

1. (20%) Let  $A, B$  be two arrays of length  $n$  where  $A[1] < A[2] < \dots < A[n]$  and  $B[1] > B[2] > \dots > B[n]$ .
  - (a) Give an  $O(1)$ -time algorithm that can find the index  $k$  so that  $A[k] - B[k]$  is minimized.
  - (b) Prove that any algorithm that can find the index  $k$  so that  $A[k] + B[k]$  is minimized requires  $\Omega(n)$  time.
2. (30%) Give asymptotic upper bounds for the following  $T(n)$ , and justify your answers. You may assume that  $T(1) = 1$ .
  - (a)  $T(n) = T(\lceil n/3 \rceil) + T(\lceil n/4 \rceil) + O(n)$ .
  - (b)  $T(n) = T(\lceil n/3 \rceil + 5) + T(\lceil n/4 \rceil + 7) + O(n)$ .
  - (c)  $T(n) = T(\lfloor n/2 \rfloor) + T(\lceil n/2 \rceil) + O(\log n)$ .
3. (10%) Let  $Q$  be the convex polygon illustrated in Figure 1. Answer the following questions by calculating cross products as we did in the lecture.

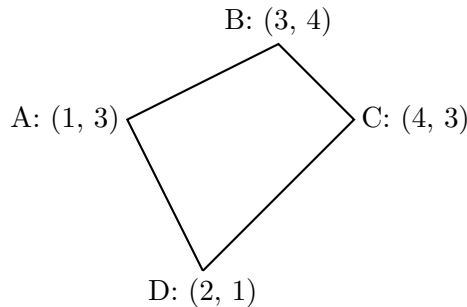


Figure 1: The convex polygon  $Q$ .

- (a) Let  $p = (4, 2)$ . Is  $p \in Q$ ? By  $p \in Q$ , we mean that  $p$  is on the boundary of  $Q$  or in its interior. If not, what is the convex hull of  $Q \cup \{p\}$ ?
- (b) Let  $p = (3, 2)$ . Is  $p \in Q$ ? If not, what is the convex hull of  $Q \cup \{p\}$ ?

4. (15%) Given an  $n$  by  $n$  matrix  $A \in \mathbb{R}^{n \times n}$ . Finding the *monotonic*<sup>1</sup> path from  $A[1][1]$  to  $A[n][n]$  so that the sum of values on the cells visited by the path is minimized, i.e. the shortest monotonic path.

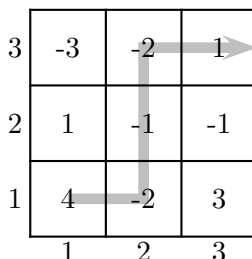


Figure 2: An illustration of the monotonic path.

- (a) Devise a *polynomial-time* algorithm that can output the length of the shortest monotonic path. By polynomial-time algorithms, we mean those algorithms whose running time is  $O(n^c)$  for some constant  $c \geq 0$ . Give the pseudocode of your algorithm.
  - (b) Devise a polynomial-time algorithm that can output the cells visited by the shortest monotonic path.
  - (c) Explain why your algorithms are correct and analyze their running time.
5. (10%) Give an array  $A$  of  $n$  real numbers. Devise an  $O(n \log n)$ -time algorithm that can output the longest *bitonic* subsequence  $S$  of  $A$ . We say a subsequence  $S$  is bitonic if there exists an index  $k$  so that  $S[i] < S[j]$  for every  $i < j, j \leq k$  and  $S[i] > S[j]$  for every  $i < j, i \geq k$ .
- (a) Give the pseudocode of your algorithm.
  - (b) Explain why your algorithm is correct and analyze its running time.
6. (15%) Give a bag and  $n$  stones where the  $i$ -th stone has weight  $w_i$  and value  $v_i$ . We would like to place some of the  $n$  stones into the bag so that the total value of the selected stones is maximized and the total weight of the selected stones does not exceed  $m$ , a given parameter.

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<sup>1</sup>By monotonic path, we mean a path that goes only upward and rightward.

- (a) Devise an  $O(mn)$ -time algorithm that can answer whether there exists a subset of the  $n$  stones whose total weight is  $k$ , for any  $k \leq m$ . Give the pseudocode of your algorithm.
- (b) Devise an  $O(mn)$ -time algorithm that can calculate the maximum value of the stones that you can pack into the bag. Give the pseudocode of your algorithm.
- (c) Explain why your algorithms are correct.