

Content of Course

- Prime and Hailstone sequences
- RSA encryption
- Chaotic Dynamical Systems.

Part I Primes and Hailstone Sequences.

Prime Numbers

- ~~A number that is divisible only by itself and one~~
- ~~A positive integer greater than 1 that is divisible only by itself and one.~~

Twin Primes

- ~~A pair of primes separated by 2~~
- ~~Eg. {3, 5}, {17, 19}~~

- Prime Number - A positive integer greater than that is divisible by itself and one.

- Twin Prime - A pair of primes separated by 2.
 - Ex {3, 5}, {17, 19}.

- Theorem - There are infinitely many primes.

• Euclid's Proof - Suppose that there are finitely many primes (p_1, p_2, \dots, p_n). Define $a = p_1, p_2, \dots, p_n + 1$.

'a' is not in the list, thus 'a' is not a prime.

Thus there is a p_i from the list that divides a.

Now consider $a - p_1, p_2, \dots, p_n = 1$, but p_i divides 'a' and p_1, p_2, \dots, p_n ; thus p_i divides $(a - p_1, p_2, \dots, p_n)$. Therefore p_i divides 1 (non-prime), thus there are infinitely many primes. (by contradiction).

Algebra

→ Goldbach Conjecture - any even number can be written as the sum of two primes.

$$- 4 = 2+2, \quad 6 = 3+3, \quad 12 = 5+7,$$

$$\blacksquare 8 = 5+3, \quad 10 = 5+5, \quad 14 = 7+7,$$

• Conjecture - no proof, no counter example.

• Fermat's Last Theorem - The equation $x^n + y^n = z^n$ has no solutions in positive integers for $n > 2$.

• Fermat-Wiles Last - $1^n + 2^n = 3^n, \quad 3^2 + 4^2 = 5^2, \quad n \neq 3, 4, \dots$

• Fermat's-Wiles Last Theorem - Wiles contributed the proof to Fermat's Conjecture.

• How to check for prime number?

- Divide said number by all of the primes leading up to it. (based on \Rightarrow)

• Fundamental Theorem of Arithmetic.

- Any positive integer, $a \geq 1$, can be written in a unique way as a product of primes.

- $a = p_1^{e_1} \cdots p_k^{e_k}$, p_i prime, e_i positive integer

- Ex. $12 = 2^2 \cdot 3^1$

• Let $n = ab$, and claim that $a \leq \sqrt{n}$ or $b \leq \sqrt{n}$. If not (by way of contradiction) then $a > \sqrt{n}$ and $b > \sqrt{n}$. This means that $ab > \sqrt{n} \cdot \sqrt{n} \Rightarrow n > n$, nonsense. Therefore $a \leq \sqrt{n}$ or $b \leq \sqrt{n}$. Therefore it is sufficient to check divisions up to \sqrt{n} to determine whether or not n is prime.

Hailstone Sequences

• Ex. 4, 2, 1; length=2

Ex. 10, 5, 16, 8, 4, 2, 1; length=6

• Hailstone Sequence - ① take any positive integer, a

② if a is even; divide by 2 to get next term

if a is odd; multiply by 3 and add 1 to get next

③ Repeat till 1 and stop.

Seminar

• Number Theory (\mathbb{Z})

- specifically positive integers.

• Prime Numbers

- twin primes
- Conjecture & theorem
- proof & counter-example.

• Hailstone Sequences.

- for $n=6$: 6, 3, 10, 5, 16, 8, 4, 2, 1. ~~for $n=7$~~



Programming

Loops and Conditional Controls

• A-Loops Repetitive operations

For ()

(statements)

Next

• Eg:- Dim i, Count As Integer ' two variables are declared of type integers

Count = 0 ' Value 0 is assigned to count.

For i = 1 to 3 ' 1 is assigned to count.

Count = Count + 1 ' ~~Value~~ Vary with:

Next

Text ~~box~~ 1. Text = Count

Eg 2

Dim i, Sum As Integer
For i = 1 To 3

Sum += i ' Sum = Sum + i

Next

TextBox1.Text = Sum.

