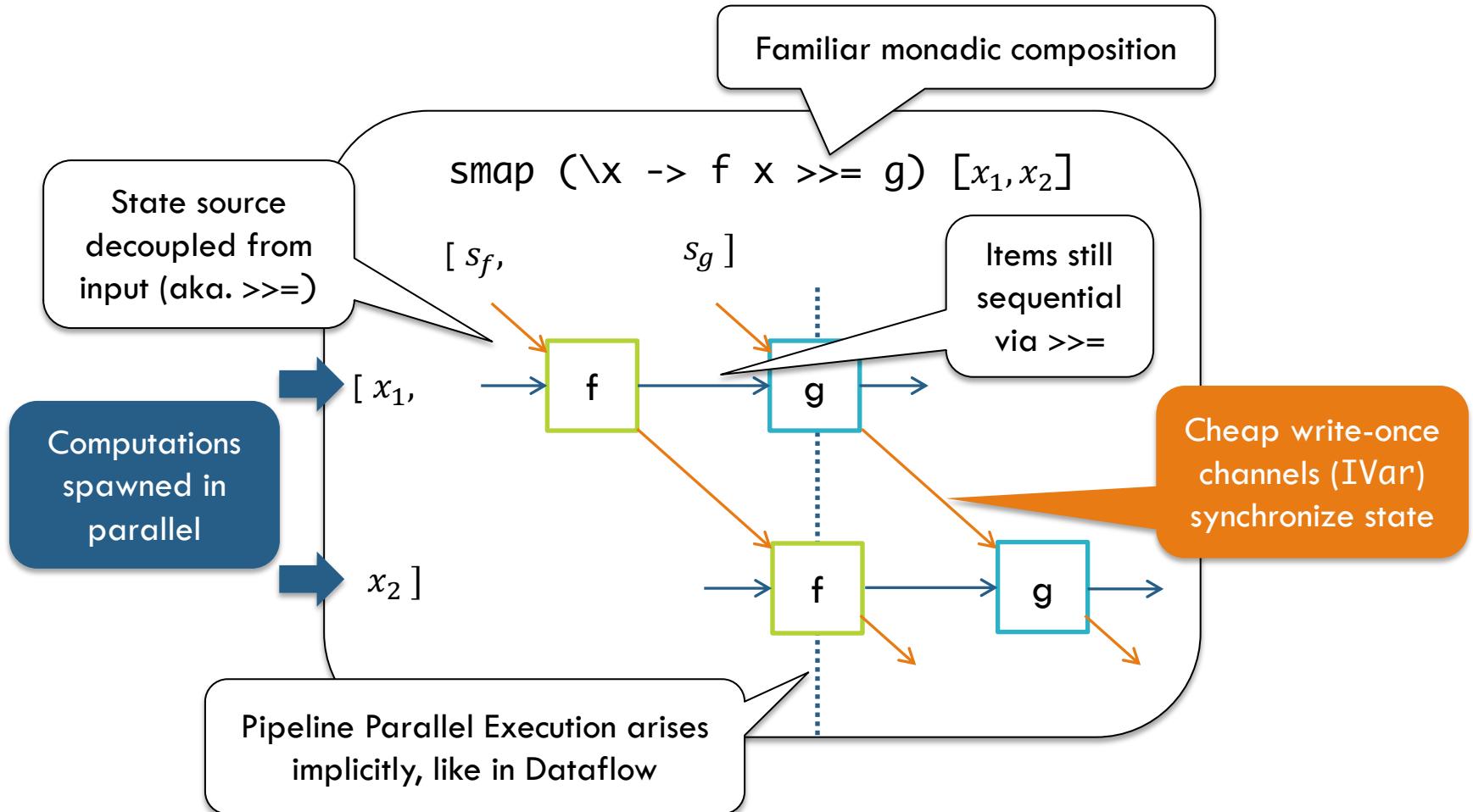
1. Wadler, Philip. "The essence of functional programming." *POPL*. Vol. 92. No. 37. 1992.
2. Launchbury, John, and Simon L. Peyton Jones. "Lazy functional state threads." *ACM SIGPLAN Notices* 29.6 (1994): 24-35.

State in Haskell



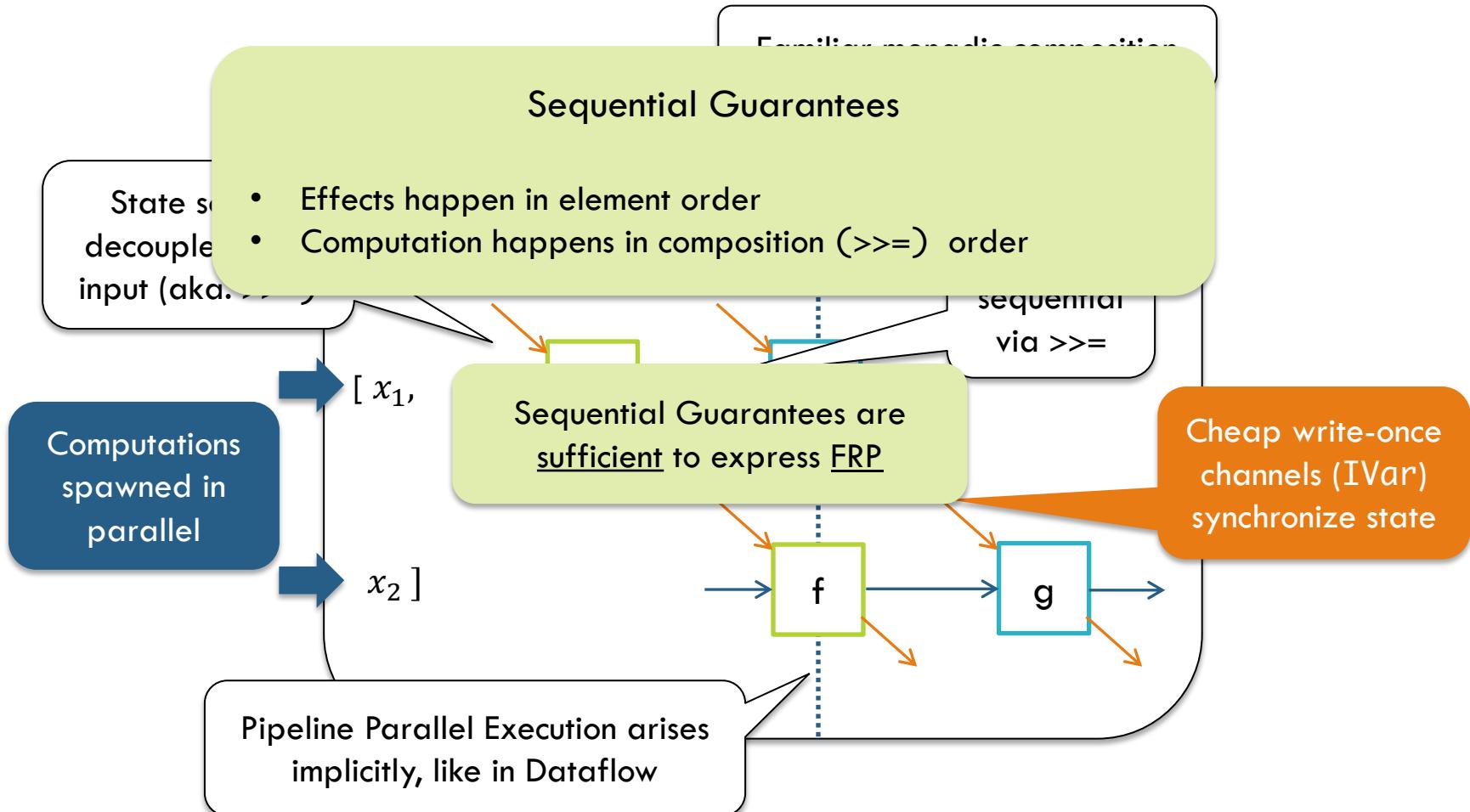
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Smap (De)Construction



