

Programming and Algorithms II

Dynamic Programming

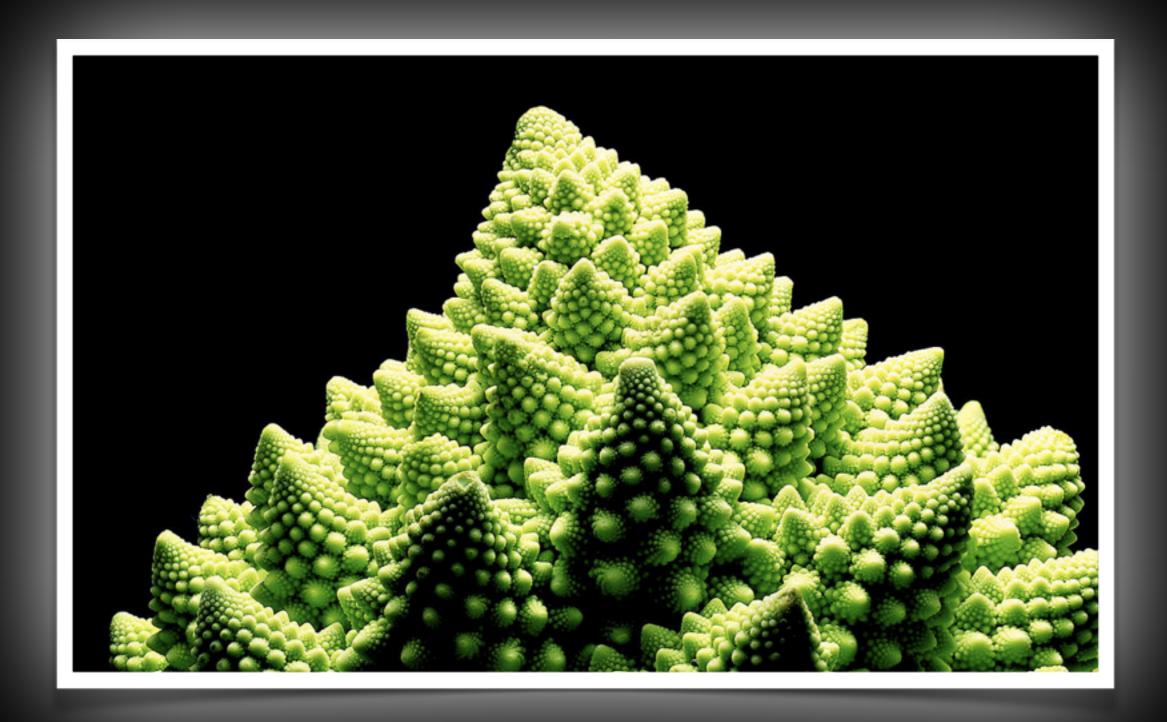
Nicolas Wu

nicolas.wu@bristol.ac.uk

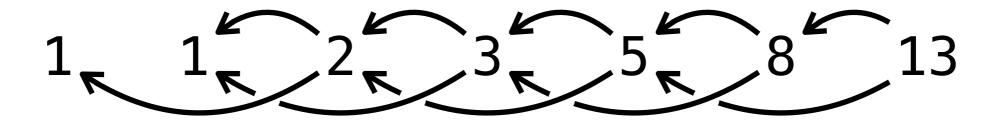
Department of Computer Science University of Bristol

Dynamic Programming

- start with naive recursive definition
- notice repeated subproblems
- tabulate results
- lookup rather than recurse

