Verifier Core Language BNF Grammar

Jenna Wise

October 20, 2018

```
\in VAR
                                                                                                                            (variables)
 x, y, z
             \in VAL
                                                                                                                                (values)
              \in EXPR
                                                                                                                        (expressions)
              \in STMT
                                                                                                                         (statements)
             \in LOC
                                                                                                                           (object Ids)
    f
             \in FIELDNAME
                                                                                                                        (field names)
             \in METHODNAME
                                                                                                                    (method names)
    C
             \in CLASSNAME
                                                                                                                        (class names)
            ::= \overline{cls} \ s
    P
   cls
            ::= class\ C\ extends\ C\ \{\overline{field}\ \overline{A}\ \overline{method}\}
            := T f;
  field
            ::= predicate \ \alpha_C(\overline{x}) := \widetilde{\phi}
    \boldsymbol{A}
    T
            ::= int \mid C \mid \top
method ::= T m(\overline{T x}) dynamic contract static contract \{s\}
contract ::= requires \widetilde{\phi} \ ensures \widetilde{\phi}
            ::= + | - | * | \setminus
    \oplus
            ::= \neq | = | < | >
            ::= skip \mid s_1 \; ; \; s_2 \mid T \; x \mid x := e \mid if \; (x \odot y) \; \{s_1\} \; else \; \{s_2\} \mid x.f := y \mid x := new \; C
               |y := z.m(\overline{x}) | y := z.m_C(\overline{x}) | assert \phi | release \phi | hold \phi \{s\}
            ::= v \mid x \mid e \oplus e \mid e.f
            ::= result \mid id \mid old(id) \mid this
            ::= n \mid o \mid null
            := \text{true} \mid e \odot e \mid \alpha(\overline{e}) \mid acc(e.f) \mid \phi * \phi
            := \phi \mid ? * \phi
```