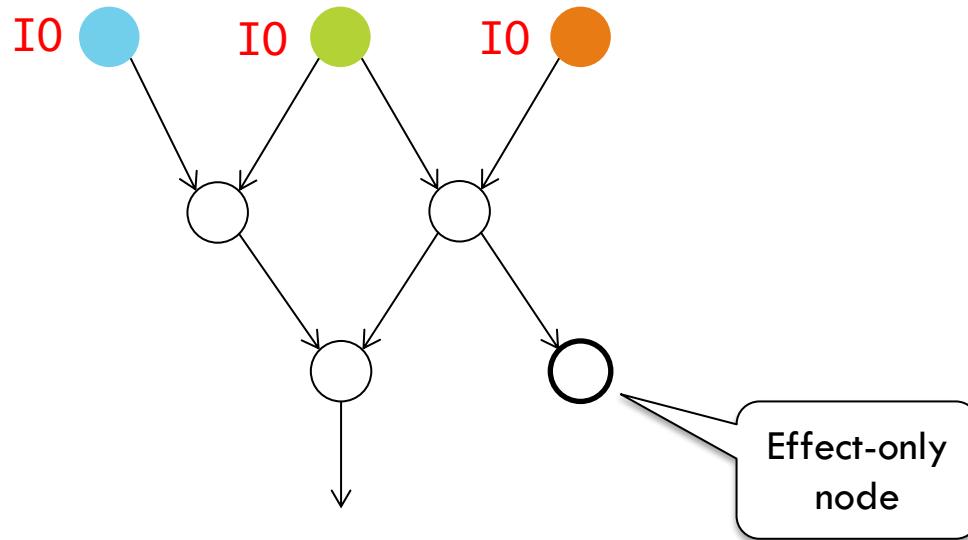
1. Simon Marlow, Ryan Newton, and Simon Peyton Jones. 2011. A monad for deterministic parallelism. In *Proceedings of the 4th ACM symposium on Haskell (Haskell '11)*. ACM, New York, NY, USA, 71-82.

