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DIGITAL CIRCUITS 1st MIDTERM EXAM (Question 1)

QUESTION 1 (30 Points):

a. A is a 4-bit and B is an 8-bit, **signed**, binary integer, which are given as follows:

$$A=(8)_{16} \text{ and } B=(FA)_{16}.$$

Perform necessary operations on **binary numbers** to compare the **absolute values** ($|A|$, $|B|$) of two numbers. Decide which number has the greater absolute value by interpreting the obtained binary result. Explain the operations and interpretation of the result.

b. Assume that the numbers A and B are **unsigned**, and answer the same question.

Solution:

Note: Base 10 values are not necessary. Digital circuits can only operate on binary numbers. All operations must be performed (as written in the question) on **binary numbers**.

a.

A= 1000 Sign extension to obtain 8-bit A= 1111 1000

B= 1111 1010 (Because Hex F=1111, Hex A=1010)

Both numbers are negative (sign =1)

To obtain absolute values we apply 2's complement operations.

2's complement of A= 1111 1000 : $0000\ 0111 + 1 = 0000\ 1000 = |A|$

2's complement of B= 1111 1010: $0000\ 0101 + 1 = 0000\ 0110 = |B|$

To compare the absolute values we perform $|A| - |B| = |A| + 2\text{'s complement of } |B|$

$$\begin{array}{r} |A| : \quad 0000\ 1000 \\ 2\text{'s complement of } |B| : + 1111\ 1010 \\ \hline 1\ 000\ 0010 \end{array}$$

Interpretation:

Absolute values are unsigned numbers. Therefore we investigate the carry (barrow).

Carry=1, that means no barrow. Consequently $|A| > |B|$

b.

A= 1000 Extension to obtain 8-bit A= 000 1000 (A is unsigned)

B= 1111 1010

As the numbers are unsigned $A = |A|$, and $B = |B|$

To compare the numbers we perform $A - B = A + 2\text{'s complement of } B$

$$\begin{array}{r} A : \quad 0000\ 1000 \\ 2\text{'s complement of } B : + 0000\ 0110 \\ \hline 0000\ 1110 \end{array}$$

Interpretation:

Absolute values are unsigned numbers. Therefore we investigate the carry (barrow).

Carry=0, that means barrow. Consequently $A < B$ and $|A| < |B|$

QUESTION 2 (35 Points):

a. Using axioms and theorems of Boolean algebra show that $(a \oplus b \oplus c) + (a \odot b \odot c) = a \oplus b \oplus c$

Note:

i. The output of the $x \oplus y$ function is 1 if the operands are different ($x \neq y$).

ii. The output of the $x \odot y$ is 1 if the operands are same (equal) ($x = y$).

b. $A = (a_1 a_0)$ and $B = (b_1 b_0)$ are two 2-bit, binary integers.

Find the logical expression F , that outputs 1 if $A = B$ by using bitwise comparison.

Implement the expression using **only** 2-input NAND gates.

Solution:

a) **Note:**

$a \oplus b = a\bar{b} + \bar{a}b$ output is 1 if the operands are different ($x \neq y$).

$a \odot b = ab + \bar{a}\bar{b}$ output is 1 if the operands are same (equal) ($x = y$).

Also:

$$\overline{(a \oplus b)} = a \odot b$$

$$\overline{(a \odot b)} = a \oplus b$$

$$(a \oplus b \oplus c) + (a \odot b \odot c) = a \oplus b \oplus c \quad ?$$

$$= \left(a\overline{(b \oplus c)} + \bar{a}(b \oplus c) \right) + \left(a \underbrace{(b \odot c)}_{(b \oplus c)} + \bar{a} \underbrace{(b \odot c)}_{(b \oplus c)} \right)$$