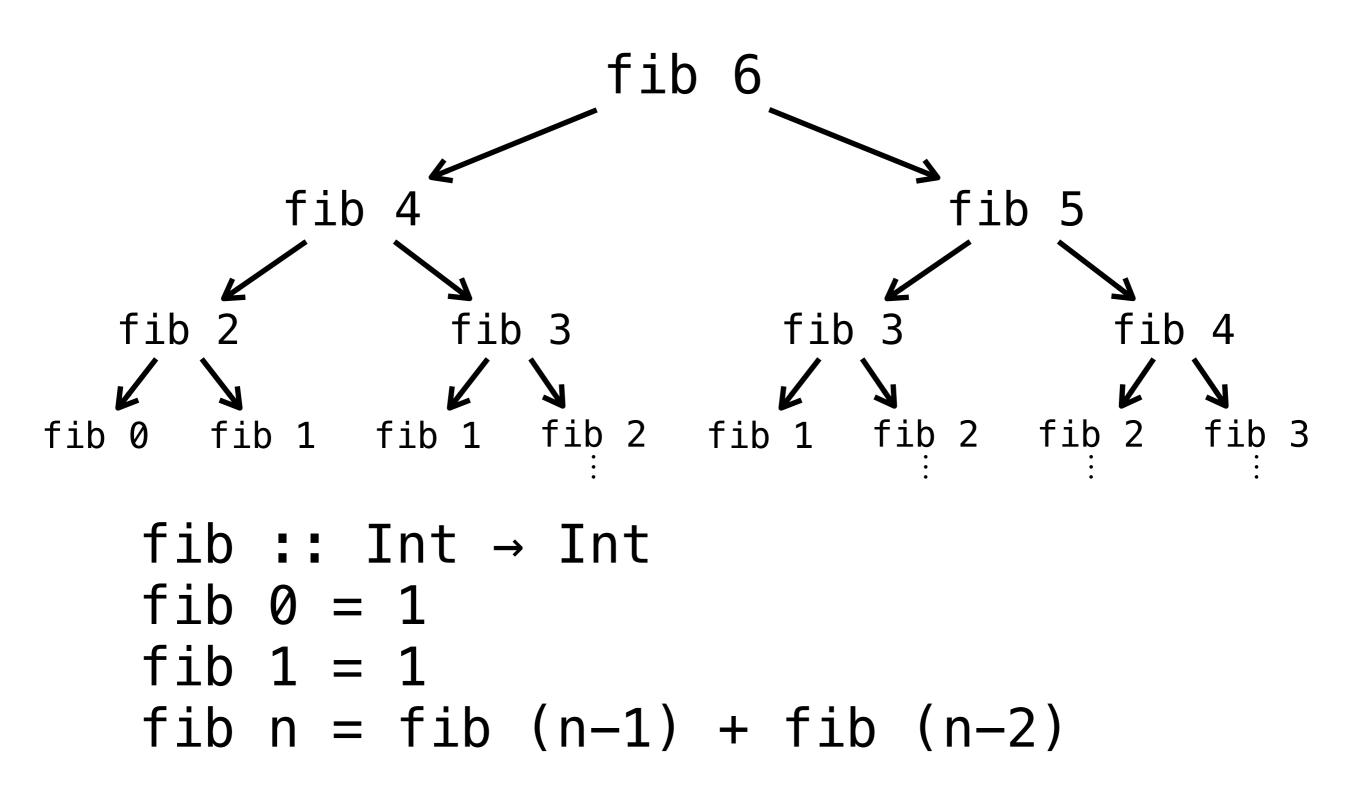


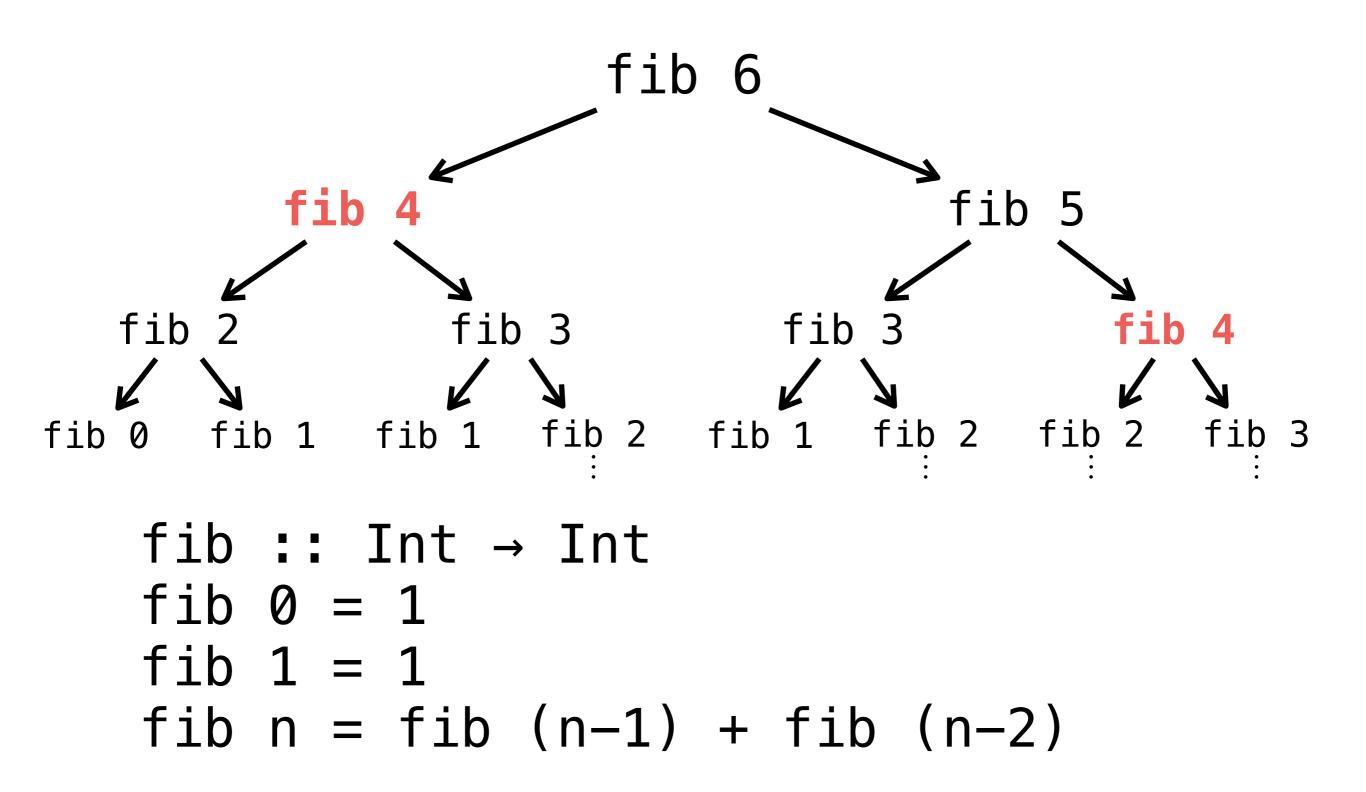
fibonacci

