

# Normalization

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# Anomalies

- If we combine information from two tables in the same, we get **anomalies**
- Insert
- Delete
- Update

<u>Id</u>	Name	<u>email</u>
111	Orlando	<u>ok@ok.com</u>
111	Orlando	<u>ok@spsu.edu</u>
222	Lina	<u>lc@spsu.edu</u>

# Functional Dependency

- $A \rightarrow B$  means if two rows have same value of A they necessarily have same value of B
  - Notice A and B are *sets* of attributes !
- Generalization of candidate key - candidate key says no two rows can have same value
- FDs are determined by **semantics**, you can't say an FD exists just by looking at data
- But can say whether it is NOT true by looking at data

# Quick Check

<b>Id</b>	<b>Name</b>	<b>Gender</b>	<b>Age</b>
1	Orlando	Male	35
2	John	Male	35
3	Jane	Female	31
4	Jane	Female	30

**Id -> Name ?**

**Age -> Gender?**

**Name-> Id?**

**Name, Age -> Id?**



# Quick Check

<b>Id</b>	<b>Name</b>	<b>Gender</b>	<b>Age</b>
1	Orlando	Male	35
2	John	Male	35
3	Jane	Female	31
4	Jane	Female	30

**Id -> Name ?      YES it \*could\* be true (no two rows with same id)**

**Age -> Gender?      YES it \*could\* be true (the two rows with same age also have same gender). Do you think it would be true in the real world ?**

**Name -> Id?      NO, 3rd and 4th row, same name (jane) but different Id**

**Name, Age -> Id?      YES, no two rows with same name and same age**

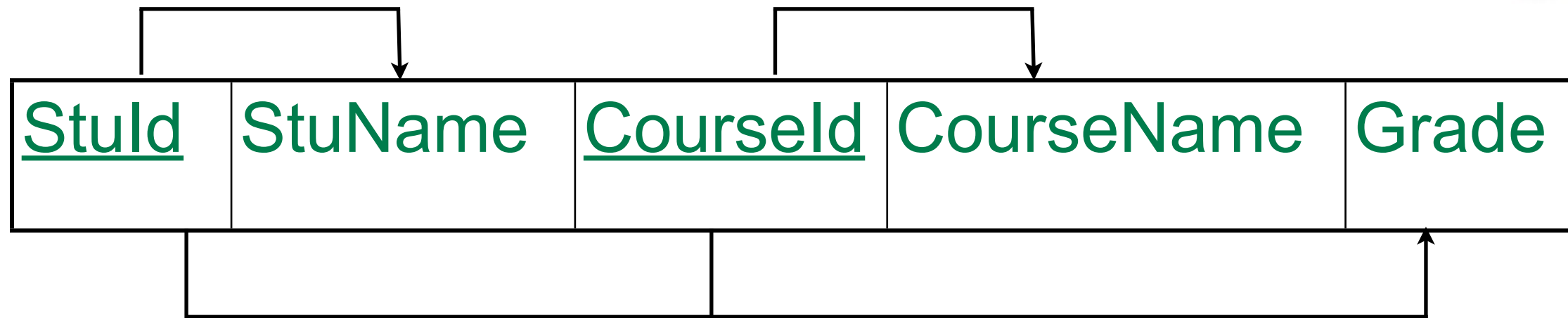
# First normal form

- Only atomic attributes (simple, single-valued)
- Given, since we can't represent tables with non-atomic attributes
- Every *relation* is in 1NF by definition

# Second normal form

- 1NF and no *partial* (functional) dependencies
- A partial dependency is when a non-key attribute depends on a *part* of the primary key
- Of course, you need a composite PK

