

Memory Management (1)

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CSE325 Principles of Operating
Systems

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Multiprogramming

- ❑ Simple uniprogramming with a single segment per process
- ❑ Uniprogramming disadvantages
 - ❑ Inefficient use of CPU time
 - ❑ Inflexibility of job scheduling
- ❑ Need multiprogramming

Memory Management Requirements

- ❑ The OS must fit multiple processes in memory
 - ❑ memory needs to be subdivided to accommodate multiple processes
 - ❑ memory needs to be allocated to ensure a reasonable supply of ready processes so that the CPU is never idle
 - ❑ memory management is an **optimization** task under **constraints**

Memory Management Wish-list

☐ Sharing

- ☐ multiple processes **coexist** in main memory

☐ Transparency

- ☐ Processes **are not aware** that memory is shared
- ☐ Run **regardless of number/locations** of other processes

☐ Protection

- ☐ **Cannot access** data of OS or other processes

☐ Efficiency: should have reasonable performance

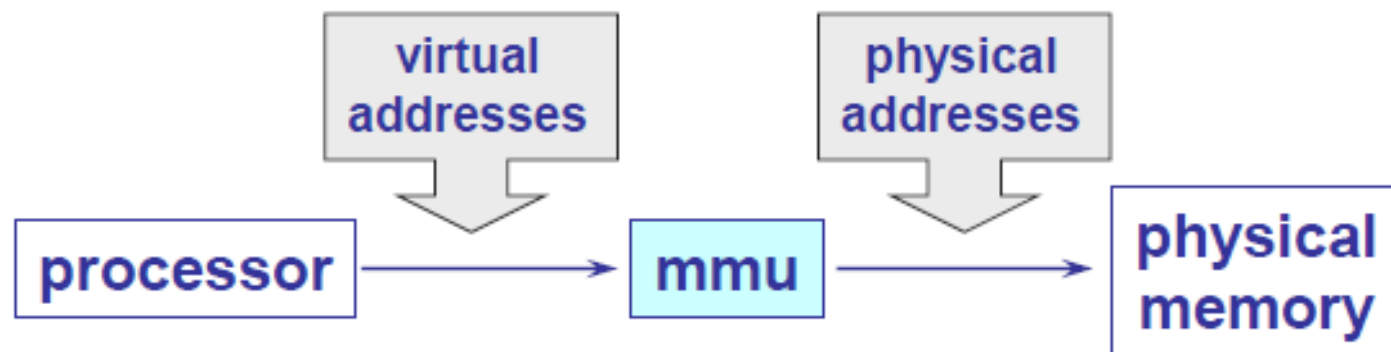
- ☐ Purpose of sharing is to increase efficiency
- ☐ **Do not waste** CPU or memory resources

Virtual Addresses for Multiprogramming

- ❑ To make it easier to manage memory of multiple processes, make processes use virtual addresses (which is not what we mean by “virtual memory” today!)
 - ❑ virtual addresses are independent of location in physical memory (RAM) where referenced data lives
 - ❑ OS determines location in physical memory
 - ❑ instructions issued by CPU reference virtual addresses
 - ❑ e.g., pointers, arguments to load/store instructions, PC ...
 - ❑ virtual addresses are translated by hardware into physical addresses (with some setup from OS)

