

## Big Ideas

- **Induction** — be familiar with establishing *base case* and *inductive step*.
  - Base Case:  $P(0)$  is true.
  - Inductive Step: Assume  $P(n)$  for a specific  $n$  and show  $P(n) \rightarrow P(n+1)$ .
- **Strong Induction** — It's like normal induction, except we potentially assume multiple base cases and need to assume  $P(0)$  through  $P(n)$  are true to show that  $P(n+1)$  is true.
- **Well-Ordering Principle** — If  $S$  is a set of one or more integers all greater than some fixed integers, then  $S$  has a least element.
- **Induction Application: Algorithms** — You can show an algorithm is “correct” by applying induction to variables that go through iterations of a loop. This involves use of a *loop invariant*, a predicate that is true before the loop and remains true after passing through the loop, and the *guard*, which is the name given to the predicate condition that keeps statements in the loop. We denote  $I(n)$  to be the  $n$ th iteration of the loop invariant through the loop.
  - Basis Step — Identical to base case. Check that the pre-condition is true before the loop, e.e. pre-condition implies  $I(0)$  is true.
  - Induction Step — Show that if loop invariant  $I(n)$  and guard  $G$  are true before the loop, then  $I(n+1)$  is true after the loop.
  - Eventual Falsity of the Guard — Show that eventually  $G$  will become false (we don't want the loop to run infinitely!).
  - Correctness of Post-Condition — Check that the post-condition is true after the loop, i.e. if  $N$  is the smallest value for which  $I(N)$  is true and the guard  $G$  becomes false, then the post-condition is also true.

I'm aware this looks like a lot, but realistically it is identical to induction with the added steps of checking the loop doesn't run infinitely and that the end result isn't contradictory.

- **Recurrence Relations** — A recurrence relation is defined as a formula for a sequence that defines elements in terms of previous elements. For example:

$$a_n = a_{n-1} + 3a_{n-2} - 7a_{n-3}$$

relates values in the sequence to the previous three values.

- **Solving Second-Order Recurrence Relations** — For a recurrence relation of the form  $a_n = Aa_{n-1} + Ba_{n-2}$ , you can find sequences that satisfy the relation by solving the characteristic equation

$$t^2 - At - B = 0$$

The solution sequence would then be  $a_n = t^n$  for  $n \geq 0$ .

For second-order recurrence relations with initial conditions  $a_0$  and  $a_1$ , you can find an explicit formula by solving for constants  $C$  and  $D$  as follows:

- If  $r$  and  $s$  are *unique* roots of the characteristic equation:

$$a_n = Cr^n + Ds^n$$

– If  $r$  is a *repeated* roots of the characteristic equation:

$$a_n = Cr^n + Dnr^n$$

- **Structural Induction** — Sets are defined recursively using a BASE group of elements, a list of RECURSION rules to create new elements, and a RESTRICTION that these are the only ways to form elements of the set. You can perform induction on the set to check if a condition is true by applying the recursion rules of the set to the base case of the induction.
- **Element Argument of Sets** — To show that  $X \subseteq Y$ , show that if you assume  $x \in X$  arbitrarily, that  $x \in Y$  as well.
- **Set Equality** — To show that sets  $X = Y$ , you must show  $X \subseteq Y$  and  $Y \subseteq X$ .
- **Element Method for the Empty Set** — To show a set  $X = \emptyset$ , use proof by contradiction - assume  $\exists x \in X$  and demonstrate this as a contradiction.
- **Power Sets** — The set of all subsets of a set  $X$  is called the *power set* of  $X$  and is denoted  $\mathcal{P}(X)$ .

## Practice Textbook Problems