# Eigengames

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#### PCA: previous approaches

- PCA task:  $\min_{\mathbf{U} \in \mathbb{R}^{d \times k}, \mathbf{U}^T \mathbf{U} = \mathbf{I}} \| \mathbf{U} \mathbf{U}^T \mathbf{X} \mathbf{X} \|_F^2$
- Full SVD of sample covariance matrix:

$$\mathcal{O}(\min\{nd^2, n^2d\})$$

- Streaming k-PCA:
  - Frequent Directions
    - Approximates the top-k subspace
    - Three steps: copy the row of A, rotate, shrink (to create a zero row)
  - Oja's algorithm
    - Approximates the top-k eigenvectors

$$w_{t+1} - w_t = \eta(w_t^T X_t) X_t; \qquad w_{t+1}^T w_{t+1} = 1,$$

#### **Algorithm 4.1.1** FrequentDirections

```
\begin{aligned} &\textbf{Input:}\ \ell, A \in R^{n \times d} \\ &B \leftarrow 0^{\ell \times d} \\ &\textbf{for}\ i \in 1, \dots, n\ \textbf{do} \\ &B_{\ell,:} \leftarrow A_{i,:} \\ &[U, \Sigma, V] \leftarrow \textbf{svd}(B) \\ &C \leftarrow \Sigma V^T \\ &\delta \leftarrow \Sigma_{\ell,\ell}^2 \\ &B \leftarrow \sqrt{\Sigma^2 - \delta I_\ell} \cdot V^T \\ &\textbf{return}\ B \end{aligned}
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$$||A - \pi_B^k(A)||_F \le (1 + \varepsilon)||A - A_k||_F$$

### Symmetric Generalized Eigenvalue Problem (SGEP)

#### Definition (SGEP)

Given matrices  $A = A^{\top}$ ,  $B = B^{\top}$ ,  $B \succ 0$ ,

$$Av = \lambda Bv, \tag{1}$$

defines Symmetric Generalized Eigenvalue Problem (SGEP)

#### Examples:

- SVD/PCA:  $A = X^{T}X, B = I$
- Graph Laplacian: L Laplacian, A = L, B = I
- CCA, ICA, etc.

# Top-k SGEP Game ( $\Gamma$ -Eigengame)

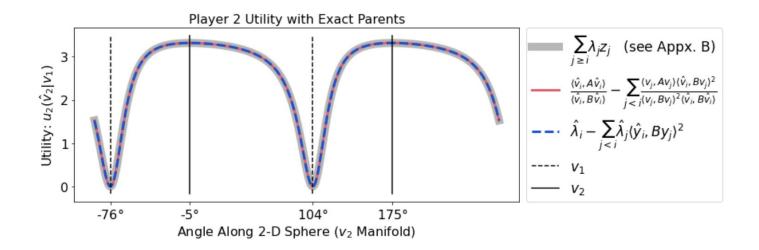
- Players: k players  $i \in \{1, \ldots, k\}$
- Strategy: generalized eigenvectors  $v_i \in \mathcal{S}^{d-1}$  associated with top-k generalized eigenvalues  $\lambda_i$
- **Utility**: *i*-th player utility conditioned on players (j < i):

$$u_{i} (\hat{v}_{i} | \hat{v}_{j < i}) = \frac{\langle \hat{v}_{i}, A \hat{v}_{i} \rangle}{\langle \hat{v}_{i}, B \hat{v}_{i} \rangle} - \sum_{j < i} \frac{\langle \hat{v}_{j}, A \hat{v}_{j} \rangle \langle \hat{v}_{i}, B \hat{v}_{j} \rangle^{2}}{\langle \hat{v}_{j}, B \hat{v}_{j} \rangle^{2} \langle \hat{v}_{i}, B \hat{v}_{i} \rangle}$$
$$= \hat{\lambda}_{i} - \sum_{j < i} \hat{\lambda}_{j} \langle \hat{y}_{i}, B \hat{y}_{j} \rangle,$$

where 
$$\hat{y}_i = \frac{\hat{v}_i}{\|\hat{v}_i\|_B}$$
,  $\hat{\lambda}_i = \frac{\langle \hat{v}_i, A \hat{v}_i \rangle}{\langle \hat{v}_i, B \hat{v}_i \rangle}$  (gen. Rayleigh quotient),  $\|z\|_B = \sqrt{\langle z, Bz \rangle}$ 

#### Utilities as periodic functions

**Proposition 1** (Utility Shape). Each player's utility is periodic in the angular deviation  $(\theta)$  along the sphere. Its shape is sinusoidal, but with its angular axis  $(\theta)$  smoothly deformed as a function of B. Most importantly, every local maximum is a global maximum.



#### Deriving algorithm

Gradient of player's utility (up to scaling constant):

$$\frac{(\hat{v}_i^{\top} B \hat{v}_i) A \hat{v}_i - (\hat{v}_i^{\top} A \hat{v}_i) B \hat{v}_i}{\langle \hat{v}_i, B \hat{v}_i \rangle^2} - \sum_{i \leq i} \frac{\hat{\lambda}_j}{\langle \hat{v}_j, B \hat{v}_j \rangle} (\hat{v}_i^{\top} B \hat{v}_j) \frac{\left[\langle \hat{v}_i, B \hat{v}_i \rangle B \hat{v}_j - \langle \hat{v}_i, B \hat{v}_j \rangle B \hat{v}_i\right]}{\langle \hat{v}_i, B \hat{v}_i \rangle^2}.$$