Logical vs. Physical Address

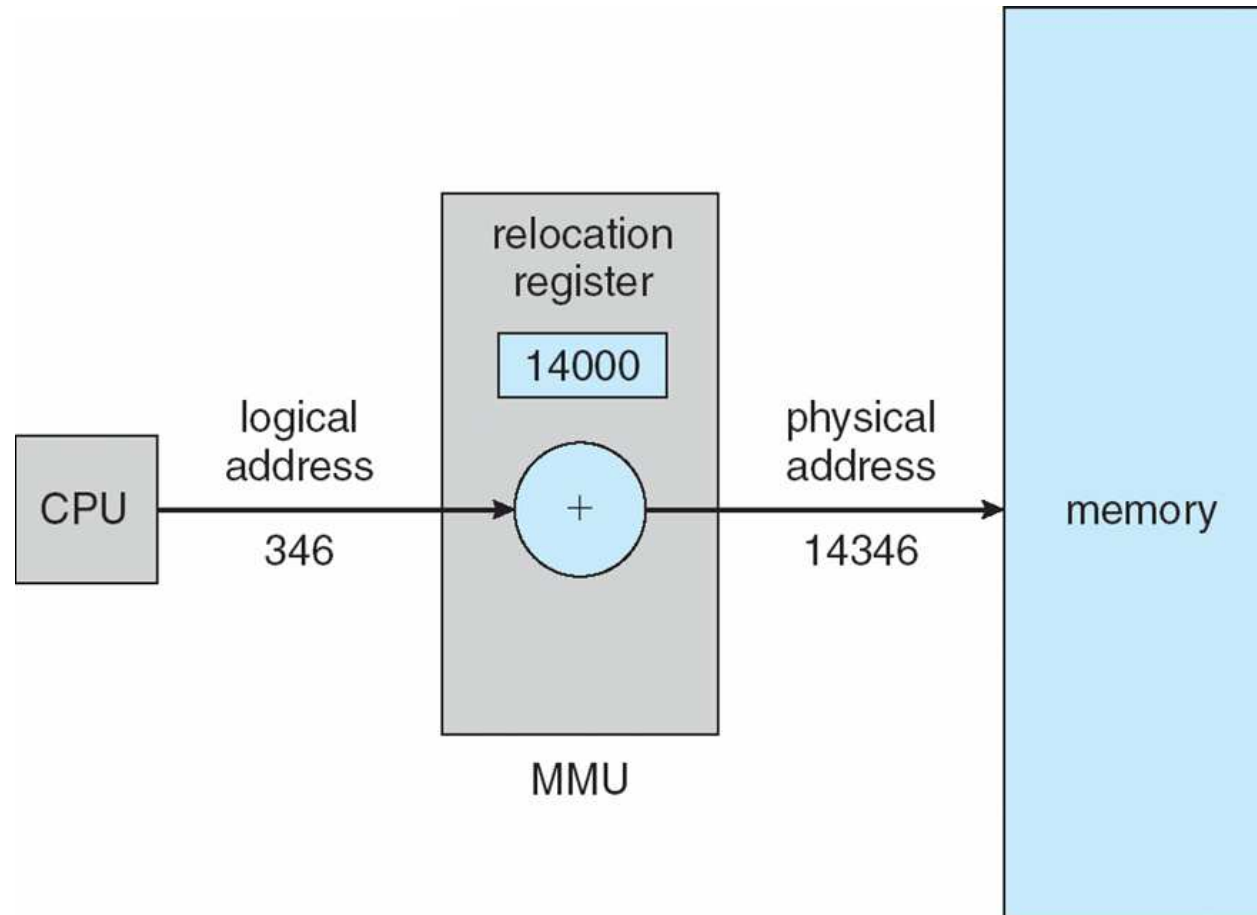
- ❑ **Logical address** – generated by the CPU; also referred to as **virtual address**
- ❑ **Physical address** – address seen by the memory unit



Old Technique #1: Fixed-Sized Partitions

- ❑ One of the simplest methods for allocating memory is to divide memory into several fixed-sized partitions
- ❑ Each partition may contain exactly one process.
- ❑ Degree of multiprogramming limited by number of partitions
- ❑ when a partition is free, a process is selected from the input queue and is loaded into the free partition.
- ❑ When the process terminates, the partition becomes available for another process.

