

# Notes on Mechanism Design

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- This study notes are mainly based on the lecture note written by Valimaki in 2018.

## 1 Single Agent

- One principal v.s. one agent.
- $a \in A$ : allocation,  $\theta \in \Theta$ : agent's private info.  $\theta \sim F(\theta)$ .  $u^P(a, \theta)$ ,  $u^A(a, \theta)$ .
- We often assume quasi-linear payoff functions:
  - $a := (x, t)$ ,  $u^P(a, \theta) := v^P(x, \theta) + t$ ,  $u^A(a, \theta) := v^A(x, \theta) - t$ .
- A mechanism is a pair  $M := (\Sigma, \phi)$ , where  $\Sigma$  is a message space and  $\phi : \Sigma \rightarrow \Delta(A)$ .
- Agent's strategy:  $\sigma : \Theta \rightarrow \Delta(\Sigma)$ . Principal commits to a mechanism  $M$ .
- Consider a social choice function  $\psi : \Theta \rightarrow A$ . We want to know whether  $\psi$  is implementable (, i.e., achievable in equilibrium,) or not.
- As for implementability, we can discuss it focusing only on direct mechanisms, assuming  $\Sigma := \Theta$ , w.l.o.g. (Revelation principle)

### 1.1 Revenue Equivalence

- In §1.1 and §1.2, we assume that the parameter space is a closed interval  $\Theta := [\underline{\theta}, \bar{\theta}] \subseteq \mathbb{R}$ .

#### 1.1.1 Milgrom and Segal (2002), Envelope Theorem

- $\Theta := [\underline{\theta}, \bar{\theta}]$ .  $f(\cdot, \theta) : X \rightarrow \mathbb{R}$ .  $\{f(\cdot, \theta)\}_{\theta \in \Theta}$ .
- $V(\theta) := \max_{x \in X} f(x, \theta)$ .  $X^*(\theta) := \operatorname{argmax}_{x \in X} f(x, \theta)$

**Def. 1.1** (Selection). *A function  $x^* : \Theta \rightarrow X$  is a selection from  $X^*$  if  $x^*(\theta) \in X^*(\theta)$  for all  $\theta \in \Theta$ .*

**Thm. 1.1** (Milgrom and Segal (2002)). *Assume the following:*

- For any  $x \in X$ ,  $f(x, \cdot) : \Theta \rightarrow \mathbb{R}$  is absolutely continuous on  $\Theta$ .
- For any  $x \in X$ ,  $f(x, \cdot) : \Theta \rightarrow \mathbb{R}$  is differentiable on  $\Theta$ .

*Then, the following holds:*

- $V$  is absolutely continuous.
- For any selection  $x^*$  from  $X^*$ ,  $V(\theta) = V(\underline{\theta}) + \int_{\underline{\theta}}^{\theta} f_{\theta}(x^*(s), s) ds$ .

*Proof.* Note that the absolute continuity of  $f(x, \theta)$  implies that  $f_{\theta}(x, \theta) \in L^1(\Theta)$  for any  $x \in X$ .

(i)  $V$  is **absolutely continuous**. It is sufficient to show that  $V$  is Lipschitz continuous. Fix any  $\theta', \theta$ . Since any integrable function is bounded, for any  $x$  there exists  $L > 0$  s.t.  $|f_\theta(x, \theta)| \leq L$  for almost all  $\theta \in \Theta$ .

$$\begin{aligned} |V(\theta') - V(\theta)| &= \left| \max_{x'} f(x', \theta') - \max_x f(x, \theta) \right| \\ &\leq \max_x |f(x, \theta') - f(x, \theta)| = \max_x \left| \int_{\theta'}^{\theta} f_\theta(x, s) ds \right| \\ &\leq L \cdot |\theta' - \theta| \end{aligned}$$

(ii) Fix any selection  $x^*$  from  $X^*$ . By the result of (i),

$$V(\theta) = V(\underline{\theta}) + \int_{\underline{\theta}}^{\theta} V'(s) ds$$

Fix any selection  $x^*$  and  $\theta', \theta$  such that  $\theta' > \theta$ . By the definition of  $V$  and  $x^*$ ,

$$\begin{aligned} V(\theta) &= f(x^*(\theta), \theta) \geq f(x^*(\theta'), \theta) \\ V(\theta') &= f(x^*(\theta'), \theta') \geq f(x^*(\theta), \theta') \end{aligned}$$

Hence,

$$\begin{aligned} V(\theta') - V(\theta) &\leq f(x^*(\theta'), \theta') - f(x^*(\theta'), \theta) \\ \frac{V(\theta') - V(\theta)}{\theta' - \theta} &\leq \frac{f(x^*(\theta'), \theta') - f(x^*(\theta'), \theta)}{\theta' - \theta} \end{aligned}$$