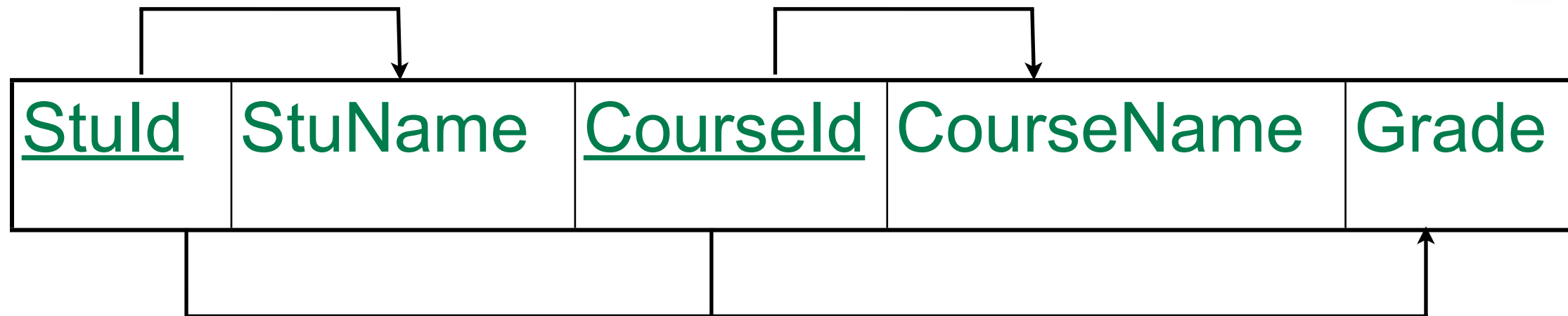
# 2NF Example



- Can represent FDs with arrow as above or
  - StuId->StuName,
  - CourseId->CourseName
  - StuId,CourseId -> Grade (and StuName, CourseName)
- Any partial FDs ?

