## File Systems (1)

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### File Systems

- ☐ Most visible aspect of an OS
- ☐ Implement an abstraction (files) for secondary storage
  - ☐ A file is <u>a named collection of related information</u> that is recorded on <u>secondary storage</u>.
  - □ Data can NOT be written to secondary storage <u>unless</u> <u>they are within a file</u>.
- ☐ Organize files logically (directories)
- ☐ Permit sharing of data between processes, users, and machines
- ☐ Protect data from unwanted access (security).

#### File Structure

- ☐ A file has a certain defined **structure** which depends on its types:
  - ☐ A <u>text</u> file is a sequence of characters organized into lines.
  - ☐ A <u>source</u> file is a sequence of subroutines and function.
  - ☐ An <u>object</u> file is a sequence of bytes organized into blocks understandable by the system's linker.
  - ☐ An <u>executable</u> file is a series of code sections that the loader can bring into memory and execute.

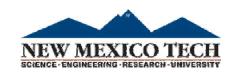


#### File Attributes

- □ Name only information kept in human-readable form
   □ Identifier unique tag (number) identifies file within
  - file system
- □ **Type** needed for systems that support different types
- ☐ **Location** pointer to file location on device
- □ **Size** current file size
- □ **Protection** controls who can do reading, writing, executing
- ☐ **Time, date, and user identification** data for protection, security, and usage monitoring
- ☐ Information about files are kept in the directory structure, which is maintained on the disk

#### File Operations

- ☐ Create
- ☐ Write at write pointer location
- Read at read pointer location
- **☐** Reposition within file seek
- Delete
- Truncate
- □  $Open(F_i)$  search the directory structure on disk for entry  $F_i$ , and move the content of entry to memory
- □  $Close(F_i)$  move the content of entry  $F_i$  in memory to directory structure on disk

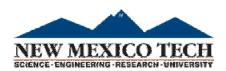


#### **Access Methods**

☐ Sequential access ☐ Read all bytes/records from the beginning ☐ Cannot jump around ☐ May rewind or back up, however ☐ Convenient when medium was magnetic tape ☐ Often useful when whole file is needed ☐ Random access ☐ Bytes (or records) read in any order ☐ Essential for database systems Read can be ... Move file marker (seek), then read or ... Read and then move file marker

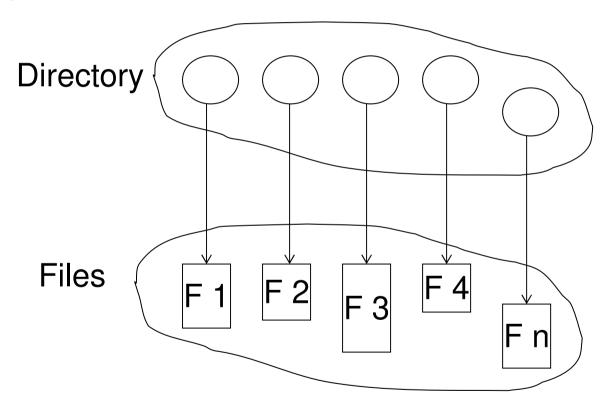
#### **Directories**

- ☐ Naming is nice, but limited
- ☐ Humans like to group things together for convenience
- ☐ File systems allow this to be done with *directories*



#### **Directory Structure**

A collection of nodes containing information about all files



Both the directory structure and the files reside on disk

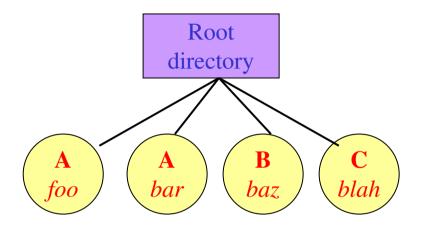


#### **Directory Organization**

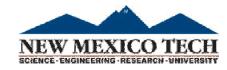
The directory is organized logically to obtain

- ☐ Efficiency locating a file quickly
- Naming convenient to users
  - ☐ Two users can have same name for different files
  - ☐ The same file can have several different names
- ☐ Grouping logical grouping of files by properties, (e.g., all Java programs, all games, ...)