Fixed-Sized Partitions

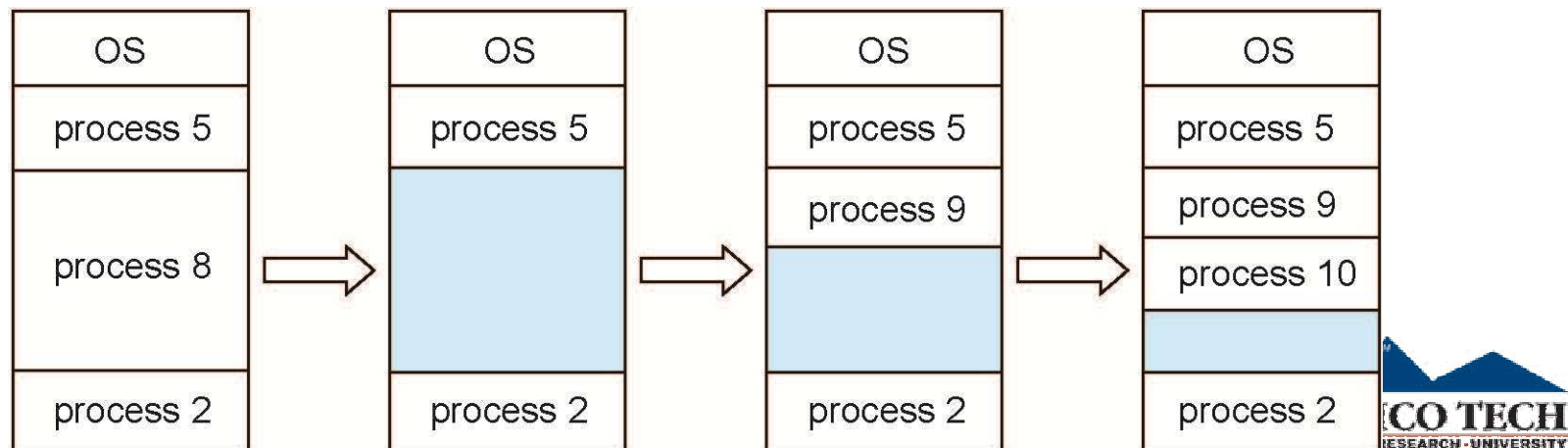


Memory Allocation

- ❑ Main memory must support both OS and user processes
- ❑ Limited resource, must allocate efficiently
- ❑ **Contiguous allocation** is one early method
- ❑ Main memory usually into two **partitions**:
 - ❑ Resident operating system, usually held in low memory with interrupt vector
 - ❑ User processes then held in high memory
 - ❑ Each process contained in single contiguous section of memory

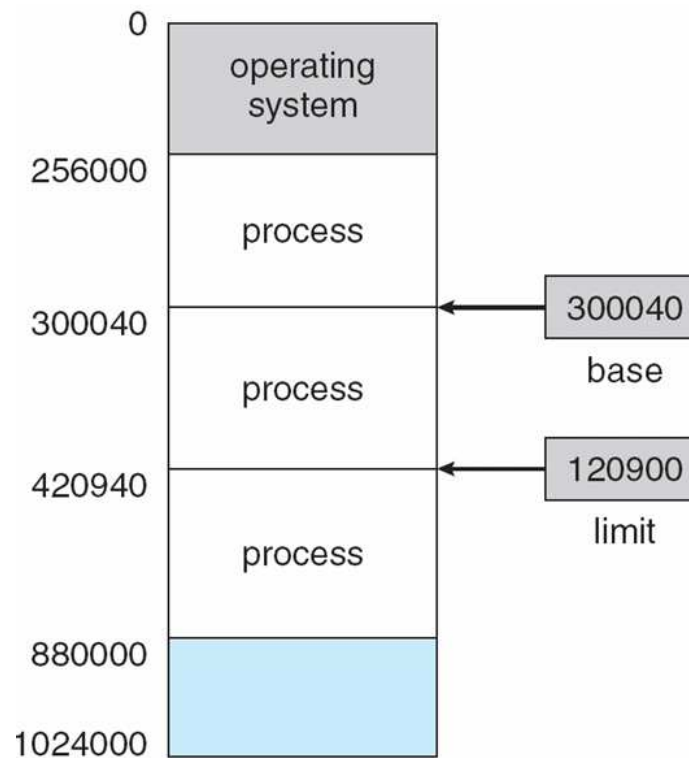
Old Technique #2: Variable Partitions

- ❑ **Variable-partition** sizes for efficiency (sized to a given process' needs)
- ❑ **Hole** – block of available memory; holes of various size are scattered throughout memory
- ❑ When a process arrives, it is allocated memory from a hole large enough to accommodate it
- ❑ Process exiting frees its partition, adjacent free partitions combined
- ❑ Operating system maintains information about:
 - a) allocated partitions b) free partitions (hole)

