AVL Trees

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open import Relation.Binary
 open import Relation. Binary. Propositional Equality
 open import Level using (Lift; lift; _⊔_)
 open import Data.Nat as N using (N; suc; zero; pred)
 open import Data.Product
 open import Data. Unit renaming (T to 1)
 open import Data. Maybe
 open import Function
 open import Data.Bool
 open import Data. Empty renaming (1 to 0)
 module AVL
     \{k \ r\} \ (Key : \mathsf{Set} \ k)
     \{\_<\_: Rel \ \mathit{Key} \ r\}
     (isStrictTotalOrder: IsStrictTotalOrder \_ \equiv \_ \_ < \_)
     where
     open IsStrictTotalOrder isStrictTotalOrder
 infix 5 [_]
 data \mathsf{T}: \mathsf{Set}\ k where
     \begin{array}{lll} \inf & 4 & \stackrel{\top}{ } < \\ & \stackrel{\top}{ } < \vdots & \stackrel{\top}{ } \rightarrow \stackrel{\top}{ } \rightarrow \text{Set } r \\ \bot & \stackrel{\top}{ } < \bot & = \text{Lift } r \ 0 \\ \bot & \stackrel{\top}{ } < \top & = \text{Lift } r \ 1 \\ \bot & \stackrel{\top}{ } < [ \ \_ \ ] = \text{Lift } r \ 1 \\ \top & \stackrel{\top}{ } < [ \ \_ \ ] & = \text{Lift } r \ 0 \\ [ \ \_ \ ] & \stackrel{\top}{ } < \bot & = \text{Lift } r \ 0 \\ [ \ \_ \ ] & \stackrel{\top}{ } < \top & = \text{Lift } r \ 1 \\ [ \ x \ ] & \stackrel{\top}{ } < [ \ y \ ] & = x < y \\ \end{array}
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 $\mathsf{rot}\mathsf{-r}:\ \forall\ \{\mathit{lb}\ \mathit{ub}\ \mathit{rh}\ \mathit{v}\}\ \{\mathit{V}:\mathit{Key}\to\mathsf{Set}\ \mathit{v}\}$

