Topological Data Analysis

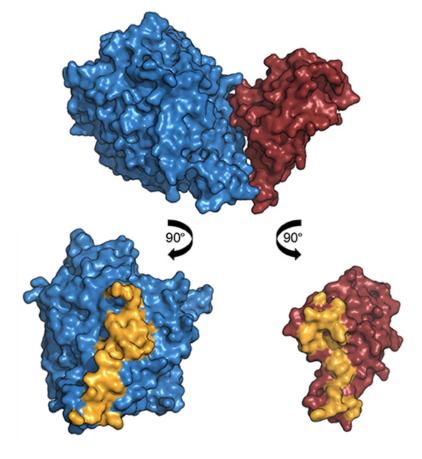
Lecture 6

Higher-order Network Science

Complex networks



Social and collaboration networks



Protein interaction networks



Functional brain networks

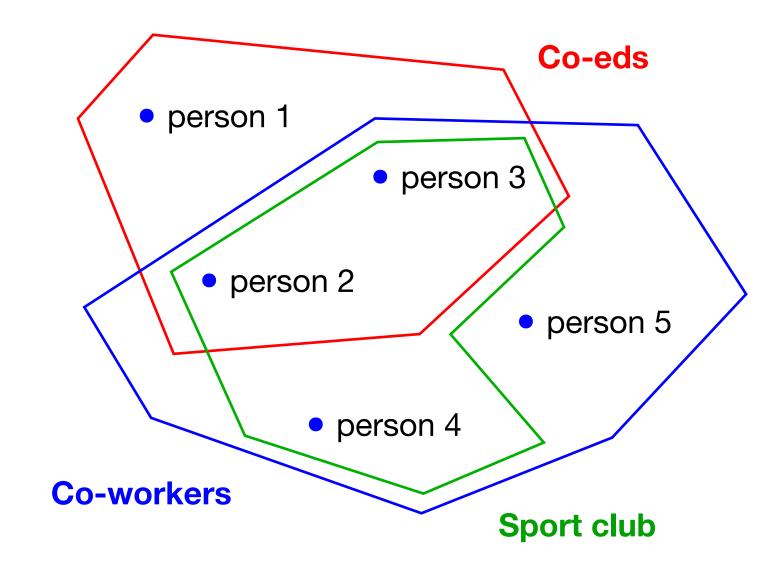


Financial networks

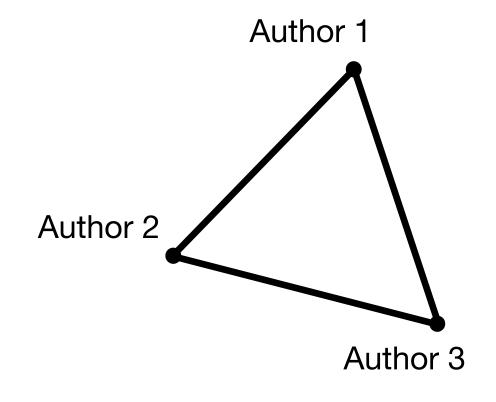
Higher-order complex networks

Collaboration networks

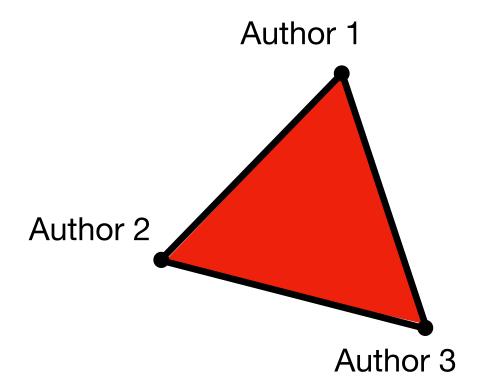
Social network



Collaboration network



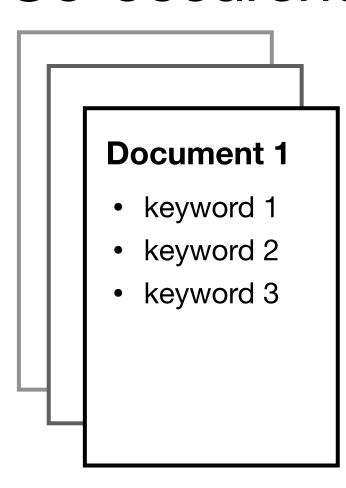
- paper 1 (authors 1-2)
- paper 2 (authors 1-3)
- paper 3 (authors 2-3)



paper 4 (Authors 1-2-3)