Smap (De)Construction

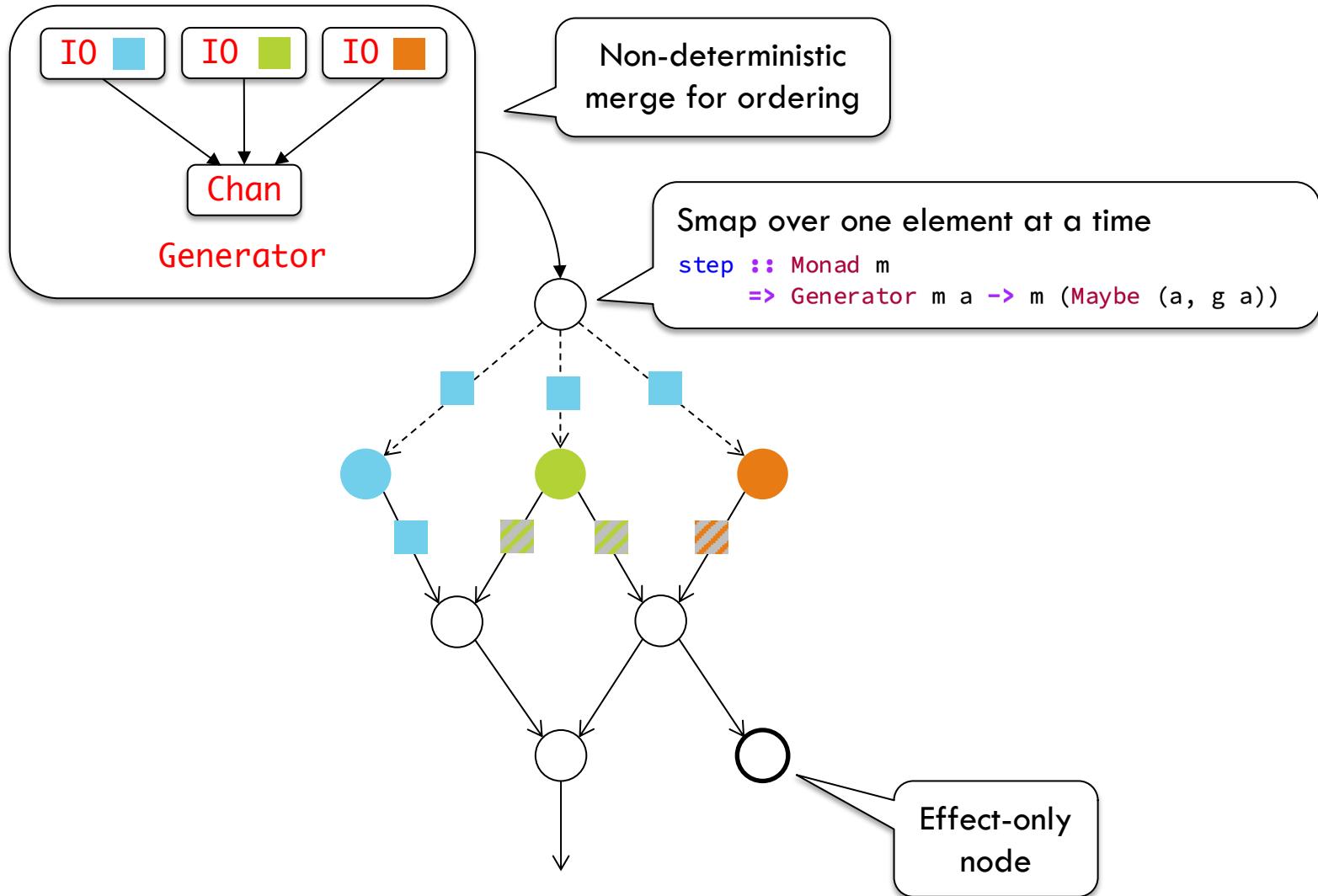


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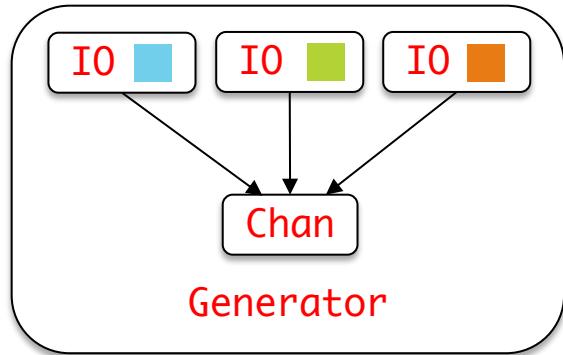
Functional Reactive Programming



Functional Reactive Programming



Functional Reactive Programming



Non-deterministic merge for ordering

Smap over one element at a time

```
step :: Monad m  
=> Generator m a -> m (Maybe (a, g a))
```

All events dispatch to all source nodes

Only matching source updates its value

Other sources repeat old value
⇒ No glitches

Using a filter avoids recompute. Returns last value instead (using state)

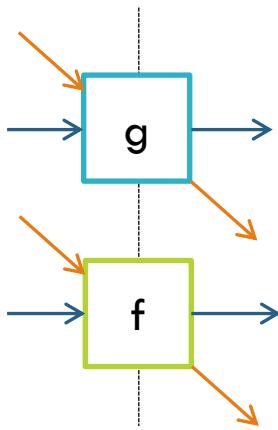
- Trivial (~50 LOC)

Synchronizes via state

Effect-only node

Other Forms for Parallelism

Task Parallelism



Concисely expressed with the `<*>` operator

1. Spawned off

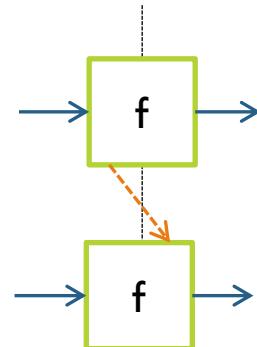
2. Compute this

 $\lambda x \rightarrow f\ x <*> g\ x$

3. Collect results & apply

Can be done automatically with the ApplicativeDo GHC Extension¹

Data Parallelism



Being lazy in a allows s to be sent on first

If s unchanged,
virtually free

Next instance of f can start before this point

`~(a,s) <- run f`

evaluate s
send s

evaluate a

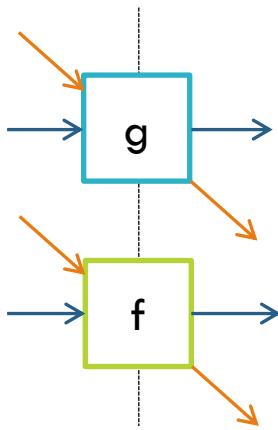
Node Implementation

Works for both read-only and unused state

1. Simon Marlow, Simon Peyton Jones, Edward Kmett, and Andrey Mokhov. 2016. Desugaring Haskell's do-notation into applicative operations. In *Proceedings of the 9th International Symposium on Haskell (Haskell 2016)*. ACM, New York, NY, USA, 92-104.

Other Sources for Parallelism

Task



- We formalized the model using category theory
- Yields a formal explanation for the forms of extracted parallelism

2. Sebastian Ertel, Justus Adam, Norman A. Rink, Andrés Goens, Jeronimo Castrillon. 2019. Category-Theoretic Foundations of "STCLang: State Thread Composition as a Foundation for Monadic Dataflow Parallelism". <https://arxiv.org/abs/1906.12098>

2. Compute this

$$\lambda x \rightarrow f\ x <*> g\ x$$

3. Collect results & apply

Can be done automatically with the ApplicativeDo GHC Extension¹

lism

g lazy in a allows s
be sent on first

If S unchanged,
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Next instance of f
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$\sim(a, s) \leftarrow \text{run } f$

