

# Cloud Networking Network Virtualization

## Case Study: VL2

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# Key Needs

## Agility

Location independent addressing

Performance uniformity

Security

Network semantics

# Case Study

## VL2: A Scalable and Flexible Data Center Network

Albert Greenberg  
Srikanth Kandula  
David A. Maltz

James R. Hamilton  
Changhoon Kim  
Parveen Patel

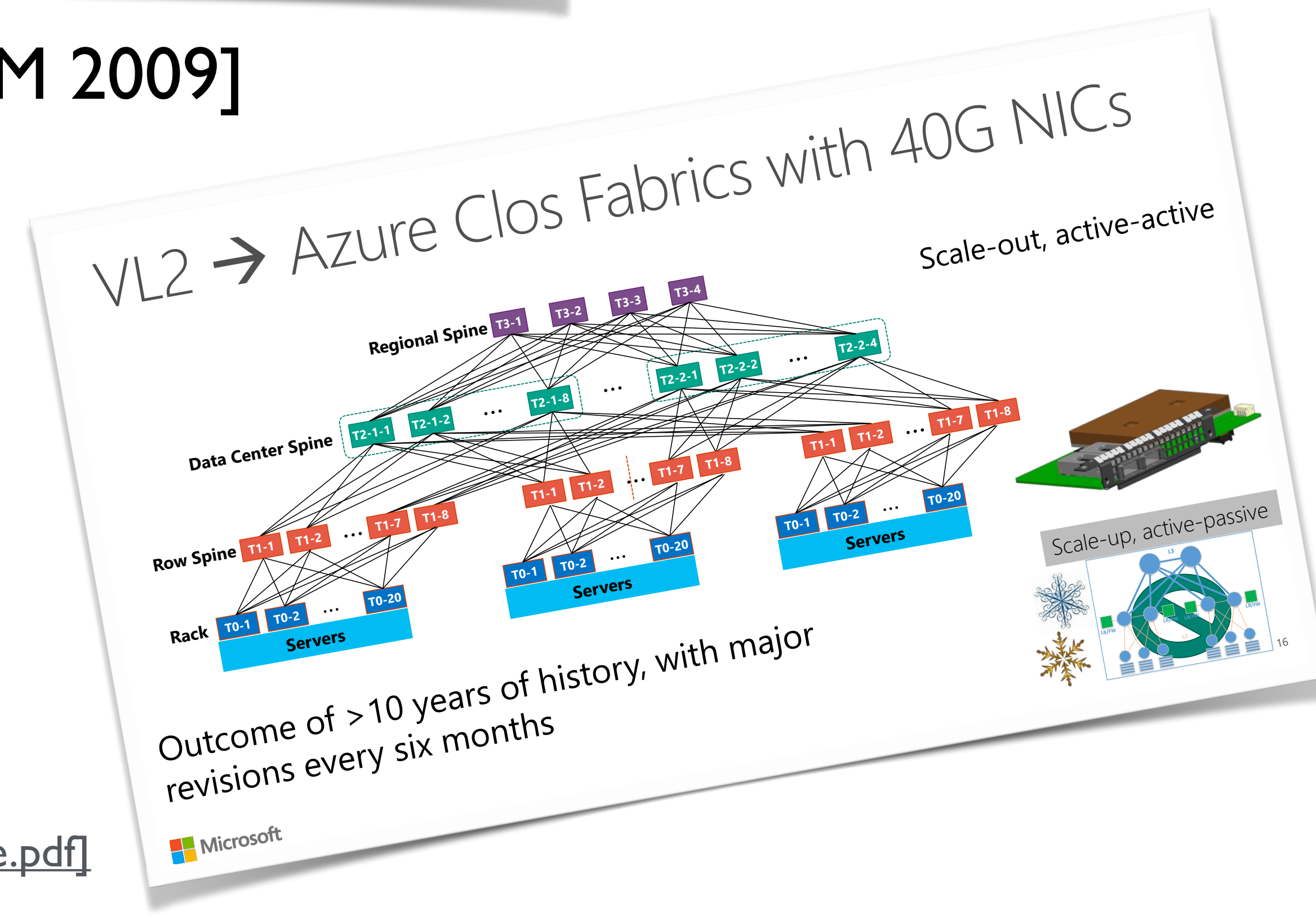
Navendu Jain  
Parantap Lahiri  
Sudipta Sengupta

Microsoft Research

[ACM SIGCOMM 2009]

Influenced architecture of  
Microsoft Azure

[From Albert Greenberg keynote at SIGCOMM 2015:  
<http://conferences.sigcomm.org/sigcomm/2015/pdf/papers/keynote.pdf>]

