Week 7 Tutorial 7

(7.1) Least square solution of a linear system: We wish to solve a system of linear equations

$$Ax = b$$
,

where $A \in \mathbb{R}^{m \times n}$ with rk A = m, $x \in \mathbb{R}^n$ and $b \in \mathbb{R}^m$. In many applications

- it is often not possible to calculate the inverse of A,
- or we have more equations than unknowns, that is m > n, and therefore can not expect a unique solution.

In the later case one can calculate the least square solution, that is the vector x that minimises the squared Euclidean: distance

$$\min_{m{x}} \frac{1}{2} \| m{A} m{x} - m{b} \|^2$$

(we multiply with $\frac{1}{2}$ to make the gradient nicer). Show that

1. the gradient of $f = \|\mathbf{A}\mathbf{x} - \mathbf{b}\|^2$ is given by

$$\nabla_{\boldsymbol{x}} f = \boldsymbol{A}^T (\boldsymbol{A} \boldsymbol{x} - \boldsymbol{b})$$

2. Calculate the critical point

$$\boldsymbol{x}^* = (\boldsymbol{A}^T \boldsymbol{A})^{-1} \boldsymbol{A}^T \boldsymbol{b}.$$

Is this point unique? Explain why $A^T A$ is invertible.

Solution: We can write the squared Euclidean distance as

$$f(\boldsymbol{x}) = \frac{1}{2}(\boldsymbol{A}\boldsymbol{x} - \boldsymbol{b})^T(\boldsymbol{A}\boldsymbol{x} - \boldsymbol{b}) = \frac{1}{2}\left(\boldsymbol{x}^T\boldsymbol{A}^T\boldsymbol{A}\boldsymbol{x} - 2\boldsymbol{x}^T\boldsymbol{A}^T\boldsymbol{b} + \boldsymbol{b}^T\boldsymbol{b}\right)$$

We note that A^TA is always symmetric, therefore we can use the example 2.5 in the lecture notes and conclude that

$$\nabla f = \boldsymbol{A}^T (\boldsymbol{A}\boldsymbol{x} - \boldsymbol{b})$$

The critical point can be found by setting $\nabla f = \mathbf{0}$, giving the stated solution.

The solution is unique since the problem is quadratic (we see that the Hessian is the matrix A^TA which is always positive definite. Furthermore A^TA is a square matrix in $\mathbb{R}^{m\times m}$ which has rank m. This follows that if $A^TAx = 0$ then

$$0 = \boldsymbol{x}^T \boldsymbol{A}^T \boldsymbol{A} \boldsymbol{x} = (\boldsymbol{A} \boldsymbol{x})^T (\boldsymbol{A} \boldsymbol{x}) = \langle \boldsymbol{A} \boldsymbol{x}, \boldsymbol{A} \boldsymbol{x} \rangle = \|\boldsymbol{A} \boldsymbol{x}\|^2.$$

Which implies that Ax = 0 and since A has full rank, we conclude that x = 0.

(7.2) Norms: Let $f:[0,1]\to\mathbb{R}$ be continuous. Show that $\|f\|_{\infty}=\sup_x |f(x)|$ is indeed a norm.

Solution:

- 1. If $||f||_{\infty} = 0$, then $\sup |f(x)| = 0$ which implies that $f \equiv 0$.
- 2. $\|\alpha f\|_{\infty} = \sup |\alpha f(x)| = |\alpha| \sup |f| = |\alpha| \|f\|_{\infty}$.
- 3. $||f+g||_{\infty} = \sup_{x} |f(x)+g(x)| \le \sup_{x} |f(x)| + |g(x)| = ||f||_{\infty} + ||g||_{\infty}$.
- (7.3) Interpolation I: Given n+1 data points (x_i, y_i) with $x_i, y_i \in \mathbb{R}$ we wish to find a polynomial, which fits these points exactly.

One can for example consider a polynomial of degree n (then we have the same number of equations and unknowns):

$$p_n(x) = a_0 + a_1 x + a_2 x^2 + \dots a_n x^n.$$

Its coefficients are determined by setting $y_i = p(x_i)$ for $i = 1, \dots n + 1$.

