

## 1. Declaration of program and validity of syntax

$sT = \{c \mid c \text{ is a character}\}$

$sR = \{sT \mid sT \text{ is a statement}\}$

Let  $\text{syntaxRules} = \{sR \mid sR \text{ is a rule}\}$

Assume semantic of a language

Let  $L = (\text{syntax}, \text{semantic})$ , where  $\text{syntax} = \text{syntaxRules}$

$\text{compile}(\text{expression}, \text{syntaxRules}) \rightarrow \begin{cases} 1, & \text{if } \text{expression} \in \text{syntaxRules} \\ 0, & \text{otherwise} \end{cases}$

Let  $\text{exp} = \{c \mid c \text{ is a character}\}$

Let  $\text{expression} = \{\text{exp} \mid \text{exp is a chain of characters}\}$

$\text{correctSyntax}(\text{expressions}) = \forall \text{expression} \in \text{expressions} \mid \text{compile}(\text{expression}, \text{syntax})$

$\text{isSyntaxValid}(P) \rightarrow \begin{cases} 1, & \text{if } \text{syntax}(\text{expressions}) \\ 0, & \text{otherwise} \end{cases}$

$\text{SyntaxValid}(P) \text{ iff } \text{isSyntaxValid}(P) = 1$

$P = \{\text{expressions}\}$

## 2. Correctness of the program by using unit tests

Let  $\text{inputs}_{ft} = \{v \mid v \text{ is a value}\}$

Let  $\text{func}_{ft}: \text{inputs}_{ft} \rightarrow \text{output}_{ft}$ , where  $\text{func}_{ft}$  is a function

Let  $t_j = (\text{func}_{ft}, \text{inputs}_{ft}, \text{output}_{ft})$

Let  $ts_i = \{t_j \mid t_j \text{ is a test}\}$

Let  $\text{inputs}_{f_{pb}} = \{v \mid v \text{ is a value}\}$

Let  $\text{func}_{f_{pb}}: \text{inputs}_{f_{pb}} \rightarrow \text{output}_{f_{pb}}$ , where  $\text{func}_{f_{pb}}$  is a function

Let  $pb_i = (\text{func}_{f_{pb}}, \text{inputs}_{f_{pb}}, \text{output}_{f_{pb}})$

Let  $P = \text{a program to test}$

Let  $\text{Blocks}(P) = \{pb_i \mid pb_i \text{ is a program block of } P\}$

Let  $TS_{pb_i} = \{ts_i \mid ts_i \text{ is a related set of tests for } pb_i\}$

Let  $TS_P = \bigcup_{pb_i \in \text{Blocks}(P)} TS_{pb_i}$

$\text{success}(t_j, pb_i) \rightarrow \begin{cases} 1, & \text{if } [t_j]_{\text{output}_{ft}} = [pb_i]_{\text{output}_{f_{pb}}} \\ 0, & \text{otherwise} \end{cases}$

Let  $p = \text{boolean variable s.t } p \rightarrow \begin{cases} 1, & \text{if True} \\ 0, & \text{otherwise} \end{cases}$

Let  $\varphi = p \wedge \neg p$

$$correct(ts_i, pb_i) = \forall t_j \in ts_i \mid success(t_j, pb_i) \text{ and } success(t_j, pb_i) \models \varphi$$

$$correctBlock(TS_{pb_i}, pb_i) = \forall ts_i \in TS_{pb_i} \mid correct(ts_i, pb_i)$$

$$correctProgram(TS_P, Blocks(P)) = \forall pb_i \in Blocks(P), \forall TS_{pb_i} \in TS_P \mid correctBlock(TS_{pb_i}, pb_i)$$

### 3. Requirements Satisfaction

$$pre - condition = \{v \mid v \text{ is a value}\}$$

$$post - condition = \text{is a value}$$

$$\text{Let } r_i = (pre - condition, post - condition)$$

$$\text{Let } RS_{pb_i} = \{r_i \mid r_i \text{ is a requirement}\}$$

$$\text{Let } RS_P = \{RS_{pb_i} \mid RS_{pb_i} \text{ is a set of requirements for a program block } i\}$$

$$conditionInput = [t_j]_{input_{ft}} = [pb_i]_{input_{f_{pb}}} = [r_i]_{pre-condition}$$

$$conditionOutput = success(t_j, pb_i) = 1 \text{ and } [t_j]_{output_{ft}} = [pb_i]_{output_{f_{pb}}} = [r_i]_{post-condition}$$

$$satisfy(success(t_j, pb_i), r_i) \rightarrow \begin{cases} 1, & \text{if } conditionOutput \wedge conditionInput \\ 0, & \text{otherwise} \end{cases}$$

$$satisfaction(pb_i, ts_i, r_i) = \forall t_j \in ts_i \mid satisfy(success(t_j, pb_i), r_i)$$

$$satisfyBlock(TS_{pb_i}, pb_i, RS_{pb_i}) = \forall ts_i \in TS_{pb_i}, r_i \in RS_{pb_i} \mid satisfaction(pb_i, ts_i, r_i)$$

$$satisfyP(TS_P, Blocks(P), RS_P) = \forall pb_i \in Blocks(P), \forall TS_{pb_i} \in TS_P, \forall RS_{pb_i} \in RS_P \mid satisfyBlock(TS_{pb_i}, pb_i, RS_{pb_i})$$