Laboratory 1

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September 16, 2020

3.3.1 Periodic tasks in Ada

Format:

$$\tau(\phi, T, C, D)$$

Tasks:

$$\tau_1(T_1 = 900, 300)$$

Task 2,3 and 5 are released at t = 1.000, task 2 finishes first, then task 6 so:

$$P_2 > P_6 > P_5$$

Task 3 is released at t = 1.200, pre-empting task 5.

$$P_3 > P_5$$

Task 2,6,3 and 5 seem to have a higher priority than task 4 since it finishes at t = 1.396 despite being scheduled to start at t = 0.900.

$$P_2 > P_6 > P_5 > P_4$$

Tasks 3 and 2 are scheduled to start at the same time at t = 2.000. Obviously task 3 has a higher priority than task 2.

$$P_3 > P_2 > P_6 > P_5 > P_4$$

Tasks 2 and 1 are scheduled to be released at the same time at t = 3.500. Task 2 seems to have higher priority since task 1 has to wait. Also task 1 has to wait for task 3 which starts at t = 3.600.

$$P_2 > P_1$$

$$P_3 > P_1$$

Task 1 and 6 are both scheduled to start at t = 6.200. In this case, task 1 starts first due to higher priority.

$$P_1 > P_6$$

Which leads to:

$$P_3 > P_2 > P_1 > P_6 > P_5 > P_4$$

3.3.2 The rate monotonic schedule.

1) Calculating utilization

The utilization of a periodic task $\tau_i(T_i, C_i)$ is defined as:

$$u_i = \frac{C_i}{T_i}$$

The total utilization for a set of tasks is defined as:

$$U = \sum_{i=1}^{n} u_i$$

The 3 tasks given are:

$$\tau_i(\phi_i, T_i, C_i, D_i)$$

$$\tau_1(100, 300, 100, 300)$$

$$\tau_2(100, 400, 100, 400)$$

$$\tau_3(100, 600, 100, 600)$$

Their respective utilizations are:

$$u_1 = \frac{C_1}{D_1} = \frac{100}{300} = \frac{1}{3}$$

$$u_2 = \frac{C_2}{D_2} = \frac{100}{400} = \frac{1}{4}$$

$$u_3 = \frac{C_3}{D_3} = \frac{100}{600} = \frac{1}{6}$$

The total utilization U can thus be calculated:

$$U = \frac{1}{3} + \frac{1}{4} + \frac{1}{6} = \frac{9}{12} = \frac{3}{4} = 0.75$$

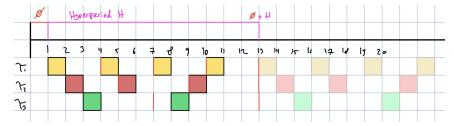
The hyperperiod H is defined as the least common multiplier from the set of task periods:

$$H = lcm(T_1, T_2, T_3) = lcm(300, 400, 600) = 1200$$

According to RMS the priorities should therefore be:

$$P_1 > P_2 > P_3$$

For one hyperperiod of 1200ms the schedule should look like this:

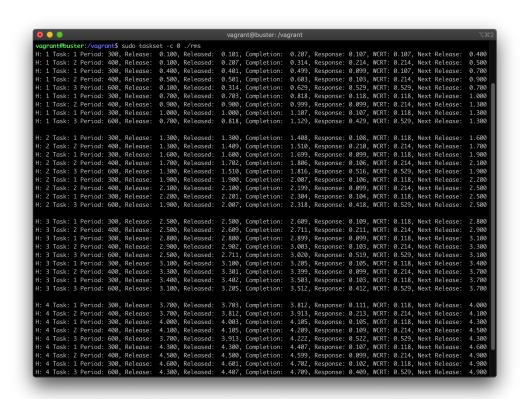


2) Calibrating execution time

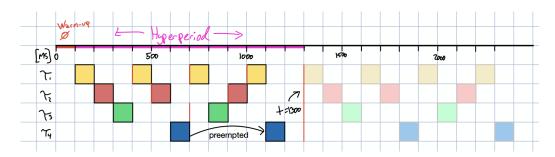
Constraining the execution of the program to use only a single CPU core using taskset was made possible on a macOS system by using a virtual machine. The calibration parameter was set to **1208** which produced the following output:

3) Implementing the periodic task set Γ_1

Running the rms on a single core produces the following output:



4) Adding the additional task τ_4



RMS2 output:

