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1 Review of Propositional Logic

Task: Recall enough propositional logic to see how it matches up with set theory.

Definition: A <u>proposition</u> is any declarative sentence that is either true or false.

1.1 Connectives

| $\underline{\text{Connectives}}$ | | Notation in Maths | |
|----------------------------------|-------------------|-----------------------------|-------------------|
| and | \wedge | | |
| or | \vee | "Inclusive or" | |
| not | \neg | Sometimes denoted \sim | |
| implies | \rightarrow | if/then; called implication | \Rightarrow |
| if and only if | \leftrightarrow | Called equivalence | \Leftrightarrow |

1.1.1 Truth Table of the Connectives

Let P, Q be propositions:

| Р | Q | $P \wedge Q$ | Р | Q | |
|---|---|--------------|---|---|---|
| F | F | F | F | F | Ī |
| F | Τ | F | F | Т | Ī |
| Т | F | F | Т | F | |
| Τ | Т | Т | Т | Т | L |

| F T | P | $\neg P$ |
|-------|---|----------|
| TF | F | Т |
| * * | Т | F |

NB In some textbooks, T is denoted by 1, and F is denoted by 0.

| Р | Q | $P \rightarrow Q$ |
|---|---|-------------------|
| F | F | Т |
| F | Т | Т |
| Т | F | F |
| Т | Τ | Т |

NB Note that the only instance when an implication (if/then statement) denoted by $P \to Q$ is false is when the hypothesis (P) is true, but the conclusion (Q) is false.

| P | Q | $P \leftrightarrow Q$ |
|---|---|-----------------------|
| F | F | Т |
| F | Τ | F |
| Т | F | F |
| Т | Т | Т |

NB The truth table for the equivalence says that both P and Q must have the same truth value, i.e. both be true or both be false for the equivalence to be true.

Priority of the Connectives

Highest to Lowest: $\neg, \land, \lor, \rightarrow, \leftrightarrow$

1.2 Important Tautologies

$$\begin{array}{cccc} (P \to Q) & \leftrightarrow & (\neg P \vee Q) \\ (P \leftrightarrow Q) & \leftrightarrow & [(P \to Q) \wedge (Q \to P)] \\ \neg (P \wedge Q) & \leftrightarrow & (\neg P \vee \neg Q) \\ \neg (P \vee Q) & \leftrightarrow & (\neg P \wedge \neg Q) \end{array} \right\} \ \, \begin{array}{c} \text{De Morgan Laws} \\ \text{(these have parallels in in} \\ \text{set theory)} \end{array}$$

As a result, \neg and \lor together can be used to represent all of \neg , \land , \lor , \rightarrow , \leftrightarrow .

Less obvious: One connective called the Sheffer stroke P|Q (which stands for "not both P and Q" or "P nand Q") can be used to represent all of \neg , \wedge , \vee , \rightarrow , \leftrightarrow since $\neg P \leftrightarrow P|P$ and $P \vee Q \leftrightarrow (P|P) \mid (Q|Q)$.

Recall that if $P \rightarrow Q$ is a given implication, then $Q \rightarrow P$ is called the <u>converse</u> of $P \rightarrow Q$, while $\neg Q \rightarrow \neg P$ is called the contrapositive of $P \rightarrow Q$.

1.3 Indirect Arguments/Proofs by Contradiction/Reductio ad absurdum

Based on the tautology (P \rightarrow Q) \leftrightarrow (\neg Q \rightarrow \neg P)

Example: Famous argument that $\sqrt{2}$ is irrational.

Proof:

Suppose $\sqrt{2}$ is rational, then it can be expressed in fraction form as $\frac{a}{b}$ with a and b integers, $b \neq 0$. Let us **assume** that our fraction is reduced, **i.e.** the only common divisor of the numerator a and denominator b is 1.

Then,

$$\sqrt{2} = \frac{a}{b}$$

Squaring both sides, we have

$$2 = \frac{a^2}{h^2}$$

Multiplying both sides by b^2 yields

$$2b^2 = a^2$$

Therefore, 2 divides a^2 , i.e. a^2 is even. If a^2 is even, then a is also even, namely a=2k for some integer k.

Substituting the value of 2k for a, we have $2b^2 = (2k)^2$ which means that $2b^2 = 4k^2$. Dividing both sides by 2, we have $b^2 = 2k^2$. That means 2 divides b^2 , so b is even.

divides b^2 , so b is even.

This implies that both a and b are even, which means that both the numerator and the denominator of our fraction are divisible by 2. This contradicts our **assumption** that the numerator a and the denominator b have no common divisor except 1. Since we found a contradiction, our assumption that $\sqrt{2}$ is rational must be false. Hence the theorem is true.

qed