1 Syntax

```
program ::= \overline{cls} \ \overline{s}
                 ::= class \ C \ \{\overline{field} \ \overline{method}\}
cls
field
               ::=T f;
method ::= T \ m(T \ x) \ contract \ \{\overline{s}\}
contract ::= requires \phi; ensures \phi;
T
                 ::= int \mid C
                 ::=x.f:=y; \mid x:=e; \mid x:=\text{new }C; \mid x:=y.m(z); \mid \text{return }x; \mid \text{assert }\phi; \mid \text{release }\phi;
s
                 ::=\mathtt{true}\ |\ e=e\ |\ e\neq e\ |\ \mathtt{acc}(x.f)\ |\ \phi*\phi
\phi
                 := v \mid x \mid e.f
e
                 ::= this | result | \langle other \rangle
Γ
                 ::=(x\mapsto T)
                ::= (o \mapsto (C, \overline{(f \mapsto v)}))
H
                ::=(x\mapsto v)
\rho
            ::= \overline{(x,f)}::= \overline{(o,f)}
A_s
                ::= (\rho, A_d, \overline{s}) \cdot S \mid nil
```

2 Static semantics

2.1 Expressions $(A_s \vdash_{\mathtt{sfrm}} e)$

$$\frac{}{A \vdash_{\mathtt{sfrm}} x} \mathsf{WFVAR}$$

$$\frac{}{A \vdash_{\mathtt{sfrm}} v} \text{ WFValue}$$

$$\frac{(x,f) \in A}{A \vdash_{\mathtt{sfrm}} x.f} \text{ WFField}$$

2.2 Formulas $(A_s \vdash_{\mathtt{sfrm}} \phi)$

$$\frac{}{A \vdash_{\mathtt{sfrm}} \mathtt{true}} \mathtt{WFTRUE}$$

$$\frac{A \vdash_{\mathtt{sfrm}} e_1}{A \vdash_{\mathtt{sfrm}} e_1 = e_2} \text{ WFEQUAL}$$

$$\frac{A \vdash_{\mathtt{sfrm}} e_1 \qquad A \vdash_{\mathtt{sfrm}} e_2}{A \vdash_{\mathtt{sfrm}} e_1 \neq e_2} \text{ WFNEQUAL}$$

$$\frac{}{A \vdash_{\mathtt{sfrm}} \mathtt{acc}(x, f)} \mathsf{WFAcc}$$

$$\frac{A_s \vdash_{\mathtt{sfrm}} \phi_1 \qquad A_s \cup \mathtt{static-footprint}(\phi_1) \vdash_{\mathtt{sfrm}} \phi_2}{A_s \vdash_{\mathtt{sfrm}} \phi_1 * \phi_2} \quad \text{WF-SepOp}$$

2.3 Footprint (static-footprint(ϕ) = A_s)

$$\begin{array}{lll} \mathtt{static\text{-}footprint}(\mathtt{true}) &= \emptyset \\ \mathtt{static\text{-}footprint}(e_1 = e_2) &= \emptyset \\ \mathtt{static\text{-}footprint}(e_1 \neq e_2) &= \emptyset \\ \mathtt{static\text{-}footprint}(\mathtt{acc}(x.f)) &= \{(x,f)\} \\ \mathtt{static\text{-}footprint}(\phi_1 * \phi_2) &= \mathtt{static\text{-}footprint}(\phi_1) \cup \mathtt{static\text{-}footprint}(\phi_2) \end{array}$$

2.4 Hoare $(\Gamma \vdash \{\phi\}\overline{s}\{\phi\})$

$$\frac{\Gamma \vdash \{\phi_p\}s_1\{\phi_{q1}\} \quad \phi_{q1} \implies \phi_{q2} \quad \Gamma \vdash \{\phi_{q2}\}s_2\{\phi_r\}}{\Gamma \vdash \{\phi_p\}s_1; s_2\{\phi_r\}} \quad \text{H-Sec}$$

$$\frac{\Gamma(x) = C \qquad \text{fields}(C) = \{\overline{f_i}\}}{\Gamma \vdash \{\phi\}x := \text{new } C\{\overline{\texttt{acc}(x.f_i)} * x \neq \texttt{null} * \phi\}} \quad \text{H-NewObj}$$

$$\frac{\phi \implies \mathtt{acc}(x.f) * x \neq \mathtt{null}}{\Gamma \vdash \{\phi\} x.f := y \{\phi * x.f = y\}} \quad \text{H-FieldAssign}$$