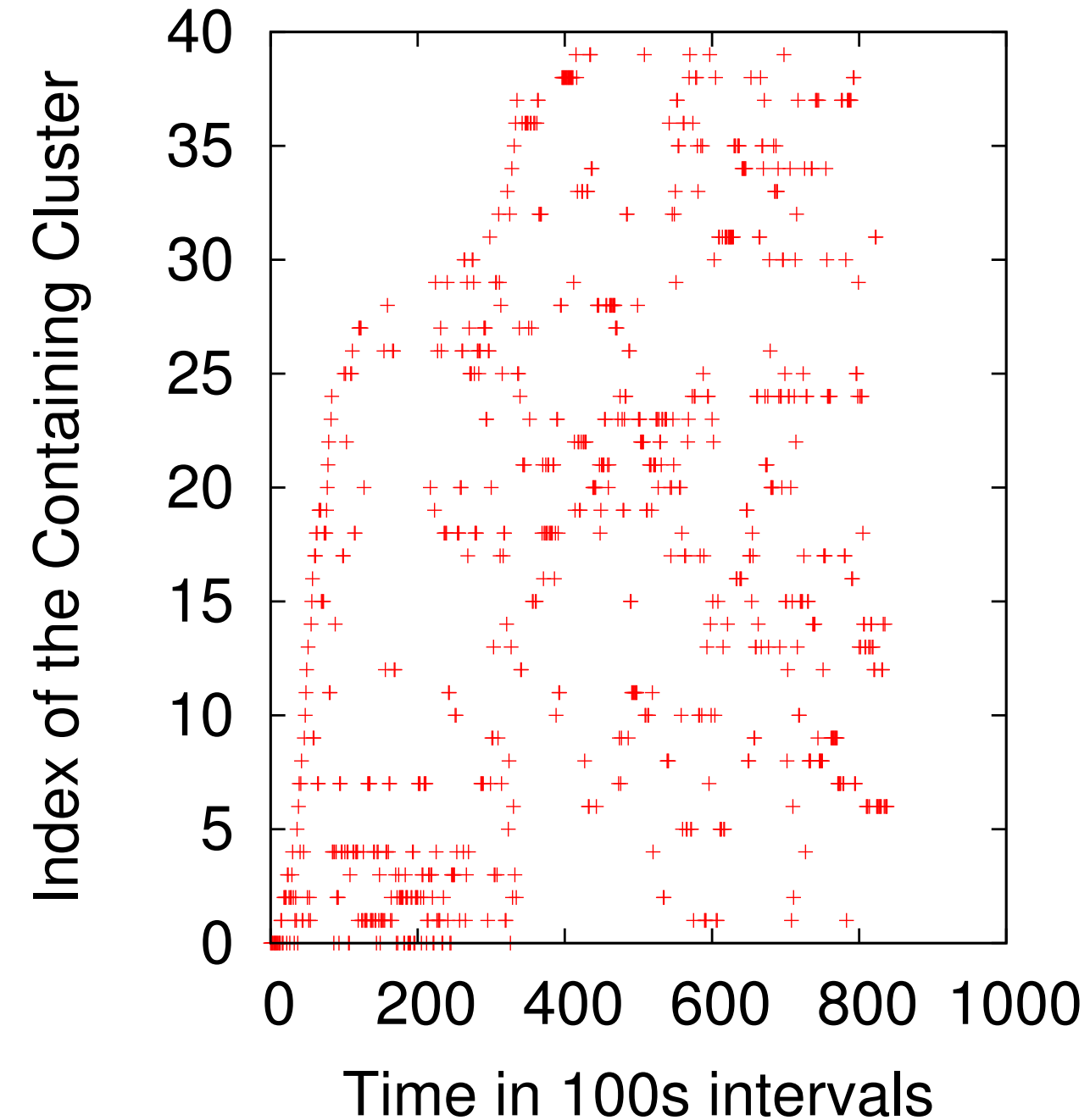


# Motivating Environmental Characteristics

Increasing internal traffic is a bottleneck

- Traffic volume between servers is 4x external traffic

Unpredictable, rapidly-changing traffic matrices (TMs)



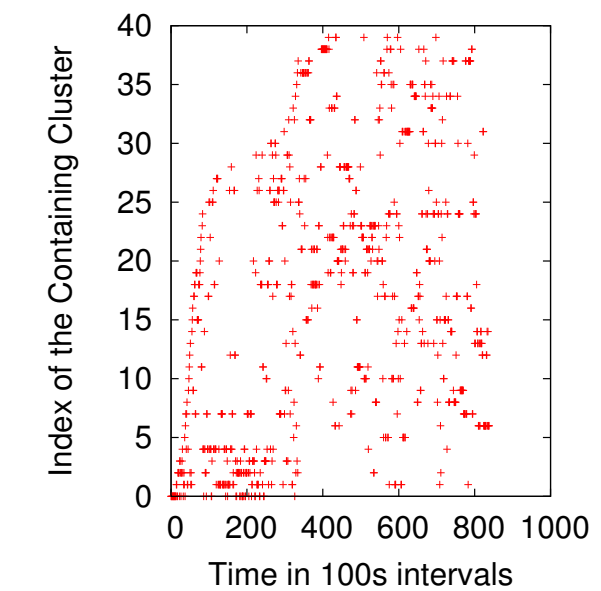
[Greenberg et al.]

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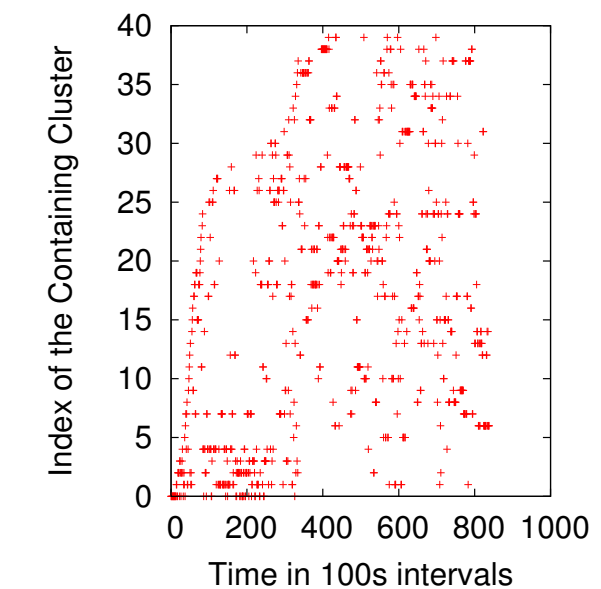


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[Greenberg et al.]