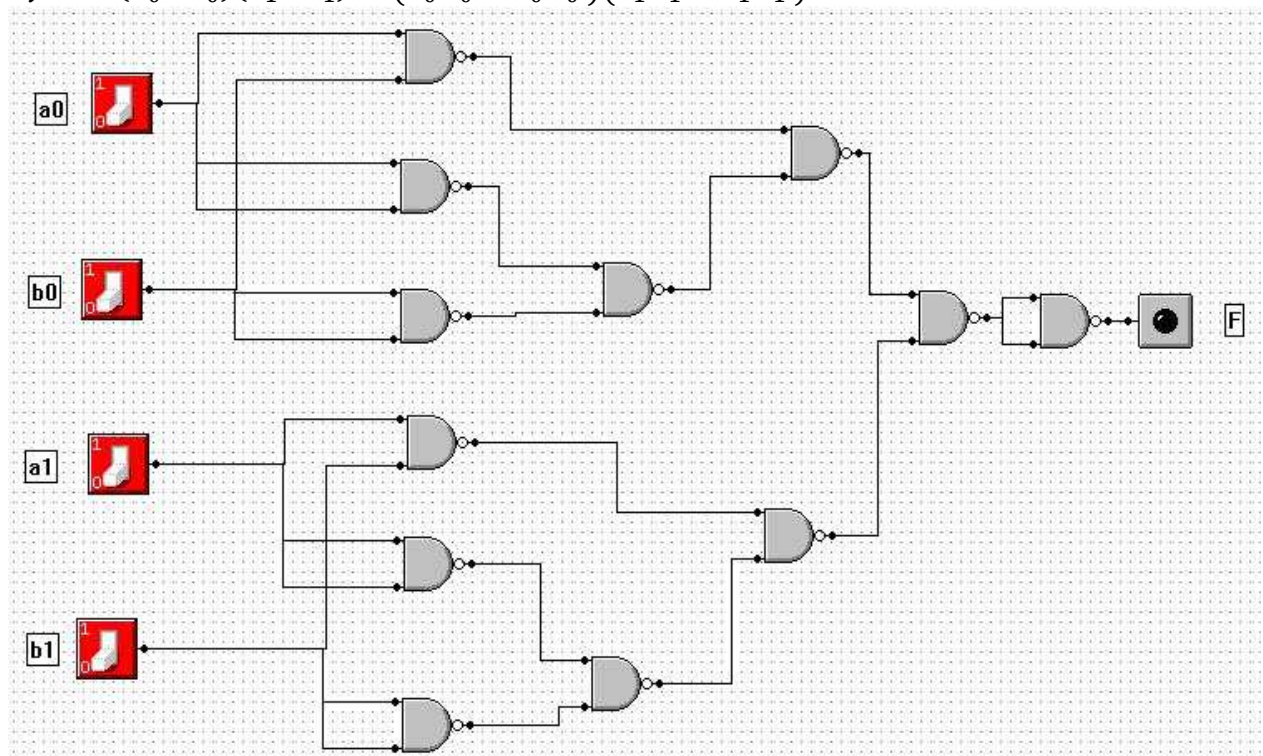
$$= a\overline{(b \oplus c)} + \bar{a}(b \oplus c) + a\overline{(b \oplus c)} + \bar{a}(b \oplus c)$$

$$= \underbrace{a\overline{(b \oplus c)} + a\overline{(b \oplus c)}}_{a\overline{(b \oplus c)}} + \underbrace{\bar{a}(b \oplus c) + \bar{a}(b \oplus c)}_{\bar{a}(b \oplus c)}$$

$$= a\overline{(b \oplus c)} + \bar{a}(b \oplus c) = a \oplus b \oplus c \quad \checkmark$$

Note: There are also other possible solutions.

$$b) F = (a_0 \odot b_0)(a_1 \odot b_1) = (a_0 b_0 + \bar{a}_0 \bar{b}_0)(a_1 b_1 + \bar{a}_1 \bar{b}_1)$$



QUESTION 3 (35 Points):

Consider the logical function given below:

$$f(a,b,c,d) = U_1(0,1,3,6,9) + U_\phi(4,7,8,13,15)$$

- Draw the Karnaugh map of the function $f(a,b,c,d)$ and find all prime implicants.
- Assuming 2 units cost for each variable and 1 unit cost for each complement, find the expression of $f(a,b,c,d)$ with the lowest cost.
- Implement the expression (found in b.) with **only** 2-input NAND gates.

a)

cd \ ab	00	01	11	10
00	1	1	1	
01	0		0	1
11		0	0	
10	0	1		

Prime Implicants	Covered points	Cost
A $\bar{b}\bar{c}$	0, 1, 9	6
B $\bar{a}bc$	6	7
C $\bar{a}cd$	3	7
D $\bar{a}b\bar{d}$	6	8
E $a\bar{c}d$	9	7
F $\bar{a}\bar{c}\bar{d}$	0	9
G $\bar{a}\bar{b}d$	1, 3	8