$$\frac{\phi' = \phi[e/x] \qquad \emptyset \vdash_{\mathtt{sfrm}} \phi' \qquad \mathtt{static-footprint}(\phi') \vdash_{\mathtt{sfrm}} e}{\Gamma \vdash \{\phi'\}x := e\{\phi\}} \qquad \text{H-VarAssign}$$

$$\frac{}{\Gamma \vdash \{\phi\} \text{return } x \{\phi * \text{result} = x\}} \quad \text{H-Return}$$

$$\Gamma(y) = C \qquad \phi \implies y \neq null * \phi_p * \phi_r \qquad \phi_p = \texttt{mpre}(C, m)[y, \overline{z}/\texttt{this}, \overline{X}] \qquad \phi_q = \texttt{mpost}(C, m)[y, \overline{z}, x/\texttt{this}, \overline{X}] \qquad \phi$$

$$\frac{\phi \implies \phi'}{\Gamma \vdash \{\phi\} \text{assert } \phi'\{\phi\}} \quad \text{H-Assert}$$

$$\frac{\phi \implies \phi' * \phi_r \qquad \emptyset \vdash_{\mathtt{sfrm}} \phi_r}{\Gamma \vdash \{\phi\}\mathtt{release} \ \phi' \{\phi_r\}} \quad \text{H-Release}$$

3 Dynamic semantics

3.1 Expressions $(H, \rho \vdash e \Downarrow v)$

$$H, \rho \vdash x \Downarrow \rho(x)$$
 EE-Var

$$\overline{H, \rho \vdash v \Downarrow v}$$
 EE-Value

$$\frac{H, \rho \vdash x \Downarrow o}{H, \rho \vdash x.f \Downarrow H(o)(f)} \quad \text{EE-Acc}$$

3.2 Formulas $(H, \rho, A \vDash \phi)$

$$H, \rho, A \vDash \mathsf{true}$$
 EA-True

$$\frac{H, \rho \vdash e_1 \Downarrow v_1 \qquad H, \rho \vdash e_2 \Downarrow v_2 \qquad v_1 = v_2}{H, \rho, A \vDash e_1 = e_2} \quad \text{EA-Equal}$$

$$\frac{H, \rho \vdash e_1 \Downarrow v_1 \qquad H, \rho \vdash e_2 \Downarrow v_2 \qquad v_1 = v_2}{H, \rho, A \vDash e_1 = e_2} \quad \text{EA-NEqual}$$

$$\frac{H, \rho \vdash x \Downarrow o \quad (o, f) \in A}{H, \rho, A \vDash \mathsf{acc}(x.f)} \quad \text{EA-Acc}$$

$$\frac{A_1 = A \setminus A_2 \qquad H, \rho, A_1 \vDash \phi_1 \qquad H, \rho, A_2 \vDash \phi_2}{H, \rho, A \vDash \phi_1 * \phi_2} \quad \text{EA-SepOp}$$

3.3 Footprint (footprint_{H,ρ}(ϕ) = A_d)

$$\begin{array}{ll} \operatorname{footprint}_{H,\rho}(\operatorname{true}) &= \emptyset \\ \operatorname{footprint}_{H,\rho}(e_1 = e_2) &= \emptyset \\ \operatorname{footprint}_{H,\rho}(e_1 \neq e_2) &= \emptyset \\ \operatorname{footprint}_{H,\rho}(\operatorname{acc}(e.f)) &= \{(o,f)\} \text{ where } H, \rho \vdash e \Downarrow o \\ \operatorname{footprint}_{H,\rho}(\phi_1 * \phi_2) &= \operatorname{footprint}_{H,\rho}(\phi_1) \cup \operatorname{footprint}_{H,\rho}(\phi_2) \end{array}$$

3.4 Small-step $((H,S) \rightarrow (H,S))$

$$\frac{H, \rho \vdash x \Downarrow o \quad H, \rho \vdash y \Downarrow v_y \quad (o, f) \in A \quad H' = H[o \mapsto [f \mapsto v_y]]}{(H, (\rho, A, x.f := y; \overline{s}) \cdot S) \to (H', (\rho, A, \overline{s}) \cdot S)} \text{ ESFIELDASSIGN}$$

$$\frac{H, \rho \vdash e \Downarrow v \qquad \rho' = \rho[x \mapsto v]}{(H, (\rho, A, x := e; \overline{s}) \cdot S) \to (H, (\rho', A, \overline{s}) \cdot S)} \text{ ESVARASSIGN}$$