- (i)  $\hat{\lambda}_j \langle \hat{v}_i, B \hat{v}_j \rangle = \langle \hat{v}_i, A \hat{v}_j \rangle$  if player i's parents match their true solutions, i.e.,  $\hat{v}_{j < i} = v_{j < i}$ ,
- (ii)  $\sqrt{\langle \hat{v}_j, B \hat{v}_j \rangle} = ||\hat{v}_j||_B$  is strictly positive and real-valued because  $B \succ 0$ ,

Simplified update:

$$\tilde{\nabla}_{i} = \overbrace{(\hat{v}_{i}^{\top} B \hat{v}_{i}) A \hat{v}_{i} - (\hat{v}_{i}^{\top} A \hat{v}_{i}) B \hat{v}_{i}}^{\text{reward}} - \sum_{i \leq i} \overbrace{(\hat{v}_{i}^{\top} A \hat{y}_{j}) \left[ \langle \hat{v}_{i}, B \hat{v}_{i} \rangle B \hat{y}_{j} - \langle \hat{v}_{i}, B \hat{y}_{j} \rangle B \hat{v}_{i} \right]}^{\text{penalty}}. \tag{7}$$

**Lemma 1** (Well-posed Utilities). Given exact parents and assuming the top-k eigenvalues of  $B^{-1}$  are distinct and positive, the maximizer of player i's utility is the unique generalized eigenvector i (up to sign, i.e.,  $-v_i$  is also valid).

**Theorem 1** (Nash Property). Assuming the top-k generalized eigenvalues of the generalized eigenvalue problem  $Av = \lambda Bv$  are positive and distinct, their corresponding generalized eigenvectors form the unique, strict Nash equilibrium of  $\Gamma$ -EigenGame.

**Lemma 2.** The direction  $\tilde{\nabla}_i$  defined in equation (7) is a steepest ascent direction on utility  $u_i(\hat{v}_i|\hat{v}_{j< i})$  given exact parents  $\hat{v}_{j< i} = v_{j< i}$ .

**Theorem 2** (Deterministic / Full-batch Global Convergence). Given a symmetric matrix A and symmetric positive definite matrix B where the top-k eigengaps of  $B^{-1}A$  are positive along with a square-summable, not summable step size sequence  $\eta_t$  (e.g., 1/t), Algorithm I converges to the top-k generalized eigenvectors asymptotically ( $\lim_{T\to\infty}$ ) with probability 1.

#### **Algorithm 1** Deterministic / Full-batch $\gamma$ -EigenGame

- 1: Given:  $A \in \mathbb{R}^{d \times d}$  and  $B \in \mathbb{R}^{d \times d}$ , step size sequence  $\eta_t$ , and number of iterations T.
- 2:  $\hat{v}_i \sim \mathcal{S}^{d-1}$ , i.e.,  $\hat{v}_i \sim \mathcal{N}(\mathbf{0}_d, \mathbf{I}_d)$ ;  $\hat{v}_i \leftarrow \hat{v}_i / ||\hat{v}_i||$  for all i
- 3: **for** t = 1 : T **do**
- parfor i = 1 : k do
- $\hat{y}_j = rac{\hat{v}_j}{\sqrt{\langle \hat{v}_i, B \hat{v}_i 
  angle}}$ 5:
- rewards  $\leftarrow (\hat{v}_i^{\top} B \hat{v}_i) A \hat{v}_i (\hat{v}_i^{\top} A \hat{v}_i) B \hat{v}_i$ 6:
- penalties  $\leftarrow \sum_{i < i} (\hat{v}_i^ op A \hat{y}_j) ig[ \langle \hat{v}_i, B \hat{v}_i 
  angle B \hat{y}_j \langle \hat{v}_i, B \hat{y}_j 
  angle B \hat{v}_i ig]$ 7:
- $\nabla_i \leftarrow \text{rewards} \text{penalties}$ 8:
- $\hat{v}_i' \leftarrow \hat{v}_i + \eta_t \nabla_i$
- $\hat{v}_i \leftarrow rac{\hat{v}_i'}{||\hat{v}_i'||}$ 10:
- end parfor 11:
- 12: end for
- 13: return all  $\hat{v}_i$

#### **Algorithm 2** Stochastic $\gamma$ -EigenGame