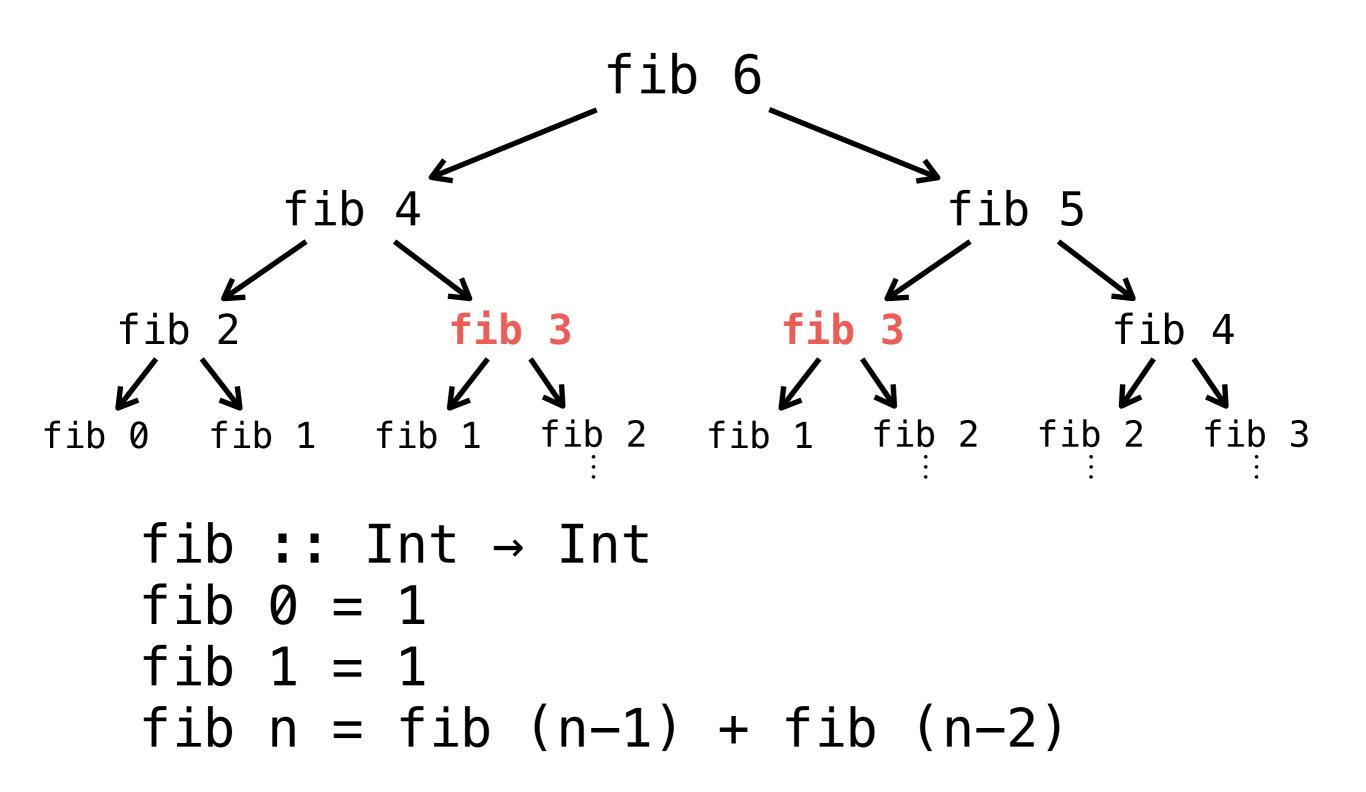


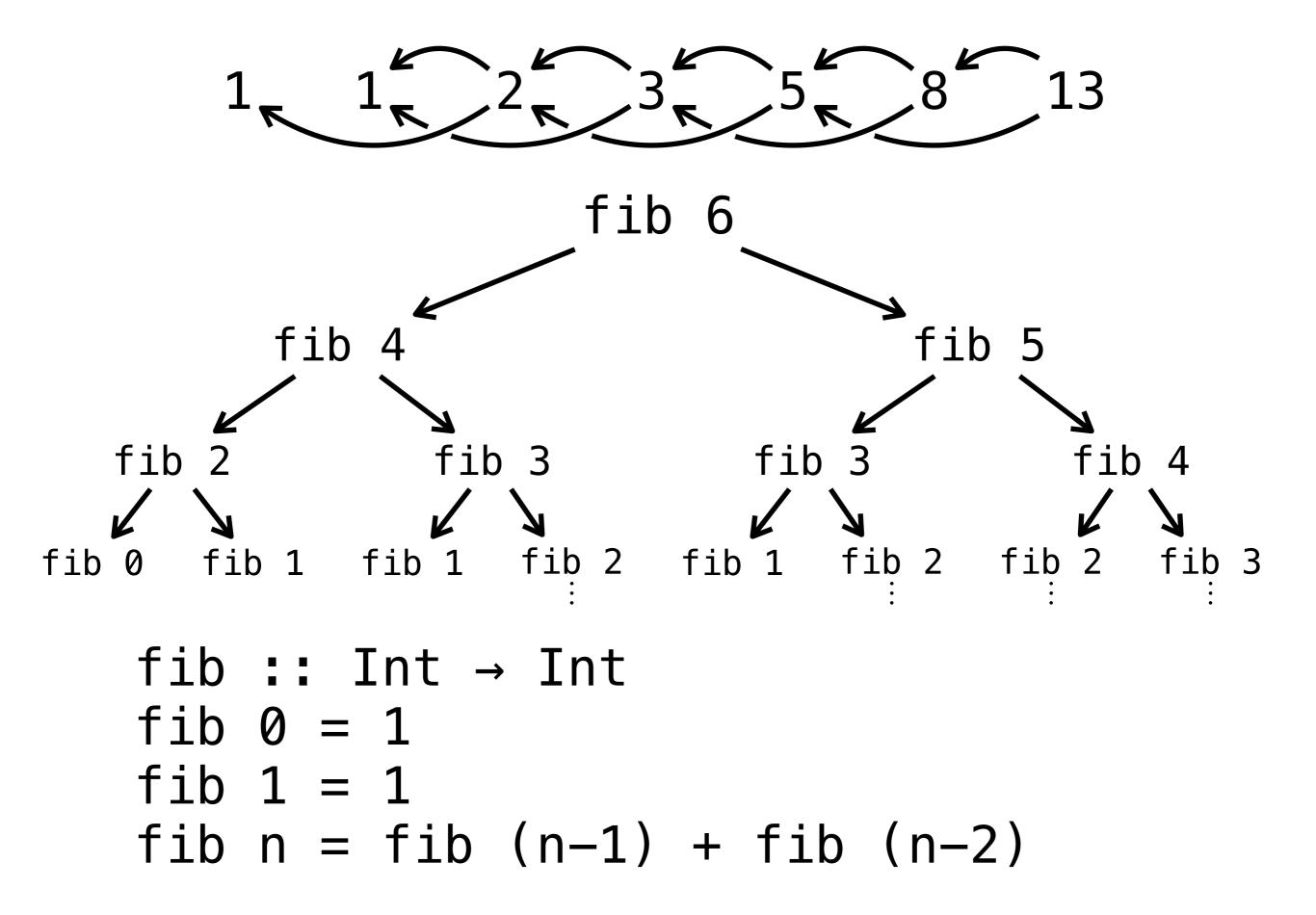
$$1 \frac{2}{5} \frac{3}{5} \frac{8}{13}$$

