## Church-Turing Thesis

All "reasonable" model of computation are equivalent to Turing machines.

A model of computation A is <u>equivalent</u> to a model of computation B if we can

- 1. simulate every A-machine with a B-machine and
- 2. simulate every B-machine with an A-machine.

Examples we've seen already...

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- 3. PDAs requiring empty stacks and PDAs that don't require empty stacks

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For example, most programming languages are Turing complete.

# Can we "upgrade" the Turing machine?

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- Adding more heads
- Adding non-determinism
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We do not know if there is a real model of computation that recognizes languages that are not Turing-recognizable.