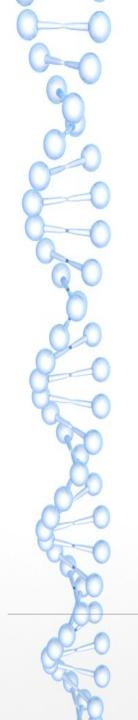
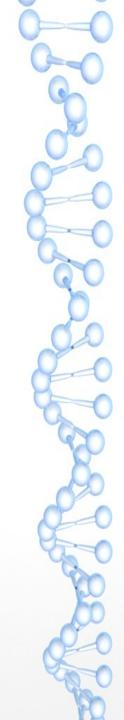


Jacopo Mondi March 2014



### **Summary**

- Power Management Overview
- Frequency Scaling in Android/Linux
- Linux Power Management
- Android Specificities



## Power Management Overview

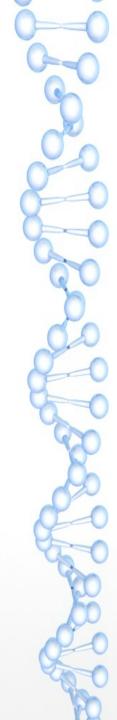
### Run-Time power consumption reduction

- Linux & Android:
- Dynamic Frequency Scaling with CPUFreq

#### **Power Management**

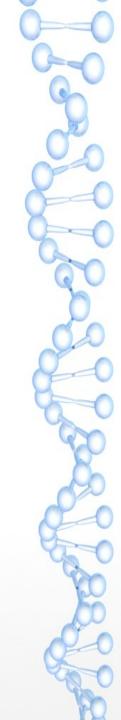
- Linux & Android
- Kernel Power Management
- Linux:
- CPUIdle
- Runtime-PM
- Android
- Wakelocks
- Earlysuspend/Lateresume

Efficiently reducing system performances



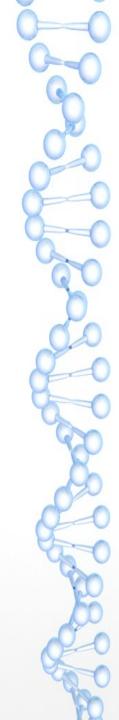
### Standard Linux framework adopted by Android

- Separates scaling logic (governor) from implementation
- Governors implement scaling logic
  - Platform specific code performs actual scaling
- OPP = (Frequency, Voltage)
  - Governor sets frequency, platform code sets OPP
- Extensive debug interface through sysfs



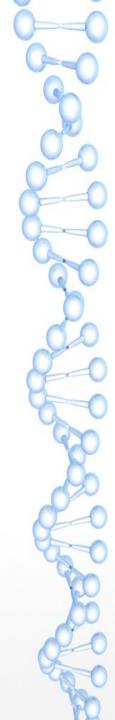
### Governors can be switched at run-time

- Different governor are usually available
- Performances
- Conservative
- On-Demand
- . . . .
- Governors have parameters that modify their behavior
- Thresholds
- Available frequencies
- Timing



#### **Android uses "Interactive"**

- Designed to up-scale faster then on-demand to
- increase system responsiveness to user-interactions
  - Instead of sampling CPU each x msecs, ramps up after idle
  - Designed for latency critical environment, such as UI interactions



#### Hands on:

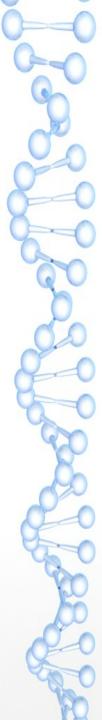
1500000:

```
# cd /sys/devices/system/cpu/cpu0/cpufreq
# cat scaling governor
Interactive
# cat scaling available frequencies
396800 800000 1100000 1300000 1500000
# cat stats/trans table
  From :
              To
              396800
                        800000
                                 1100000
                                           1300000
                                                      1500000
  396800:
                           323
                                       0
  800000:
                 315
                                      60
 1100000:
                                                 35
                            27
 1300000:
                                                           31
```

16

0

# cat stats/time\_in\_states
396800 145499
800000 7876
1100000 789
1300000 382
1500000 2285



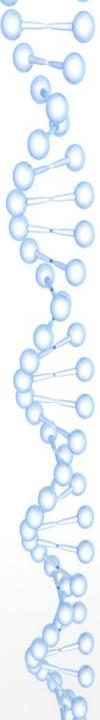
#### Hands on:

```
# cat scaling_cur_freq
396800
```