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\rightarrow (k: Key)
       \rightarrow V k
       \rightarrow Tree V lb [k] (suc (suc <math>rh))
       \rightarrow Tree V [k] ub rh
       \rightarrow Inserted V lb ub (suc (suc rh))
rot-r u uc (node v vc 	imes ta tb) tc = same (node v vc 	imes ta (node u uc 	imes tb tc))
rot-r u uc (node v vc - ta tb) tc = chng (node v vc > ta (node u uc < tb tc))
rot-r u uc (node v vc > ta (node w wc bw tb tc)) td =
   same (node w \ wc - (node \ v \ vc \ ( \searrow \implies \angle bw) \ ta \ tb) \ (node \ u \ uc \ ( \swarrow \implies \searrow bw) \ tc \ td))
rot-I: \forall \{lb \ ub \ lh \ v\} \{V : Key \rightarrow \mathsf{Set} \ v\}
       \rightarrow (k: Key)
       \rightarrow V k
       \rightarrow Tree V lb [k] lh
       \rightarrow Tree V [k] ub (suc (suc <math>lh))
       \rightarrow Inserted V lb ub (suc (suc lh))
rot-l u\ uc\ tc\ (\mathsf{node}\ v\ vc\ \cdot \ tb\ ta) = \mathsf{same}\ (\mathsf{node}\ v\ vc\ \cdot \ (\mathsf{node}\ u\ uc\ \cdot \ tc\ tb)\ ta)
rot-l u uc tc (node v vc - tb ta) = chng (node v vc \checkmark (node u uc \gt tc tb) ta)
rot-l u uc td (node v vc \checkmark (node w wc bw tc tb) ta) =
   same (node w \ wc - (\text{node } u \ uc \ (\searrow \Longrightarrow \swarrow bw) \ td \ tc) \ (\text{node } v \ vc \ (\swarrow \Longrightarrow \searrow bw) \ tb \ ta))
     insert : \forall \{l \ u \ h \ v\} \{V : Key \rightarrow Set \ v\} (k : Key)
              \rightarrow (V k \rightarrow V k \rightarrow V k)
              \rightarrow Tree V l u h
              \rightarrow l < k < u
              \rightarrow Inserted V l u h
     insert v \ vc \ f (leaf l < u) (l, u) = chng (node v \ vc \ne (leaf l) (leaf u))
     insert v \ vc \ f (node k \ kc \ bl \ tl \ tr) prf with compare v \ k
     insert v \ vc \ f \ (node \ k \ kc \ bl \ tl \ tr) \ (l \ , \ )
     ... | chng tl' with bl
     ... | = \text{rot-r } k \ kc \ tl' \ tr
              | \cdot | = \text{chng (node } k \, kc \, \angle \, t \, l' \, tr)
              | \times | = \text{same (node } k \ kc + t l' \ tr)
     insert v \ vc \ f (node k \ kc \ bl \ tl \ tr)
       \mid tri\approx _ refl _ = same (node k (f vc kc) bl tl tr)
     insert v \ vc \ f (node k \ kc \ bl \ tl \ tr) (_ , u)
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| \text{ tri>} \_\_ c \text{ with insert } v \ vc \ f \ tr \ (c \ , \ u) \\ ... \ | \text{ same } tr' = \text{ same (node } k \ kc \ bl \ tl \ tr')
                 ... | chng tr' with bl
                  ... | = same (node \ kc - tl \ tr')
                                               |\cdot| = \mathsf{chng} \; (\mathsf{node} \; k \; kc \; \cdot \; tl \; tr')
                                          | \cdot | = \text{rot-l } k \ kc \ tl \ tr'
       lookup: (k: Key)
                                             \rightarrow \forall \{l \ u \ s \ v\} \{V : Key \rightarrow \mathsf{Set} \ v\}
                                             \rightarrow Tree V \ l \ u \ s
                                             \rightarrow Maybe (V k)
        lookup k (leaf l < u) = nothing
        lookup k (node v\ vc\ \_\ tl\ tr) with compare k\ v
       \dots \mid \mathsf{tri} < \_ \_ \_ = \overline{\mathsf{lookup}} \; k \; tl
        ... | tripprox _ refl _ = just vc
         \dots \mid \mathsf{tri} > \underline{\phantom{a}} = \mathsf{lookup} \ k \ tr
module DependantMap where
        data Map \{v\} (V: Key \rightarrow Set \ v): Set (k \sqcup v \sqcup r) where
                 tree : \forall \{h\} \rightarrow \mathsf{Bounded}.\mathsf{Tree}\ V \perp \top\ h \rightarrow \mathsf{Map}\ V
       \mathsf{insertWith}: \ \forall \ \{v\} \ \{\mathit{V}: \mathit{Key} \to \mathsf{Set} \ \mathit{v}\} \ (\mathit{k}: \mathit{Key})
                                                              \rightarrow V k
                                                              \rightarrow (V k \rightarrow V k \rightarrow V k)
                                                              \rightarrow Map V
                                                               \rightarrow Map V
       insertWith k \ v \ f (tree tr) =
                  tree (proj<sub>2</sub> (Bounded.insert k \ v \ f \ tr (lift tt , lift tt)))
        insert : \forall \{v\} \{V : Key \rightarrow Set \ v\} (k : Key) \rightarrow V \ k \rightarrow Map \ V \rightarrow Map \ V
        insert k v = \text{insertWith } k v \text{ const}
       \mathsf{lookup}: (k: Key) \to \forall \{v\} \{V: Key \to \mathsf{Set}\ v\} \to \mathsf{Map}\ V \to \mathsf{Maybe}\ (V\ k)
        lookup \ k \ (tree \ tr) = Bounded.lookup \ k \ tr
module Map where
         data Map \{v\} (V : \mathsf{Set}\ v) : \mathsf{Set}\ (k \sqcup v \sqcup r) where
                  tree : \forall \{h\} \rightarrow \mathsf{Bounded}.\mathsf{Tree} \; (\mathsf{const} \; V) \perp \top \; h \rightarrow \mathsf{Map} \; V
        \mathsf{insertWith}: \ \forall \ \{v\} \ \{\mathit{V}: \mathsf{Set} \ \mathit{v}\} \ (\mathit{k}: \ \mathit{Key}) \rightarrow \ \mathit{V} \rightarrow \ (\mathit{V} \rightarrow \ \mathit{V} \rightarrow \ \mathit{V}) \rightarrow \ \mathsf{Map} \ \ \mathit{V} \rightarrow \ \mathsf{Map} \ \ \ \mathit{V} \rightarrow \ \mathsf{Map} \ \ \mathsf{V} \rightarrow \ \mathsf{V} \rightarrow
        insertWith k \ v \ f (tree tr) =
                  tree (proj_2 (Bounded.insert k \ v \ f \ tr (lift tt , lift tt)))
       insert : \forall \{v\} \{V : \mathsf{Set}\ v\} (k : Key) \to V \to \mathsf{Map}\ V \to \mathsf{Map}\ V
       insert k v = insertWith k v const
       \mathsf{lookup} : (k : Key) \to \forall \{v\} \{V : \mathsf{Set}\ v\} \to \mathsf{Map}\ V \to \mathsf{Maybe}\ V
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lookup k (tree tr) = Bounded.lookup k tr

module Sets where
data \langle \mathsf{Set} \rangle: Set (k \sqcup r) where
tree: \forall \{h\} \rightarrow \mathsf{Bounded.Tree} \ (\mathsf{const} \ 1) \perp \top h \rightarrow \langle \mathsf{Set} \rangle
insert: Key \rightarrow \langle \mathsf{Set} \rangle \rightarrow \langle \mathsf{Set} \rangle
insert k (tree tr) =
tree (proj<sub>2</sub> (Bounded.insert k tt const tr (lift tt , lift tt)))

member: Key \rightarrow \langle \mathsf{Set} \rangle \rightarrow \mathsf{Bool}
member k (tree tr) = is-just (Bounded.lookup k tr)
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