Verifier Core Language BNF Grammar

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\begin{array}{rcll}
x, y, z & \in & VAR & \textit{(variables)} \\
v & \in & VAL & \textit{(values)} \\
e & \in & EXPR & \textit{(expressions)} \\
s & \in & STMT & \textit{(statements)} \\
o & \in & LOC & \textit{(object Ids)} \\
f & \in & FIELDNAME & \textit{(field names)} \\
m & \in & METHODNAME & \textit{(method names)} \\
C & \in & CLASSNAME & \textit{(class names)} \\
P & ::= & \overline{cls}\ s & \\
cls & ::= & class\ C\ extends\ C\ \{\overline{field}\ \overline{method}\} & \\
field & ::= & T\ f; & \\
T & ::= & int \mid C \mid \top & \\
method & ::= & T\ m(\overline{T\ x})\ contract\ \{s\} & \\
contract & ::= & requires\ \widetilde{\phi}\ ensures\ \widetilde{\phi} & \\
\oplus & ::= & + \mid - \mid * \mid \backslash & \\
\odot & ::= & \neq \mid = \mid < \mid > & \\
s & ::= & skip \mid s_1\ ;\ s_2 \mid T\ x \mid x := e \mid if\ (x \odot y)\ \{s\}\ else\ \{s\} \mid x.f := y \mid x := new\ C & \\
 & & \mid y := z.m(\overline{x}) \mid assert\ \phi \mid release\ \phi \mid hold\ \phi\ \{s\} & \\
e & ::= & v \mid x \mid e \oplus e \mid e.f & \\
x & ::= & result \mid id \mid old(id) \mid this & \\
v & ::= & n \mid o \mid null & \\
\phi & ::= & \text{true} \mid e \odot e \mid acc(e.f) \mid \phi * \phi & \\
\widetilde{\phi} & ::= & \phi \mid ? * \phi &
\end{array}
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