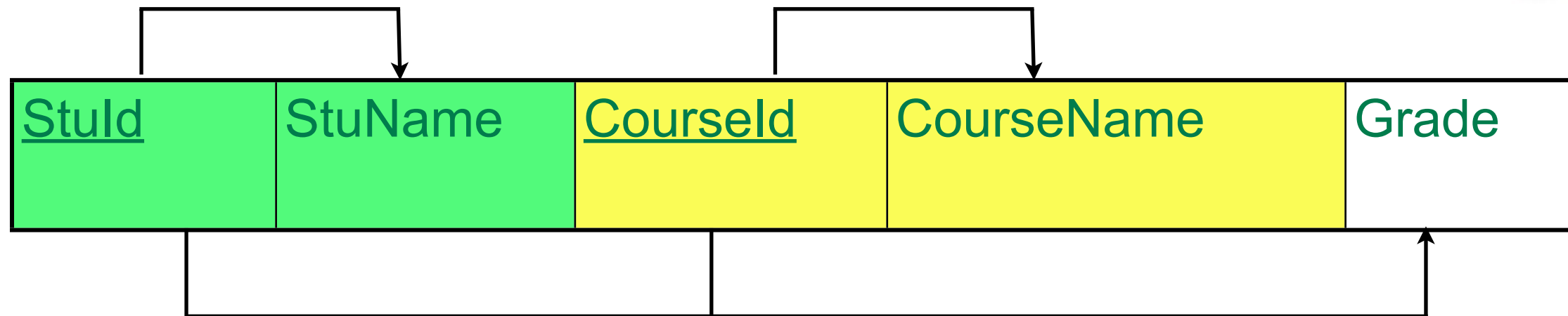


# 2NF Example (again)

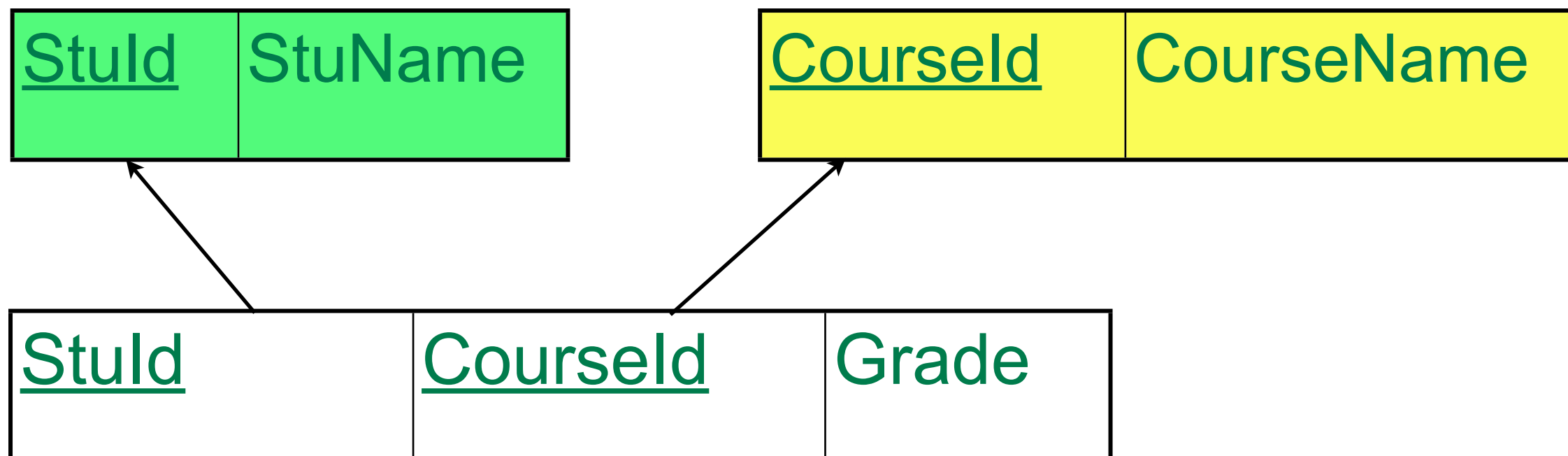


- Can represent FDs with arrow as above or
  - **Stuld->StuName** **PARTIAL**
  - **CourseId->CourseName** **PARTIAL**
  - **Stuld, CourseId -> Grade** (and StuName, CourseName)
- How do we convert the partial dependencies into normal ones ? By breaking into more tables

# 2NF: Normalizing



Becomes ... (notice above arrows mean functional dependency, below they represent foreign key constraints)



# You try it now

<u>SeriesId</u>	<u>EpisodeId</u>	SeriesTitle	EpisodeTitle	AiringDate
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- FDs
  - SeriesId -> SeriesTitle
  - SeriesId,EpisodeId -> SeriesTitle, EpisodeTitle, AiringDate
- Any partial dependencies ? How do we eliminate them ?

# Solution

<u>SeriesId</u>	<u>EpisodeId</u>	SeriesTitle	EpisodeTitle	AiringDate
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- FDs
  - SeriesId -> SeriesTitle PARTIAL
  - SeriesId,EpisodeId -> SeriesTitle, EpisodeTitle, AiringDate
- How do we eliminate the partial FD? Make new table !

<u>SeriesId</u>	<u>EpisodeId</u>	EpisodeTitle	AiringDate
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<u>SeriesId</u>	SeriesTitle
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# Third normal form

- 2NF and no *transitive* (functional) dependencies
- A *transitive* dependency is when a non-key attribute depends on another non-key attribute

# 3NF Example

<u>Course</u>	<u>SectNum</u>	Classroom	Capacity
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- Classroom->Capacity
- Any partial FDs ?
- Any transitive FDs ?