Similarly,

$$\begin{aligned} V(\theta) - V(\theta') &\leq f(x^*(\theta'), \theta) - f(x^*(\theta'), \theta') \\ \frac{V(\theta) - V(\theta')}{\theta - \theta'} &\geq \frac{f(x^*(\theta'), \theta) - f(x^*(\theta'), \theta')}{\theta - \theta'} \end{aligned}$$

Note that by assumption  $f(x, \cdot)$  is differentiable at all  $\theta \in \Theta$ . Therefore, if  $V$  is differentiable at  $\theta$ , we have  $V'(\theta) = f_\theta(x^*(\theta), \theta)$ .  $\square$

### 1.1.2 RET

- Focus on the agent's utility:  $u := u^A$ .
- $A := \phi(\Theta)$ .  $V(\theta) := \max_{a \in A} u(a, \theta)$ .  $A^*(\theta) := \operatorname{argmax}_{a \in A} u(a, \theta)$ .
- Assume that  $u(a, \cdot)$  is absolutely continuous and differentiable on  $\Theta$  for all  $a \in A$ .
- By incentive compatibility,  $\phi(\theta) \in A^*(\theta)$  for all  $\theta \in \Theta$ :  $\phi$  is a selection from  $A^*$ .

**Thm. 1.2** (Revenue Equivalence Theorem).

$$V(\theta) = V(\underline{\theta}) + \int_{\underline{\theta}}^{\theta} u_\theta(\phi(s), s) ds$$

In particular, under quasi-linear utility,

$$\begin{aligned} V(\theta) &= V(\underline{\theta}) + \int_{\underline{\theta}}^{\theta} v_\theta(x(s), s) ds \\ t(\theta) &= v(x(\theta), \theta) - V(\underline{\theta}) - \int_{\underline{\theta}}^{\theta} v_\theta(x(s), s) ds \end{aligned}$$

*Proof.* Milgrom and Segal. As for quasi-linear cases, the results follow from

$$V(\theta) = v(x(\theta), \theta) - t(\theta)$$

$\square$

- RET states that under any IC mechanism, except for the constant  $V(\underline{\theta})$ , the transfer from the agent to the principal is uniquely determined once the allocation rule  $x$  is fixed.

## 1.2 Characterization of IC

### 1.2.1 Monotone Comparative Statics

This subsection is based on the lecture slides by John K.-H. Quah:

<http://www.johnquah.com/lecture-slides.html>

- Consider parameterized optimization problems.
- We often want to know how optimizers and optimal values change according to the changes in parameters.
- comparative statics = Sensitivity analysis
- Implicit function theorem: Not only the direction of changes but also the rate of change. Many assumptions are required.
- Monotone comparative statics: Only the direction of changes. Fewer assumptions.
- $\Theta \subseteq \mathbb{R}$ . Two functions  $g : \Theta \rightarrow \mathbb{R}$  and  $f : \Theta \rightarrow \mathbb{R}$ .

**Def. 1.2** (Single Crossing).  $g$  dominates  $f$  by single crossing property (SCP),  $g \succsim_{SC} f$ , if for all  $x'' > x'$ ,

- $f(x'') - f(x') \geq 0 \implies g(x'') - g(x') \geq 0$
- $f(x'') - f(x') > 0 \implies g(x'') - g(x') > 0$

$\{f(\cdot, \theta)\}_{\theta \in \Theta}$  is an SCP family if

$$\forall \theta'' > \theta'; f(\cdot, \theta'') \succsim_{SC} f(\cdot, \theta')$$

**Def. 1.3** (Increasing Differences).  $g$  dominates  $f$  by increasing differences,  $g \succsim_{IN} f$ , if for all  $x'' > x'$ ,

$$g(x'') - g(x') \geq f(x'') - f(x').$$

$\{f(\cdot, \theta)\}_{\theta \in \Theta}$  satisfies increasing differences if

$$\forall \theta'' > \theta'; f(\cdot, \theta'') \succsim_{IN} f(\cdot, \theta')$$

**Def. 1.4** (Strictly Increasing Differences).  $g$  dominates  $f$  by strictly increasing differences,  $g \succsim_{SID} f$ , if for all  $x'' > x'$ ,

$$g(x'') - g(x') > f(x'') - f(x').$$

$\{f(\cdot, \theta)\}_{\theta \in \Theta}$  satisfies strictly increasing differences (SID) if

$$\forall \theta'' > \theta'; f(\cdot, \theta'') \succ_{SID} f(\cdot, \theta')$$

- Note that  $g \succsim_{IN} f$  implies  $g \succsim_{SC} f$ .