Finger scroll the Android configuration settings page, make system 'active'....

```
# cat scaling_cur_freq
1500000
```

Yes, it's quite 'interactive'



Nice... Now show me the code (You know, talk is cheap blah blah..)

```
# ls -1 drivers/cpufreq/cpufreq*.c (governors)
drivers/cpufreq/cpufreq.c
drivers/cpufreq/cpufreq conservative.c
drivers/cpufreq/cpufreq interactive.c
drivers/cpufreq/cpufreq-nforce2.c
drivers/cpufreq/cpufreq ondemand.c
drivers/cpufreq/cpufreq performance.c
drivers/cpufreq/cpufreq powersave.c
drivers/cpufreq/cpufreq stats.c
drivers/cpufreq/cpufreq userspace.c
# ls -1 drivers/cpufreq/*omap*.c (omap adaption layer)
Drivers/cpufreq/omap-cpufreq.c
# ls -1 arch/arm/mach-omap2/ (an author selection of)
arch/arm/mach-omap2/dvfs.c
arch/arm/mach-omap2/voltage.c
arch/arm/mach-omap2/abb.c
arch/arm/mach-omap2/frequency.c
arch/arm/mach-omap2/voltagedomains44xx data.c
LOT OF NASTY STUFF HERE!!
```



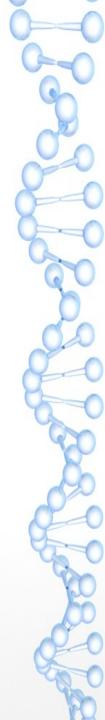
Nice... Now show me the code

Follow the (pseudo)code...



Nice... Now show me the code

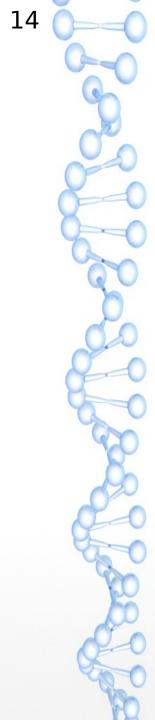
Follow the (pseudo)code...



Nice... Now show me the code

Follow the (pseudo)code...

```
# omap cpufreq.c
static struct cpufreq driver omap driver = {
      .flags
                        = CPUFREQ STICKY,
      .verify
                        = omap verify speed,
      .target
                        = omap target,
                        = omap getspeed,
      .get
      .init
                        = omap cpu init,
      .exit
                        = omap cpu exit,
      .name
                        = "omap",
      .attr
                        = omap cpufreq attr,
};
static int omap cpufreq scale(struct cpufreq policy *policy,
                        unsigned int target freq, unsigned int cur freq,
                        unsigned int relation)
     mutex lock(&omap cpufreq lock);
      current target freq = freq table[i].frequency;
      /* This dives into platform code! */
      ret = omap cpufreq scale(policy, current target freq, policy->cur,
                        relation);
     mutex unlock(&omap cpufreq lock);
```



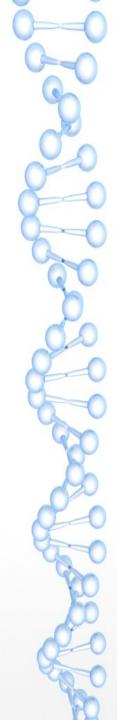
## Suspend

...sleeping never been so hard...

## Diaspora

# Android and standard Linux have diverged long time ago

- Different target devices
- Linux has to be as generic as possible
- Android has well defined use-cases
- Linux-only solutions were not ready at the time Android was designed
- Android developers introduced their very own solutions
- As well as Linux kernel community
- Two different approaches to sleep state management
- Idle-based
- Opportunistic



## Diaspora

Linux

**Frameworks** 

Runtime PM

CPUIdle

Device driver PM

**Strategies** 

Suspend when Idle

Idle devices can reduce their power consumption

User space triggers suspend execution

Android

Wakelocks

Earlysuspend/Lateresume

**CPUIdle** 

Device driver PM

Opportunistic suspend

Always tries to sleep

Kernel-based mechanism



### Two power management models for device drivers

### System Sleep:

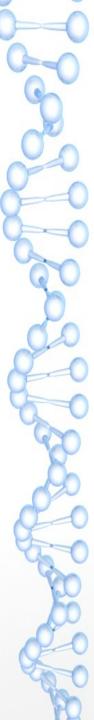
Drivers enter low-power state during system-wide suspend operation.