Eg 3

Dim i, Sum As Integer
For i = 1 To 3

Sum += i

Next

TextBox1.Text = Sum

Eg 4

Dim i, j As Integer

Dim Count As Long

Count = 0

{ For i = 1 To 2

 For j = 2 To 4

 Count += 1

 Next

Next,

TextBox1.Text = Count

* If

B - Conditional Controls

- IF (condition) Then

(statements)

End If

- OR IF (condition) Then

(statements)

Else

(statements)

End If.

Eg 3

Dim i, Count As Integer

i=5

Count=3

If i > 4 Then

Count *= 2

End If.

Textbox1.Text = Count.

Lecture Goal

- Write a program that gives the length of a Hailstone Sequence
 - Reminder: The Collatz Conjecture - let a_1 be a positive integer. Then the HS for which a_1 is the first term ends at 1 after finitely many steps.

- Eg. Take $a_1 = 10$, then the HS for 10 is

$$a_1 = 10 \quad a_2 = 4$$

$$a_2 = 5 \quad a_3 = 2$$

$$a_3 = 16 \quad a_4 = 1, \text{ length} = 6.$$

$$a_4 = 8$$

- Operation Mod.

- Method n gives the remainder when m is divided by n .

- $16 \bmod 5 = 1$.

- $23 \bmod 4 = 3$.

- Programming Tips:- Work incrementally.

- Code then test then code then test

- Given a_n , display a_{n+1}

Dim x As Integer.

$x = \text{Textbox1.Text}$.

If $x \bmod 2 = 0$ Then

$\text{Textbox2.Text} = \underline{x}$.

Else

$\text{Textbox2.Text} = 3 * x + 1$

End If

- b) Given a_n , determine & display a_{n+1}

Dim x As Integer

$x = \text{Textbox1.Text}$

If $x \bmod 2 = 0$ Then

~~$x = \frac{x}{2}$~~

Else

$x = 3 * x + 1$

End If

$\text{Textbox2.Text} = x$.

• c) Complete Code Idea
loop → repeat code (b)
loop { code (b)
until we reach 1

• Do While loops.

- When we don't know how many times we have to repeat but we know "until a condition is satisfied"
- Do While (condition)
(statements).

Loop.

- Bgs. Dim i As Integer
i = 1.

Do While i < 5
i = ~~i + 1~~ 2.

i | i
1 | 1, 3, 5.
~~2~~

Avoid infinite
loops.

• c) Complete Code Idea.

Dim seed As Integer
Seed = TextBox1.Text
Do While Seed > 1
 If Seed Mod 2 = 0 Then
 Seed = Seed / 2
 Else
 Seed = 3 * Seed + 1
 End If

Loop
TextBox2.Text = Seed.

Dim Seed, Length As Integer
Seed = TextBox1.Text
Length = 0
Do While Seed > 1
 If Seed Mod 2 = 0 Then
 Seed = Seed / 2
 Else
 Seed = 3 * Seed + 1
 End If
 Length = Length + 1
Loop.
TextBox2.Text = Length.

Lecture Goal

- To create a program to determine the maximum length of a hailstone for $n \leq 10,000$.
- Needs:
 - code of length of the HS for $n = \text{seed}$ [Review (last Wednesday) code]
 - for seed = 1 to 10,000.
 - max length to be compared with new length
- Code:

Dim Seed, i, Length, Upperbound, Maxlength, Maxseed As Long

Upperbound = Textbox1.Text 'eg. 10001

Maxlength = 0

Maxseed = 1.