Design result: Nonblocking fabric

- High throughput for *any* TM that respects server NIC rates

# Motivating Environmental Characteristics

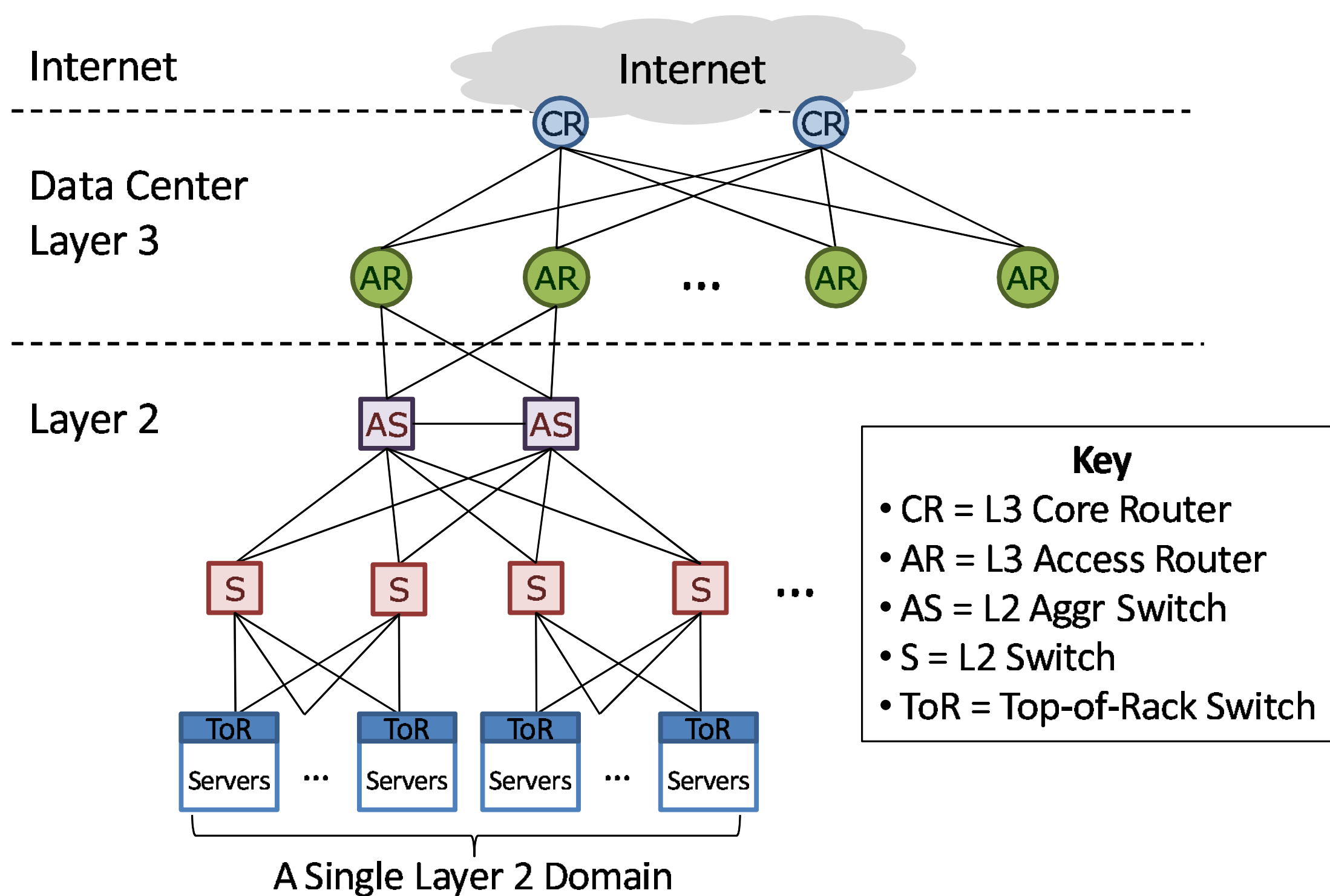
## Failure characteristics

- Analyzed 300K alarm tickets, 36M error events
- 0.4% of failures were resolved in over one day
- 0.3% of failures eliminated all redundancy in a device group (e.g. both uplinks)

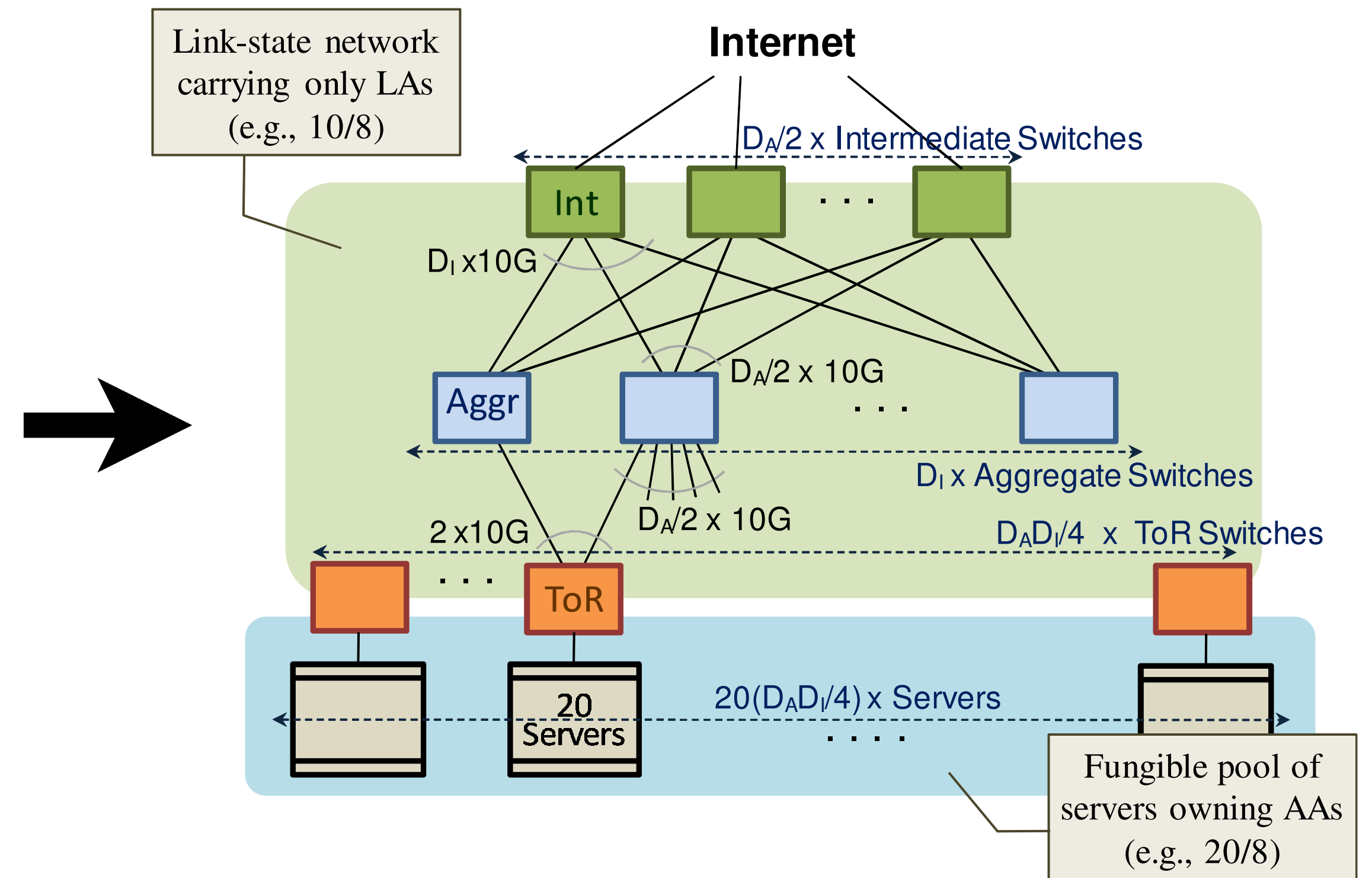
## Design result: Clos topology

- “Scale out” instead of “scale up”