Higher-order complex networks

Co-occurence networks



Word-document matrix

	Word 1	Word 2	Word 3
Doc 1		1	
Doc 2	1	1	
Doc 3			1
Doc 4		1	

Basket 1

• item 1, item 2, item 3

Basket 2

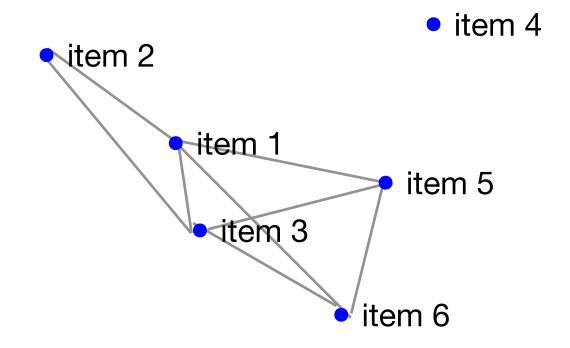
• item 1, item 3, item 5, item 6

Item-basket matrix

	Item 1		Item 6
Basket 1	1	•••	
Basket 2	1		1

Graph

Homogenous



Context 1

quick brown **fox** jumps over

Context 2

brown fox **jumps** over lazy

Context 3

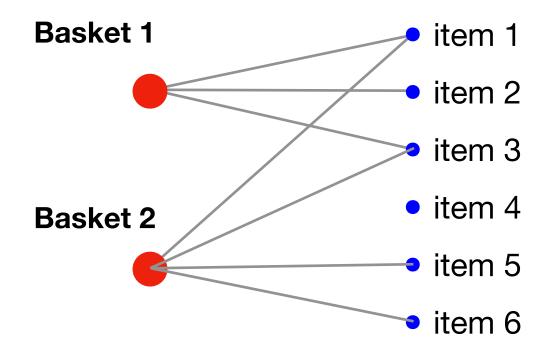
fox jumps **over** lazy dog

Word-context matrix

	Word 1	Word 2	Word 3
Context 1	1		
Context 2	1	1	
Context 3		1	
Context 4	1		1

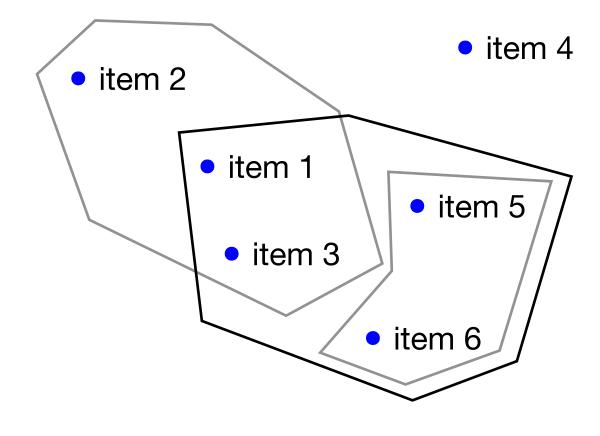
Bipartite graph

Heterogenous



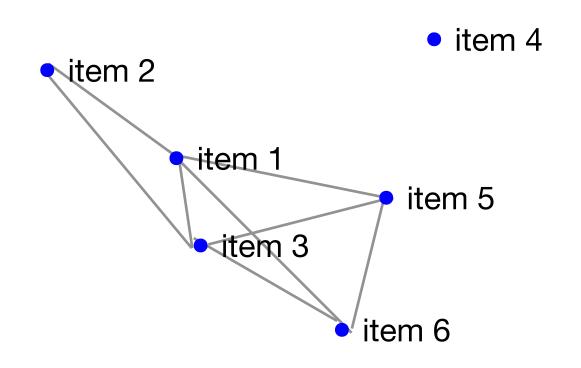
Hypergraph

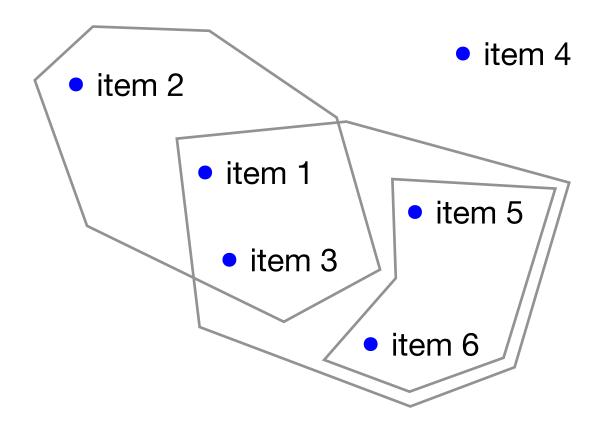
Homogenous

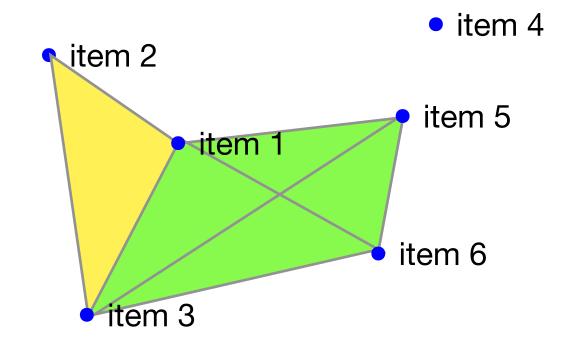


Higher-order complex networks