for i = 1 To Upperbound

Length = 0

Seed = i

Do While Seed > 1

If Seed Mod 2 = 0 Then

Seed = Seed / 2

Else

Seed = 3 * Seed + 1

End If

Length = length + 1

Loop

If Length > Maxlength Then

Maxlength = length

Maxseed = i

End If

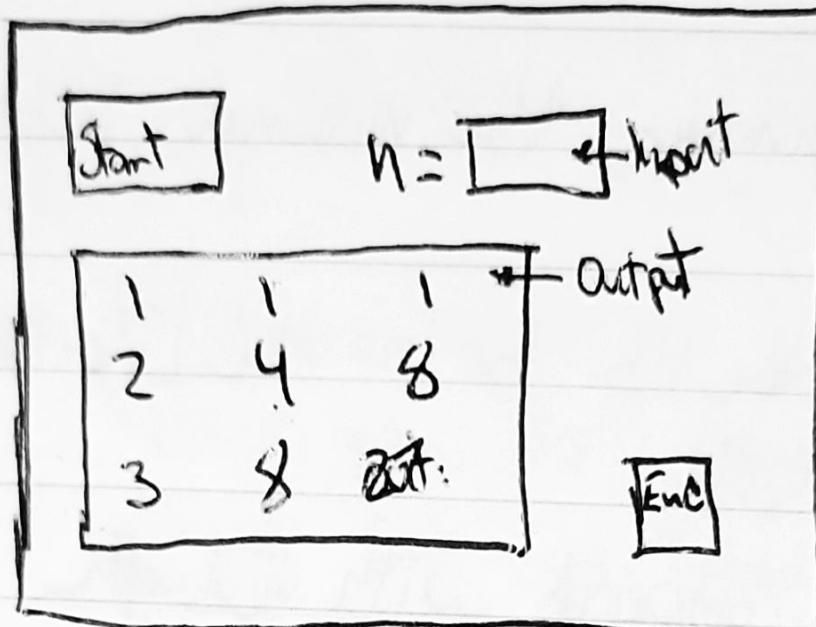
Next.

Textbox2.Text = Maxlength

Textbox3.Text = Maxseed

Interface to Show Many Output in a Table

Eg.



Output Box → Multiline = True.

Scrollbar = Vertical.

Part II RSA Encryption

- RSA \rightarrow Rivest-Shamir-Adleman (1971).
- Have to make an unbreakable code that anyone can encode ~~but~~ but only one person can decode ("trap door cipher")
- Algebra Background

- What is the remainder of $18/5$?

$$- 18 \bmod 5 = 3$$

$$\rightarrow \text{take } x = 23, y = 1004.$$

$$- x \bmod 5 = 3$$

$$y \bmod 5 = 4$$

$$- x+y = 23+1004$$

$$= 1027 \rightarrow 1027 \bmod 5 = 2$$

$$- xy = 23 \cdot 1004$$

$$= 23092 \rightarrow 23092 \bmod 5 = 2$$

$$x \bmod 5 + y \bmod 5 = 3+4$$

$$= 7 \rightarrow 7 \bmod 5 = 2.$$

$$(x \bmod 5)(y \bmod 5) = 12$$

$$= 12 \rightarrow 12 \bmod 5 = 2.$$

• $\bmod 5$

+	0	1	2	3	4
0	0	1	2	3	4
1	1	2	3	4	0
2	2	3	4	0	1
3	3	4	0	1	2
4	4	0	1	2	3

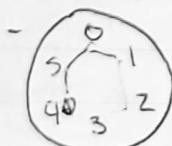
*	0	1	2	3	4
0	0	0	0	0	0
1	0	1	2	3	4
2	0	2	4	1	3
3	0	3	1	4	2
4	0	4	3	2	1

• Arithmetic in the integers modulus 5

- $\mathbb{Z}_5 = \{0, 1, 2, 3, 4\}$, together with $+, \cdot$

• Define - $\mathbb{Z}_n = \{0, 1, 2, \dots, n-1\}$, together with $+, \cdot$

"Clock" arithmetic



\mathbb{Z}_6

$$\textcircled{1} \quad 4+3 \bmod 6 = 1$$

$$\textcircled{2} \quad 4 \cdot 3 \bmod 6 = 0$$

• \mathbb{Z}_6

	0	1	2	3	4	5
0	0	0	0	0	0	0
1	0	1	2	3	4	5
2	0	2	4	1	3	5
3	0	3	0	2	4	1
4	0	4	1	3	5	2
5	0	5	4	3	2	1

$$xy = 16$$

For prime no. small integer factor \rightarrow no zeros.