# VL2 physical topology



Traditional



VL2



# Routing in VL2

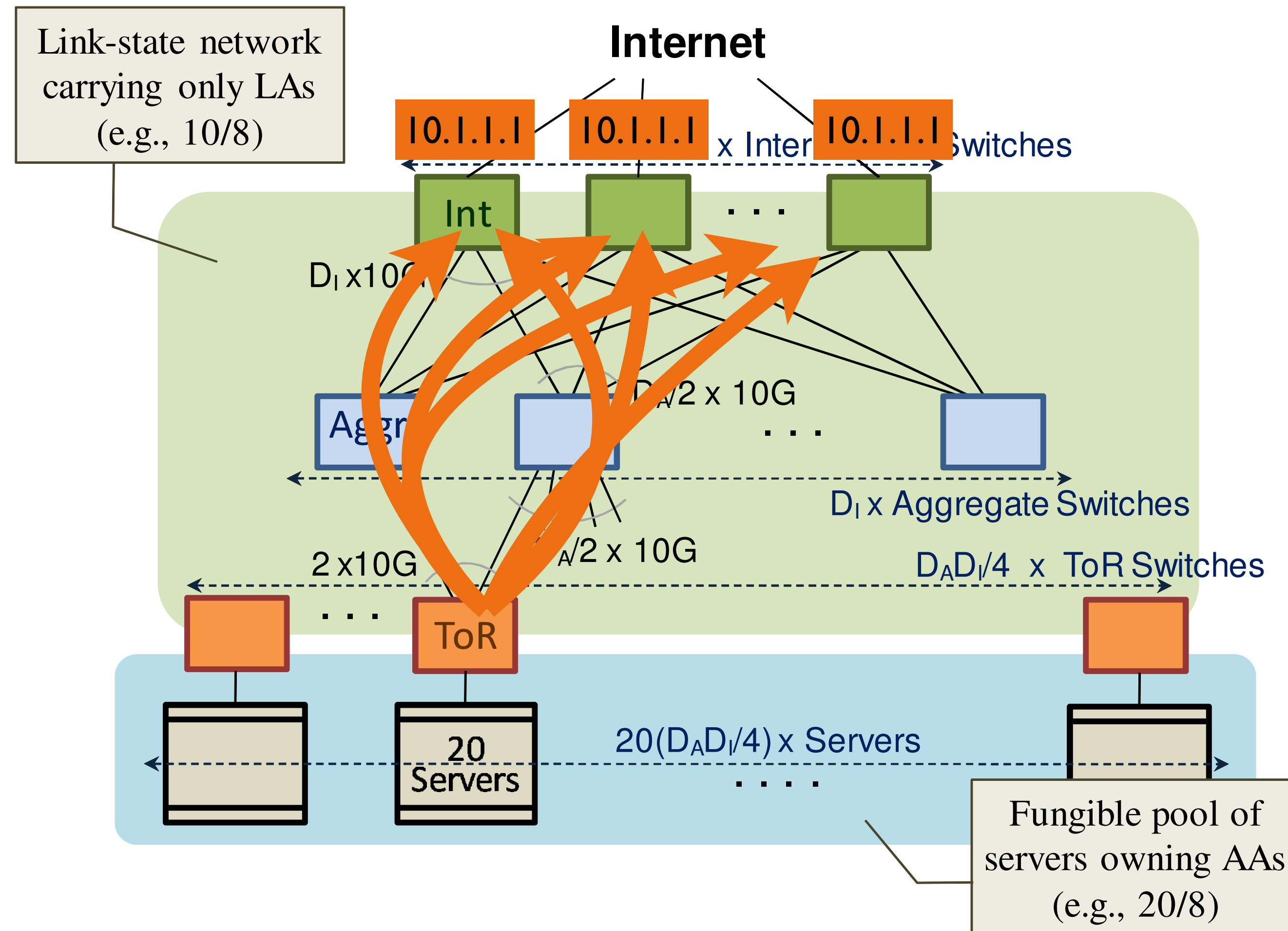
## Unpredictable traffic

- Difficult to adapt

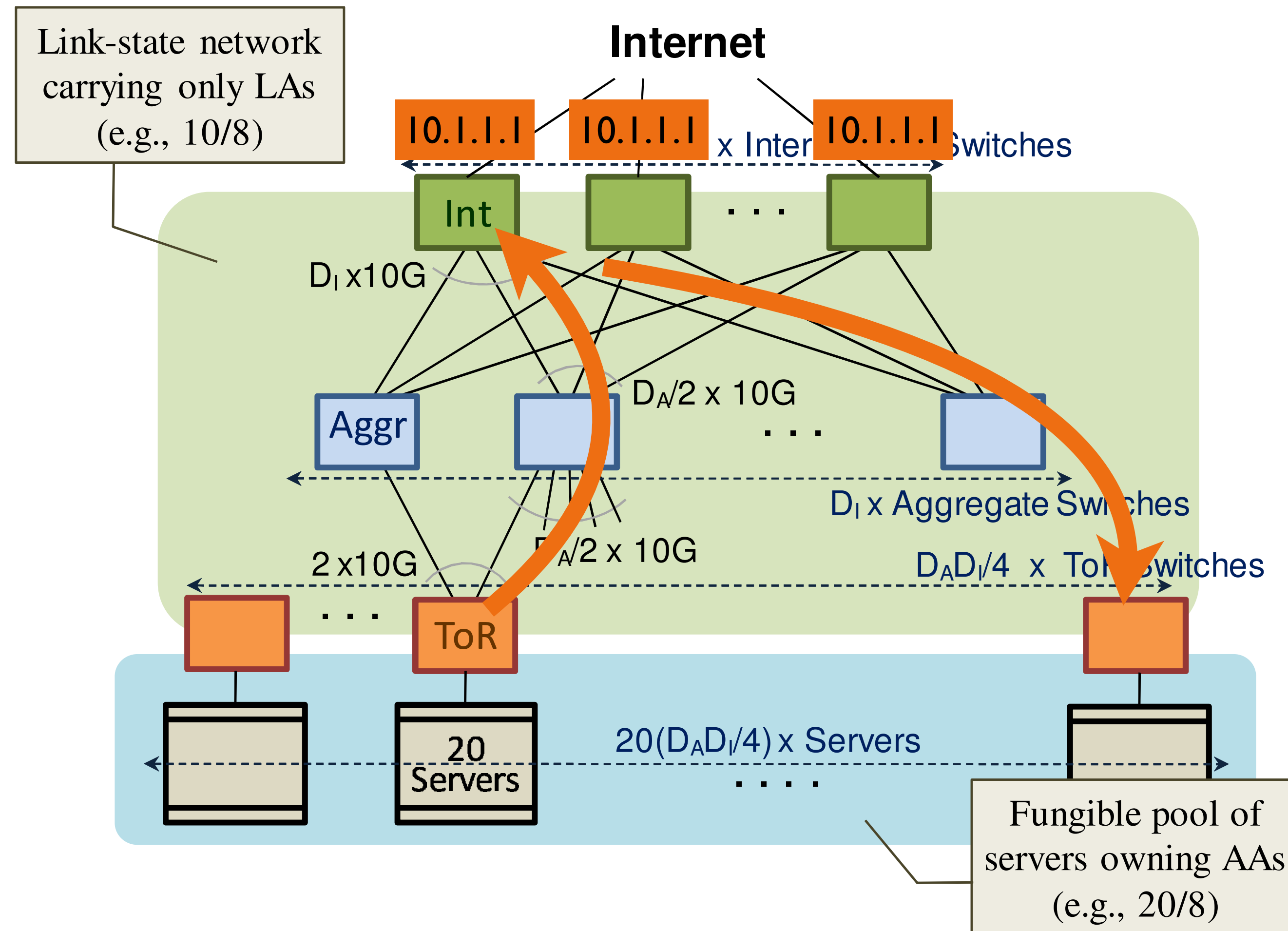
## Design result: “Valiant Load Balancing”

- Route traffic independent of current traffic matrix
- Spreads arbitrary traffic pattern so it's uniform among top layer switches