Models







Graph

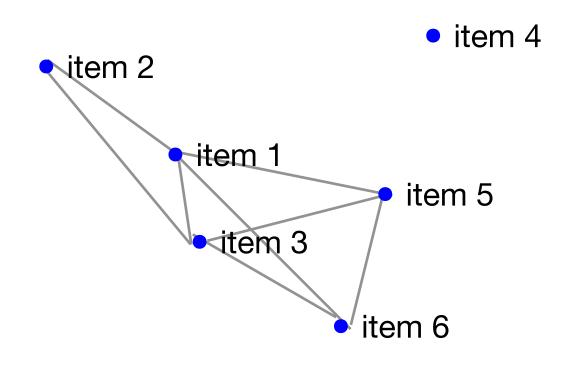
Hypergraph

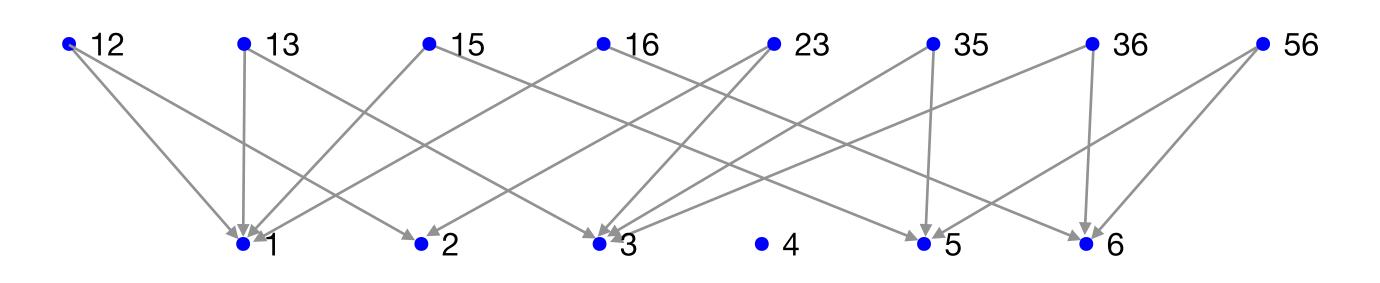
Simplicial complex

Closed under inclusion Edges of different dimensions

Hasse diagram

Graph

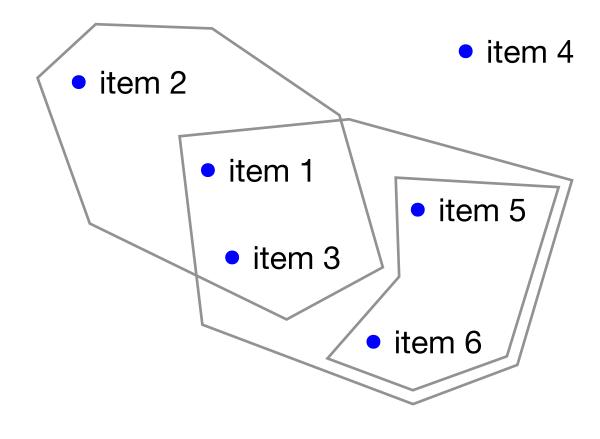


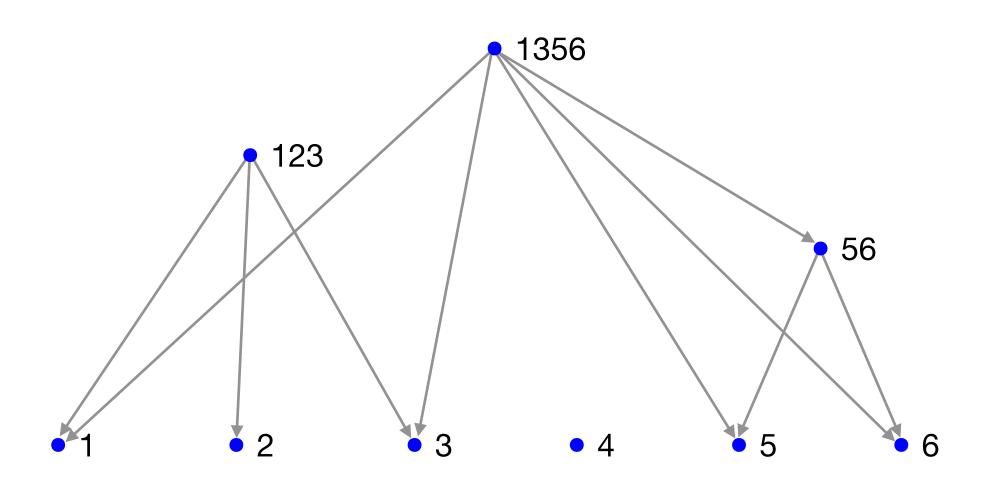


Graph Hasse diagram

Hasse diagram

Hypergraph

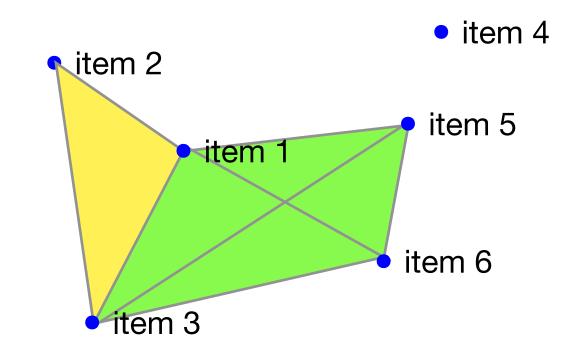




Hypergraph Hasse diagram

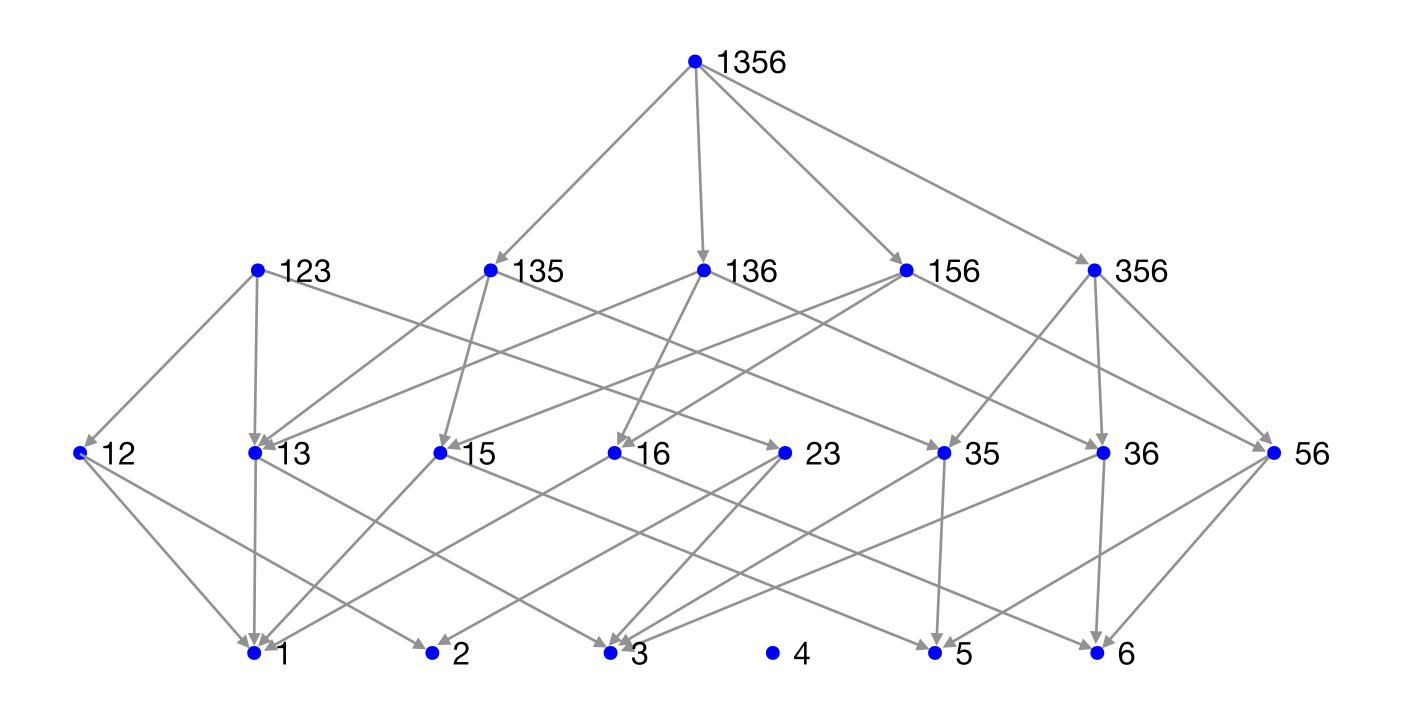
Hasse diagram

Simplicial complex



Simplicial complex

Closed under inclusion Edges of different dimensions



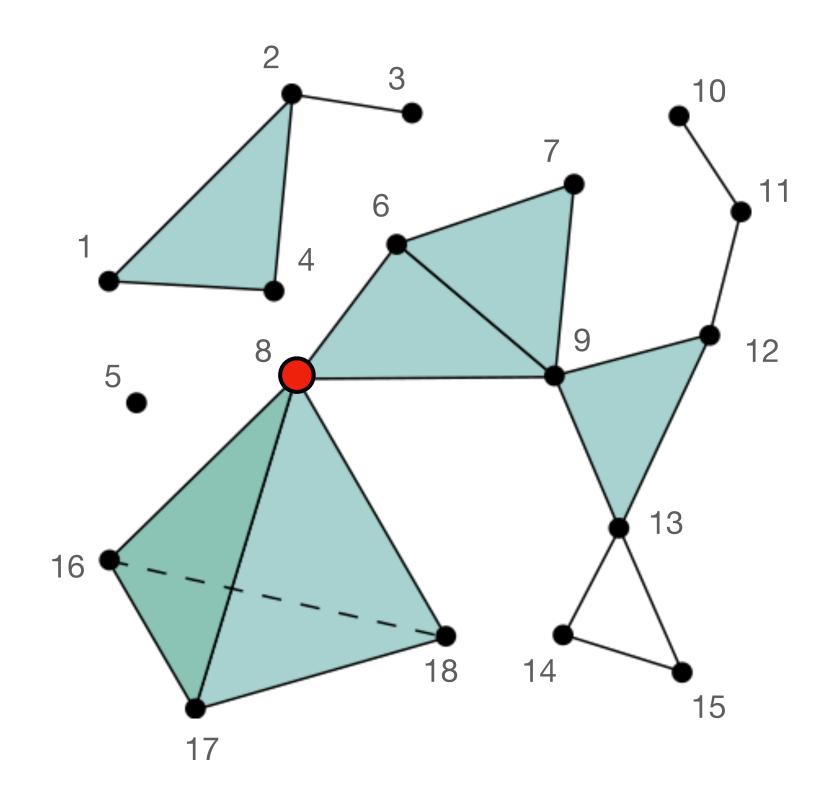
Hasse diagram

Incidence

For a pair of simplices $\tau \subseteq \sigma$

- au is a face of σ ,
- σ is a coface of au.

A p-simplex $\sigma_1^{(p)}$ is q-incident to a q-simplex $\sigma_2^{(q)}$, denoted $\sigma_1^{(p)} \to_q \sigma_2^{(q)}$ if $p \neq q$ and $\sigma_1^{(p)}$ is a face or coface of $\sigma_2^{(q)}$.