# 3NF Example

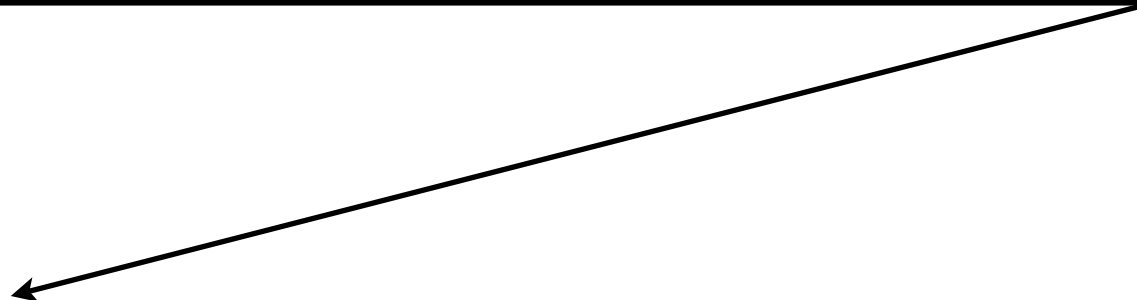
<u>Course</u>	<u>SectNum</u>	Classroom	Capacity
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- Classroom->Capacity      TRANSITIVE
- Any partial FDs ?      NO
- Any transitive FDs ?      Yes ! how do we eliminate ? by breaking into its own table

# 3NF Normalization

<u>Course</u>	<u>SectNum</u>	Classroom
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<u>Classroom</u>	Capacity
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# You try it

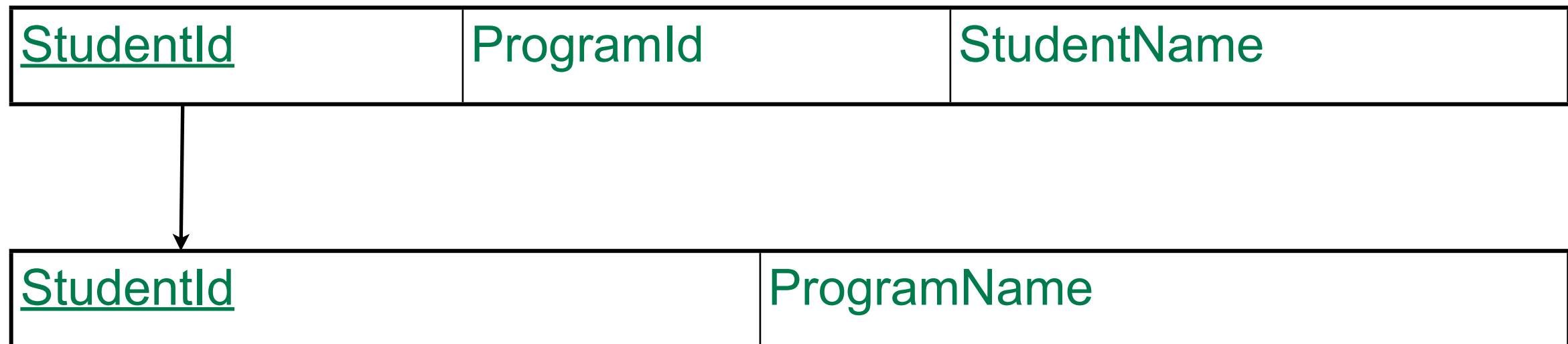
<u>StudentId</u>	ProgramId	StudentName	ProgramName
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- ProgramId -> ProgramName
- Any Partial FDs ?
- Any Transitive FDs?

# Solution

<u>StudentId</u>	ProgramId	StudentName	ProgramName
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- ProgramId -> ProgramName TRANSITIVE
- Any Partial FDs ? NO
- How to eliminate transitive FD ? New table !



# Further Normalization

- Boyce-Codd Normal form (BCNF)
  - Slight difference with 3NF
  - To be in 3NF but not in BNF, needs two composite candidate keys, with one attribute of one key depending on one attribute of the other
  - Not very common :)
- Fourth Normal Form (4NF)
  - To break it, need to have two independent multivalued attributes in the same table
- Usually, if you're in 3NF you're in BCNF, 4NF, ...