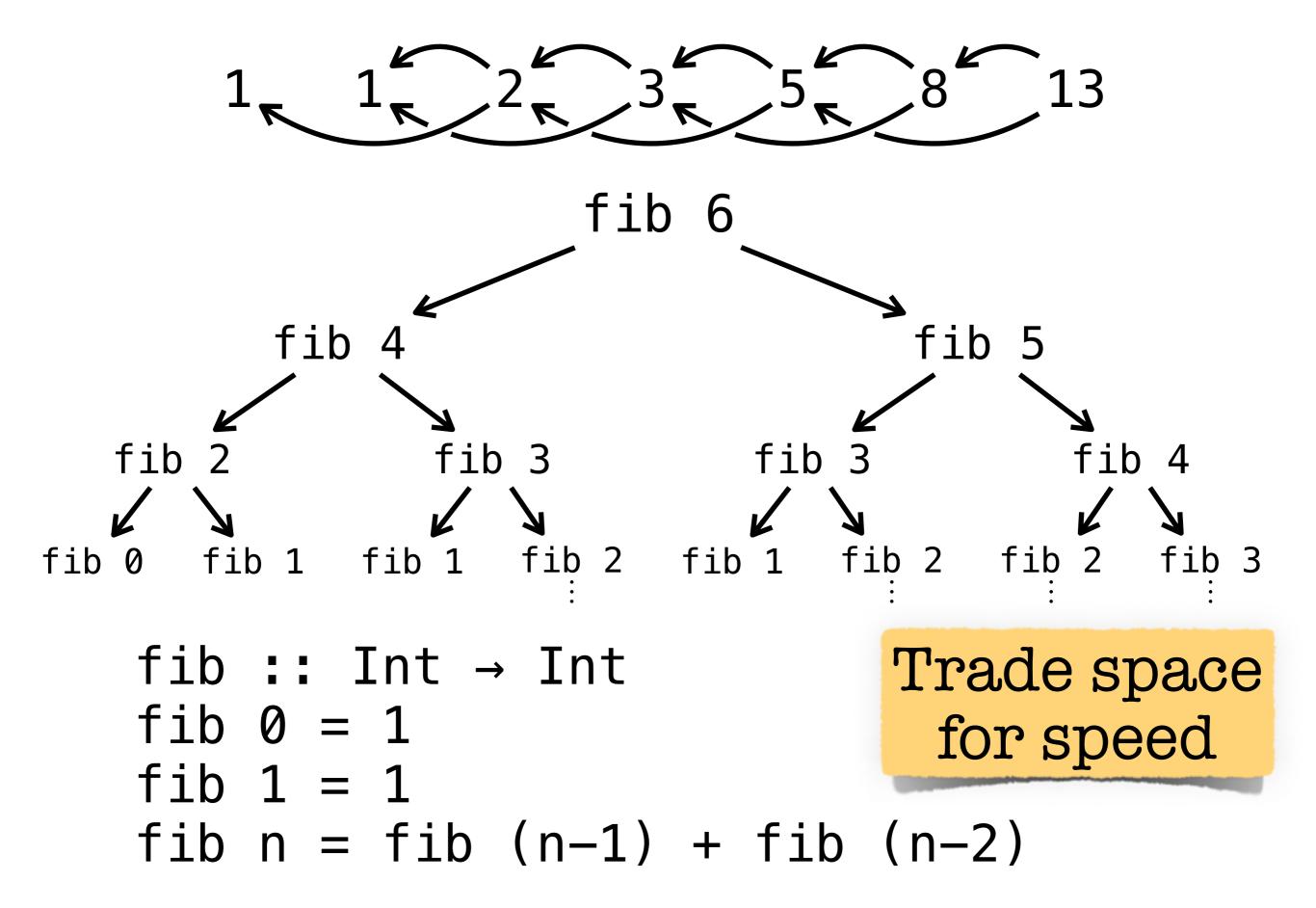
```
fib :: Int → Int
fib 0 = 1
fib 1 = 1
fib n = fib (n-1) + fib (n-2)
```











```
fib n = table!n where

,
,
table! table!
```

```
fib :: Int \rightarrow Int

fib 0 = 1

fib 1 = 1

fib n = fib (n-1) + fib (n-2)
```

```
fib :: Int → Int
fib n = table!n where
  table = tabulate (0,n) fib'
  fib' 0 = 1
  fib' 1 = 1
  fib' n = table!(n-1) + table!(n-2)
```

```
fib :: Int \rightarrow Int

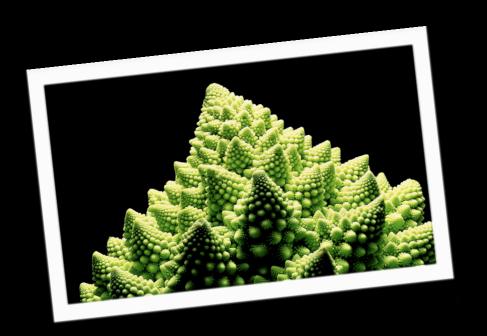
fib 0 = 1

fib 1 = 1

fib n = fib (n-1) + fib (n-2)
```

```
fib :: Int → Int
fib n = table!n where
  table = tabulate (0,n) fib'
  fib' 0 = 1
  fib' 1 = 1
  fib' n = table!(n-1) + table!(n-2)
```

```
tabulate :: Ix i \Rightarrow (i,i) \rightarrow (i \rightarrow a) \rightarrow Array i a tabulate (m,n) f = array (m,n) [ (i, f i) | i \leftarrow [m.n]] array :: Ix i \Rightarrow (i,i) \rightarrow [(i, a)] \rightarrow Array i a
```



