

1. Proof the correctness of program based on unit testing

1.1 Assumptions

Let P be a program to test. Assume that the program is syntactically valid.

1.2 Proof that a program P is correct for a set of tests

1.2.1 Declaration of a program block

Let $inputs_{f_{pb}} \equiv \{v \mid v \text{ is a value}\}$

Let $func_{f_{pb}}: inputs_{f_{pb}} \rightarrow output_{f_{pb}}$, where $func_{f_{pb}}$ is a function

Let $pb_i \equiv (func_{f_{pb}}, inputs_{f_{pb}}, output_{f_{pb}})$

1.2.2 Declaration of all program blocks

Let $Blocks(P) \equiv \{pb_i \mid pb_i \text{ is a program block of } P\}$

1.2.3 Declaration of a test

Let $inputs_{f_t} \equiv \{v \mid v \text{ is a value}\}$

Let $func_{f_t}: inputs_{f_t} \rightarrow output_{f_t}$, where $func_{f_t}$ is a function

Let $t_j \equiv (func_{f_t}, inputs_{f_t}, output_{f_t})$

1.2.4 Declaration of all test sets for P

Let $TA_{pb_i} \equiv [t_j \mid t_j \text{ is a test}]$

Let $TS_P \equiv \bigcup_{pb_i \in Blocks(P)} TA_{pb_i}$

1.2.5 Show the validity and success of a test on a program block

$success(t_j, pb_i) \equiv \begin{cases} 1, & \text{if } [t_j]_{output_{f_t}} = [pb_i]_{output_{f_{pb}}} \\ 0, & \text{otherwise} \end{cases}$

$valid(t_j, pb_i) \equiv \begin{cases} 1, & \text{if } [t_j]_{func_{f_t}} = [pb_i]_{func_{f_{pb}}} \text{ and } [t_j]_{inputs_{f_t}} = [pb_i]_{inputs_{f_{pb}}} \\ 0, & \text{otherwise} \end{cases}$

$valid(t_j, pb_i) \rightarrow success(t_j, pb_i) = 1$

1.2.6 Proof that the whole program is correct

$correctBlock(TA_{pb_i}, pb_i) \equiv \forall t_j \in TA_{pb_i} \mid valid(t_j, pb_i)$

$correctP(TS_P, Blocks(P)) \equiv$

$\forall pb_i \in Blocks(P), \forall TA_{pb_i} \in TS_P \mid correctBlock(TA_{pb_i}, pb_i)$

2. Proof correctness of program from requirements point of view.

2.1 Considerate the previous parts 1.2.1 to 1.2.4 include as the first part of this proof.

2.2 Declaration of requirements

Let $condition \equiv [v \mid v \text{ is a value}]$

Let $conditionsPre \equiv \{condition \mid condition \text{ is a set of values}\}$

Let $conditionsPost \equiv \{condition \mid condition \text{ is a set of values}\}$

Let $pre - condition([t_j]_{inputs_{ft}}) \equiv$
 $\begin{cases} 1, \forall input \in [t_j]_{inputs_{ft}}, \exists condition \in conditionsPre \mid input \in condition \\ 0, otherwise \end{cases}$

Let $post - condition([t_j]_{output_{ft}}) \equiv$
 $\begin{cases} 1, \forall output \in [t_j]_{output_{ft}}, \exists condition \in conditionsPost \mid output \in condition \\ 0, otherwise \end{cases}$

Let $r_i \equiv (pre - condition([t_j]_{inputs_{ft}}), post - condition([t_j]_{output_{ft}}))$

Let $RA_{pb_i} \equiv [r_i \mid r_i \text{ is a requirement}]$

Let $RS_P \equiv \bigcup_{pb_i \in Blocks(P)} RS_{pb_i}$

2.3 Test satisfaction of a requirement

$inputsEquality(t_j, pb_i) \equiv \begin{cases} 1, if [t_j]_{inputs_{ft}} = [pb_i]_{inputs_{f_{pb}}} \\ 0, otherwise \end{cases}$

$outputsEquality(t_j, pb_i) \equiv \begin{cases} 1, if [t_j]_{output_{ft}} = [pb_i]_{output_{f_{pb}}} \\ 0, otherwise \end{cases}$

$conditionInput(t_j, pb_i, r_i) \equiv \begin{cases} 1, if inputsEquality(t_j, pb_i) \mid [r_i]_{pre-condition}([pb_i]_{inputs_{f_{pb}}}) \\ 0, otherwise \end{cases}$

$conditionOutput(t_j, pb_i, r_i) \equiv \begin{cases} 1, if outputsEquality(t_j, pb_i) \mid [r_i]_{post-condition}([pb_i]_{output_{f_{pb}}}) \\ 0, otherwise \end{cases}$

$satisfy(t_j, pb_i, r_i) \equiv$
 $\begin{cases} 1, if conditionInput(t_j, pb_i, r_i) \text{ and } conditionOutput(t_j, pb_i, r_i) \text{ and } valid(t_j, pb_i) \\ 0, otherwise \end{cases}$

2.4 A specific requirement is satisfied for a program block

Let $satisfactionR(TS_{pb_i}, pb_i, r_i) \equiv$
 $\begin{cases} 1, if \forall t_j \in TS_{pb_i} \mid satisfy(t_j, pb_i, r_i) \rightarrow (\exists t_j \in TS_{pb_i} \mid satisfy(t_j, pb_i, r_i) = 1) \\ 0, otherwise \end{cases}$

2.5 All requirements are satisfied for a program block

Let $satisfactionSetR(TS_{pb_i}, pb_i, RS_{pb_i}) \equiv$
 $\begin{cases} 1, if \forall t_j \in TS_{pb_i}, \forall r_i \in RS_{pb_i} \mid satisfactionR(TS_{pb_i}, pb_i, r_i) = 1 \\ 0, otherwise \end{cases}$

2.6 All requirements are satisfied for a program

Let $satisfactionALL \equiv \forall pb_i \in Blocks(P), \forall TS_{pb_i} \in TS_P, RS_{pb_i} \in RS_P \mid satisfactionSetR(TS_{pb_i}, pb_i, RS_{pb_i}) = 1$