Simplicial complex K

Incidence

For a pair of simplices $\tau \subseteq \sigma$

- au is a face of σ ,
- σ is a coface of τ .

A p-simplex $\sigma_1^{(p)}$ is *q-incident* to a q-simplex $\sigma_2^{(q)}$, denoted $\sigma_1^{(p)} \to_q \sigma_2^{(q)}$ if $p \neq q$ and $\sigma_1^{(p)}$ is a face or coface of $\sigma_2^{(p)}$.

Upper incidence p < q

$$\{8\} \rightarrow_1 \{6,8\}$$

$$\{8\} \rightarrow_1 \{6,8\} \qquad \{8\} \rightarrow_2 \{8,17,18\}$$

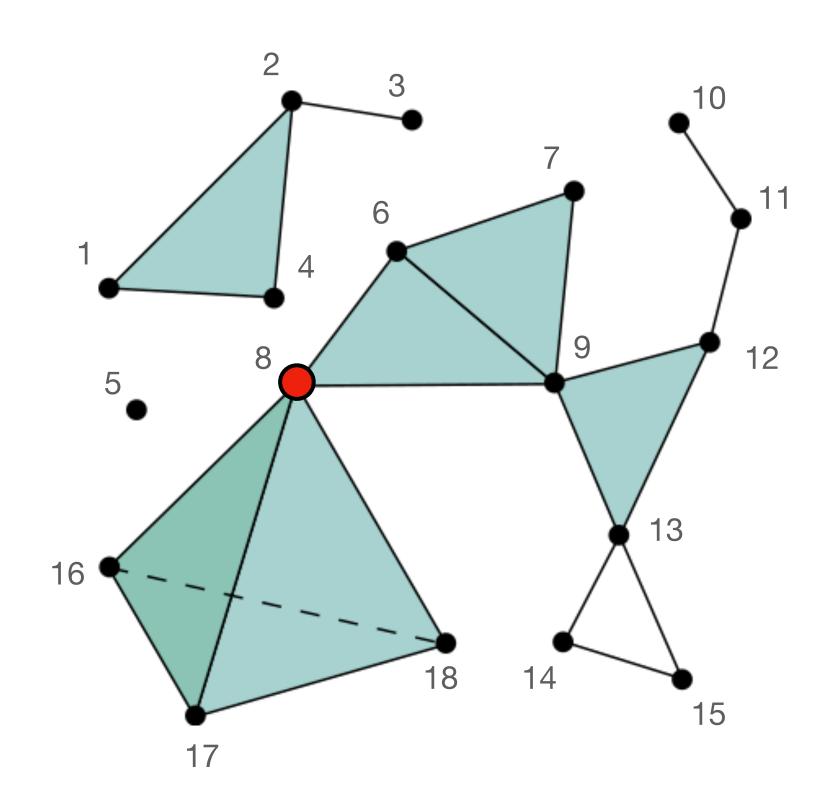
$$\{8\} \rightarrow_3 \{8,16,17,18\}$$

$$\{8,16\} \rightarrow_3 \{8,16,17,18\}$$

Lower incidence p < q

$$\{6,7,9\} \rightarrow_1 \{6,9\}$$

$$\{8,16,17,18\} \rightarrow_1 \{16,18\}$$

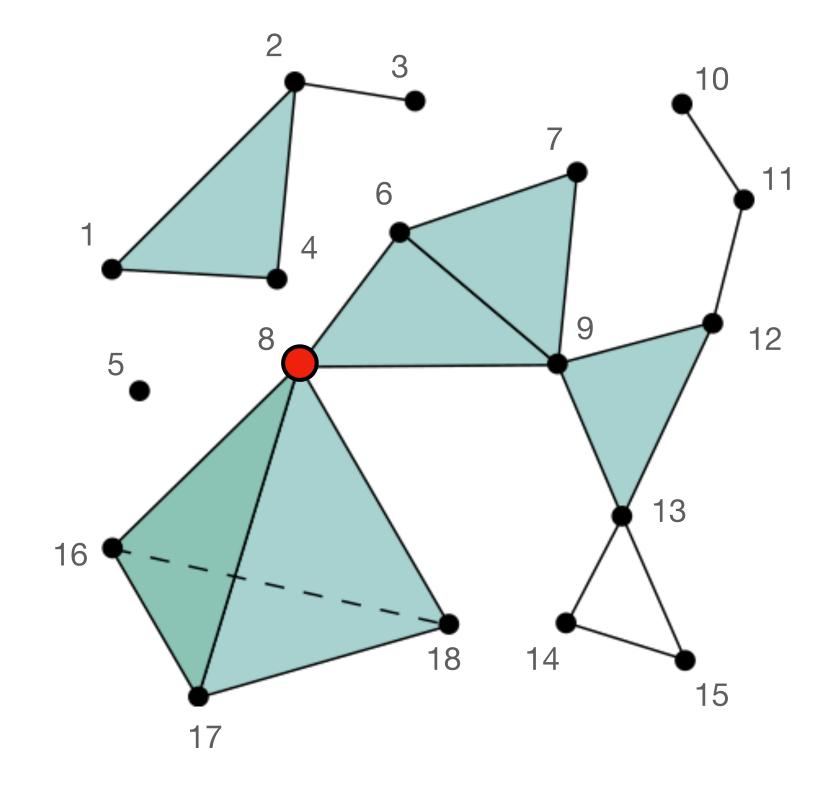


Simplicial complex K

Incidence

(p, q)-incidence of a p-simplex is the # q-simplices q-incident to it.