evaluate s
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evaluate a

Node Implementation

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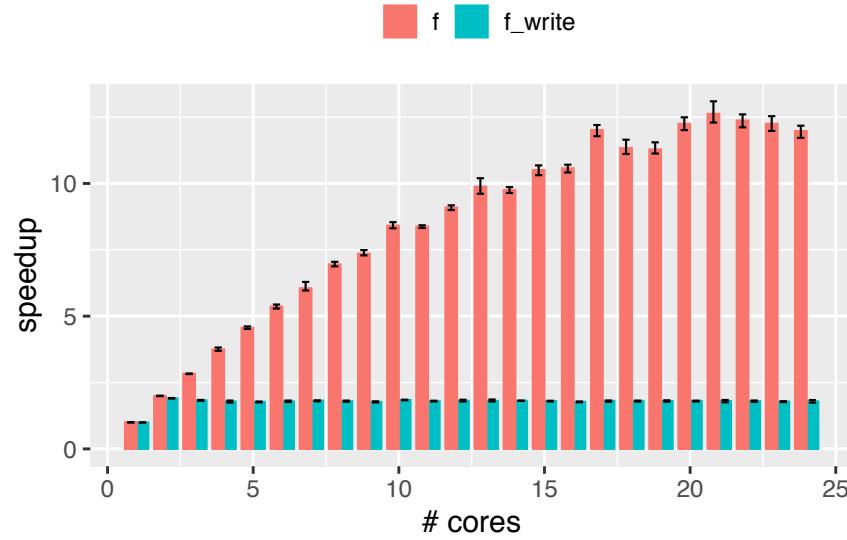
Data Parallelism

Benchmark executes State Thread (`f` or `f_write`) twice, as a two stage pipeline

```
f i = (i +) <$> get
```

```
f_write i = do
    r <- f i
    put r
    pure r
```

Same as `f` but also writes state

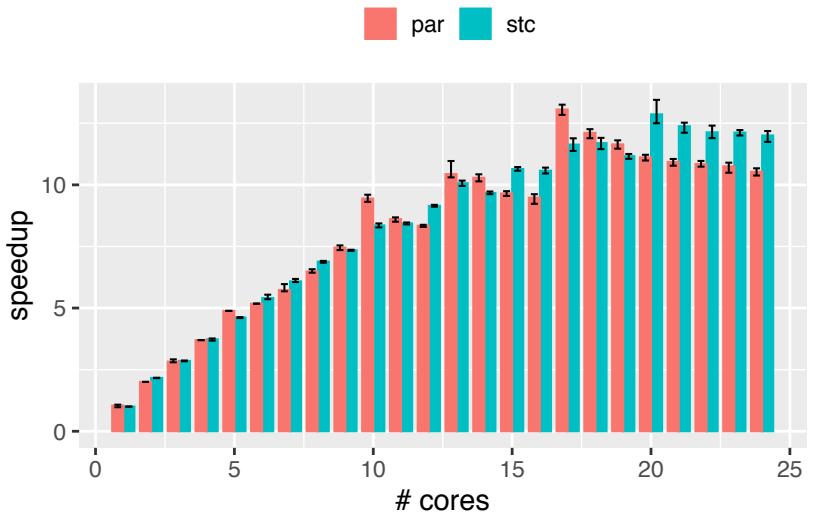


`f_write` exploits pipeline parallelism (~2x)

`f` additionally exploits data parallelism (linear scaling)

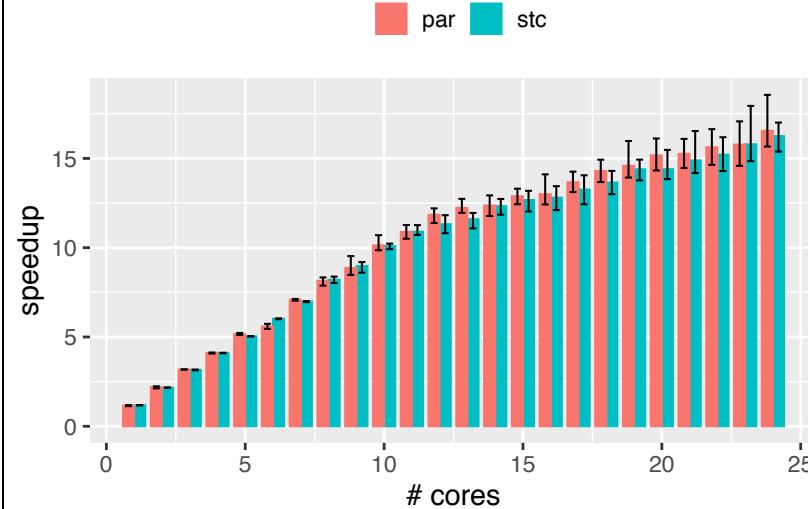
Microbenchmarks

Simple Compositions



stc w = smap (w => w)
par w = parMap (w . w)

Benchmarks from [1]