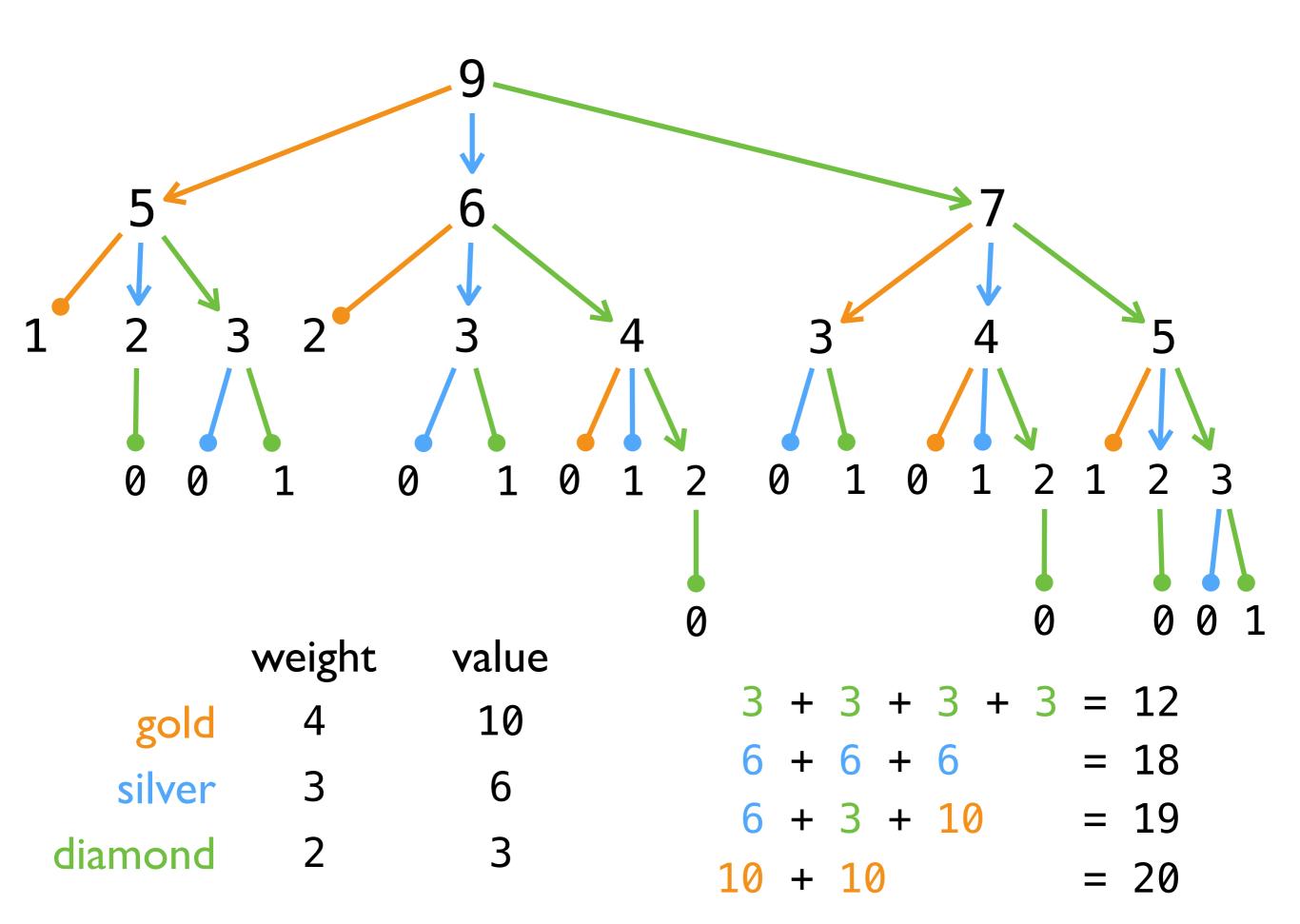
Recap

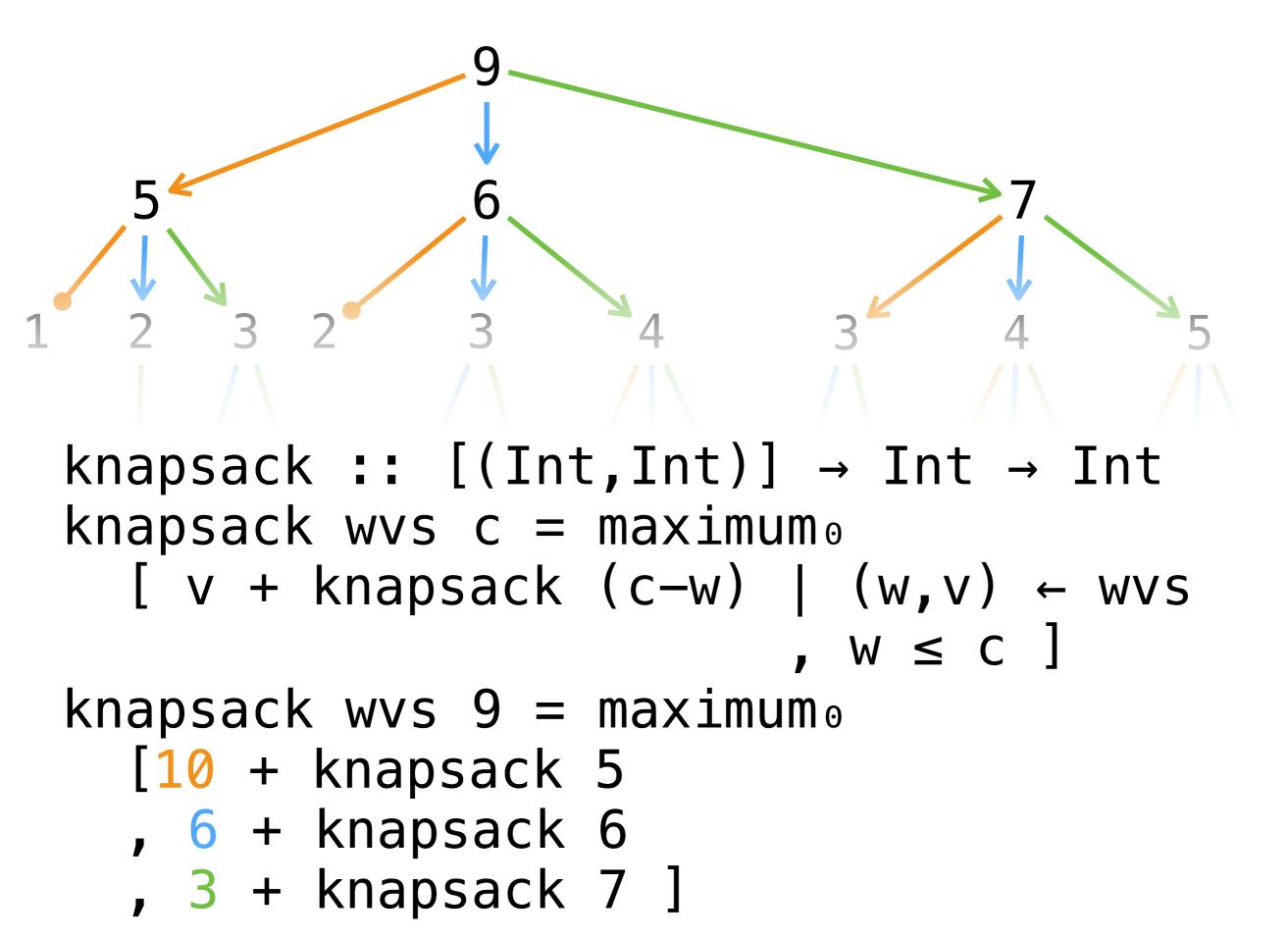
- start with naive recursive definition
- notice repeated subproblems
- tabulate results
- lookup rather than recurse