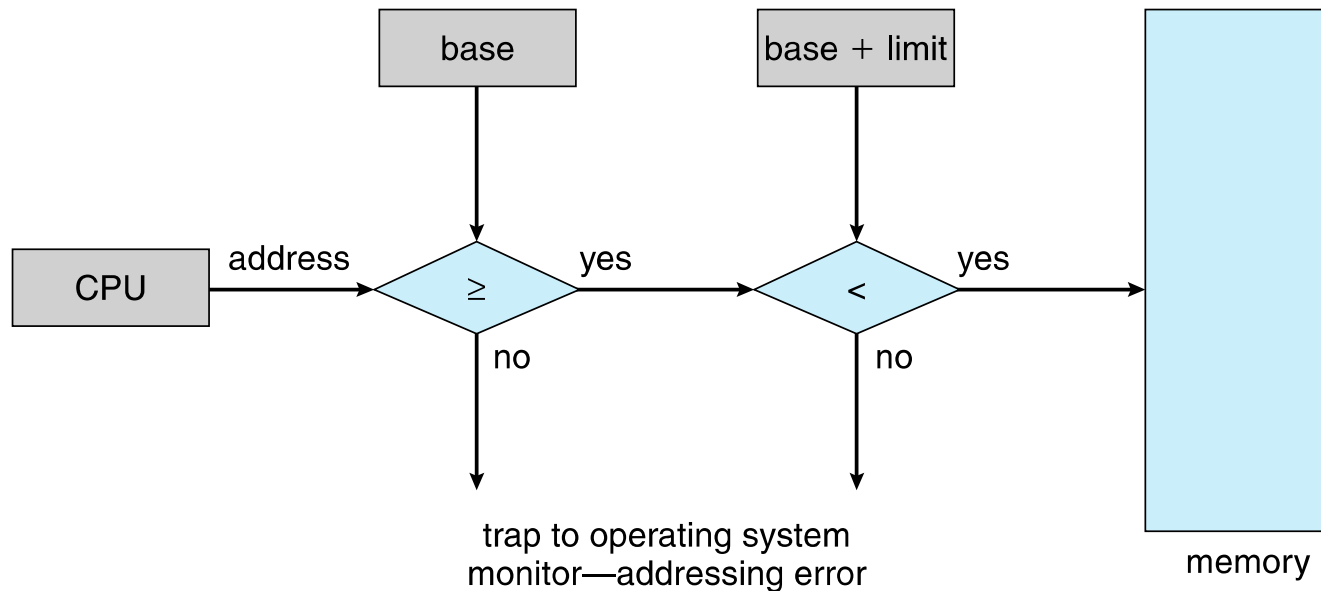


Variable Partitions

- ❑ A pair of registers provide address protection between processes:
 - ❑ **base register:** smallest legal address
 - ❑ **limit register:** size of the legal range



Hardware Address Protection



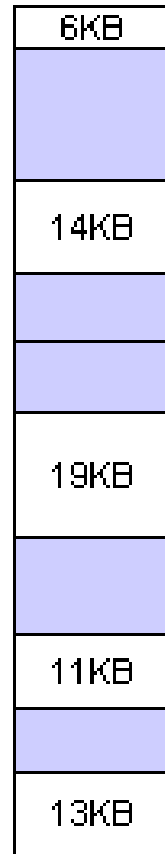
Dynamic Memory Allocation Problem

How to satisfy a request of size n from a list of free holes?

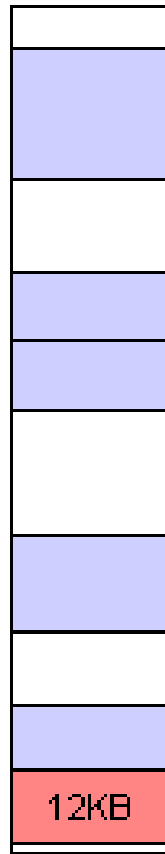
- ❑ **First-fit:** Allocate the *first* hole that is big enough
- ❑ **Best-fit:** Allocate the *smallest* hole that is big enough; must search entire list, unless ordered by size
 - ❑ Produces the smallest leftover hole
- ❑ **Worst-fit:** Allocate the *largest* hole; must also search entire list
 - ❑ Produces the largest leftover hole

First-fit and best-fit better than worst-fit in terms of speed and storage utilization

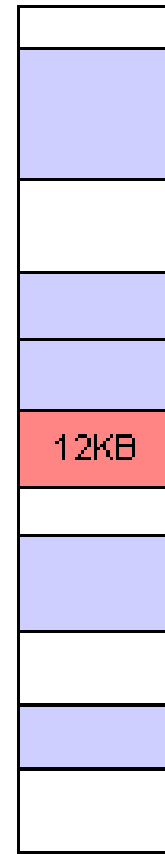
Allocation Example



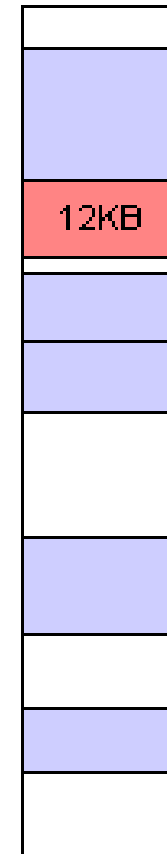
Memory



Best-fit



Worst-fit



First-fit