Black-Scholes²

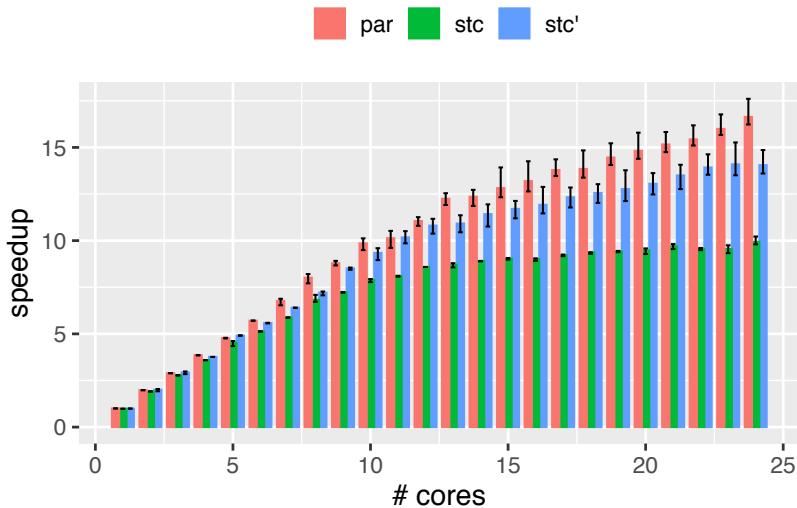
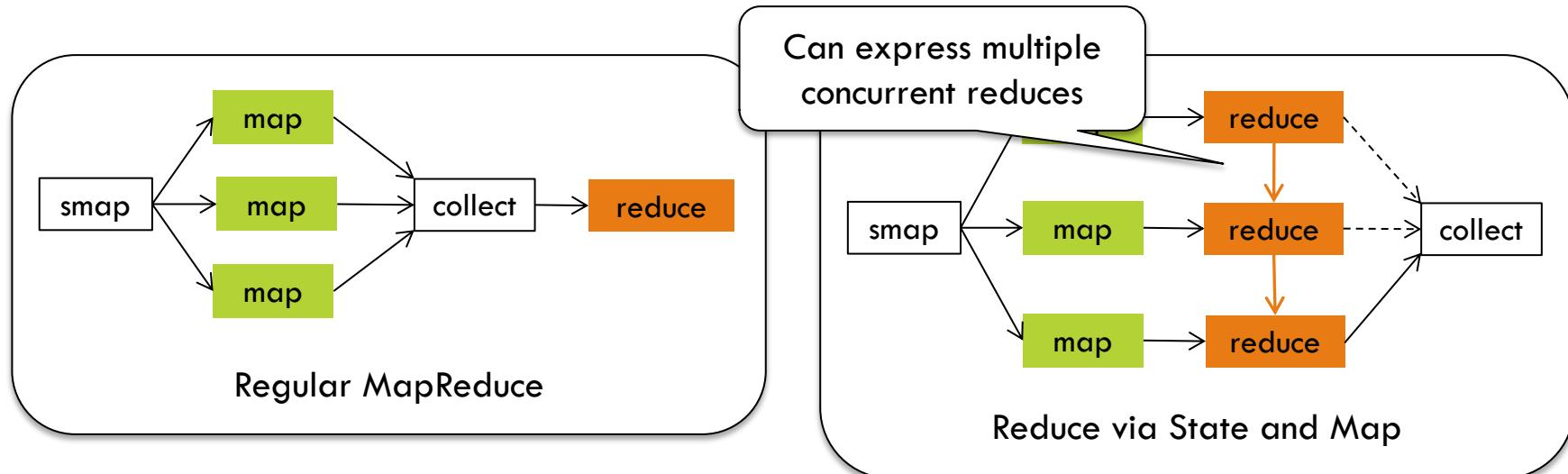
STCLang is implemented
using the monad-par library

Overhead over manual monad-par
implementations is mostly negligible

1. For simplicities sake only includes the more favorable monad-par measurement (`par2` in the paper)

2. Marlow, Simon, Ryan Newton, and Simon Peyton Jones. "A monad for deterministic parallelism." ACM SIGPLAN Notices. Vol. 46. No. 12. ACM, 2011.

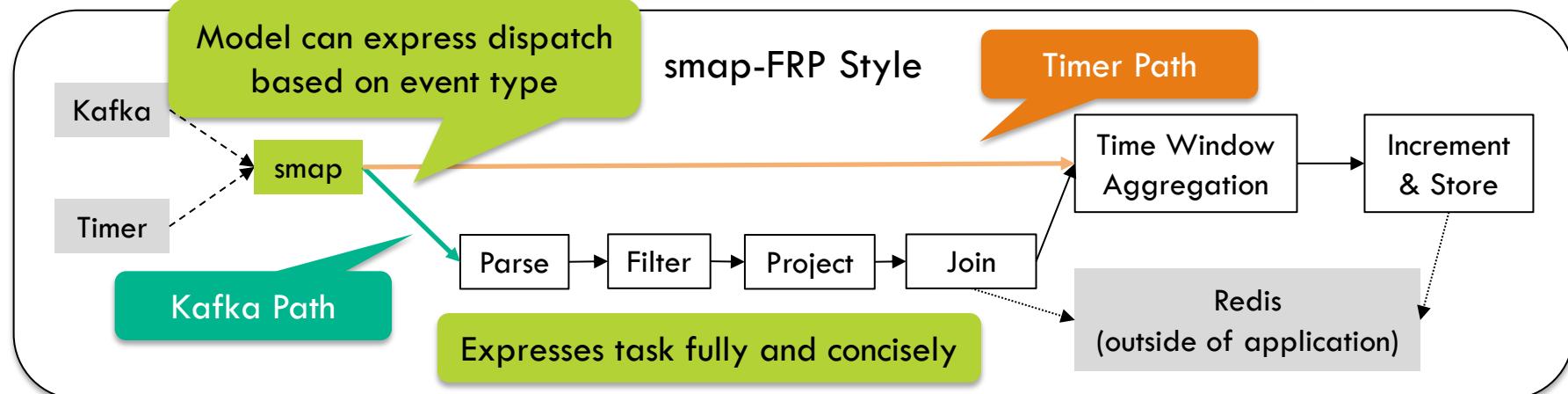
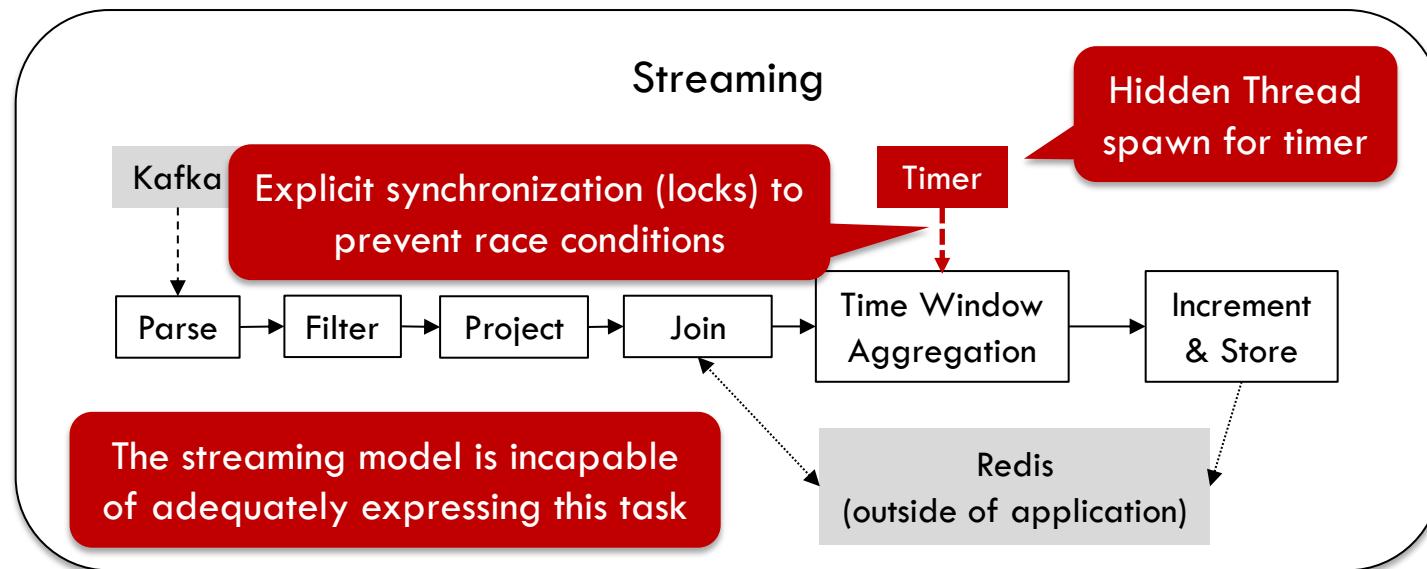
MapReduce Benchmark