These operations are usually device specific, and device drivers provide a set of standard callbacks to PM core, invoked during suspend operations

#### Runtime PM (Linux only)

Drivers provide an interface to reduce runtime power consumption. In order to guarantee correctness of operations, the whole device tree is affected by runtime PM.

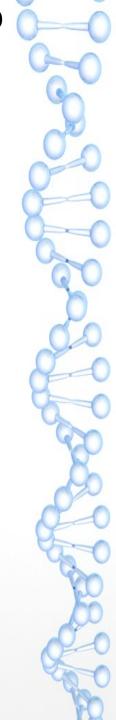
Linux-only mechanism, alternative to Android mechanisms.



#### Power management operations (include/linux/pm.h):

Drivers provides callbacks invoked by PM core during suspend and (eventually) runtime power management

```
struct dev pm ops {
     int (*prepare)(struct device *dev);
     void (*complete)(struct device *dev);
     int (*suspend)(struct device *dev);
      int (*resume)(struct device *dev);
      int (*freeze)(struct device *dev);
      int (*thaw)(struct device *dev);
      int (*poweroff)(struct device *dev);
      int (*restore)(struct device *dev);
      int (*suspend late)(struct device *dev);
      int (*resume early)(struct device *dev);
      int (*freeze late)(struct device *dev);
      int (*thaw early)(struct device *dev);
      int (*poweroff late)(struct device *dev);
      int (*restore early)(struct device *dev);
      int (*suspend noirg)(struct device *dev);
      int (*resume noirg)(struct device *dev);
      int (*freeze noirg)(struct device *dev);
      int (*thaw noirg)(struct device *dev);
      int (*poweroff noirg)(struct device *dev);
      int (*restore noirg)(struct device *dev);
      int (*runtime suspend)(struct device *dev);
      int (*runtime resume)(struct device *dev);
      int (*runtime idle)(struct device *dev);
Run-Time only operations
```

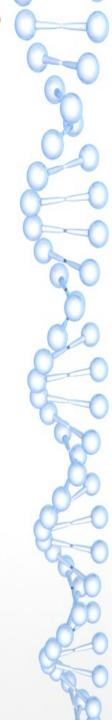


#### **Power management phases**

- During system suspend/resume different phases are traversed by the system
- Each phase consists in the invocation of PM-related callbacks
- PM Callbacks have precedence compared to other frameworks' callbacks

Once phase commences when all the callbacks invoked in the preceding one ends

- The (exclusive) order of callback invocation is:
  - Device Driver
  - Device Type
  - Device Bus
  - Device Class

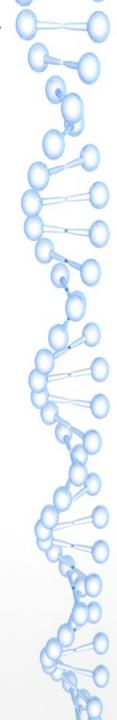


#### **Suspend Phases**

- Power management phases: *Prepare*
- Prevent new devices from being registered
- Device should not enter low power, but eventually prepare
- Power management phases: Suspend
- Quiesce device to prevent new I/O
- Enter Low Power state
- Power management phases: Suspend\_late
- Eventually store device state after I/O has been disabled
- Power management phases: Suspend\_noirq
- Final callback, with interrupt disabled

Not all callbacks are mandatory

If any of the above fails, system suspend is aborted

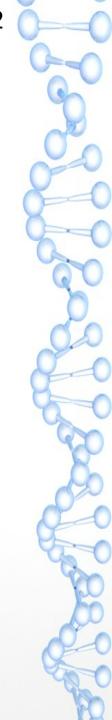


#### **Resume Phases**

- Power management phases: Resume noirg
- Actions to be executed before interrupts for the device is enabled
- Power management phases: Resume\_early
- Prepare for resume
- Dual of suspend\_late
- Power management phases: Resume
- Restore normal activity, enable I/O
- Power management phases: Complete
- Re-enable registration of child devices (for bus and classes)

Hardware can be reset during suspend, driver should handle this.

Errors returned from these callbacks are ignored, system resume is never aborted



### **Power Domain Management**

#### **Power Domain**

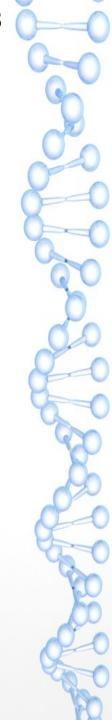
set of hw components sharing a resource, such as clock, power sources, voltage regulators.

#### Domains have policy manager associated:

Eg.

