

System Programming

PC Assembly Language

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Topics

PC Assembly Language

- x86 Processors
- Instructions
- Directives
- System Calls

Assembly and C

- Subroutines
- Calling Conventions
- C from Assembly
- Assembly from C

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x86 Processors

- ▶ very similar from the programming standpoint
- ▶ 8086: 16-bit processor, real mode
- ▶ 80386: 32-bit processor, protected mode

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Segments

- ▶ **code** segment: read-only parts
 - ▶ instructions
 - ▶ constants
- ▶ **data** segment
 - ▶ initialized data
- ▶ **bss** segment
 - ▶ uninitialized data
- ▶ **stack** segment

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8086 Registers

- ▶ 4 general purpose data registers
- ▶ 2 index registers
- ▶ 2 pointer registers
- ▶ 4 segment registers
- ▶ 2 control registers

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Data Registers

- ▶ AX: accumulator register
- ▶ BX: base register
 - ▶ used to address data in memory
- ▶ CX: counter register
 - ▶ used as repetition counter in loop operations
- ▶ DX: data register
 - ▶ used in multiplication and division operations
- ▶ high and low halves can be accessed as 8-bit registers: AH-AL, BH-BL, CH-CL, DH-DL

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Index and Pointer Registers

- ▶ index registers:
 - ▶ DI: data index
 - ▶ SI: stack index
 - ▶ they can be used like general purpose registers
- ▶ pointer registers:
 - ▶ BP: base pointer
 - ▶ SP: stack pointer

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Segment Registers

- ▶ CS: code segment register
- ▶ DS: data segment register
- ▶ SS: stack segment register
- ▶ ES: extra segment register

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Control Registers

- ▶ IP: instruction pointer
 - ▶ CS + IP: address of next instruction
- ▶ FLAGS: status conditions
 - ▶ ZF (zero), OF (overflow), SF (sign), CF (carry), PF (parity)

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80386

- ▶ 32 bit registers:
EAX EBX ECX EDX ESI EDI EBP ESP
EIP
- ▶ AX, BX, ..., BP, SP are still valid (lower 16 bits)
- ▶ AH, AL, ..., DH, DL are still valid

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Operand Types

- ▶ register
- ▶ memory
 - ▶ offset from beginning of segment
- ▶ immediate
 - ▶ listed in the instruction itself
- ▶ implied
 - ▶ not explicitly specified

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Basic Instructions

<code>mov dest, src</code>	move src to dest
<code>add dest, src</code>	add src to dest
<code>adc dest, src</code>	add src to dest with carry
<code>sub dest, src</code>	subtract src from dest
<code>sbb dest, src</code>	subtract src from dest with borrow
<code>inc dest</code>	increment dest
<code>dec dest</code>	decrement dest
<code>mul src</code>	multiply eax with src, result in edx:eax
<code>div src</code>	divide edx:eax by src, result in eax and edx

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Bitwise Instructions

<code>not dest</code>	bitwise not (one's complement)
<code>and dest, src</code>	bitwise and
<code>or dest, src</code>	bitwise or
<code>xor dest, src</code>	bitwise xor
<code>neg dest</code>	negate (two's complement)
<code>shl dest, amount</code>	logical shift left
<code>shr dest, amount</code>	logical shift right
<code>asl dest, amount</code>	arithmetic shift left
<code>asr dest, amount</code>	arithmetic shift right
<code>rol dest, amount</code>	rotate left
<code>ror dest, amount</code>	rotate right
<code>rcl dest, amount</code>	rotate left with carry
<code>rcr dest, amount</code>	rotate right with carry

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Branching Instructions

<code>jmp</code>	unconditional
<code>jz</code>	if ZF is set
<code>jnz</code>	if ZF is unset
<code>jo</code>	if OF is set
<code>jno</code>	if OF is unset
<code>js</code>	if SF is set
<code>jns</code>	if SF is unset
<code>jc</code>	if CF is set
<code>jnc</code>	if CF is unset
<code>jp</code>	if PF is set
<code>jnp</code>	if PF is unset

