



Hacettepe University
Computer Engineering Department

BBM204 Software Practicum II - 2024 Spring

Programming Assignment 1

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1 Problem Definition

We've conducted a thorough analysis of sorting and searching algorithms, evaluating their efficiency across diverse datasets. Our study focused on three sorting algorithms (insertion, merge, and counting) and two searching algorithms (linear and binary).

Using milliseconds for sorting algorithms and nanoseconds for searching algorithms, we measured their execution times across different data sizes. Additionally, we analyzed the computational and auxiliary space complexities associated with each algorithm.

To present our findings effectively, we created graphs for each dataset, offering a visual representation of algorithmic efficiency. Our study provides valuable insights into the performance and scalability of these algorithms, aiding in informed decision-making for practical applications and algorithmic optimizations.

2 Solution Implementation

2.1 Insertion Sort Algorithm

```
1 public void sort(int[] arr) {  
2     for (int j = 1; j < arr.length; j++) {  
3         int key = arr[j];  
4         int i = j - 1;  
5         while (i >= 0 && arr[i] > key) {  
6             arr[i + 1] = arr[i];  
7             i--;  
8         }  
9         arr[i + 1] = key;  
10    }  
11 }
```

2.2 MergeSort Algorithm

```
13 public void sort(int[] arr) {
14     mergeSort(arr);
15 }
16
17 public static int[] mergeSort(int[] A) {
18     int n = A.length;
19     if (n <= 1) {
20         return A;
21     }
22     int mid = n / 2;
23     int[] left = Arrays.copyOfRange(A, 0, mid);
24     int[] right = Arrays.copyOfRange(A, mid, n);
25
26     left = mergeSort(left);
27     right = mergeSort(right);
28
29     return merge(left, right);
30 }
31
32 private static int[] merge(int[] A, int[] B) {
33     int[] C = new int[A.length + B.length];
34     int i = 0, j = 0, k = 0;
35
36     while (i < A.length && j < B.length) {
37         if (A[i] <= B[j]) {
38             C[k] = A[i];
39             i++;
40         } else {
41             C[k] = B[j];
42             j++;
43         }
44         k++;
45     }
46     while (i < A.length) {
47         C[k] = A[i];
48         i++;
49         k++;
50     }
51     while (j < B.length) {
52         C[k] = B[j];
53         j++;
54         k++;
55     }
56     return C;
57 }
```

2.3 Counting Sort Algorithm

```
60 public void sort(int[] arr) {  
61     countingSort(arr, Arrays.stream(arr).max().orElse(0));  
62 }  
63  
64 public static int[] countingSort(int[] A, int k) {  
65     int[] count = new int[(k + 1)];  
66     int[] output = new int[A.length];  
67     int size = A.length;  
68  
69     for (int i = 0; i < size; i++) {  
70         count[A[i]]++;  
71     }  
72  
73     for (int i = 1; i <= k; i++) {  
74         count[i] += count[i - 1];  
75     }  
76  
77     for (int i = size - 1; i >= 0; i--) {  
78         output[(count[A[i]] - 1)] = A[i];  
79         count[A[i]]--;  
80     }  
81  
82     return output;  
83 }
```

2.4 Linear Search Algorithm

```
86 public int Search(int[] arr, int value) {  
87     int size = arr.length;  
88     for (int i = 0; i < size; i++) {  
89         if (arr[i] == value) {  
90             return i;  
91         }  
92     }  
93     return -1;  
94 }
```

2.5 Binary Search Algorithm

```
96 public int Search(int[] arr, int value) {  
97     int lowIndex = 0;  
98     int highIndex = arr.length - 1;  
99  
100     while (highIndex - lowIndex > 1){  
101         int mid = (highIndex + lowIndex) / 2;  
102         if(arr[mid] < value) {  
103             lowIndex = mid + 1;  
104         }  
105         else {  
106             highIndex = mid;  
107         }  
108     }  
109  
110     if(arr[lowIndex] == value)  
111         return lowIndex;  
112  
113     if(arr[highIndex] == value)  
114         return highIndex;  
115  
116     return -1;  
117 }
```

3 Results, Analysis, Discussion

Table 1: Results of the running time tests performed for varying input sizes (in ms).