A clock manager within a domain takes into consideration joint constraints of the domain sub-components.

Components in a power domain have to be put in low power state together, as well as woke up at the same time



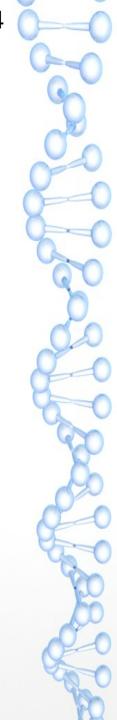
### **Power Domain Management**

#### **Domain Managers**

The struct device holds a pointer to a dev\_pm\_domain object, which provides a set of domain-wide power management callbacks (as it happens for drivers and classes or bus)

The power management domain callbacks, if defined for the given device, always take precedence over the callbacks provided by the device's subsystem (e.g. bus type)

Relevant for SoCs and embedded platforms



#### **CPU Idle states**

Different idle states between active device and full suspend

```
C1 = CPU0/1 ON; MPU ON; CORE ON
```

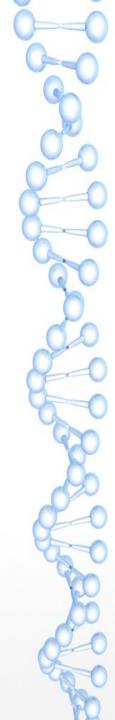
- C2 = CPU0/1 OFF; MPU INA; CORE INA
- C3 = CPU0/1 OFF; MPU CSWR; CORE CSWR
- C4 = CPU0/1 OFF; MPU OSWR; CORE OSWR
- Where:

OSWR: Open Switch Retention

Logic is lost for all the modules in the power domain except for the ones with Built in retention Flip Flops

**CSWR: Closed Switch Retention** 

Logic is preserved for all modules



CPUIdle

Framework to manage Idle states

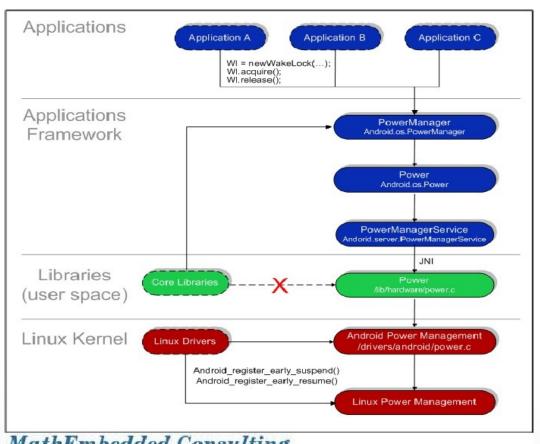
Separates logic from implementation, as CPUFreq

Platform specific driver in arch/arm/mach-omap2/cpuidle-44xx.c

Generic driver code in
#ls drivers/cpuidle/\*.c -1
drivers/cpuidle/coupled.c
drivers/cpuidle/cpuidle.c
drivers/cpuidle/driver.c
drivers/cpuidle/governor.c
drivers/cpuidle/sysfs.c

### **Android**

Introducing locks, which once locked, lock the system from exiting wake state



#### MathEmbedded Consulting

Embedded and Open Source Software Consulting, Development and Training

#### PowerManager class

Provide power management features to applications and services

- Part of the Android framework, mediates between userspace and kernel interface, usually based on *sysfs*
- Application can require wakelocks from power manager:
- Very dangerous: badly written application can drain the battery
- A combination of flags determine which system components have to be kept "awake"

FLAGS	CPU	SCREEN	KEYBOARD
PARTIAL_WAKE_LOCK	On	Off	Off
SCREEN_DIM_WAKE_LOCK	On	Dim	Off
SCREEN_BRIGHT_WAKE_LOCK	On	Bright	Off
SCREEN_BRIGHT_WAKE_LOCK	On	Bright	Bright



#### The power HAL

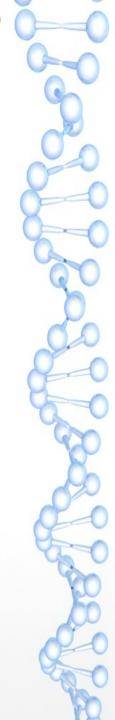
#### HAL for power management

```
$find . -name power.c
./libhardware_legacy/power/power.c
./qcom/power/power.c
./libhardware/modules/power/power.c
```