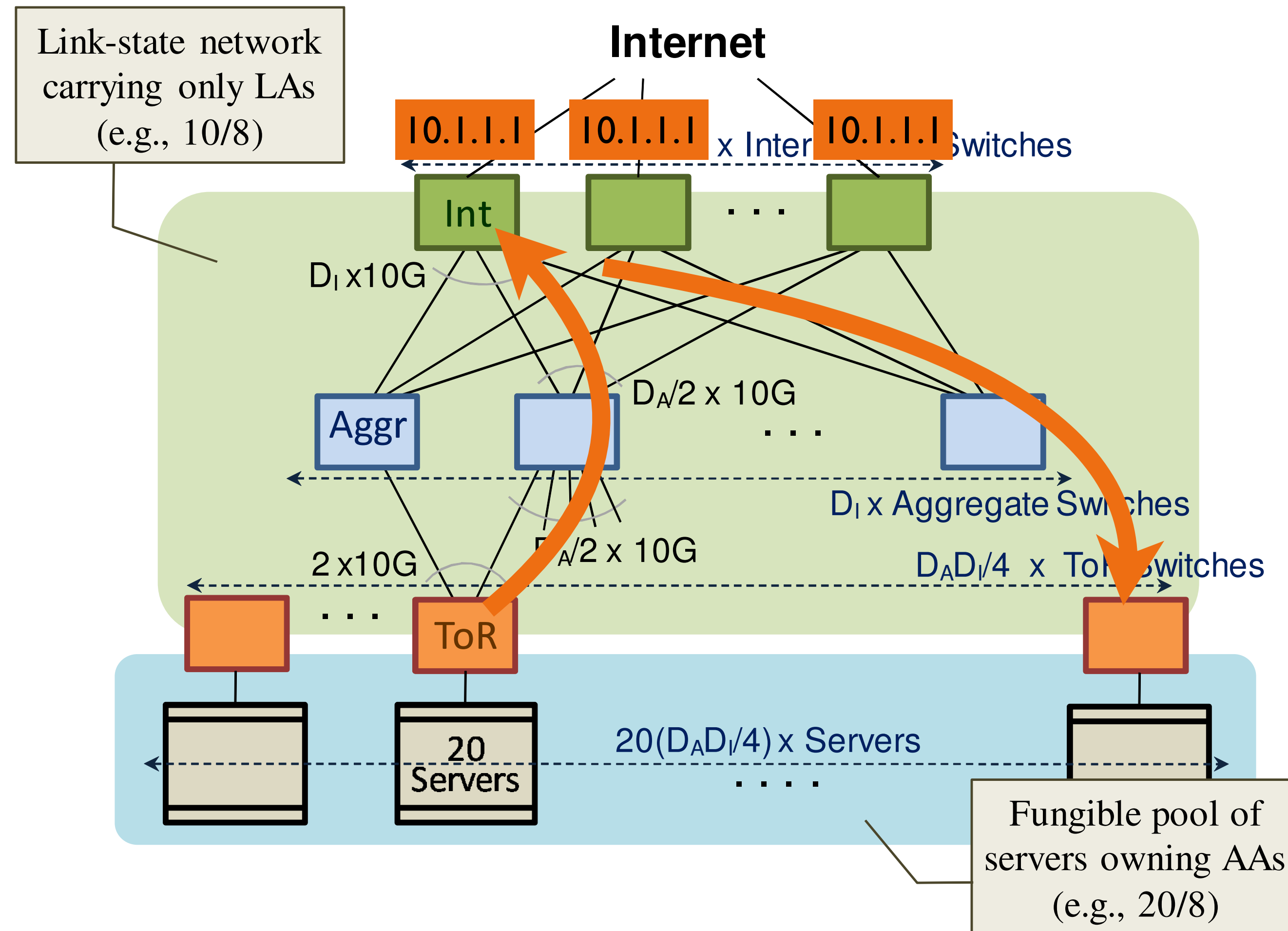
# Routing Implementation



# Routing Implementation



# Routing Implementation



Similar effect to ECMP  
to each rack

Smaller forwarding  
tables at most switches

# Virtualization

“All problems in computer science can be solved by another level of indirection.”

– *David Wheeler*

App / Tenant layer

- Application Addresses (AAs): Location independent
- Illusion of a single big Layer 2 switch connecting the app

Virtualization layer

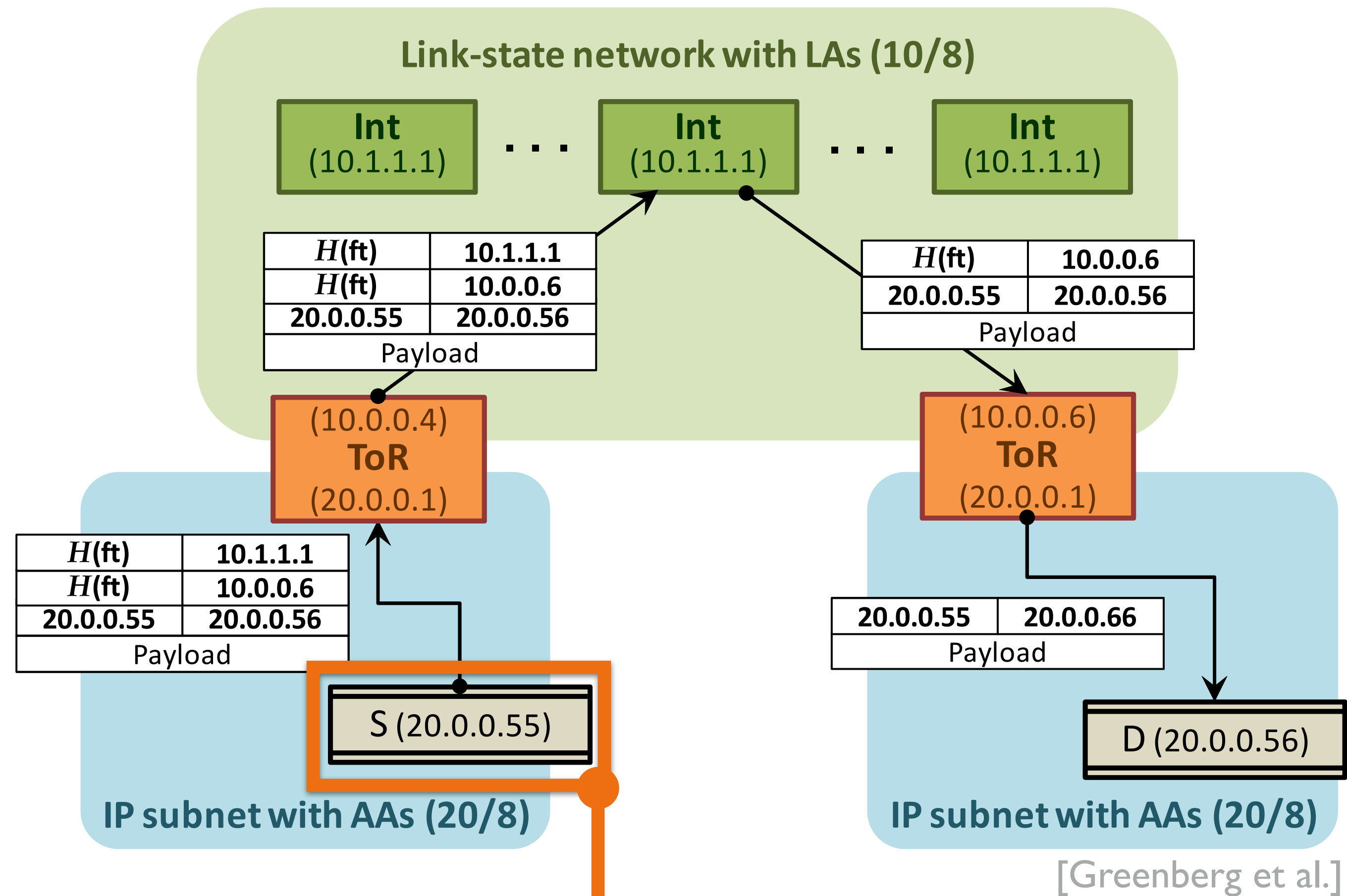
- Directory server: Maintain AA to LA mapping
- Server agent: Query server, wrap AAs in outer LA header