* Inverse Element - find where the remainder is 1.

\mathbb{Z}_5	
0	none
1	1
2	3
3	2
4	4

\mathbb{Z}_6	
0	none
1	1
2	none
3	none
4	none
5	5

* \mathbb{Z}_6 - integers mod 6.

+	0	1	2	3	4	5
0	0	1	2	3	4	5
1	1	2	3	4	5	0
2	2	3	4	5	0	1
3	3	4	5	0	1	2
4	4	5	0	1	2	3
5	5	0	1	2	3	4

*	0	1	2	3	4	5
0	0	0	0	0	0	0
1	0	1	2	3	4	5
2	0	2	4	0	2	4
3	0	3	0	3	0	3
4	0	4	2	0	4	2
5	0	5	4	3	2	1

- $\mathbb{Z}_6 = \{0, 1, 2, 3, 4, 5\}$; together $+ \cdot$, the set of all ~~other~~ remainders when divided by 6.

- "Clock arithmetic."

- $26 \bmod 6 = 2$.

- $5 \cdot 5 \bmod 6 = 1$, thus 5 is the inverse of itself.

• ← Let's look at all the inverses in \mathbb{Z}_7 .

\mathbb{Z}_7 inverses	
0	none
1	1
2	4
3	5
4	2
5	3
6	6

* Observation - Similar to \mathbb{Z}_5 , every non-zero ~~non-zero~~ element in \mathbb{Z}_7 has an inverse.

* Definition - We say that a is congruent to b mod n , denoted by $a \equiv b \pmod{n}$, provided that $a \equiv b$ have the same ~~remainder~~ after divided by n .

• Note this is equivalent to say when $n \nmid (b-a)$
 "n divides evenly into $(b-a)$ "
 \Leftrightarrow no remainder for $\frac{b-a}{n}$.

- Eg ① $6 \equiv 8 \pmod{2}$.
 Since $2 \nmid (8-6)$

Or

Since 6 has a remainder of 0 when divided by 2 and so does 8.

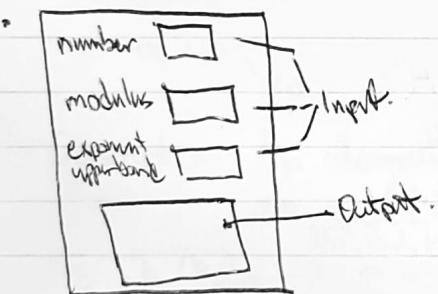
- Eg ② $6 \not\equiv 8 \pmod{3}$
 Since $3 \nmid (8-6) \pmod{2}$

- Eg ③ $24 \equiv 39 \pmod{5}$
 Since $5 \nmid (39-24)$

- Eg ④ $2^8 \equiv 1 \pmod{5}$
 $(2^4)^2 \equiv 1 \pmod{5}$.

$76 \equiv 1 \pmod{5} \checkmark$, so $2^4 \equiv 1 \pmod{5}$.

- Eg ⑤ $2^{1002} \equiv ? \pmod{5}$
 $(2^{4 \cdot 250}) \cdot 2^2 \equiv (1)^{250} \cdot 2^2 \pmod{5}$.
 Thus $2^{1002} \equiv 4 \pmod{5}$.



• Idea -

$3^0 = 1$	"power"	\rightarrow power = 1
$3^1 = 3$		for i = 1 to upperbound
$3^2 = 9 \equiv 4 \pmod{5}$		\rightarrow power = power * number Mod 5
$3^3 = 27 \equiv 2 \pmod{5}$		Next

add in the table for output

• Theorem - If $a \equiv b \pmod{n}$ and $c \equiv d \pmod{n}$ then $(a+c) \equiv (b+d) \pmod{n}$.

