

## 1. Proof the correctness of program based on unit testing

### 1.1 Assumptions

Let  $P$  be a program to test. Assume that the program is syntactically and semantically valid.

### 1.2 Proof that a block of $P$ is correct on a test

#### 1.2.1 Declaration of a program block

Let  $inputs_{f_{pb}} = \{v \mid v \text{ is a value}\}$

Let  $func_{f_{pb}}: inputs_{f_{pb}} \rightarrow output_{f_{pb}}$ , where  $func_{f_{pb}}$  is a function

Let  $pb_i = (func_{f_{pb}}, inputs_{f_{pb}}, output_{f_{pb}})$

#### 1.2.2 Declaration of all program blocks

Let  $Blocks(P) = \{pb_i \mid pb_i \text{ is a program block of } P\}$

#### 1.2.3 Declaration of a test

Let  $inputs_{f_t} = \{v \mid v \text{ is a value}\}$

Let  $func_{f_t}: inputs_{f_t} \rightarrow output_{f_t}$ , where  $func_{f_t}$  is a function

Let  $t_j = (func_{f_t}, inputs_{f_t}, output_{f_t})$

#### 1.2.4 Declaration of all test sets for $P$

Let  $TS_{pb_i} = \{t_j \mid t_j \text{ is a test}\}$

Let  $TS_P = \bigcup_{pb_i \in Blocks(P)} TS_{pb_i}$

#### 1.2.5 Show the validity and success of a test on a program block

$$success(t_j, pb_i) = \begin{cases} 1, & \text{if } [t_j]_{output_{f_t}} = [pb_i]_{output_{f_{pb}}} \\ 0, & \text{otherwise} \end{cases}$$

$$valid(t_j, pb_i) = \begin{cases} 1, & \text{if } [t_j]_{func_{f_t}} = [pb_i]_{func_{f_{pb}}} \text{ and } [t_j]_{inputs_{f_t}} = [pb_i]_{inputs_{f_{pb}}} \\ 0, & \text{otherwise} \end{cases}$$

$$valid(t_j, pb_i) \rightarrow success(t_j, pb_i) = 1$$

#### 1.2.6 Proof that the whole program is correct

$$correctBlock(TS_{pb_i}, pb_i) = \forall t_j \in TS_{pb_i} \mid valid(t_j, pb_i)$$

$$correctP(TS_P, Blocks(P)) = \forall pb_i \in Blocks(P), \forall TS_{pb_i} \in TS_P \mid correctBlock(TS_{pb_i}, pb_i)$$

## 2. Proof correctness of program from requirements point of view.

### 2.1 Considerate the previous parts 1.2.1 to 1.2.4 include as the first part of this proof.

### 2.2 Declaration of requirements

Let  $pre - condition = \{v \mid v \text{ is a value}\}$

Let  $post - condition = \{v \mid v \text{ is a value}\}$

Let  $r_i = (pre - condition, post - condition)$

Let  $RS_{pb_i} = \{r_i \mid r_i \text{ is a requirement}\}$

Let  $RS_P = \bigcup_{pb_i \in Blocks(P)} RS_{pb_i}$

### 2.3 Test satisfaction of a requirement

$$\text{inputsEquality}(t_j, pb_i) = \begin{cases} 1, & \text{if } [t_j]_{\text{inputs}_{ft}} = [pb_i]_{\text{inputs}_{f_{pb}}} \\ 0, & \text{otherwise} \end{cases}$$

$$\text{outputsEquality}(t_j, pb_i) = \begin{cases} 1, & \text{if } [t_j]_{\text{output}_{ft}} = [pb_i]_{\text{output}_{f_{pb}}} \\ 0, & \text{otherwise} \end{cases}$$

$$\text{conditionInput}(t_j, pb_i, r_i) = \begin{cases} 1, & \text{if } \text{inputsEquality}(t_j, pb_i) \text{ and } \forall v \in [pb_i]_{\text{inputs}_{f_{pb}}} \mid v \in [r_i]_{\text{pre-condition}} \\ 0, & \text{otherwise} \end{cases}$$

$$\text{conditionOutput}(t_j, pb_i, r_i) = \begin{cases} 1, & \text{if } \text{outputsEquality}(t_j, pb_i) \text{ and } \forall v \in [pb_i]_{\text{output}_{f_{pb}}} \mid v \in [r_i]_{\text{post-condition}} \\ 0, & \text{otherwise} \end{cases}$$

$$\text{satisfy}(t_j, pb_i, r_i) = \begin{cases} 1, & \text{if } \text{conditionInput}(t_j, pb_i, r_i) \text{ and } \text{conditionOutput}(t_j, pb_i, r_i) \text{ and } \text{valid}(t_j, pb_i) \\ 0, & \text{otherwise} \end{cases}$$

### 2.4 A specific requirement is satisfied for a program block

$$\text{Let } \text{satisfactionR}(TS_{pb_i}, pb_i, r_i) = \begin{cases} 1, & \text{if } \forall t_j \in TS_{pb_i} \mid \text{satisfy}(t_j, pb_i, r_i) \rightarrow (\exists t_j \in TS_{pb_i} \mid \text{satisfy}(t_j, pb_i, r_i) = 1) \\ 0, & \text{otherwise} \end{cases}$$

### 2.5 All requirements are satisfied for a program block

$$\text{Let } \text{satisfactionSetR}(TS_{pb_i}, pb_i, RS_{pb_i}) = \begin{cases} 1, & \text{if } \forall t_j \in TS_{pb_i}, \forall r_i \in RS_{pb_i} \mid \text{satisfactionR}(TS_{pb_i}, pb_i, r_i) = 1 \\ 0, & \text{otherwise} \end{cases}$$

### 2.6 All requirements are satisfied for a program

$$\text{Let } \text{satisfactionALL} = \forall pb_i \in Blocks(P), \forall TS_{pb_i} \in TS_P, RS_{pb_i} \in RS_P \mid \text{satisfactionSetR}(TS_{pb_i}, pb_i, RS_{pb_i}) = 1$$