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Branching Instructions

- `cmp vleft, vright`: compare vleft and vright

condition	signed	unsigned
vleft = vright	<code>je</code>	<code>je</code>
vleft \neq vright	<code>jne</code>	<code>jne</code>
vleft < vright	<code>jl</code>	<code>jb</code>
vleft \nless vright	<code>jnl</code>	<code>jnb</code>
vleft \leq vright	<code>jle</code>	<code>jbe</code>
vleft \nless vright	<code>jnle</code>	<code>jnbe</code>
vleft > vright	<code>jg</code>	<code>ja</code>
vleft \nless vright	<code>jng</code>	<code>jna</code>
vleft \geq vright	<code>jge</code>	<code>jae</code>
vleft \nless vright	<code>jnge</code>	<code>jnae</code>

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Directives

- needed by the assembler
- not part of the instruction set
- labels
 - mark points in code and data
 - entry labels have to be marked `global`
- segments
- data definition
- named constants: `equ`
 - no memory allocated

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Code Template

```
segment .data
; initialized data definitions

segment .bss
; uninitialized data definitions

segment .text
global _start

_start:
; entry point
```

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Data Definition

type	initialized	uninitialized
byte	<code>db</code>	<code>resb</code>
word	<code>dw</code>	<code>resw</code>
dword	<code>dd</code>	<code>resd</code>
qword	<code>dq</code>	<code>resq</code>
tword	<code>dt</code>	<code>rest</code>

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Data Definition Examples

Example

```
L1 db 0
L2 dw 1000
L3 dd 1A92h
L4 db 0, 1, 2, 3
L5 db "w", "o", "r", "d", 0
L6 db 'word', 0
L7 times 100 db 0
L8 resb 1
L9 resw 100
```

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Dereferencing

- ▶ plain label:
address of memory

Example

```
mov eax, L1
```

- ▶ label in brackets:
contents of memory

Example

```
mov eax, [L1]
```

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System Calls

- ▶ system calls are implemented using software interrupt 80h

system call setup

```
eax ← system call number
ebx ← first argument
ecx ← second argument
edx ← third argument
int 80h
```

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System Call Examples

- ▶ exit system call number: 1
- ▶ arg. 1: return status
 - ▶ 0: success, 1: failure
- ▶ read system call number: 3
- ▶ arg. 1: input descriptor
 - ▶ 0: stdin, 1: stdout, 2: stderr
- ▶ arg. 2: start of input buffer
- ▶ arg. 3: length of input
- ▶ write system call number: 4
- ▶ arg. 1: output descriptor
 - ▶ 0: stdout, 1: stderr, 2: stderr
- ▶ arg. 2: start of output buffer
- ▶ arg. 3: length of output

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Example: Hello, world!

```
segment .data
msg db "Hello, world!", 10
len equ 14

segment .text
global _start

_start:
    mov eax, 4
    mov ebx, 1
    mov ecx, msg
    mov edx, len
    int 80h

    mov eax, 1
    mov ebx, 0
    int 80h
```

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References

Required Reading: Carter

- ▶ Chapter 1: Introduction
 - ▶ 1.2. Computer Organization
 - ▶ 1.3. Assembly Language

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Stack

- ▶ the stack is accessed in 4-byte units

push

push operand

- ▶ subtract 4 from esp
- ▶ store operand to address [esp]

pop

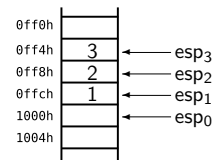
pop register

- ▶ store operand at address [esp] to register
- ▶ add 4 to esp

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Stack Example

Example



```
push dword 1
push dword 2
push dword 3
pop  eax
pop  ebx
pop  ecx
```

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Subroutine Call

call

call target

- ▶ push address of next instruction
- ▶ jump to target

ret

ret

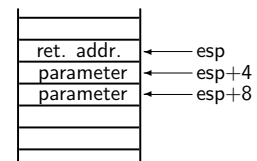
- ▶ pop return address
- ▶ jump to return address