Algorithm	500	1000	2000	4000	8000	16000	32000	64000	128000	250000
Random Input Data Timing Results in ms										
Insertion sort	0.5	0.3	0.7	2.0	7.0	33.0	108.3	387.4	1684.0	6986.8
Merge sort	0.2	0.1	0.3	0.3	0.8	1.5	3.1	7.6	20.4	34.9
Counting sort	296.1	168.0	149.2	108.3	104.2	103.8	105.3	107.0	110.8	120.6
Sorted Input Data Timing Results in ms										
Insertion sort	0.0	0.0	0.1	0.0	0.0	0.1	0.0	0.1	0.3	0.6
Merge sort	0.1	0.0	0.1	0.2	0.4	1.0	1.9	3.8	7.7	23.6
Counting sort	113.5	100.7	96.0	96.7	101.3	104.8	104.8	109.3	107.9	111.0
Reversely Sorted Input Data Timing Results in ms										
Insertion sort	0.0	0.3	0.8	3.2	12.7	50.6	198.9	772.5	3233.2	12653.8
Merge sort	0.0	0.1	0.1	0.2	0.4	0.9	2.0	5.3	14.9	15.8
Counting sort	100.0	104.3	107.3	112.8	107.3	111.7	112.5	112.9	110.5	116.1

Table 2: Results of the running time tests of search algorithms of varying sizes (in ns).

Algorithm	500	1000	2000	4000	8000	16000	32000	64000	128000	250000
Linear search (random data)	160.8	259.4	396.3	676.7	1253.1	1901.6	3779.3	7611.0	14126.9	26240.5
Linear search (sorted data)	1675.0	2748.4	454.1	835.0	1543.5	2781.2	5771.2	9749.6	23151.8	47888.9
Binary search (sorted data)	298.6	178.0	212.8	124.4	127.7	140.9	157.9	183.2	267.2	366.2

Table 3: Computational complexity comparison of the given algorithms.

Algorithm	Best Case	Average Case	Worst Case
Insertion sort	$O(n)$	$O(n^2)$	$O(n^2)$
Merge sort	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$
Counting Sort	$O(n + k)$	$O(n + k)$	$O(n + k)$
Linear Search	$O(1)$	$O(n)$	$O(n)$
Binary Search	$O(1)$	$O(\log n)$	$O(\log n)$

Table 4: Auxiliary space complexity of the given algorithms.

Algorithm	Auxiliary Space Complexity
Insertion sort	$O(1)$
Merge sort	$O(n)$
Counting Sort	$O(n + k)$
Linear Search	$O(1)$
Binary Search	$O(1)$