```
1: Given: paired data streams X_t \in \mathbb{R}^{b \times d_x} and Y_t \in \mathbb{R}^{b \times d_y}, number of parallel machines M
        per player (minibatch size per machine b'=\frac{b}{M}), step size sequences \eta_t and \gamma_t, scalar \rho lower
       bounding \sigma_{min}(B), and number of iterations T.
  2: \hat{v}_i \sim \mathcal{S}^{d-1}, i.e., \hat{v}_i \sim \mathcal{N}(\mathbf{0}_d, \mathbf{I}_d); \hat{v}_i \leftarrow \hat{v}_i / ||\hat{v}_i|| for all i
  3: [B\hat{v}]_i \leftarrow \hat{v}_i^0 for all i
  4: for t = 1 : T do
            parfor i=1:k do
  5:
                 parfor m=1:M do
 6:
                     Construct A_{tm} and B_{tm} (*unbiased estimates using independent data batches)
  7:
                     \hat{y}_j = rac{\hat{v}_j}{\sqrt{\max(\langle \hat{v}_j, [B\hat{v}]_j 
angle, oldsymbol{
ho})}}
  8:
                     [B\hat{y}]_j = \frac{[B\hat{v}]_j}{\sqrt{\max(\langle \hat{v}_i, [B\hat{v}]_i \rangle, \rho)}}
 9:
                     rewards \leftarrow (\hat{v}_i^{\top} B_{tm} \hat{v}_i) A_{tm} \hat{v}_i - (\hat{v}_i^{\top} A_{tm} \hat{v}_i) B_{tm} \hat{v}_i
10:
                     \text{penalties} \leftarrow \textstyle \sum_{j < i} (\hat{v}_i^\top A_{tm} \hat{y}_j) \big[ \langle \hat{v}_i, B_{tm} \hat{v}_i \rangle [B \hat{y}]_j - \langle \hat{v}_i, [B \hat{y}]_j \rangle B_{tm} \hat{v}_i \big]
11:
                     	ilde{
abla}_{im} \leftarrow 	ext{rewards} - 	ext{penalties}
12:
                     \nabla_{im}^{Bv} = (B_{tm}\hat{v}_i - [B\hat{v}]_i)
13:
                end parfor
14:
                \tilde{\nabla}_i \leftarrow \frac{1}{M} \sum_m [\tilde{\nabla}_{im}]
15:

\hat{v}_i' \leftarrow \hat{v}_i + \eta_t \tilde{\nabla}_i \\
\hat{v}_i \leftarrow \frac{\hat{v}_i'}{||\hat{v}_i'||}