1. Show that the coefficients $\mathbf{a} = (a_0, a_1, \dots a_n)$ can be found y solving the system Va = y, where V is the so called Vandermonde matrix

$$V = \begin{pmatrix} 1 & x_1 & x_1^2 & x_1^3 & \dots & x_1^n \\ 1 & x_1 & x_2^2 & x_2^3 & \dots & x_2^n \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 1 & x_n & x_n^2 & x_n^3 & \dots & x_n^n \end{pmatrix}$$

and $y = (y_1, ..., y_n)$.

2. The determinant of V is given by det $V = \prod_{0 \le i,j \le n} (x_j - x_i)$. Why can this become problematic?

Solution: Since we require that $p_n(x_i) = y_i$ we obtain n equation of the form

$$a_0 + a_1 x_i + a_2 x_i^2 + \dots + a_n x_i^2 = y_i$$

for $i = 1, \dots n$. These equations can be written exactly in the above form.

The condition of the Vandermonde matrix may deteriorate for a large number of points. In particular if points get very close, the determinant can quite small. Furthermore finding the interpolation polynomial requires that the solution of a linear system, which is also computationally costly.

(7.4) Interpolation II Lagrangian interpolation is an attractive alternative to the Vandermonde interpolation, because it does not involve solving a system to find the interpolating polynomial. It is based on writing the polynomial in a different way, for example instead of writing $y = x^2 - 3x + 2$ we can also write y = (x - 2)(x - 3).

Given a data set $(x_i, y_i) \in \mathbb{R}^2$, with $0 \le i \le n$. Then the Lagrange basis polynomials are given by

$$\ell_i(x) = \frac{\prod_{k \neq i} (x - x_k)}{\prod_{k \neq i} (x_i - x_k)} \quad i = 0, \dots n$$

· Show that

$$\ell_i(x_j) = \begin{cases} 1 & j = i \\ 0 & j \neq i, \end{cases}$$

and that the Lagrange interpolating polynomial through those data points

$$p_n(x) = \sum_{k=0}^{n} y_k \ell_k(x)$$

satisfies $p_n(x_i) = y_i$ for every $i = 0, \dots n$.

• Consider the function $f(x)=\frac{1}{x}$. Given the function values at $x_0=2$, $x_1=2.5$ and $x_2=4$, that is given the data pairs $(2,\frac{1}{2})$, $(\frac{5}{2},\frac{2}{5})$ and $(4,\frac{1}{4})$, calculate the corresponding interpolating Lagrange basis polynomials $\{\ell_0,\ell_1,\ell_2\}$ and the interpolating Lagrange polynomial

$$p_2(x) = \sum_{k=0}^{2} y_k \ell_k(x).$$

Sketch the solution and the function f.

Solution: For the basis polynomials we see that the numerator and the denominator are the same for i=j and that for $i\neq j$ the terms in the numerator are 0, showing the first statement. This implies the second statement (since ℓ either takes the value 0 or 1). The Lagrange basis polynomials are

$$\ell_0(x) = \frac{(x - x_1)(x - x_2)}{(x_0 - x_1)(x_0 - x_2)} = (x - 6.5)x + 10$$

$$\ell_1(x) = \frac{(x - x_0)(x - x_2)}{(x_1 - x_0)(x_1 - x_2)} = \frac{1}{3}((-4x + 24)x - 32)$$

$$\ell_2(x) = \frac{1}{3}((x - 4.5)x + 5)$$

The respective Lagrange interpolating polynomial is then

$$p_2(x) = \sum_{k=0}^{2} y_k \ell_k(x) = (0.05x - 0.425)x + 1.15.$$

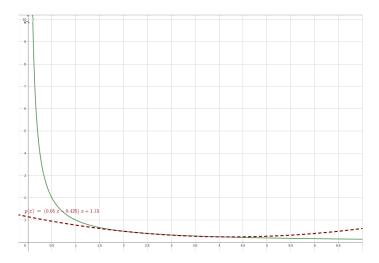


Figure 1: Lagrange interpolating polynomial vs. function f