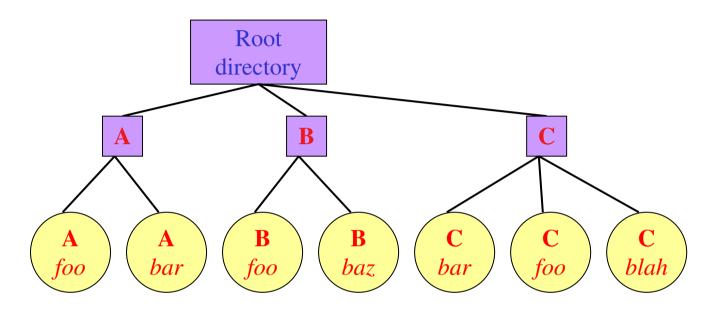
### Single-Level Directory



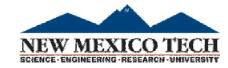
- ☐ One directory in the file system
- ☐ Example directory
  - □ Contains 4 files (foo, bar, baz, blah)
  - □ owned by 3 different people: A, B, and C (owners shown in red)
- ☐ Problem: what if user B wants to create a file called *foo*?
- ☐ Naming problem, grouping problem



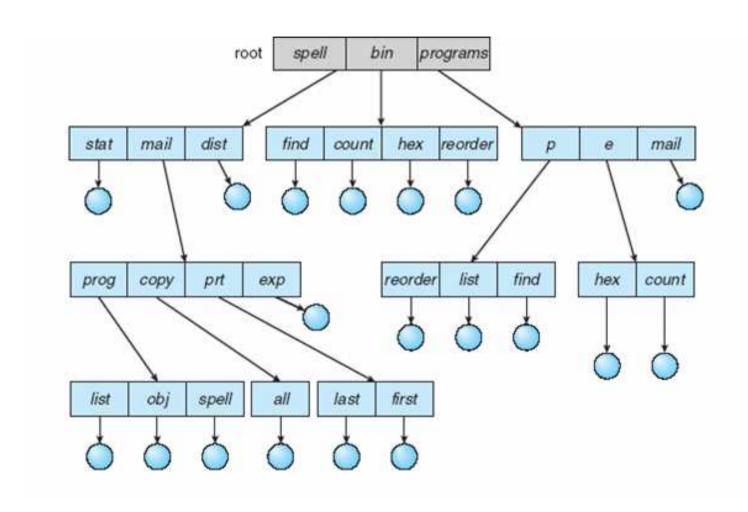
#### **Two-Level Directory**



- □ Solves naming problem: each user has her own directory
- ☐ Multiple users can use the same file name
- ☐ By default, users access files in their own directories
- ☐ Extension: allow users to access files in others' directories
- ☐ No grouping capability



#### **Tree-Structured Directories**





#### **Tree-Structured Directories**

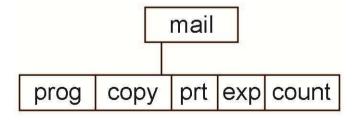
- ☐ **Absolute** or **relative** path name
- ☐ Creating a new file is done in current directory
- ☐ Delete a file

#### rm <file-name>

☐ Creating a new subdirectory is done in current directory

■ Example: if in current directory /mail

#### mkdir count



Deleting "mail"  $\Rightarrow$  deleting the entire subtree rooted by "mail"

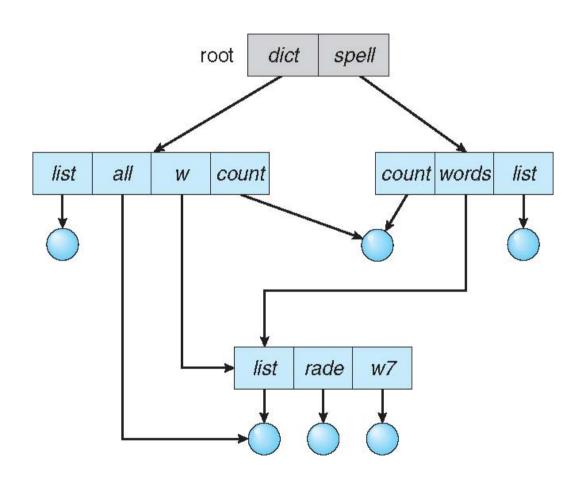
#### **Tree-Structured Directories**

- ☐ Efficient searching
- ☐ Grouping Capability
- ☐ Current directory (working directory)



## **Acyclic-Graph Directories**

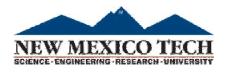
Have shared subdirectories and files





### **Acyclic-Graph Directories**

- ☐ Two different names (aliasing)
- ☐ New directory entry type
  - ☐ Link another name (pointer) to an existing file
  - □ Resolve the link follow pointer to locate the file
- $\Box$  If *dict* deletes *count*  $\Rightarrow$  dangling pointer
  - ☐ How to solve this problem?



# Solutions to Dangling Pointer Problem

- ☐ Backpointers, so we can delete all pointers Variable size records a problem
- ☐ Entry-hold-count solution