**Thm. 1.3** (Milgrom and Shannon (1994)).  $X \subseteq \mathbb{R}$ .  $f, g : X \rightarrow \mathbb{R}$ .

$$[\forall Y \subseteq X; \operatorname{argmax}_{x \in Y} g(x) \geq \operatorname{argmax}_{x \in Y} f(x)] \iff g \succsim_{SC} f$$

Note that, for  $Y, Z \subseteq \mathbb{R}$ ,

$$Y \geq Z \stackrel{\Delta}{\iff} [y \in Y, z \in Z \implies y \vee x \in Y, y \wedge z \in Z.]$$

Proof. .

$\Rightarrow$ ) We show contrapositive. Suppose that  $g \not\prec_{SC} f$ . There exist  $x'', x'$  such that  $x'' > x'$  and at least one of the following holds:

$$f(x'') \geq f(x'), g(x'') < g(x') \quad (1)$$

or

$$f(x'') > f(x'), g(x'') \leq g(x') \quad (2)$$

Let  $Y := \{x', x''\}$ ,  $G_Y := \operatorname{argmax}_{x \in Y} g(x)$  and  $F_Y := \operatorname{argmax}_{x \in Y} f(x)$ . In case of (1),  $x' \vee x'' \notin G_Y$ . In case of (2),  $x' \wedge x'' \notin F_Y$ .

$\Leftarrow$ ) Fix any  $Y \subseteq X$  and  $x'', x' \in Y$  such that  $x' \in G_Y$  and  $x'' \in F_Y$ . We need to show that  $x' \vee x'' \in G_Y$  and  $x' \wedge x'' \in F_Y$ . First, since  $x'' \in F_Y$ , we have  $f(x'') \geq f(x')$ . By assumption,  $g(x'') \geq g(x')$ . Since  $x' \in G_Y$ , we have  $x'' \in G_Y$  and  $x' \vee x'' \in G_Y$ .

Next, we show  $f(x'') = f(x')$ . Note that this implies that  $x' \wedge x'' \in F_Y$ . Suppose toward contradiction that  $f(x'') > f(x')$ . Then, since  $g \succsim_{SC} f$ , we have  $g(x'') > g(x')$ . This contradicts  $x' \in G_Y$ .  $\square$

### 1.2.2 Characterization of IC

- Consider quasi-linear utility cases. Assume that  $v(x, \theta)$  is absolutely continuous and differentiable on  $\Theta$  for all  $x$ .

**Lem. 1.1.** Let  $V(\theta) := v(x(\theta), \theta) - t(\theta)$ . If a mechanism  $(x, t)$  is IC, then

$$V(\theta) = V(\underline{\theta}) + \int_{\underline{\theta}}^{\theta} v_{\theta}(x(s), s) ds \quad (\text{LIC})$$

*Proof.* RET.  $\square$

**Lem. 1.2.** If a mechanism  $(x, t)$  is IC and  $\{v(\cdot, \theta)\}_{\theta \in \Theta}$  satisfies SID, then

$$x(\theta) \text{ is non-decreasing in } \theta. \quad (\text{M})$$

*Proof.* Fix  $\theta'', \theta'$  such that  $\theta'' > \theta'$ . Since  $\{v(\cdot, \theta)\}_{\theta \in \Theta}$  satisfies SID,  $v(\cdot, \theta'') \succsim_{SID} v(\cdot, \theta')$ . Suppose toward contradiction that  $x(\theta'') < x(\theta')$ . Since  $v(\cdot, \theta'') \succsim_{SID} v(\cdot, \theta')$ ,

$$v(x(\theta'), \theta'') - v(x(\theta''), \theta'') > v(x(\theta'), \theta') - v(x(\theta''), \theta') \geq 0$$

This violates IC. A contradiction.  $\square$

- The lemmas above shows that, assuming  $\{v(\cdot, \theta)\}_{\theta \in \Theta}$  satisfies SID, IC of  $(x, t)$  implies (LIC) and (M).
- We can show that the converse also holds.

**Lem. 1.3.** Assume that  $\{v(\cdot, \theta)\}_{\theta \in \Theta}$  satisfies SID. If the conditions (LIC) and (M) hold, then  $(x, t)$  is IC.

*Proof.* Fix any  $\theta, \theta'$ . We need to show that  $v(x(\theta), \theta) - t(\theta) \geq v(x(\theta'), \theta) - t(\theta')$ . Note that, by (LIC), we have

$$t(\theta) = v(x(\theta), \theta) - V(\underline{\theta}) - \int_{\underline{\theta}}^{\theta} v_{\theta}(x(s), s) ds$$