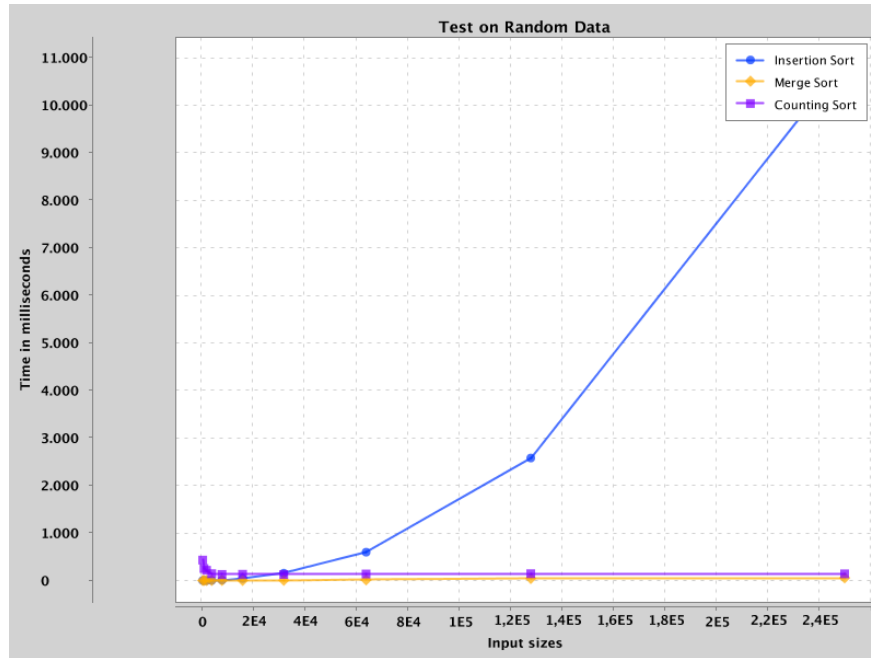


Figure 1: Algorithms Running Times - Random Data Set

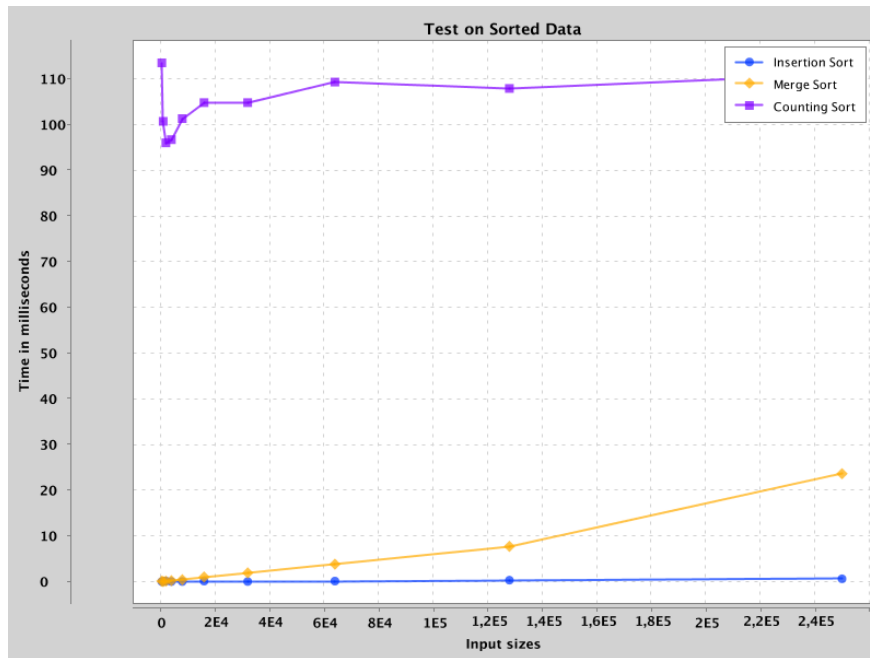


Figure 2: Algorithms Running Times - Sorted Data Set

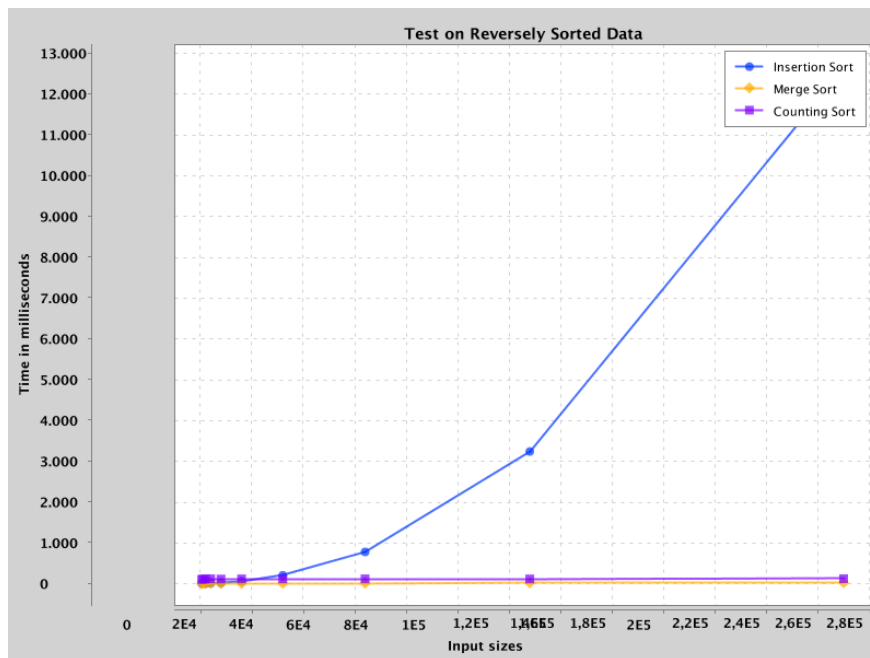


Figure 3: Algorithms Running Times - Reverse Sorted Data Set

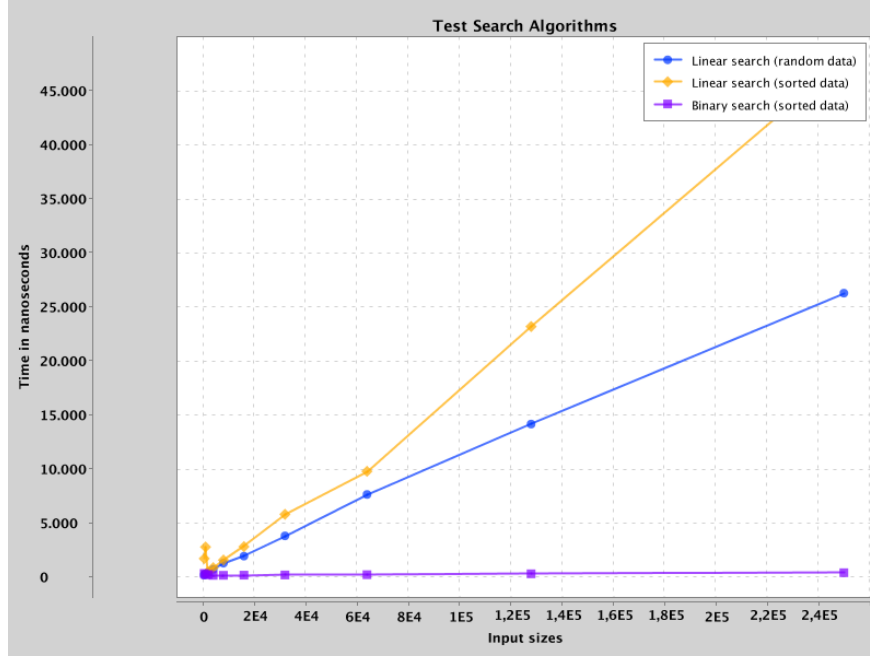


Figure 4: Search Algorithms Running Times - Random&Sorted Data Set

As the input size increases, the time taken also increases for all three sorting algorithms. The increase rate is highest for insertion sort and lowest for counting sort. The difference is not very pronounced in small datasets, but becomes more significant as the dataset grows.

For insertion sort, the best-case scenario appears to be when the data is already sorted because in this case only $n-1$ comparisons are needed, resulting in $O(n)$ time complexity in the best case. Insertion sort exhibits a quadratic curve with random and reverse-sorted data, fitting into $O(n^2)$ time complexity.

For merge sort, the best-case scenario seems to be when the data is already sorted. Merge sort shows more controlled increases compared to insertion sort and exhibits similar behavior with random and reverse-sorted data, fitting into $O(n \log n)$ time complexity.

For counting sort, the best-case scenario appears to be when the data is already sorted because it almost instantly stops sorting in the sorted set. Similarly, counting sort exhibits similar behavior with random and reverse-sorted data, fitting into $O(n + k)$ time complexity.

In terms of search algorithms, the best-case scenario for both search algorithms is under a sorted dataset, which results in $O(1)$ time complexity. Binary search and linear search almost instantly stop searching in a random dataset, resulting in $O(n)$ and $O(\log n)$ time complexity, respectively.

References

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