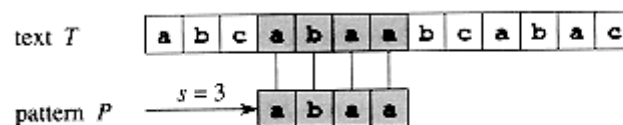


Pattern Matching

In this chapter, we will primarily talk about two string searching algorithms, the **Knuth Morris Pratt** algorithm and **Z - Algorithm**.

Introduction

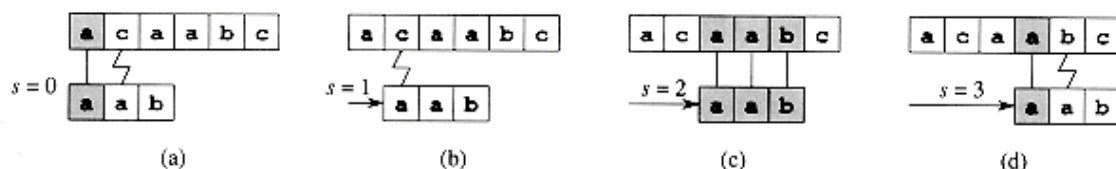
Suppose you are given a string text of length n and a string pattern of length m . You need to find all the occurrences of the pattern in the text or report that no such instance exists.



In the above example, the pattern “abaa” appears at position 3 (0 indexed) in the text “abcabaabcbac”.

Naive Algorithm

A naive way to implement this pattern searching is to move from each position in the text and start matching the pattern from that position until we encounter a mismatch between the characters or say that the current position is a valid one.



In the given picture, The length of the text is 5 and the length of the pattern is 3. For each position from 0 to 3, we choose it as the starting position and then try to match the next 3 positions with the pattern.

Naive pattern matching

- For each i from 0 to $N - M$
- For each j from 0 to $M - 1$, try to match the j th character of the pattern with $(i + j)$ th character of the string text.
- If a mismatch occurs, skip this instance and continue to the next iteration.
- Else output this position as a matching position.

```
function NaivePatternSearch(text, pattern)
    // iterate for each candidate position
    for i from 0 to text.length - pattern.length

        // boolean variable to check if any mismatch occurs
        match = True

        for j from 0 to pattern.length - 1
            // if mismatch make match = False
            if text[i + j] not equals pattern[j]
                match = False
                break

        // if no mismatch print this position
        if match == True
            print the occurrence i

    return
```

Time Complexity: $O(N \cdot M)$, where N is the length of the text and M is the length of the pattern we need to search.