Physical network layer

- Locator Addresses (LAs): Tied to topology, used to route
- Layer 3 routing via OSPF

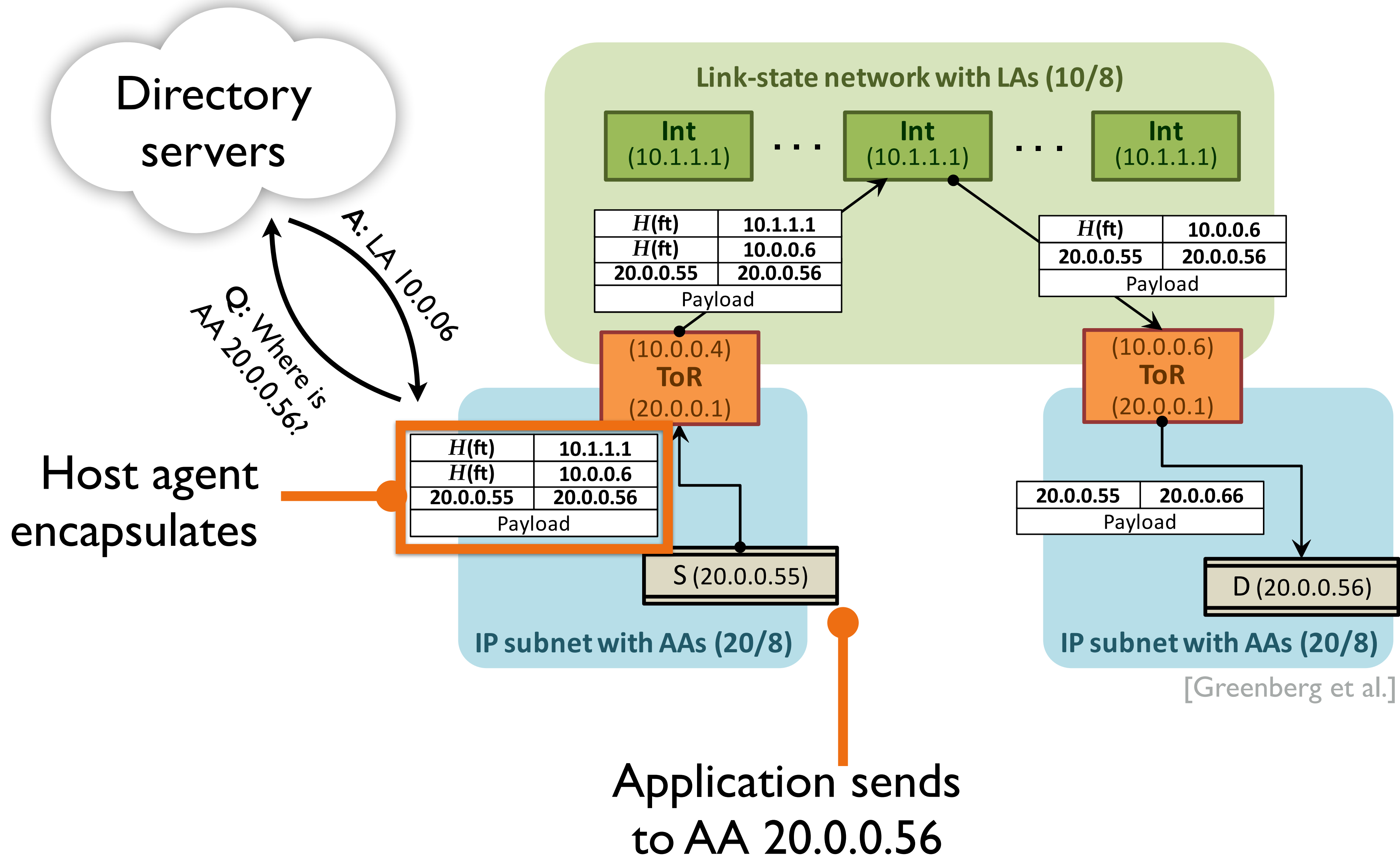


# End-to-end example

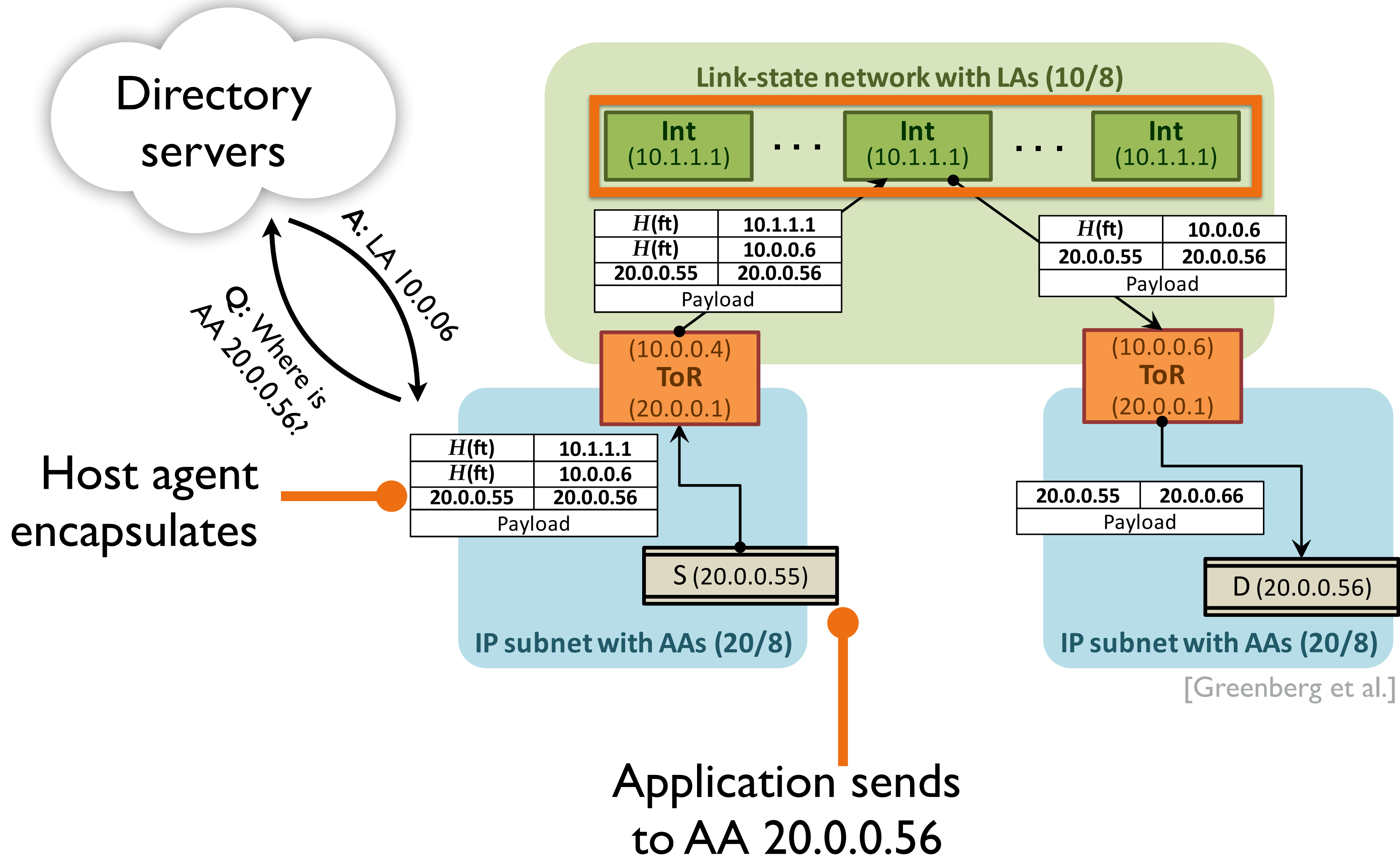


Application sends  
to AA 20.0.0.56

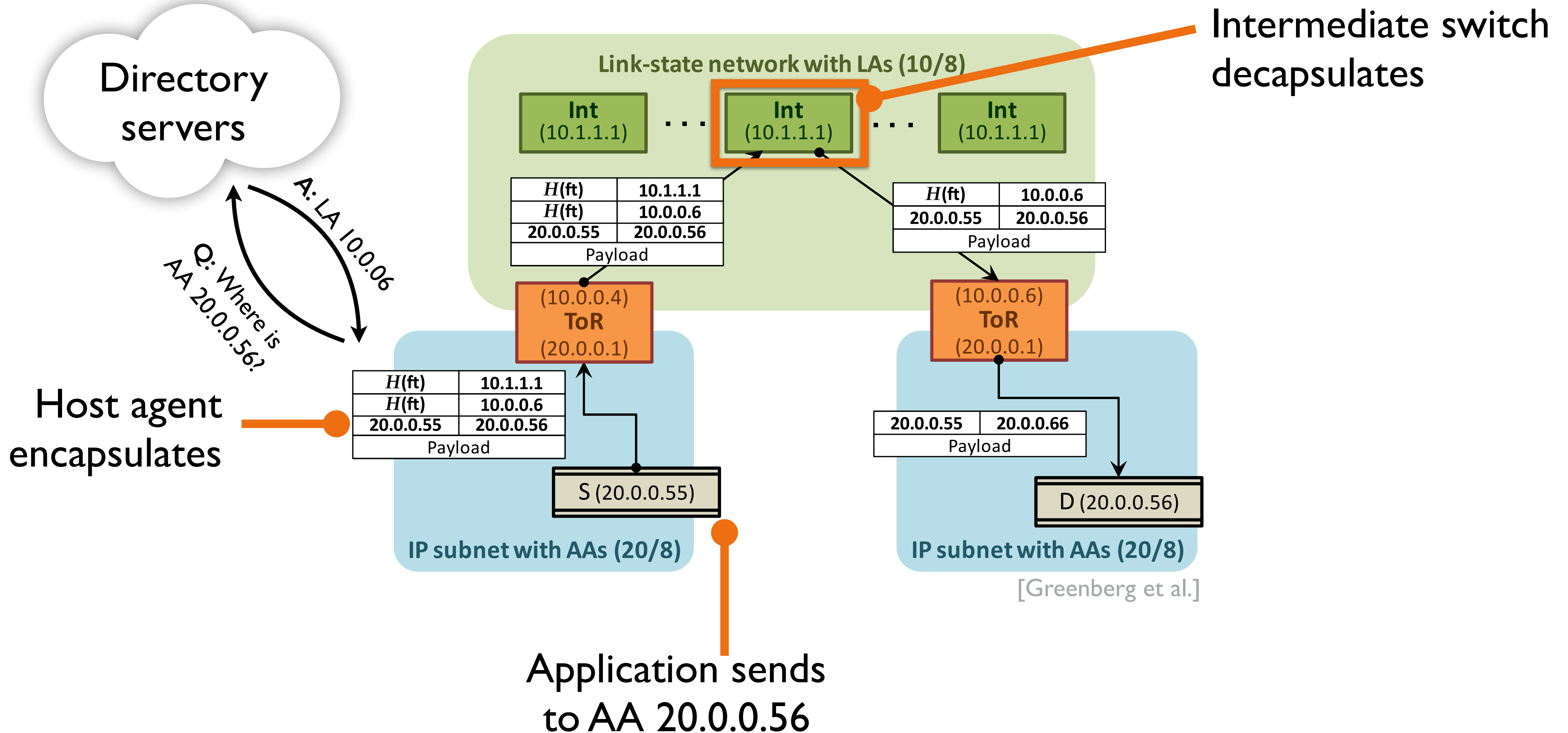
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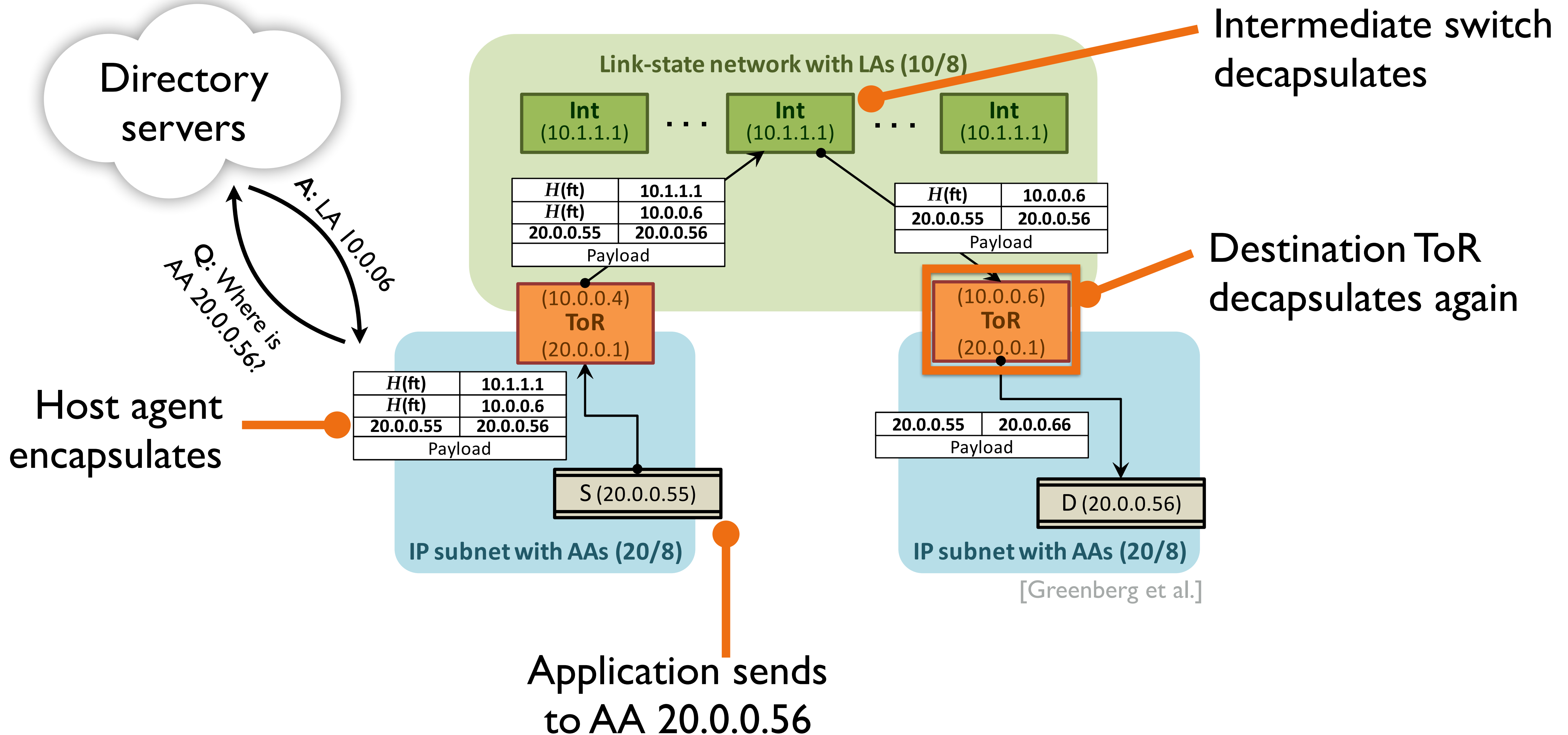
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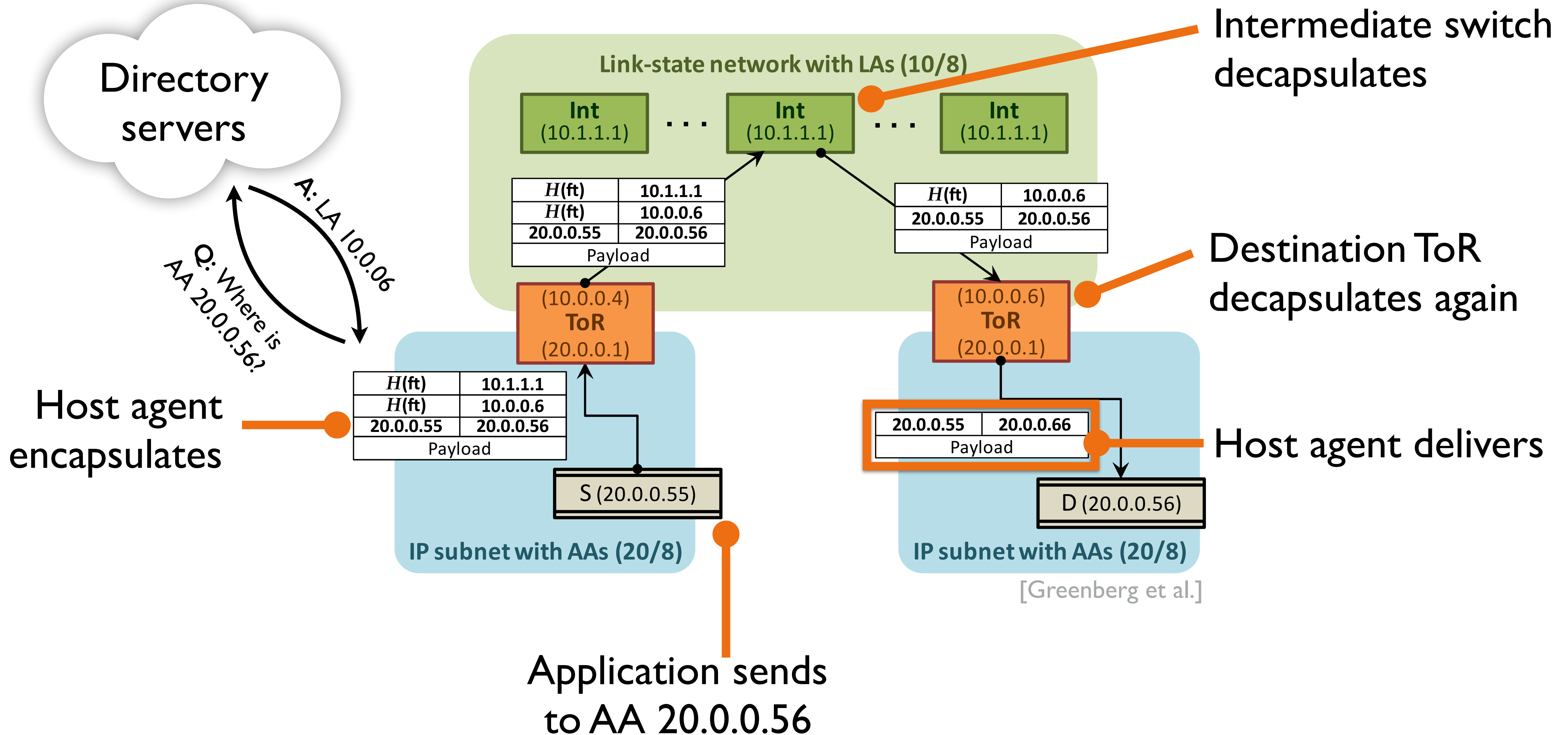


# End-to-end example





# End-to-end example



# Did we achieve agility?

## Location independent addressing

- AAs are location independent

## L2 network semantics

- Agent intercepts and handles L2 broadcast, multicast
- Both of the above require “layer 2.5” shim agent running on host; but, concept transfers to hypervisor-based virtual switch

# Did we achieve agility?

## Performance uniformity

- Clos network is nonblocking (non-oversubscribed)
- Uniform capacity everywhere
- ECMP provides good (though not perfect) load balancing
- But, performance isolation among tenants depends on TCP backing off to rate destination can receive
- Leaves open the possibility of fast load balancing

## Security

- Directory system can allow/deny connections by choosing whether to resolve an AA to a LA
- But, segmentation not explicitly enforced at hosts

# Where's the SDN?

Directory servers: Logically centralized control

- Orchestrate application locations
- Control communication policy

Host agents: dynamic “programming” of data path