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Stack Parameters

- ▶ called subroutine does not pop parameters
- ▶ accesses parameters on the stack

stack layout

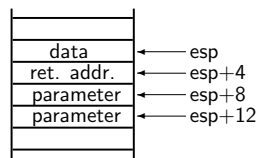


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Accessing Parameters

- ▶ offsets from esp may change

Example (after a push)



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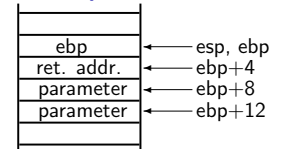
Accessing Parameters

- ▶ use ebp

subroutine template

```
push ebp
mov  ebp, esp
...
pop  ebp
ret
```

stack layout



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Example: Factorial

```
segment .bss
f resd 1

segment .text
fact:
    push ebp
    mov  ebp, esp

    mov  dword [f], 1
    mov  ecx, [ebp+8]

back:
    mov  eax, [f]
    mul  ecx
    mov  [f], eax
    dec  ecx
    cmp  ecx, 1
    jne  back

    pop  ebp
    ret
```

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Example: Calling Factorial

```
segment .data
k dd 5

segment .bss
f resd 1

segment .text
global _start

_start:
    push ebp
    mov  ebp, esp

    push dword [k]
    call fact
    add  esp, 4

    pop  ebp
    ret

fact:
    ...
```

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Calling Conventions

- ▶ how will parameters be passed?
- ▶ if using stack:
 - ▶ in what order will the parameters be pushed?
 - ▶ who will remove parameters from the stack?
- ▶ how will the result be returned?
- ▶ which registers should remain unchanged?

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C Calling Conventions

- ▶ parameters are passed via the stack
 - ▶ caller pushes parameters in reverse order
 - ▶ caller removes parameters from the stack
- ▶ result is returned over eax
- ▶ ebx, esi, edi, ebp, cs, ds, ss, es should remain unchanged

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Calling C from Assembly

- ▶ to call a C function from Assembly:
- ▶ declare function as **extern**
- ▶ push arguments in reverse order
- ▶ call function
- ▶ adjust esp

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Example: printf

```
segment .data
k dd 5
intf db "%d", 10, 0

segment .bss
f resd 1

segment .text
global main
extern printf

main:
    ...

    push dword [k]
    call fact
    add  esp, 4

    push dword [f]
    push intf
    call printf
    add  esp, 8

    ...

fact:
    ...
```

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C Variables

- ▶ global: in fixed memory locations
- ▶ static: same as global, only scope is different
- ▶ automatic: on stack
- ▶ register: in a register (if possible)
- ▶ volatile: do not optimize

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Automatic Variables

- ▶ allocation is done by subtracting from esp

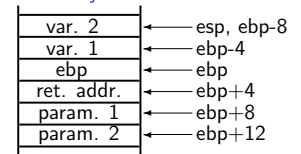
subroutine template

```
push ebp
mov  ebp, esp
sub  esp, N_BYTES

...

mov  esp, ebp
pop  ebp
ret
```

stack layout



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Example: Factorial (C)

```
int y;

void fact(int k)
{
    register int i;

    y = 1;
    for (i = k; i > 1; i--)
        y = y * i;
}
```

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Example: Factorial (C)

```
int fact(int k)
{
    int y;
    register int i;

    y = 1;
    for (i = k; i > 1; i--)
        y = y * i;
    return y;
}
```

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Example: Factorial

```
segment .text
global fact

fact:
    push ebp
    mov  ebp, esp
    sub  esp, 4

    mov  dword [ebp-4], 1
    mov  ecx, [ebp+8]

back:
    mov  eax, [ebp-4]
    mul  ecx
    mov  [ebp-4], eax
    dec  ecx
    cmp  ecx, 1
    jne  back

    mov  eax, [ebp-4]
    mov  esp, ebp
    pop  ebp
    ret
```

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Calling Assembly from C

- ▶ to call an Assembly function from C:
- ▶ in Assembly file: declare function as **global**
- ▶ in C file: declare the prototype

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Example: Factorial

```
int fact(int k);  
  
int main(void)  
{  
    int x, y;  
  
    ...  
    y = fact(x);  
    ...  
}
```

References

Required Reading: Carter

- Chapter 4: **Subprograms**