16:
17:
                \nabla_i^{Bv} \leftarrow \frac{1}{M} \sum_m [\nabla_{im}^{Bv}]
18:
               [B\hat{v}]_i \leftarrow [B\hat{v}]_i + \gamma_t \nabla_i^{Bv}
19:
            end parfor
20:
21: end for
22: return all \hat{v}_i
```

Naive Complexity:

$$\mathcal{O}(bdk^2)$$

Parallelized Complexity:

$$\mathcal{O}(dk)$$

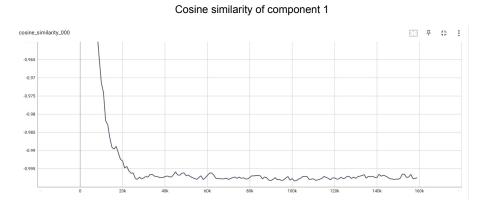
#### Experiments

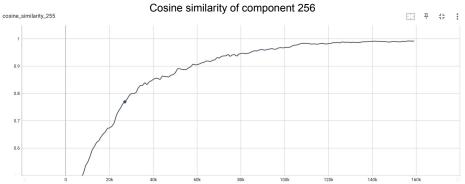
We run two kinds of experiments:

- 1. Basic synthetic data generated from normal distribution
- 2. Images from mnist dataset

We also track cosine similarity as a loss function whenever possible

#### Synthetic data: convergence

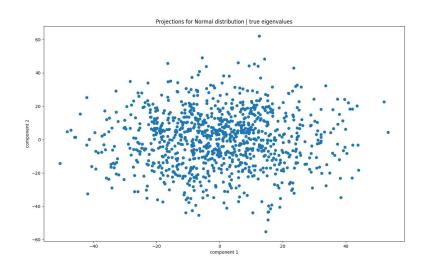


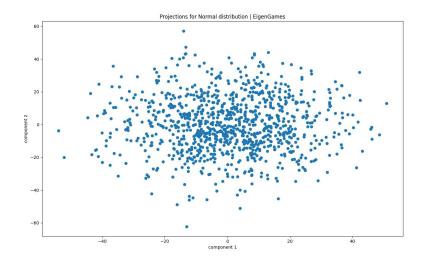


#### Valuable observations:

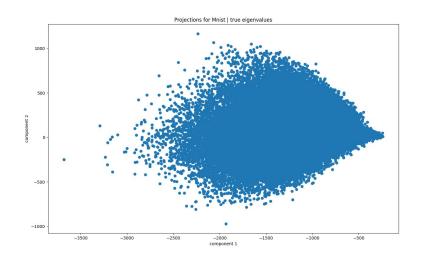
- 1. Algorithm is able to recover the true eigenvectors up to a sign
