

# Divided and Conquer

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# Divided and Conquer

- Divide the problem into sub-problems
- Conquer each sub-problem separately
- Solve the original problem by combining solutions for sub-problems

# Merge Sort

$\text{MergeSort}(A = [a_1, \dots, a_n])$

- If  $n = 1$  Then Return  $A$
- Let  $r \leftarrow \lfloor n/2 \rfloor$
- Let  $L \leftarrow [a_1, \dots, a_r]$  and  $R \leftarrow [a_{r+1}, \dots, a_n]$
- $L' \leftarrow \text{MergeSort}(L)$
- $R' \leftarrow \text{MergeSort}(R)$
- Return  $\text{Merge}(L', R')$

# Merge Arrays

Merge([ $a_1, \dots, a_\ell$ ], [ $b_1, \dots, b_r$ ])

- Set  $C \leftarrow$  empty array of length  $\ell + r$
- Set  $i \leftarrow 1$  and  $j \leftarrow 1$
- For  $k = 1$  to  $\ell + r$ 
  - If ( $a_i \leq b_j$ ) and ( $i \leq \ell$ ) Then  $C_k \leftarrow a_i$ ;  $i \leftarrow i + 1$ ;
  - Else  $C_k \leftarrow b_j$ ;  $j \leftarrow j + 1$ ;
- EndFor
- Return  $C$

# Analysis of Merge

## Lemma

Assume that  $A = [a_1, \dots, a_\ell]$  and  $B = [b_1, \dots, b_r]$  are sorted arrays. Then the output of  $\text{Merge}(A, B)$  is a sorted array of  $(A, B)$ .

The running time of Merge is  $O(\ell + r)$ .

# Correctness of MergeSort

## Theorem

*MergeSort( $A$ ) outputs a sorted array with the same elements as  $A$*

## Running Time of MergeSort

Let  $T(n)$  be the worst-case running time of Mergesort on arrays of size  $n$

An observation:  $T(n) \leq T(\lfloor n/2 \rfloor) + T(n - \lfloor n/2 \rfloor) + O(n)$

## Running Time of MergeSort

We assume that  $n = 2^k$  for now. There is a constant  $C > 1$  such that:

- $T(1) \leq C$
- $T(n) \leq 2 \cdot T(n/2) + C \cdot n$

# Solving the Recurrence

## Recurrence for General Parameters

$$f(n) = a \cdot T(n/b) + f(n)$$

# Master Theorem

## Theorem

Let  $a \geq 1, b > 1$ , and let  $T(n)$  defined by  $T(1) = \Theta(1)$  and  $T(n) = a \cdot T(n/b) + f(n)$ :

- If  $f(n) = \Theta(n^{\log_b a})$ , then  $T(n) = O(n^{\log_b a} \cdot \log n)$ ;
- If there is an  $\epsilon > 0$  with  $f(n) = O(n^{\log_b a - \epsilon})$ , then  $T(n) = O(n^{\log_b a})$ .

## Examples

- MergeSort:  $a, b = 2$  and  $f(n) = \Theta(n)$
- $T(n) = 9 \cdot T(n/3) + n \cdot \log n$

# Binary Search

## Input

- A sorted array  $A = [a_1, \dots, a_n]$
- A number  $b$

## Output

Determine whether  $b \in A$ .

Thanks!