Proof - We want to show that $n \mid (b+d) - (a+c)$, but $(b+d) - (a+c) = (b-a) + (d-c)$
 thus $n \mid (b-a) + (d-c)$

Hilary

• Theorem - If $a \equiv b \pmod{n}$, and if $c \equiv d \pmod{n}$ then

$$a \cdot b \equiv b \cdot d \pmod{n}$$

• Proof - Idea similar to previous proof

$$n \mid (a-b) \Rightarrow n \mid (a-b) \cdot c$$

$$n \mid (c-d) \Rightarrow n \mid (c-d) \cdot b$$

Thus ... - (exercise)

• What do these theorems tell us?

- A result (in addition & multiplication) that holds
in \mathbb{Z} is also true in any \mathbb{Z}_n

- Exercise generate \mathbb{Z} examples for each theorem.

• Definition - An element $a \in \mathbb{Z}_n$ is called invertible in \mathbb{Z}_n if there is an element $b \in \mathbb{Z}$ such that

$$ab \equiv 1 \pmod{n}$$

• Eye 2 is invertible in \mathbb{Z}_5 . Since $2 \cdot 3 \equiv 1 \pmod{5}$ but
 ~~2 is not invertible in \mathbb{Z}_6 since $2 \cdot 3, 2 \cdot 4, 2 \cdot 5 \equiv 1 \pmod{6}$~~

• We are interested in finding out invertible elements in \mathbb{Z}_n . Where:

\mathbb{Z}_n	Invertible Element	Non-invertible Element
1		
2	1	0
3	1, 2	0
4	1, 3	0, 2
5	1, 2, 3, 4	0
6	1, 5	0, 2, 3, 4
7	1, 2, 3, 4, 5, 6	0
8	1, 3, 5, 7	0, 2, 4, 6
9	1, 2, 4, 5, 7, 8	0, 3, 6
10	1, 3, 7, 9	0, 2, 4, 5, 6, 8
11	1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11	0
12	1, 5, 7, 11	0, 2, 3, 4, 6, 8, 9, 10

• Observations

- Every non-zero element in \mathbb{Z}_p (p prime) is invertible
(theorem).

- 0 is always non-invertible 1 is always invertible in any \mathbb{Z}_n $n \geq 2$

- $n-1$ is always invertible in \mathbb{Z}_n

- Invertible elements in \mathbb{Z}_n never share prime factors with n .

- Definition - Two integers a & b are coprime or relatively prime if they have no common prime factors.
- Note: This is equivalent to say that their greatest common divisor is 1, here $\gcd(a, b) = 1$.

- Eg. $\gcd(16, 10) = 2$,
 $\gcd(11, 52) = 1$, are relatively prime (because 11 is invertible mod 52)

$$\gcd(24, 10) = 2,$$

- How can we find the gcd of any two numbers?

- Algorithm (Euclid)

- $24 = n10 + i$

$$= 2(10) + 4.$$

$$\overbrace{10}^{= n4 + i}$$

$$= (2)4 + 2.$$

$$\overbrace{4}^{= n2 + i}$$

$$2 = (2)1 + 0.$$

↑
This is our gcd

- $\gcd(52, 64) = 4$.

$$64 = n52 + i$$

$$64 = (1)52 + 12$$

$$\overbrace{52}^{= (4)12 + 4}$$

$$12 = (3)4$$

gcd

- Observation - An element $a \in \mathbb{Z}_n$ is invertible

$\Leftrightarrow \gcd(a, n) = 1$

- Eg. $7 \in \mathbb{Z}_{20}$ is invertible since $\gcd(7, 20) = 1$. Indeed
 $7 \cdot 3 \equiv 1 \pmod{20}$

$7 \in \mathbb{Z}_{35}$ is not invertible since $\gcd(7, 35) = 7 \neq 1$. 7 & 35 are not relatively prime.