### Interact with drivers through sysfs

```
$cat ./libhardware_legacy/power/power.c
...
const char * const NEW_PATHS[] = {
    "/sys/power/wake_lock",
    "/sys/power/wake_unlock",
};
...
```

Writes to these files in order to obtain or release a wakelock



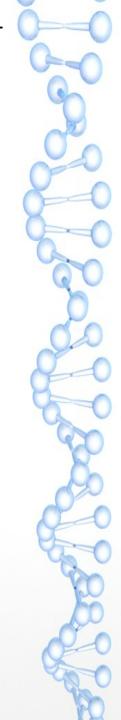
#### Wakelocks for applications

Wakelocks prevent system from going in suspend mode

- Two locking methods:
- Unlimited: wakelock has to be released explicitly
- Timed: wakelock expires after a timeout
- Two wakelock types:
  - WAKE\_LOCK\_SUSPEND: prevents a full system suspend.
  - WAKE\_LOCK\_IDLE: low-power states, which often cause large interrupt latencies or that disable a set of interrupts, will not be entered from idle until the wakelocks are released.

Locking a wakelock during a suspend operation, aborts the suspend sequence if suspend\_late has not been reach

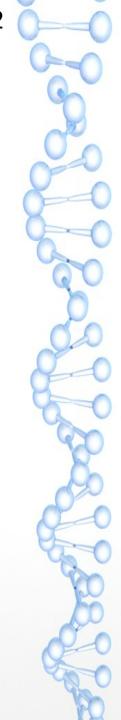
This has unconventional consequences on wakeup sequences



#### Wakelocks on kernel side

- Device drivers can hold wakelocks as well
- In fact, they often have to do that, in order to pass wake up notifications to user space
- Kernel API:

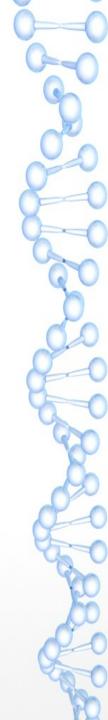
```
struct wakelock wakelock;
wake_lock_init(&wakelock, WAKE_LOCK_SUSPEND, "wakelockname");
wake_lock(&wakelock);
wake_lock_timeout(&wakelock, HZ);
wake_unlock(&wakelock);
wake_lock_destroy(&wakelock);
```



#### Wakelocks on kernel side

kernel/power/wakelock.c

- One 'suspend' task declared (workqueue item), scheduled when
  - Last locked wakelock is released
  - Last locked wakelock expires (Timed wakelock)
  - Another condition when locking that is not clear at all to me:)



#### Wakelocks on kernel side

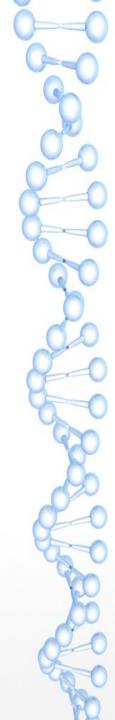
kernel/power/wakelock.c

The suspend work

```
if (has_wake_lock(WAKE_LOCK_SUSPEND)) {
    if (debug_mask & DEBUG_SUSPEND)
        pr_info("suspend: abort suspend\n");
    return;
}

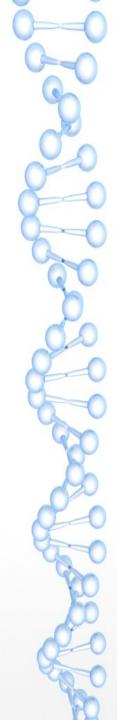
entry_event_num = current_event_num;
sys_sync();
if (debug_mask & DEBUG_SUSPEND)
        pr_info("suspend: enter suspend\n");
getnstimeofday(&ts_entry);
ret = pm_suspend(requested_suspend_state);
```

pm\_suspend calls standard linux enter\_state(PM\_SUSPEND\_MEM)
This will enter the suspend state, invoking the proper callbacks on each
device/subsystem, as explained in previous slides



#### Earlysuspend

- Introduced alongside with Wakelocks
- Basically, a series of callbacks that drivers register to the PM core
- Drivers gets notification of userspace writing to
- /sys/power/request state
- Modify the user-visible power state of the device
- EARLY SUSPEND LEVEL BLANK SCREEN
- EARLY SUSPEND LEVEL STOP DRAWING
- EARLY\_SUSPEND\_LEVEL\_DISABLE\_FB
- EARLY\_SUSPEND\_LEVEL\_STOP\_INPUT
- Applies to 'visible' devices (Display, sensors, framebuffer...)
- Dual to 'late resume' feature