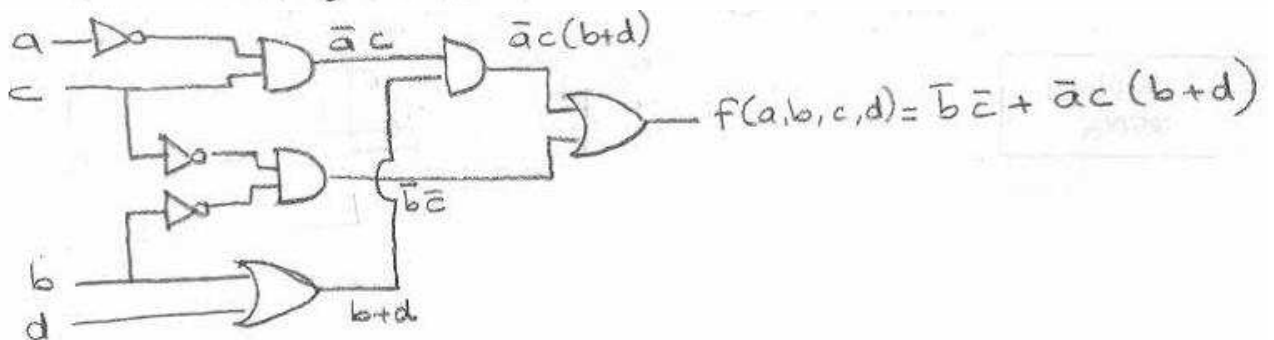
b) Prime implicant chart

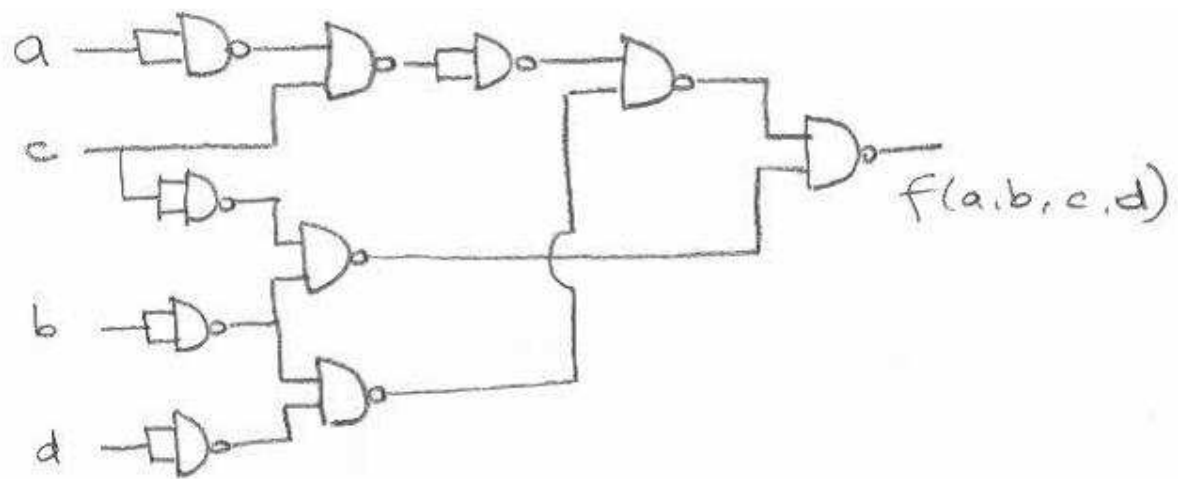
	0	1	3	6	9	Cost
A	X	X			X	6
B				X		7
C			X			7
D				X		8
E					X	7
F	X					9
G		X	X			8

- 1) A covers F and has lower cost. F is deleted.
- 2) A covers E and has lower cost. E is deleted.
- 3) B covers D and has lower cost. D is deleted.
- 4) '3' is covered by C and G. As C has lower cost, it is selected.

$$f(a,b,c,d) = \bar{b}\bar{c} + \bar{a}bc + \bar{a}cd$$

6 7 7 ← total cost = 20 units





or

$$\begin{aligned}
 f(a, b, c, d) &= \left[(\bar{b}\bar{c} + \bar{a}c(b+d))' \right]' \\
 &= \left[(\bar{b}\bar{c})' \cdot (\bar{a}c(b+d))' \right]' \\
 &= \left[(\bar{b} \downarrow \bar{c}) \cdot [(\bar{a}c)' + (b+d)'] \right]' \\
 &= \left[(\bar{b} \downarrow \bar{c}) \cdot \{[(\bar{a} \downarrow c) + (\bar{b} \downarrow d)]'\}' \right]' \\
 &= \left[(\bar{b} \downarrow \bar{c}) \cdot \{(\bar{a} \downarrow c)' \cdot (\bar{b} \downarrow d)'\}' \right]' \\
 &= \left[(\bar{b} \downarrow \bar{c}) \cdot \{(\bar{a} \downarrow c)' \downarrow (\bar{b} \downarrow d)\}' \right]' \\
 &= (\bar{b} \downarrow \bar{c}) \downarrow \{(\bar{a} \downarrow c)' \downarrow (\bar{b} \downarrow d)\}
 \end{aligned}$$