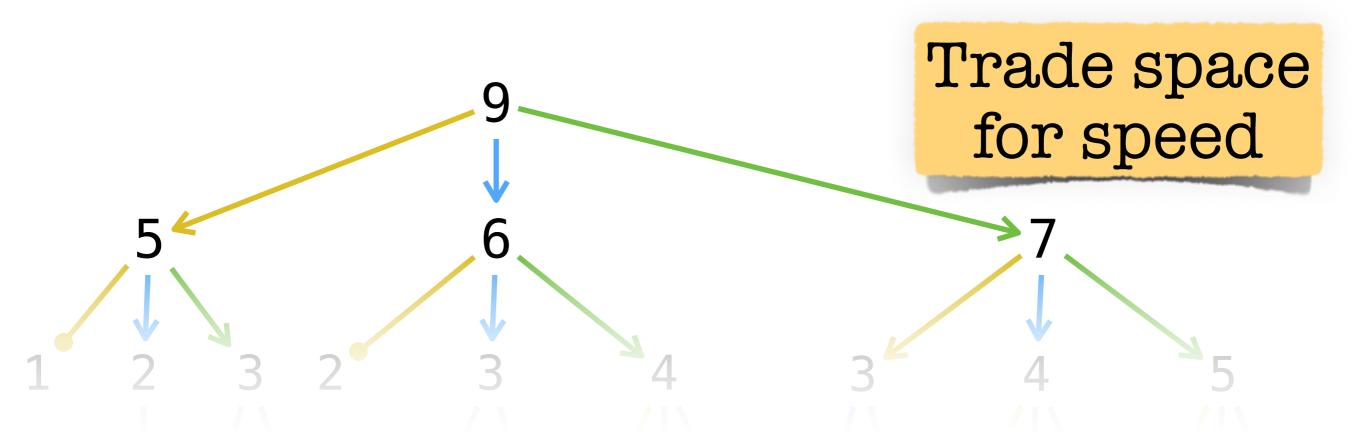


the knapsack problem

	weight	value
gold	4	10
silver	3	6
diamond	2	3







```
knapsack :: [(Int,Int)] → Int → Int
knapsack wvs c = maximum₀
  [v + knapsack wvs(c-w) | (w,v) ← wvs
, w ≤ c]
```

Trade space for speed

knapsack wvs c = table!c where

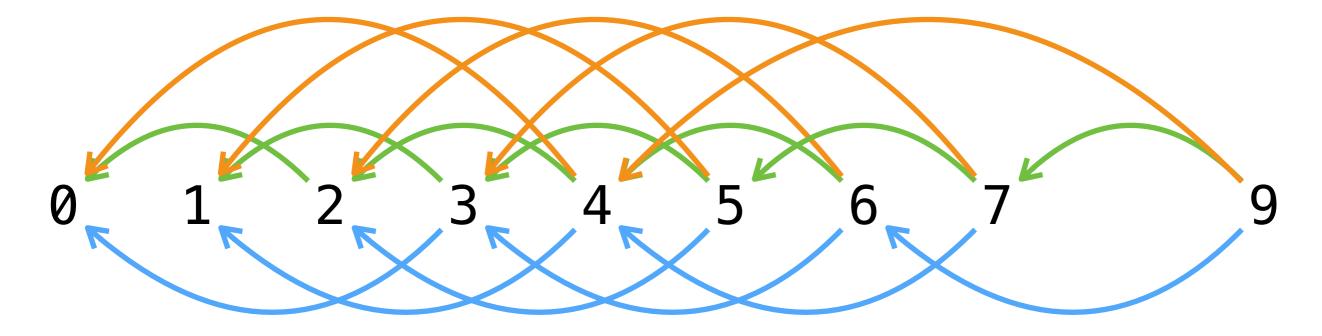
table!

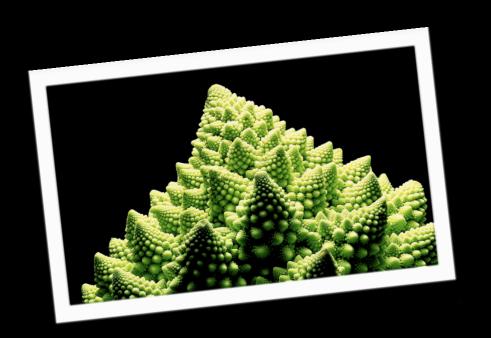
```
knapsack :: [(Int,Int)] → Int → Int
knapsack wvs c = maximum₀
  [v + knapsack (c-w) | (w,v) ← wvs
, w ≤ c]
```

```
knapsack :: [(Int,Int)] → Int → Int
knapsack wvs c = table!c where
  table = tabulate (0,c) knapsack'
  knapsack' = maximum₀
  [v + table!(c-w) | (w,v) ← wvs
  , w ≤ c]
```

```
knapsack :: [(Int,Int)] → Int → Int
knapsack wvs c = maximum₀
  [v + knapsack (c-w) | (w,v) ← wvs
, w ≤ c]
```

```
knapsack :: [(Int,Int)] → Int → Int
knapsack wvs c = table!c where
  table = tabulate (0,c) knapsack'
  knapsack' = maximum₀
  [ v + table!(c-w) | (w,v) ← wvs
  , w ≤ c ]
```





Recap



- start with naive recursive definition
- notice repeated subproblems
- tabulate results