$$i_1(\{8\}) = 5$$
 $i_2(\{8\}) = 4$ $i_3(\{8\}) = 1$



Simplicial complex K

Adjacency

A set of p-simplices $\{\sigma_1^{(p)},...,\sigma_n^{(p)}\}$ is q-adjacent via p-simplex τ^q denoted $\sigma_1^{(p)} \sim_q \sigma_2^{(p)}$ pairwise, or $\sim_q \{\sigma_1^{(p)},...,\sigma_n^{(p)}\}$ for all p-simplices if

- all p-simplices $\sigma_i^{(p)}$ are p-faces of q-simplex $au^{(q)}$, conversely
- q-simplex $au^{(q)}$ is a q-coface for all k-simplices $\sigma_i^{(p)}$.

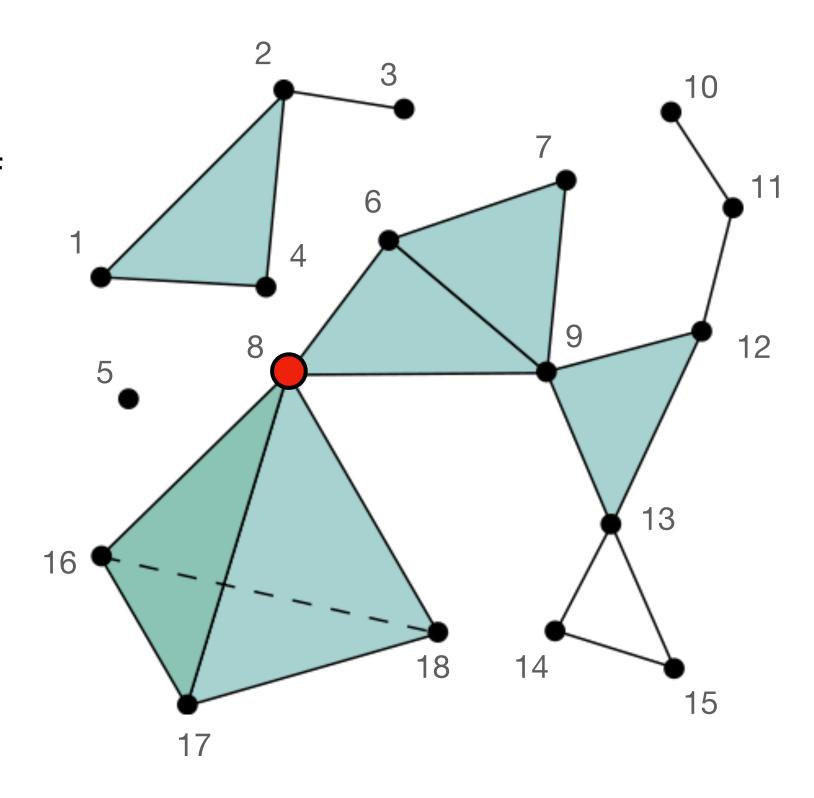
Upper adjacency p < q

$$\{6,9\} \sim_2 \{6,7\}$$
 $\{8\} \sim_3 \{18\}$

Lower adjacency p > q

$$\{8,16,17\} \sim_0 \{8,16,17,18\} \quad \{6,7,9\} \sim_1 \{6,8,9\}$$

(p, q)-degree of a p-simplex is the # of q-adjacent to it p-simplices.



Simplicial complex K

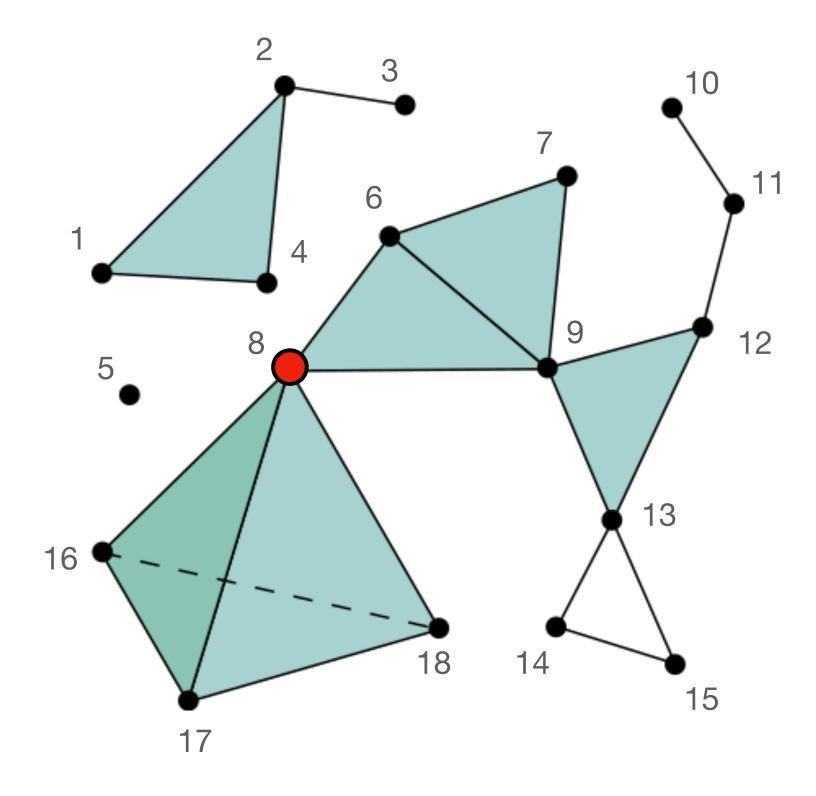
Degree

(p, q)-degree of a p-simplex is the # of q-adjacent to it p-simplices.