sum is very efficient,
(+) is too cheap for
the overhead

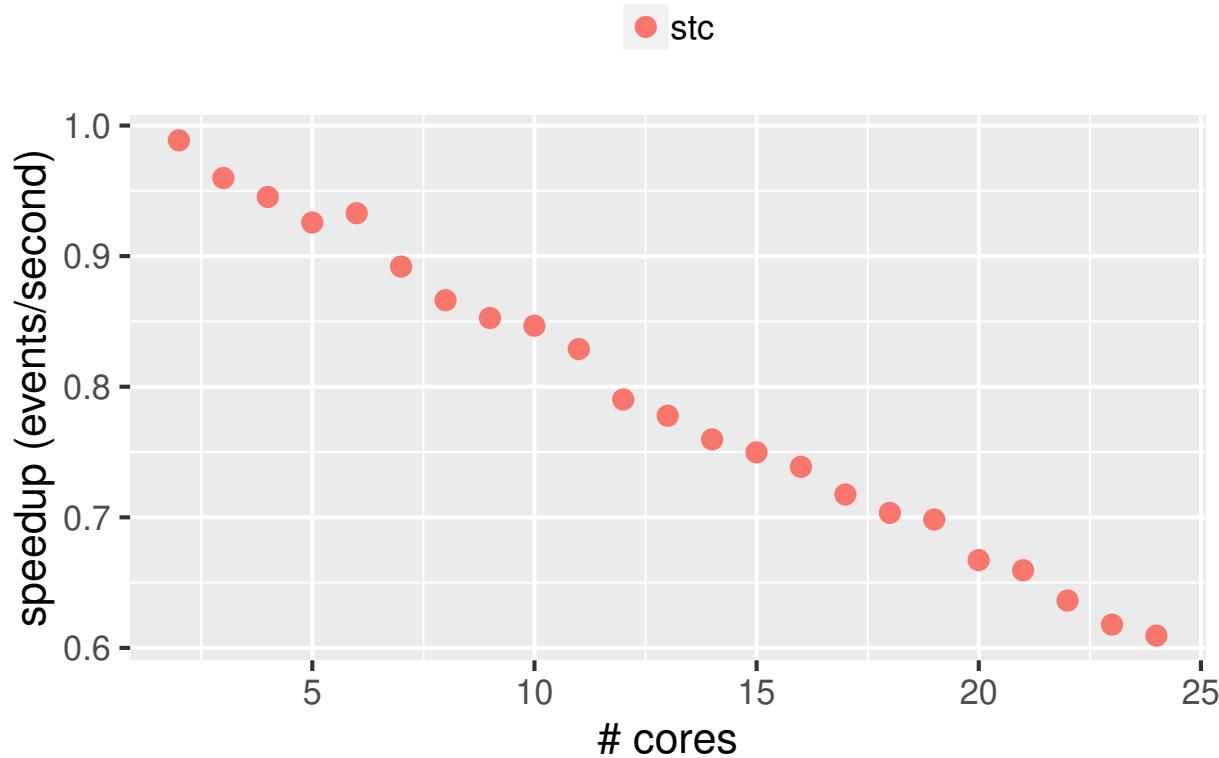
par = sum <\$> parMap euler
stc = mapReduceSTC euler (+) 0
stc' = sum <\$> smap euler

Data Streaming Benchmark



1. Sanket Chintapalli, et al. "Benchmarking streaming computation engines: Storm, flink and spark streaming." 2016 IEEE international parallel and distributed processing symposium workshops (IPDPSW). IEEE, 2016.

Data Streaming Benchmark



Performance degrades
with parallelism

Work stealing scheduler does not have a
notion of output favoured scheduling

1. I-Ting Angelina Lee, Charles E. Leiserson, Tao B. Schardl, Zhunping Zhang, and Jim Sukha. 2015. On-the-Fly Pipeline Parallelism. *ACM Trans. Parallel Comput.* 2, 3, Article 17 (September 2015), 42 pages