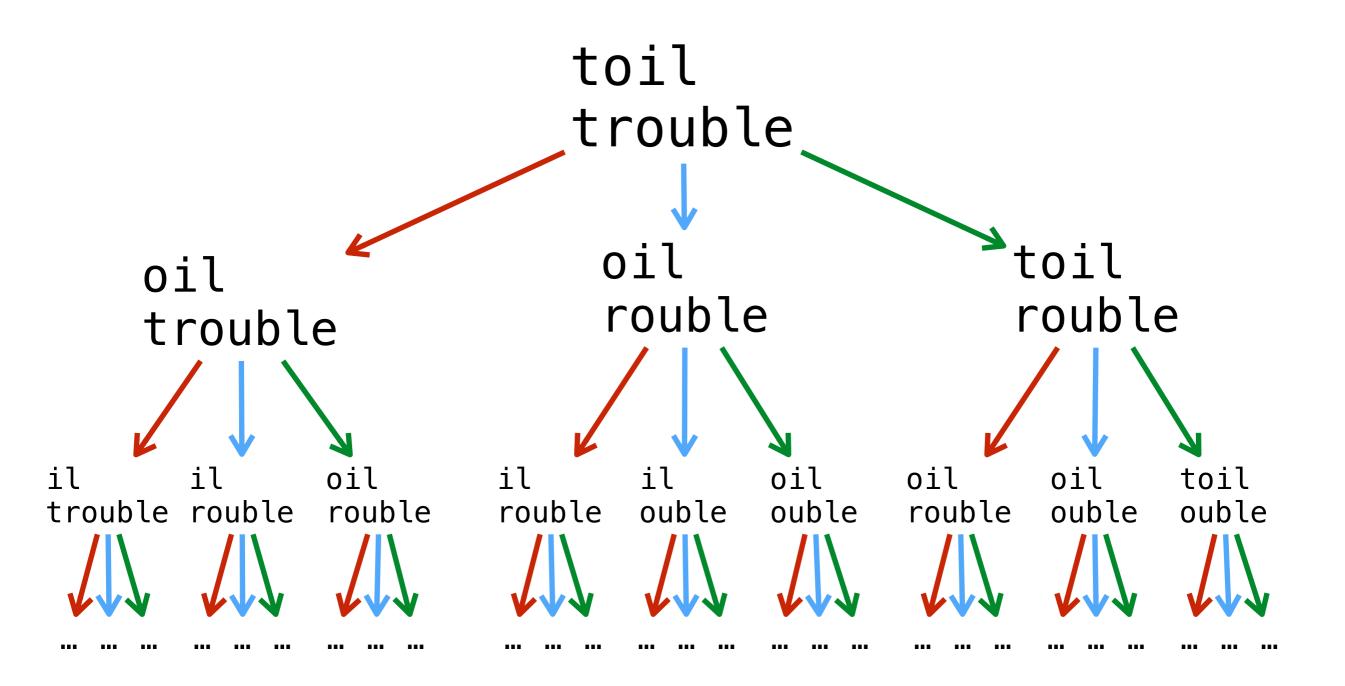
lookup rather than recurse

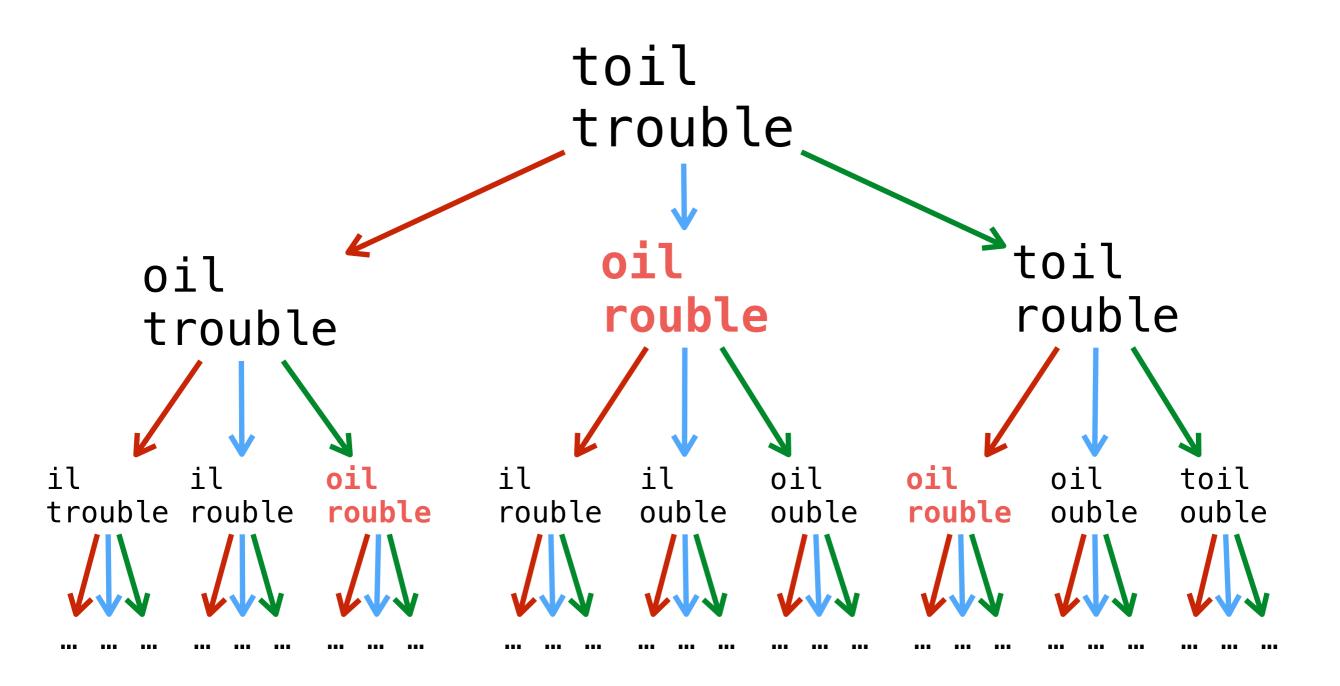


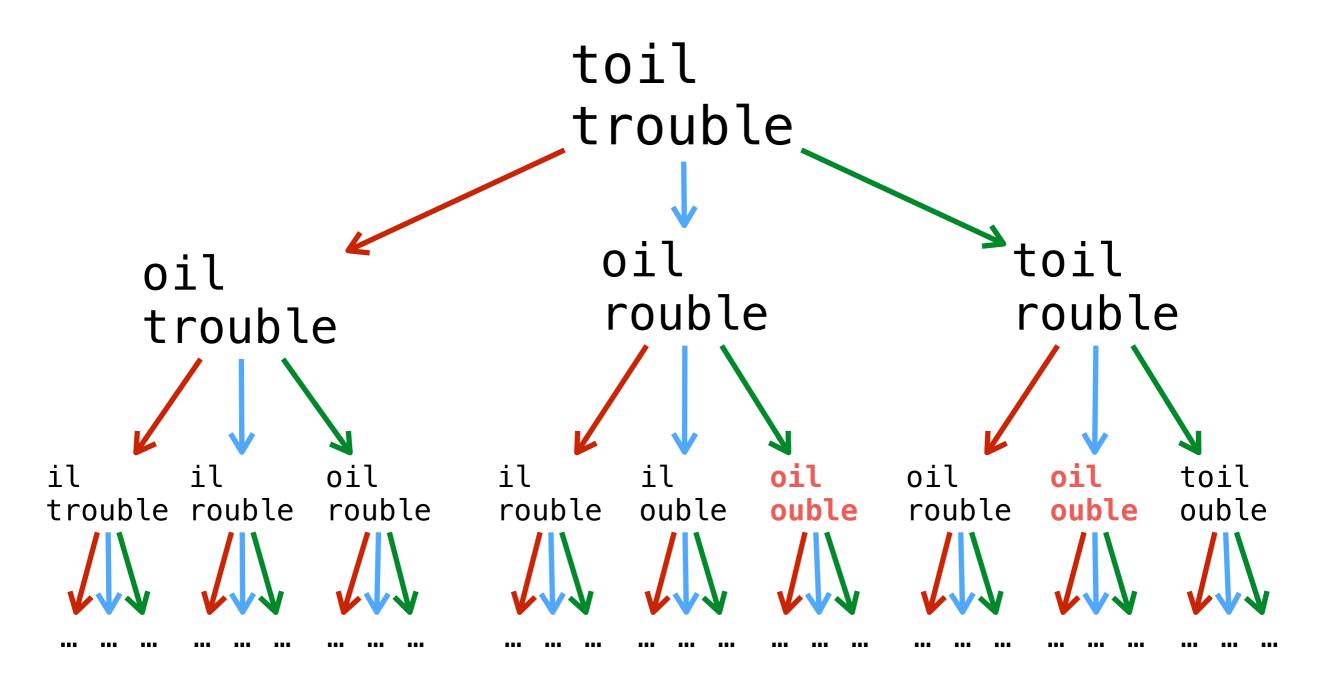
the edit-distance problem

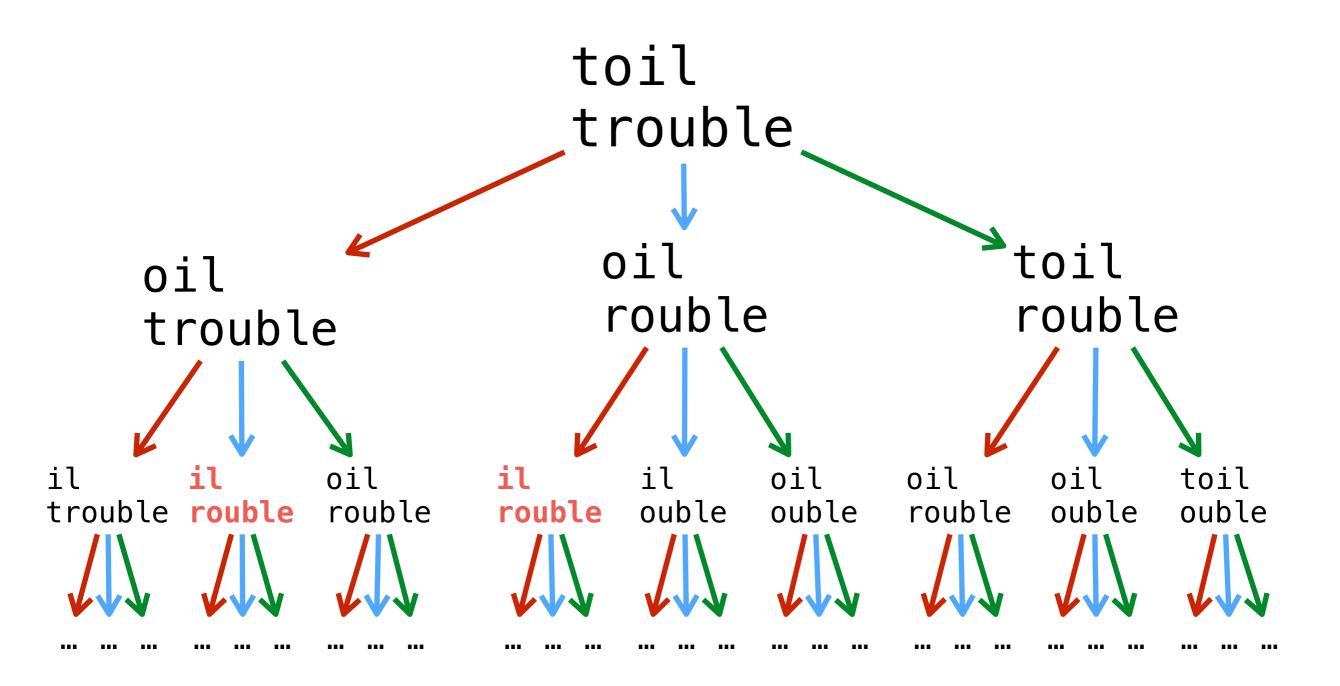
toil
troil
troul
trouble

trouble touble toile toile toil





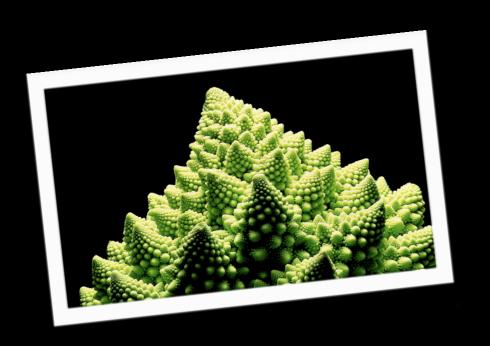




```
toil
                   trouble
                    oil
                                    toil
   oil
                    rouble
                                    rouble
   trouble
dist :: String → String → Int
dist xs
                  = length xs
dist[] ys = length ys
dist(x:xs)(y:ys) = minimum
  [ dist xs (y:ys) + 1, dist (x:xs) ys + 1
  , dist xs ys + if x == y then 0 else 1 ]
```

```
toil
                   trouble
                    oil
                                    toil
   oil
                    rouble
                                    rouble
   trouble
dist :: String → String → Int
dist xs
                  = length xs
dist[] ys = length ys
dist(x:xs)(y:ys) = minimum
  [ dist xs (y:ys) + 1, dist (x:xs) ys + 1
  , dist xs ys + if x == y then 0 else 1 ]
```

```
dist :: String → String → Int
dist xs ys = table!mn where
 mn = (length xs, length ys)
table = tabulate ((0,0), mn) dist'
 dist'i 0 = i
 dist' 0 j = j
 dist'(i+1)(j+1) = minimum
  [ table!(i,j+1) + 1, table!(i+1,j) + 1
  , table!(i,j) + if xs!i == ys!j then 0 else 1 ]
dist :: String → String → Int
dist xs [] = length xs
dist [] ys = length ys
dist (x:xs) (y:ys) = minimum
  [ dist xs (y:ys) + 1, dist (x:xs) ys + 1
  , dist xs ys + if x == y then 0 else 1 ]
```



Recap



- start with naive recursive definition
- notice repeated subproblems
- tabulate results
- lookup rather than recurse