- 2. We can observe how the prediction of eigenvectors converge hierarchically
- Algorithm takes time to converge.
   Therefore, there is no advantage on running it on simple datasets that can fit into a memory

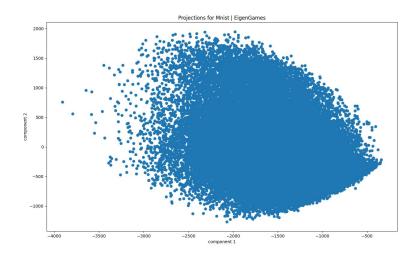
# Synthetic data: projections

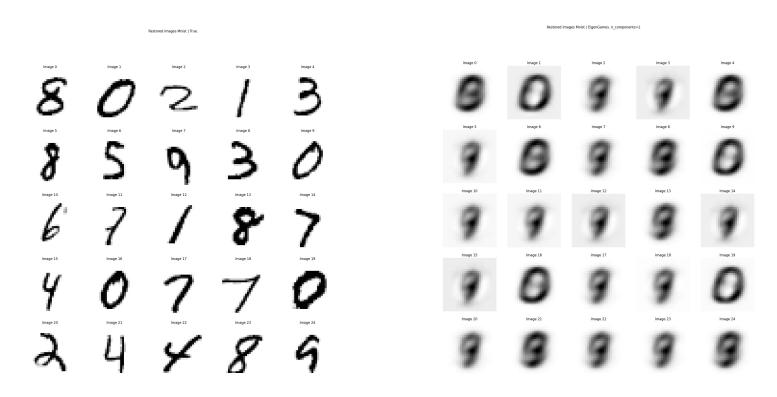


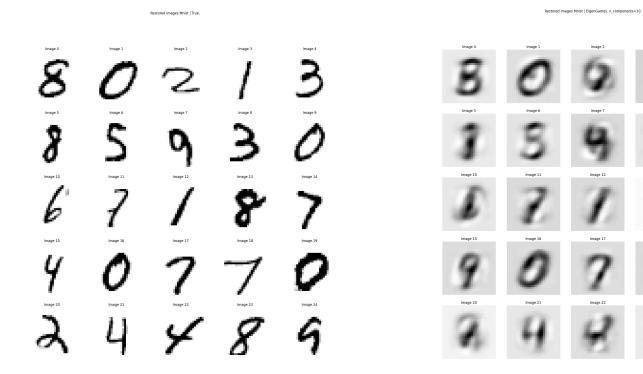


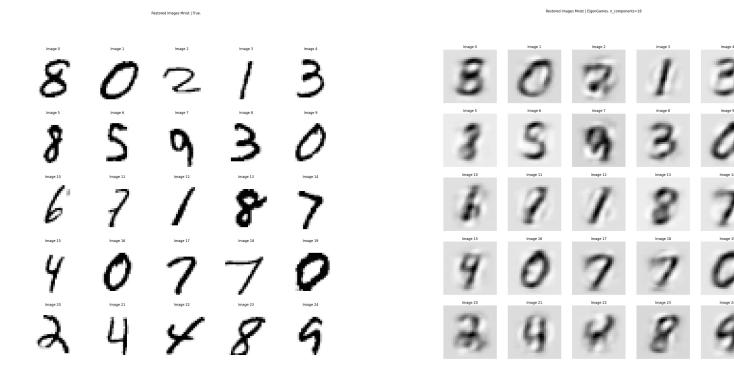
# Mnist projections

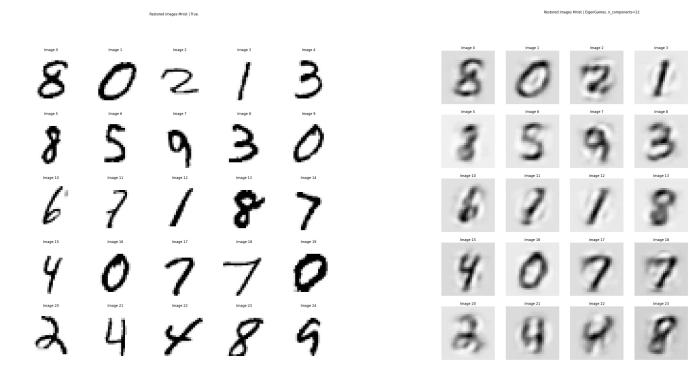


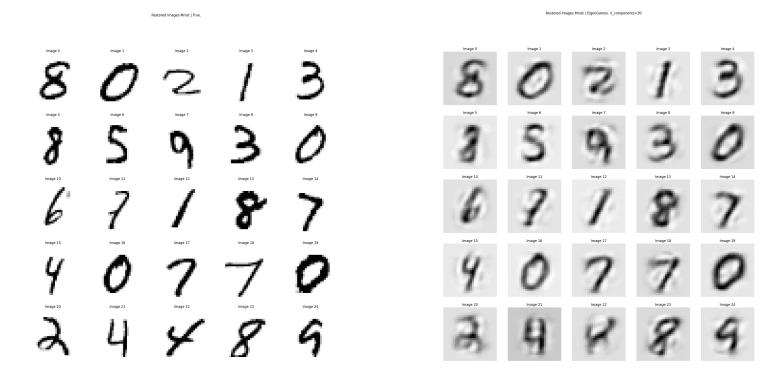














#### References

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- [Gemp2022] EigenGame Unloaded: When Playing Games is Better Than Optimizing, https://arxiv.org/abs/2102.04152
- [Gemp2023] The Symmetric Generalized Eigenvalue Problem as a Nash Equilibrium, https://arxiv.org/abs/2206.04993