$$\frac{H(o) = \bot \quad \mathtt{fields}(C) = f}{\rho' = \rho[x \mapsto o] \quad A' = A * \overline{(o, f_i)} \quad H' = H[o \mapsto [(\overline{(f_i, \mathtt{null})})]]}{(H, (\rho, A, x := \ \mathtt{new} \ C; \overline{s}) \cdot S) \rightarrow (H', (\rho', A', \overline{s}) \cdot S)} \; \mathtt{ESNEWOBJ}$$

$$\frac{H, \rho \vdash x \Downarrow v_x \qquad \rho' = \rho[\mathtt{result} \mapsto v_x]}{(H, (\rho, a, \mathtt{return} \ x; \overline{s}) \cdot S) \to (H, (\rho', a, \overline{s}) \cdot S)} \text{ ESRETURN}$$

$$\begin{split} H, \rho \vdash y \Downarrow o & H, \rho \vdash z \Downarrow v & H(o) = (C, fvf) \\ \texttt{mbody}(C, m) = \overline{r} & \texttt{mparam}(C, m) = (T, w) & \texttt{mpre}(C, m) = \phi \\ \rho' = [\texttt{this} \mapsto o, w \mapsto v] & H, \rho', A \vDash \phi & A' = \texttt{footprint}_{H, \rho'}(\phi) \\ \hline (H, (\rho, A, x := y.m(z); \overline{s}) \cdot S) \to (H, (\rho', A', \overline{r}) * (\rho, A \setminus A', x := y.m(z); \overline{s}) \cdot S) \end{split}$$
 ESAPP

$$\frac{\operatorname{mpost}(C,m) = \phi}{H,\rho',A' \vDash \phi \qquad A'' = \operatorname{footprint}_{H,\rho'}(\phi) \qquad H,\rho' \vdash \operatorname{result} \Downarrow v_r}{(H,(\rho',A',\emptyset)*(\rho,A,x := y.m(z);\overline{s}) \cdot S) \to (H,(\rho[x \mapsto v_r],A*A'',\overline{s}) \cdot S)} \text{ ESAPPFINISH}$$

$$\frac{H,\rho,A \vDash \phi}{(H,(\rho,A,\mathtt{assert}\ \phi;\overline{s})\cdot S) \to (H,(\rho,A,\overline{s})\cdot S)} \ \mathtt{ESASSERT}$$

$$\frac{H,\rho,A \vDash \phi \qquad A' = A \setminus \mathtt{footprint}_{H,\rho}(\phi)}{(H,(\rho,A,\mathtt{release}\;\phi;\overline{s})\cdot S) \to (H,(\rho,A',\overline{s})\cdot S)} \; \mathsf{ESRelease}$$

4 Theorems

Hoare preserves self-framing

$$\begin{split} \forall \; \Gamma, \phi_1, \phi_2, s : \Gamma \vdash \{\phi_1\} s \{\phi_2\} \\ &\Longrightarrow \; \mathsf{static\text{-}footprint}(\phi_1) \vdash_{\mathsf{sfrm}} \phi_1 \\ &\Longrightarrow \; \mathsf{static\text{-}footprint}(\phi_2) \vdash_{\mathsf{sfrm}} \phi_2 \end{split}$$

Hoare progress

$$\forall \Gamma, \phi_1, \phi_2, s, H_1, \rho_1, A_1 : \Gamma \vdash \{\phi_1\} s \{\phi_2\}$$

$$\Longrightarrow H_1, \rho_1, A_1 \vDash \phi_1$$

$$\Longrightarrow \exists H_2, \rho_2, A_2 : (H_1, (\rho_1, A_1, s'; \overline{s}) \cdot S) \to^* (H_2, (\rho_2, A_2, \overline{s}) \cdot S)$$

Hoare preservation

$$\begin{split} \forall \; \Gamma, \phi_1, \phi_2, s, H_1, H_2, \rho_1, \rho_2, A_1, A_2 : \Gamma \vdash \{\phi_1\} s \{\phi_2\} \\ &\Longrightarrow \; H_1, \rho_1, A_1 \vDash \phi_1 \\ &\Longrightarrow \; (H_1, (\rho_1, A_1, s'; \overline{s}) \cdot S) \to^* (H_2, (\rho_2, A_2, \overline{s}) \cdot S) \\ &\Longrightarrow \; H_2, \rho_2, A_2 \vDash \phi_2 \end{split}$$