$$d_1(\{8\}) = 5$$

$$d_1(\{8\}) = 5$$
 $d_2(\{8\}) = 5$ $d_3(\{8\}) = 3$

$$d_3(\{8\}) = 3$$



Simplicial complex K

Degree

(p, q)-degree of a p-simplex is the # of q-adjacent to it p-simplices.

$$d_1(\{8\}) = 5$$

$$d_2(\{8\}) = 5$$

$$d_1(\{8\}) = 5$$
 $d_2(\{8\}) = 5$ $d_3(\{8\}) = 3$

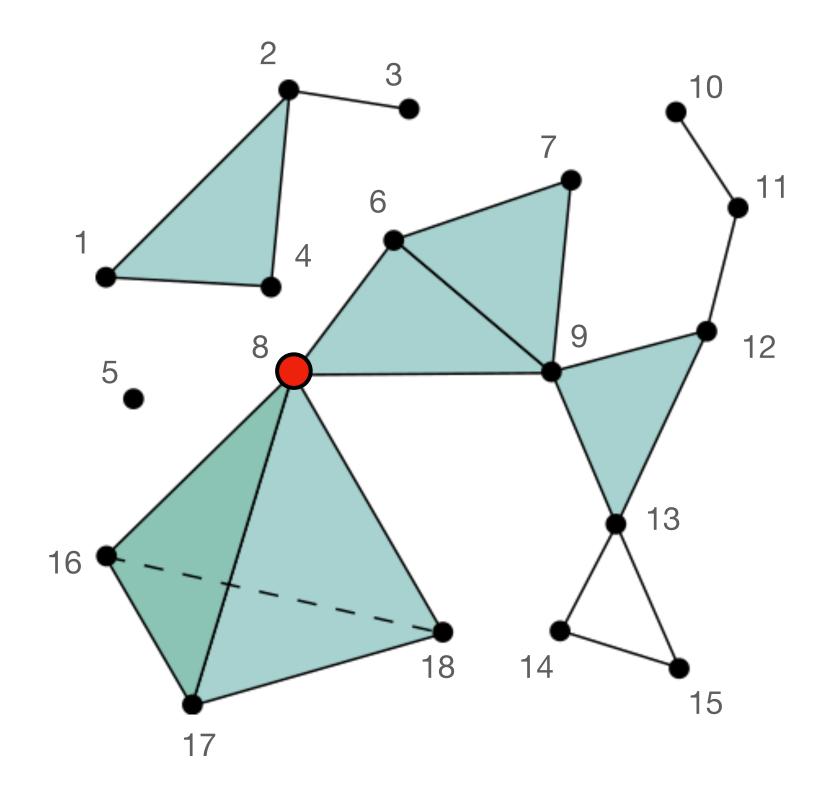
same for q = 1

different for q > 1

$$i_1(\{8\}) = 5$$

$$i_2(\{8\}) = 4$$

$$i_1(\{8\}) = 5$$
 $i_2(\{8\}) = 4$ $i_3(\{8\}) = 1$



Simplicial complex K

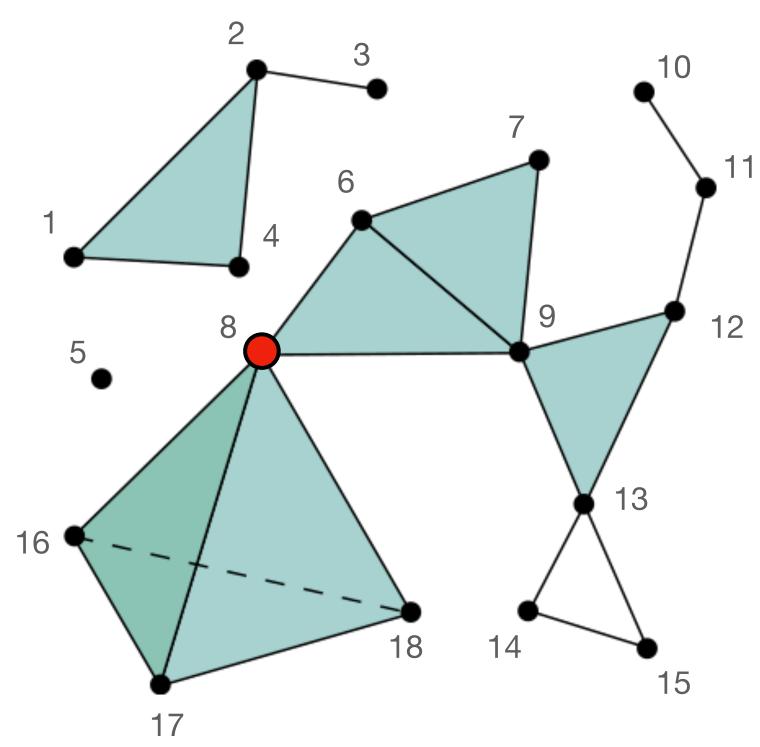
Connected components

A (p, q)-path is a sequence of simplices $\{\sigma_1^{(p)}, \sigma_2^{(q)}, \sigma_3^{(p)}, \sigma_4^{(q)}, \ldots, \sigma_{n-1}^{(q)}, \sigma_n^{(p)}\}$ such that $\sigma_i^{(p)} \to_q \sigma_{i+1}^{(p)} \to_q \sigma_{i+2}^{(p)} \ \forall i$.

Two simplices $\sigma_1^{(p)}, \sigma_2^{(p)}$ are *(p, q)-connected* if there exists (p, q)-path between them.

A (p, q)-connected component is an equivalence relation on K.