Then,

$$\begin{aligned} & [v(x(\theta), \theta) - t(\theta)] - [v(x(\theta'), \theta) - t(\theta')] \\ &= [v(x(\theta), \theta) - t(\theta)] - [v(x(\theta'), \theta) + v(x(\theta'), \theta') - v(x(\theta'), \theta') - t(\theta')] \\ &= \int_{\underline{\theta}}^{\theta} v_{\theta}(x(s), s) ds - \int_{\underline{\theta}}^{\theta'} v_{\theta}(x(s), s) ds - [v(x(\theta'), \theta) - v(x(\theta'), \theta')] \\ &= \int_{\theta'}^{\theta} v_{\theta}(x(s), s) ds - \int_{\theta'}^{\theta} v_{\theta}(x(s), \theta') ds = \int_{\theta'}^{\theta} [v_{\theta}(x(s), s) - v_{\theta}(x(s), \theta')] ds \geq 0 \end{aligned}$$

$\square$

**Thm. 1.4** (Characterization of IC). Assume that  $\{v(\cdot, \theta)\}_{\theta}$  satisfies SID. Then,

$$(x, t) \text{ is IC} \iff x \text{ is non-decreasing, and } t \text{ is calculated by (LIC)}$$

### 1.3 General Case: Rochet's Theorem and Cyclical Monotonicity

- Consider quasi-linear utility cases.
- Characterize IC mechanisms.

**Def. 1.5** (weak monotonicity). *An allocation rule  $x : \Theta \rightarrow A$  is weakly monotone if*

$$\forall \theta, \theta'; [v(x(\theta), \theta') - v(x(\theta), \theta)] + [v(x(\theta'), \theta) - v(x(\theta'), \theta')] \leq 0$$

**Prop. 1.1.** *If  $(x, t)$  is IC, then  $x$  is weakly monotone.*

**Def. 1.6** (cyclical monotonicity).

$$S := \{(\theta^1, \dots, \theta^{k+1}) \mid \forall i \in [k+1]; \theta^i \in \Theta, \theta^1 = \theta^{k+1}, k \in \mathbb{Z}^+\}$$

*An allocation rule  $x$  is cyclically monotone if, for any  $(\theta^1, \dots, \theta^{k+1}) \in S$ ,*

$$\sum_{i=1}^k [v(x^i, \theta^{i+1}) - v(x^i, \theta^i)] \leq 0, \text{ where } x^i := x(\theta^i) \quad (\text{CM})$$

**Thm. 1.5** (Rochet (1987)).

$$\exists t; (x, t) : \text{IC} \iff x \text{ is cyclically monotone.}$$

*Proof.* .

$\Rightarrow$ ) Easy.

$\Leftarrow$ ) Fix  $\theta_0 \in \Theta$ .

$$S(\theta) := \{(\theta^1, \dots, \theta^{k+1}) \mid \forall i \in [k+1]; \theta^i \in \Theta, \theta^1 = \theta_0, \theta^{k+1} = \theta, k \in \mathbb{Z}^+\}$$

$$V(\theta) := \sup_{(\theta^1, \dots, \theta^{k+1}) \in S(\theta)} \sum_{i=1}^k [v(x^i, \theta^{i+1}) - v(x^i, \theta^i)]$$

**(i)  $[V(\theta_0) = 0.]$**  By CM,  $V(\theta_0) \leq 0$ . Considering the case where  $k := 1$ , we see that  $(\theta_0, \theta_0) \in S(\theta_0)$  satisfies  $[v(x^1, \theta^2) - v(x^1, \theta^1)] = 0$ . Therefore,  $V(\theta_0) = 0$ .

**(ii)  $[V(\theta) < \infty \text{ for all } \theta \in \Theta.]$**  Fix any  $(\theta^1, \dots, \theta^{k+1}) \in S(\theta)$ .

$$\begin{aligned} 0 = V(\theta_0) &\geq \sum_{i=1}^k [v(x^i, \theta^{i+1}) - v(x^i, \theta^i)] + [v(x^{k+1}, \theta_0) - v(x^{k+1}, \theta^{k+1})] \\ &= \sum_{i=1}^k [v(x^i, \theta^{i+1}) - v(x^i, \theta^i)] + [v(x(\theta), \theta_0) - v(x(\theta), \theta)] \\ &\therefore \sum_{i=1}^k [v(x^i, \theta^{i+1}) - v(x^i, \theta^i)] \leq v(x(\theta), \theta) - v(x(\theta), \theta_0) \\ &\therefore V(\theta) \leq v(x(\theta), \theta) - v(x(\theta), \theta_0) \end{aligned}$$

**(iii) [Construct the transfer rule]** Fix any  $\theta, \theta'$ . By the same argument as in (ii), we can show that

$$V(\theta) \geq V(\theta') + v(x(\theta'), \theta) - v(x(\theta'), \theta')$$

Define  $t(\theta) := v(x(\theta), \theta) - V(\theta)$ . With this  $t$ , a mechanism  $(x, t)$  satisfies IC:

$$v(x(\theta), \theta) - t(\theta) - (v(x(\theta'), \theta) - t(\theta')) = V(\theta) - V(\theta') - v(x(\theta'), \theta) + v(x(\theta'), \theta') \geq 0$$

□