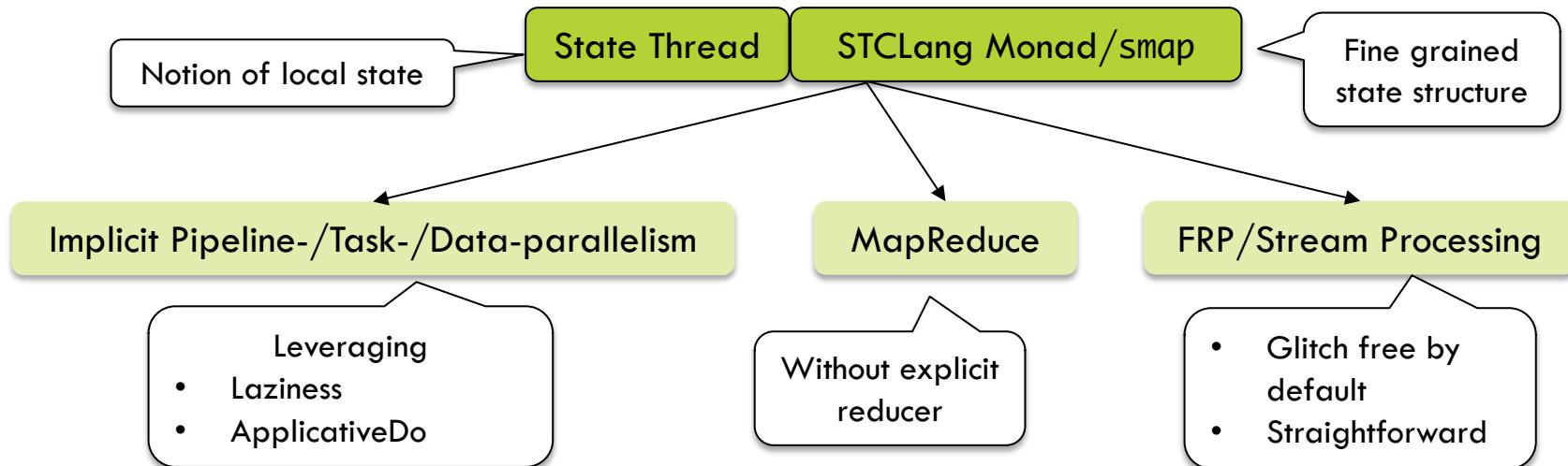
Conclusions

Monad
State ✓ Parallelism ✗ Composition ✓

Dataflow
State ✓ Parallelism ✓ Composition ✗

State Monad
+ *Dataflow Node* =

Monad Composition
+ *Dataflow Execution* =

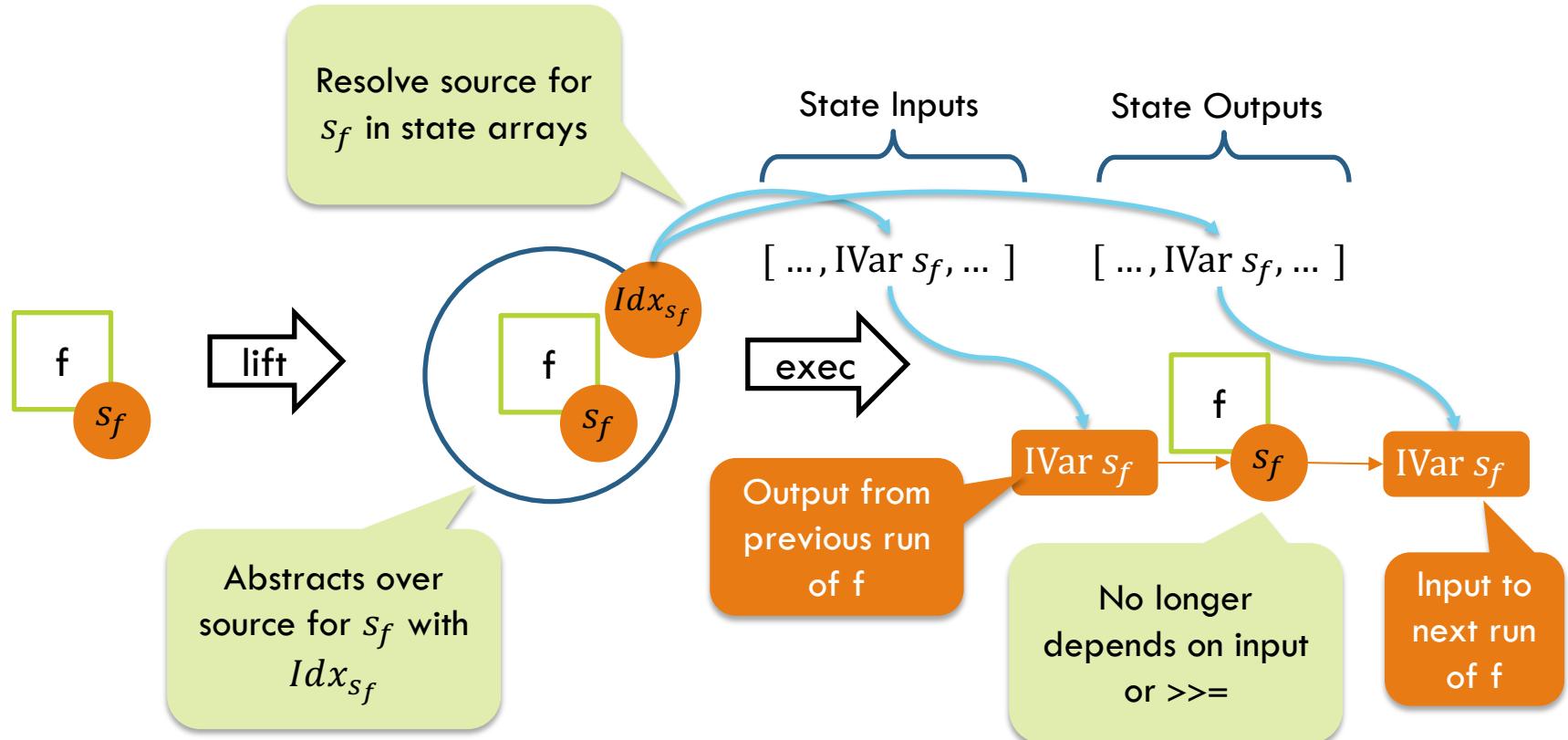


Slides <https://ohua-dev.github.io/slides/haskell-19-stclang.pptx> (.pdf for pdf)

Repo <https://github.com/ohua-dev/stc-lang>

Hackage <https://hackage.haskell.org/package/stc-lang>

The STC Monad



For Sebastian: This is my backup slide to explain our monad construction. I like it, but its also dense, so I relegated it as backup.