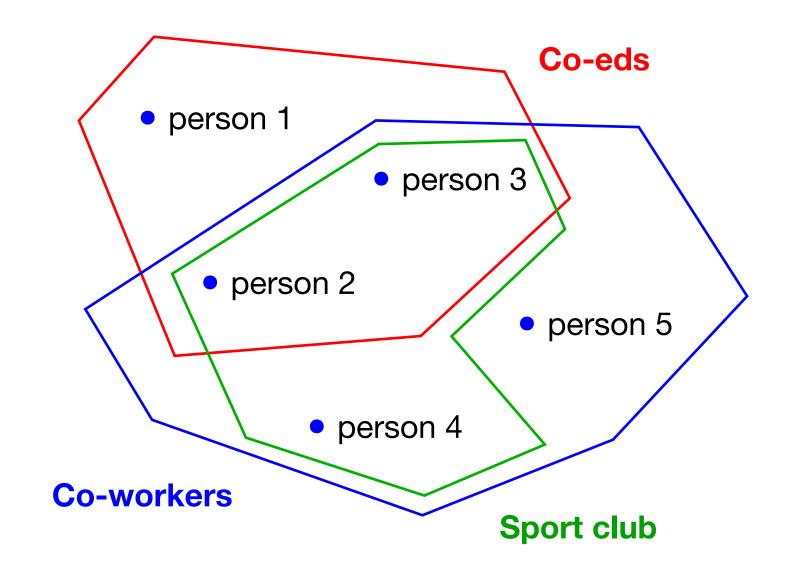
$$|C_{0,1}(K)| = 3$$
 $|C_{0,2}(K)| = 8$



Simplicial complex K

Q-analysis

Social networks



Chess networks



Binary relation

Let X and Y be sets or cardinalities m and n, respectively

$$X = \{x_1, ..., x_m\}$$
 $Y = \{y_1, ..., y_n\}$

Binary relation R is a subset of Cartesian product X x Y

$$R \subseteq X \times Y$$

$$xRy \iff (x,y) \in R$$

Matrix representation

R on sets X and Y is represented by m x n matrix A

$$a_{ij} = \begin{cases} 1, & x_i R y_j, \\ 0, & \text{otherwise}. \end{cases}$$

	y 1	y 2	Уз
X 1	1		
X 2	1	1	
X 3	1		1
X 4		1	1

Let R is a binary relation on sets X and Y of cardinalities m and n, respectively.

A Dowker complex K of a binary relation R on sets X and Y is defined

$$K = \{ \sigma^{(m)} = \{ x_{\sigma_0}, ..., x_{\sigma_m} \} \mid \exists y \text{ s.t.} x_i R y \quad \forall x_i \in \sigma^{(m)} \}$$

Analogously, a *Dowker complex L* is defined

$$L = \{ \sigma^{(n)} = \{ y_{\sigma_0}, ..., y_{\sigma_n} \} \mid \exists x \text{ s.t.} x_i R y \quad \forall y_j \in \sigma^{(n)} \}$$

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	y 1	У 2	У 3
X 1	1		
X 2	1	1	
X 3	1		1
X 4		1	1

$$\sigma^{(2)} = \{x_1, x_2, x_3\} \in K$$

Let R is a binary relation on sets X and Y of cardinalities m and n, respectively.

A Dowker complex K of a binary relation R on sets X and Y is defined

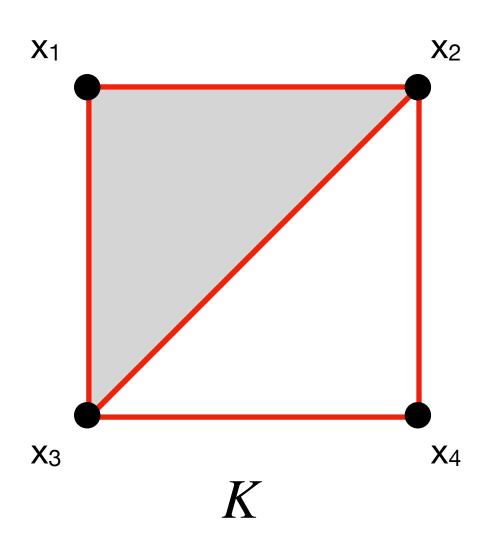
$$K = \{ \sigma^{(m)} = \{ x_{\sigma_0}, ..., x_{\sigma_m} \} \mid \exists y \text{ s.t.} x_i R y \quad \forall x_i \in \sigma^{(m)} \}$$

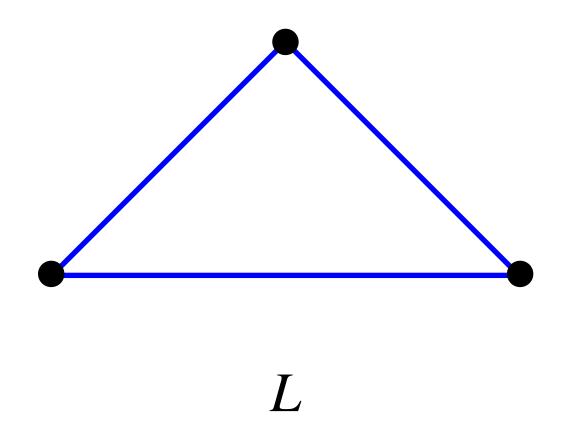
Analogously, a *Dowker complex L* is defined

$$L = \{ \sigma^{(n)} = \{ y_{\sigma_0}, ..., y_{\sigma_n} \} \mid \exists x \text{ s.t.} x_i R y \quad \forall y_j \in \sigma^{(n)} \}$$

	y 1	y 2	y 3
X ₁	1		
X 2	1	1	
X 3	1		1
X 4		1	1

$$\sigma^{(1)} = \{y_1, y_3\} \in L$$



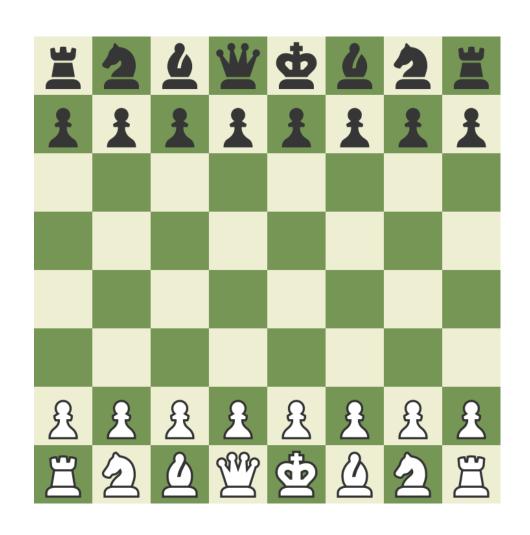


	y 1	y 2	y 3
X 1	1		
X 2	1	1	
X 3	1		1
X 4		1	1

A

Homology groups of Dowker complexes K and L are isomorphic, i.e. $H_{\bullet}(K) \simeq H_{\bullet}(L)$ [Dowker1952, Thm. 1].