The Euler φ function

- Define $\varphi(n)$ is the number of positive integers smaller than n that are relatively prime to n

Eg. $\varphi(3) = 2$	$1, \cancel{2}$
$\varphi(4) = 2$	$1, \cancel{2}, \cancel{3}$
$\varphi(5) = 4$	$1, \cancel{2}, \cancel{3}, \cancel{4}$
$\varphi(6) = 2$	$1, \cancel{2}, \cancel{3}, \cancel{4}, \cancel{5}$

- Notice: $\varphi(n)$ gives the number of invertible elements in \mathbb{Z}_n
- We will want to calculate $\varphi(n)$ for large n
- Question: if $ab = n$ then $\varphi(a)\varphi(b) = \varphi(n)$
Is φ multiplicative?

- We Explore:

a	b	ab	$\varphi(a)$	$\varphi(b)$	$\varphi(ab)$	Multiplicative?
3	3	9	2	2	4	No
2	2	14	6	1	6	Yes.
4	2	8	2	1	4	No
4	4	16	2	2	8	No
4	5	20	2	4	8	Yes
2	10	20	1	4	8	No

* Theorem: If $a \nmid b$ are relatively prime, then $\varphi(ab) = \varphi(a)\varphi(b)$

* Corollary: If $p \nmid q$ are two distinct primes, then

$$\varphi(pq) = (p-1)(q-1)$$

- Proof: $\varphi(pq) = \varphi(p)\varphi(q)$ since $\gcd(pq) = 1$
 $= (p-1)(q-1)$
since $p \nmid q$ are prime.

* Euler's Theorem - If $\gcd(a, n) = 1$, $a^{(\varphi(n))} \equiv 1 \pmod{n}$.

- Eg. $a=2, n=3$

$$\gcd(2, 3) = 1, \varphi(3) = 2, 2^2 \equiv 4 \pmod{3}$$

$$2^{(\varphi(3))} \equiv 1 \pmod{3} \Rightarrow 4 \equiv 1 \pmod{3}$$

$$\begin{aligned} \text{- Eg. } a=3, n=1,000,000,001 \\ \gcd(3, 1,000,000,001) = 1, 3^{(\varphi(1,000,000,001))} \equiv 1 \pmod{1,000,000,001} \end{aligned}$$

RSA Encryption

Rivest - Shamir - Adleman

- Alice $\xleftarrow{\text{message}}$ Bob a method so that everyone knows how to encode messages to Alice but only A knows how to decode it. "Trapdoor cipher"

• It is based on the fact that it is virtually impossible to factor very large numbers $> 10^{200}$

• Here is the procedure

• Part 1: Alice prepares her public key

1) Alice selects two distinct giant prime numbers $> 10^{100}$

$$\text{Eg. } p=23, q=41$$

2) find $n = pq$, $\Phi(n) = \Phi(pq)$

$$\text{Eg. } n = pq = 943$$

$$\Phi(n) = (p-1)(q-1) = 880$$

3) Alice chooses a number 'e' relatively prime to $\Phi(n)$

Eg. Alice selects $e=7$: when by $\gcd(7, 880)=1$

and Alice publishes the public key: (n, e) & everyone can see the public key.

Eg Alice publishes $(943, 7)$.

• Part 2: Bob sends encrypted secret message to Alice.

4) Assume the message is turned into numbers (ref: lab)

$$\text{Eg. } M=35$$

5) Bob encrypts message by calculation $M^e \bmod n = c$

~~Bob~~ $c = \text{"coded message"}$

$$\text{Eg. } 35^7 \bmod 943 = 545$$

Bob sends $c = 545$ to Alice and everyone can see it.

• Part 3: Back to Alice to decode the message received

6) Alice calculates the "privatekey" by finding d such that $e \cdot d \equiv 1 \pmod{\Phi(n)}$

Eg. Alice finds d such that

$$7d \equiv 1 \pmod{880}$$

$$d = 503 \text{ (and thus) secret}$$

7) Alice decodes message by calculating $c^d \bmod n = M$ E.g. $545^{503} \bmod 943 = 35$