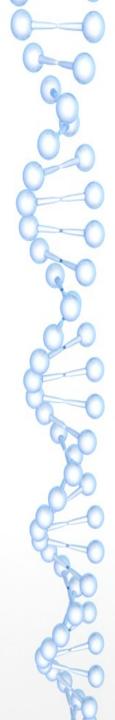


#### Earlysuspend core

kernel/power/earlysuspend.c

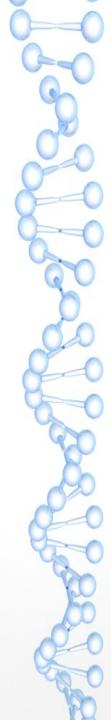
- Main function:
- request\_suspend\_state(suspend\_state\_t new\_state);
- Schedule execution of earlysuspend/lateresume work items

```
if (!old_sleep && new_state != PM_SUSPEND_ON) {
    state |= SUSPEND_REQUESTED;
    queue_work(suspend_work_queue, &early_suspend_work);
} else if (old_sleep && new_state == PM_SUSPEND_ON) {
    state &= ~SUSPEND_REQUESTED;
    wake_lock(&main_wake_lock);
    queue_work(suspend_work_queue, &late_resume_work);
}
```



early\_suspend\_work

Walk the list of callbacks registered by device drivers



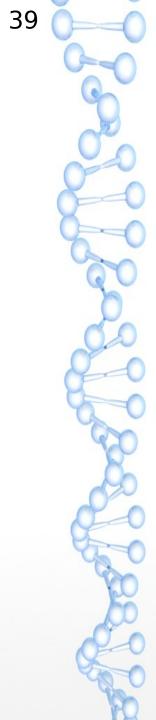
late\_resume\_work

Walk the list of callbacks registered by device drivers

**Driver: callback register** 

#### **Driver API:**

```
void register early suspend(struct early suspend *handler)
void unregister early suspend(struct early suspend *handler)
  Client Drivers: drivers/leds/ledtrig sleep.c
  static void ledtrig sleep early suspend(struct early suspend *h);
  static void ledtrig sleep early resume(struct early suspend *h);
  static struct early suspend ledtrig sleep early suspend handler = {
       .suspend = ledtrig sleep early suspend,
       .resume = ledtrig sleep early resume,
  static int init ledtrig sleep init(void)
       led trigger register simple("sleep", &ledtrig sleep);
       register pm notifier(&ledtrig sleep pm notifier);
       register early suspend(&ledtrig sleep early suspend handler);
       return 0;
```



### References

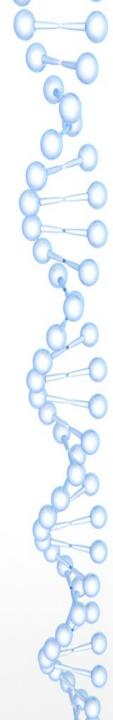
Cool stuff in here

### References

#### **CPUFreq**

http://bamboopuppy.com/changing-androids-cpu-frequency/

- http://bamboopuppy.com/android-cpu-frequency-using-interactive-governor/
- https://code.google.com/p/milestone-overclock/wiki/SmashingTheAndroidKernel
- https://www.kernel.org/doc/Documentation/cpu-freg/cpufreg-stats.txt



### References

#### **Power Management**

```
http://elinux.org/OMAP_Power_Management
http://www.docstoc.com/docs/36734923/OMAP3430-Linux-Power-Management---PowerPoint
http://elinux.org/Android_Power_Management
http://developer.android.com/reference/android/os/PowerManager.html
http://2013.efyexpo.com/wp-content/uploads/2013/03/_2_Android_Power_Management_EF
Y.pdf
http://events.linuxfoundation.org/slides/2011/linuxcon-japan/lcj2011_wysocki2.pdf
http://elinux.org/images/0/08/ELC-2010-Hilman-Runtime-PM.pdf
http://bamboopuppy.com/android-suspend-a-brief-overview/
http://forum.xda-developers.com/wiki/Wakelocks
http://www.cs.rochester.edu/~sandhya/csc256/seminars/peter_android.pdf
http://www.scribd.com/doc/22685058/android-power-management
http://www.kandroid.org/online-pdk/guide/power_management.html
http://www.kandroid.org/online-pdk/guide/early_suspend.html
https://patchwork.kernel.org/patch/58064/ (Motorola Quickwakeup)
```

And of course, the Linux kernel source code, and AOSP sources

https://community.freescale.com/thread/261901