



**Trinity College Dublin**

Coláiste na Tríonóide, Baile Átha Cliath

The University of Dublin

## Week 2

## Monads

**CS4012**

Topics in Functional Programming

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# Monads

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- **Certain patterns of computation come up over and over**
- **I want to look at a way to structure programs that can solve lots of problems for us.**
- **You've already seen this, but not from the ground up...**

# Monads

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- Last year you encountered the IO type as a mechanism for dealing with certain kinds of computations.
- Specifically...
- Functions that represent *actions which have side-effects*.
- The word “*monad*” was used to refer to this abstraction
- What’s going on with these functions, and how do they fit into my claim that we have a useful way to structure programs?

# Monads

## The problem of IO

- Let's remind ourselves of the issues around IO.
- Imagine we have some functions such as:

```
primGetChar :: Char  
primPutChar :: Char → ()
```

- For these functions to be meaningful they would have to be performing some side-effecting IO operations whenever they are evaluated.
- That's clearly going to violate the principle of referential transparency, and make us sad!

# Monads

## Referential transparency?

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- It won't take much to illustrate the problem.
- Do we know what this will do?

```
f1 = (primGetChar, primGetChar)
```

- How about this?

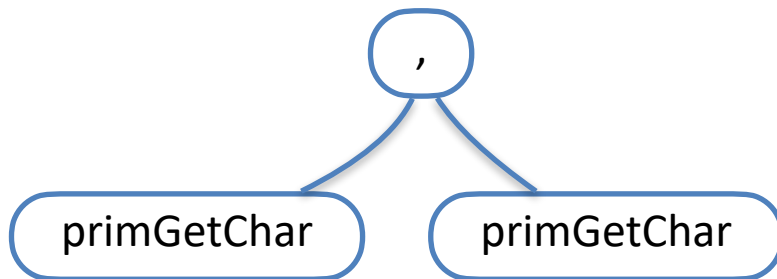
```
f2 = let x = primGetChar in (x,x)
```

# Monads

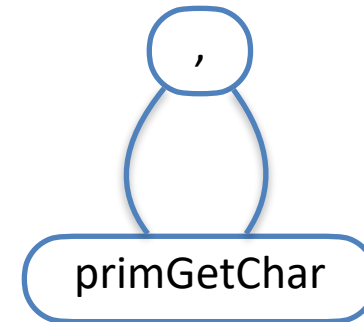
## Referential transparency?

- If we draw the graphs for the two expressions we can see the problem

```
f1 = (primGetChar, primGetChar)
```



```
f2 = let x = primGetChar in (x,x)
```



- In the right-hand case we will get only one IO action performed, while on the left there will be two. But the expressions were supposed to be the same!
- Side-effects really mess things up...

# Monads

## Token passing

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- I'm going to propose a solution.
- We want the two uses of “primGetChar” to be different somehow, so I will add an argument to the function, and use that to distinguish them.

```
getChar :: World → (Char,World)
```

- The (unsafe) primGetChar will be wrapped up in this (safe) function that takes the side-effects into account.

# Monads

## World?

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- The “World” here is a parameter which represents the state of the world from one moment to another.
- The actual details of how that is encoded, and of how `getChar` might use it, are not important right now.
- We will just assume that it is not possible to make copies of the World, nor to magic-up a World out of nowhere.



# Monads

## State passing

- With this idea in place we must write our hazardous function differently:

```
f3 :: World → ((Char, Char), World)
f3 w0 = ((ch1, ch2), w2)
      where
          (ch1, w1) = getChar w0
          (ch2, w2) = getChar w1
```

- Having to “thread” the various w- parameters forces the evaluation to happen the way we want.
- As long as we are careful never to make more than one reference to any given moment in the state of the world!

# Monads

## State passing

- We need to ensure that this sort of nonsense never happens:

```
f3 :: World → ((Char, Char), World)
f3 w = ( (ch1, ch2), w2 )
      where
          (ch1, w1) = getChar w
          (ch2, w2) = getChar w
```

- While we are at it, we would like to make it easier to write this style of function
- Manually “threading” those various “w”s around will get tedious very quickly!

# Monads

## Structured state passing

- Let's assume that all our world-mangling functions have this “shape”:

```
f :: World → (a, World)
```

- We can declare a type that captures this, which will clean up our code a bit.

```
type IO a = World → (a, World)
```

- Again, not really worrying about what `World` actually is at this time

# Monads

## Structured state passing

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- Then we get some types that will look rather familiar to you

```
getChar :: IO Char  
putChar :: Char → IO ()
```

# Monads

## Structured state passing

- The second thing we can do is hide all the “plumbing” involved in threading the state.
- For example, if I declare an (infix) function with this type:

```
(>>) :: IO a → IO b → IO b
```

- Used like this:

```
f4 = (putChar 'a') >> (putChar 'b')
```

- Read it as “do the first thing, throw away the result but keep the World, then do the second thing”.
- The results of printing are all just “( )” values so there’s nothing of interest dropped.

# Monads

## Structured state passing

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- A possible implementation for ( $\gg$ )

```
( $\gg$ )  $\lambda r = \lambda w \rightarrow \text{let } (_, w1) = \lambda w$   
       $\text{in } r\ w1$ 
```

# Monads

## Structured state passing?

- If the result is significant then instead of throwing it away we can keep it.
- Our first thought might be to have some function like this:

```
pair :: IO a → IO b → IO (a,b)
```

- But that's not really the right structure.
- For one thing, that type doesn't really say that the left thing happens *before* the right thing. If we wrote “pair” to do it in the opposite order that would still be the type!
- We will often want to do something in the second action that makes use of the result of the first action

# Monads

## Structured state passing

- So we'll really do something like this:

$$(>\!\!\asymp) :: IO\ a \rightarrow (a \rightarrow IO\ b) \rightarrow IO\ b$$

- In fact, if we have this then it's easy to write `>>`

$$(>>) \text{ l } r = \text{ l } >\!\!\asymp (\_\rightarrow r)$$



# Monads

## Structured state passing

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- To write useful programs I'll add one utility function
- A computation that does nothing (has no side effect) but which produces a value

```
return :: a → IO a
```

- This is useful if we want to do combine IO and non-IO computations.

# Monads

## Using the monad functions

- Let's see this in use
- Here's a function that reads two characters in order, then returns them as a string

```
f = getChar >>= ( \ch1 →  
    getChar >>= ( \ch2 →  
        return [ch1,ch2] ))
```

# Monads

## The IO Monad

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- What's we've seen is a possible implementation for the IO monad in Haskell.

# Monads

## Monads in general

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- A monad is an abstraction which represents a computation.
- The computations have results (reflected in the type).
- The monad provides at least two basic operations:
  - `return`, which produces a result without doing anything
  - `>>=`, which binds together two computations in the monad
- Of course a monad will also provide a collection of primitive operations (like `getChar`), in order to be useful.

# Monads

## Syntactic sugar

- One thing with this style of programming is that you can end up with long chains of `>>=` and `>>` operations.
- Haskell provides some syntactic sugar, called the “do-notation”, that allows us to write the previous program like this:

```
f = do
    ch1 ← getChar
    ch2 ← getChar
    return [ch1, ch2]
```

# Monads

## Syntactic sugar

- The notation is automatically de-sugared into the combinatory form.
- A summary of the rules:

`do x`  
`y`  $\longrightarrow$  `x >> y`

`do a ← x`  
`y`  $\longrightarrow$  `x >=> \a → do y`

`do x`  $\longrightarrow$  `x`



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# Thank you

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