The Sequential Monad

$f >>= g$ \Rightarrow

Monad

- Convenient, familiar
- Inherently sequential

$g <*> f$ \Rightarrow

Applicative

- Only task parallelism
- No flow from f to g before the effects of f

$f >>> g$ \Rightarrow

Arrow

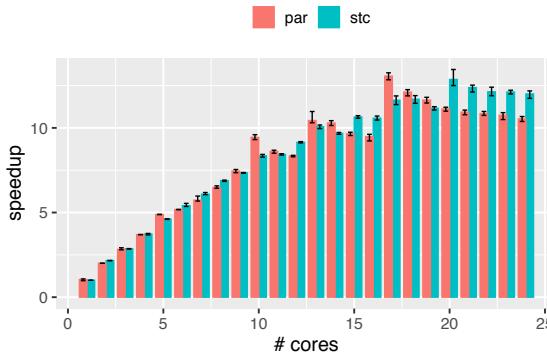
- Effect composition independent from data
- Does anyone use Arrows?
(You should though,
they're cool)

State Thread Examples

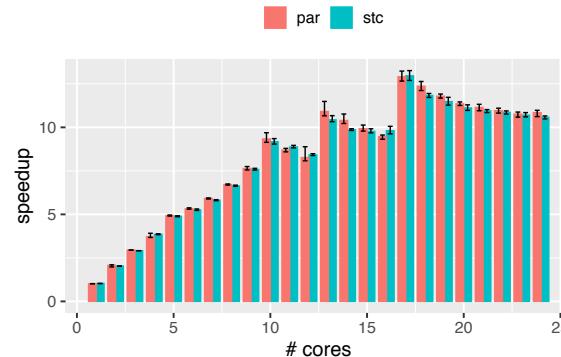
```
memoized :: Ord a
    => (a -> b) -> a
    -> State (LRUChache a b) b
memoized operation elem = do
  cache <- get
  case lookup elem cache of
    Just result -> return result
    Nothing -> do
      let res = operation elem
      put $ insert elem result cache
      return result
```

```
windowedAvg :: Int -> Float
              -> State [Float] Float
windowedAvg wsize i = do
  win <- get
  let win' = take wsize $ i : win
  put win'
  return $ sum win' `div` realToFrac wsize
```

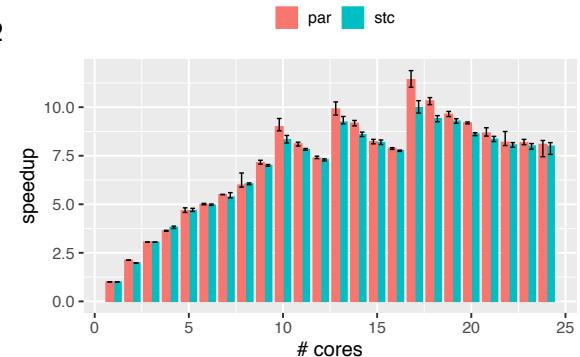
Microbenchmarks



1



2



```
stc w = smap (w >=> w)
par w = parMap (w . w)
```

```
go w = (\x ->
          (,) <$> w x
          <*> w x)
```

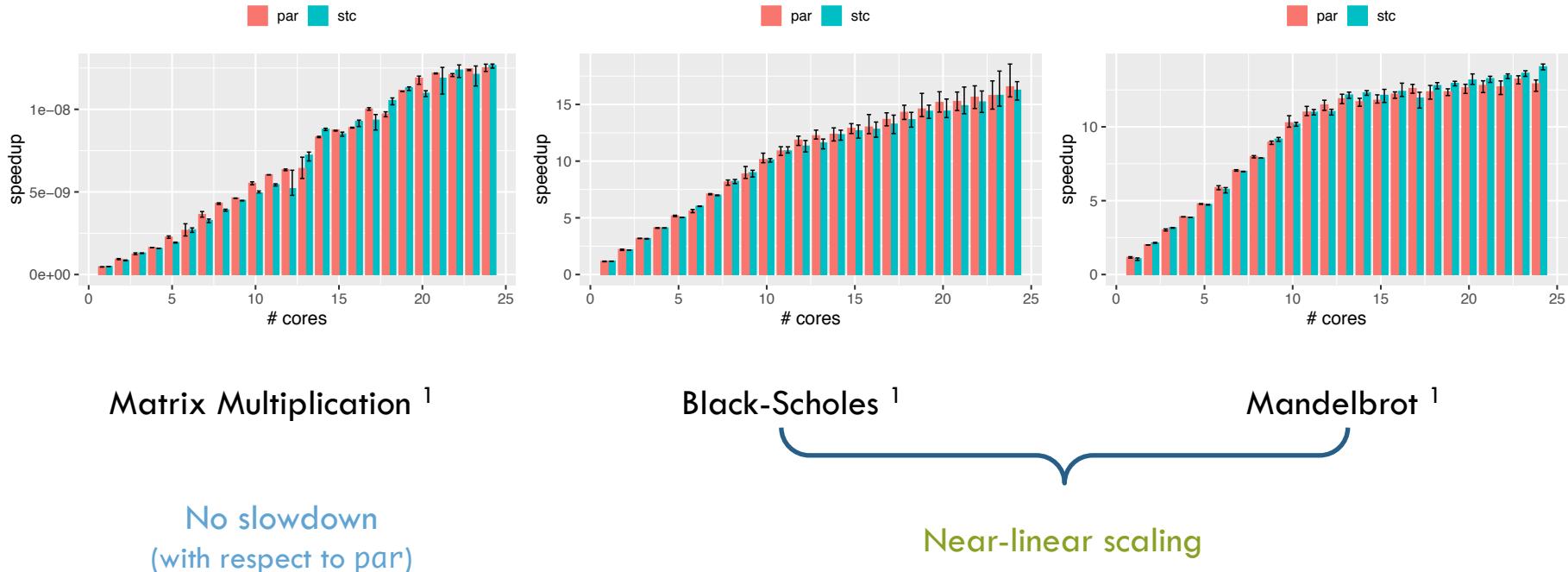
```
stc = smap . go
par = parMapM . go
```

```
go cond w x =
  if cond x
    then w x
    else w x
```

```
stc = smap . go even?
par = parMap . go even?
```

- For simplicities sake only includes the more favorable monad-par measurement (`par2` in the paper)
- NOINLINE version of the benchmark, as it is representative of what we want to test. See the paper for the complete rational

Benchmarks



1. Marlow, Simon, Ryan Newton, and Simon Peyton Jones. "A monad for deterministic parallelism." ACM SIGPLAN Notices. Vol. 46. No. 12. ACM, 2011.