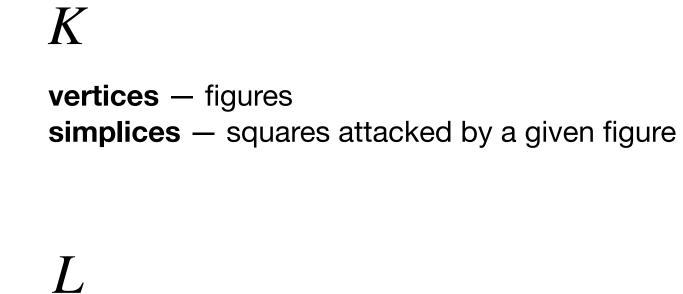
Chess network



Chess board

	f ₁	•••	f ₁₆
S ₁	1		
S ₂	1	1	
	1		1
S 64		1	1

Relation "figure attacks square"



vertices — squaressimplices — figures attacking a given square

Dowker complexes K and L

Q-analysis

Q-connected components

Given a Dowker complex K (L) Atkin's Q-analysis is to associate a Q-vector defined as a vector of #(max, q)-connected components of K (L) for all q.

Q-vector of K

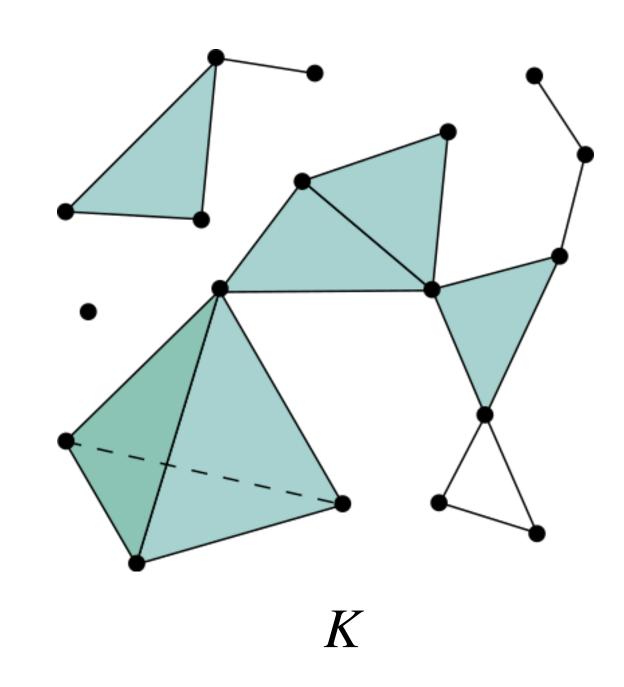
$$(3, 4+m) \tag{3}$$

0 1

(3, 6, 15)

(0, q)-vector of *K*

2 3

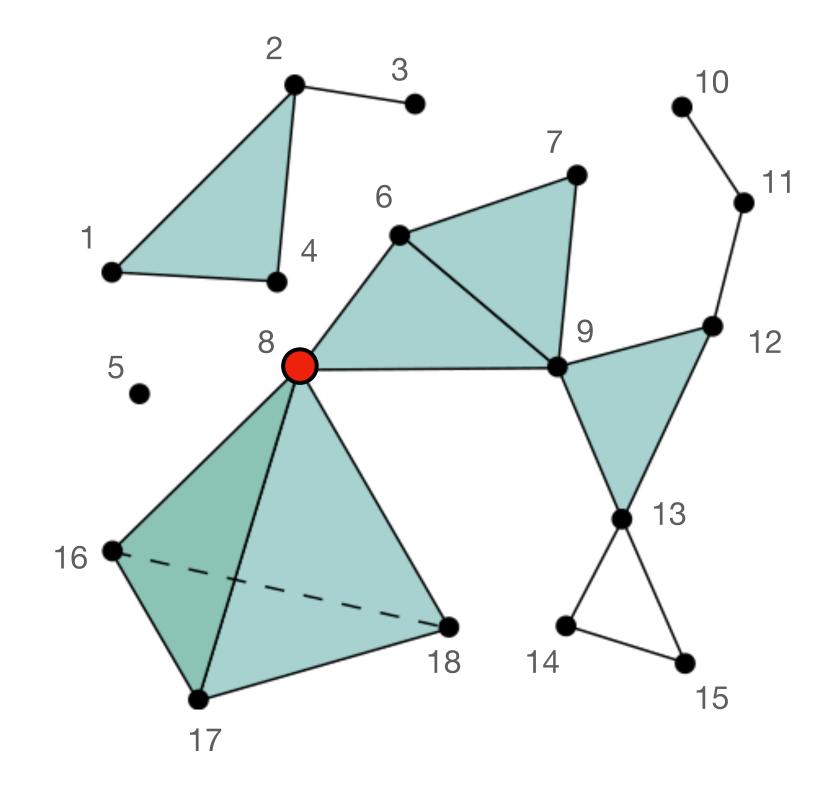


F-vector (18, 19, 8, 1)

Line graph

A (p, q)-line graph of a simplicial complex K is a graph G(V, E) where

- V consists of p-simplices of K,
- $(\sigma_1^{(p)}, \sigma_2^{(p)}) \in E_G$ if $\sigma_1^{(p)} \sim_q \sigma_2^{(p)}$.



Simplicial complex K

Centralities

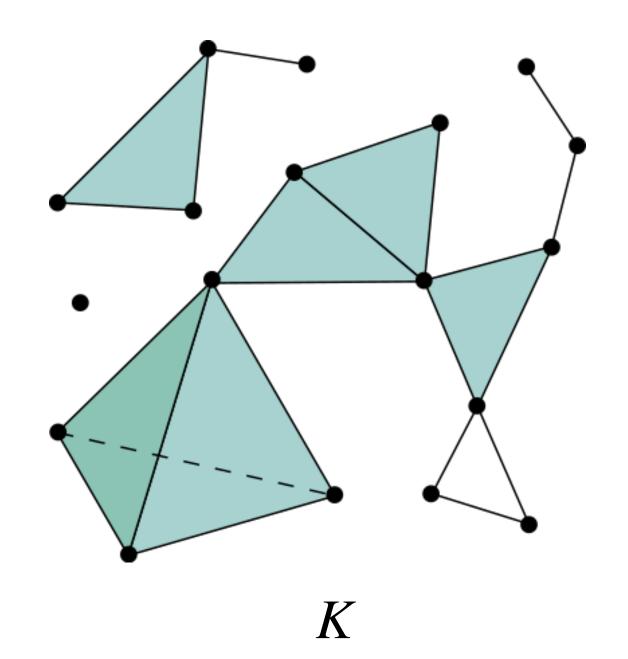
Betweenness, closeness centrality

(p, q)-betweenness centrality

shortest (p, q)-paths passing through given p-simplex

(p, q)-closeness centrality

Inverse of sum of the length of (p, q)-paths between given p-simplex and other p-simplices



Can be computed using (p, q)-line graph