

Biomeg 261 Lecture 08 2017

Flux balance analysis Cont'd.

(& constraint based analysis)

- Null spaces
  - Spans
  - Geometry & constraints
  - Optimality conditions & optimisation problems
- } Linear algebra
- } Linear programming

Null spaces (Linear algebra).

The nullspace for a matrix

A is the set of all solutions to  $\underline{A} \underline{x} = \underline{0}$

Zero is always in the null space:  $\underline{A} \underline{0} = \underline{0}$

A non-trivial nullspace is when we have non-zero solutions

$$\text{eg } \begin{bmatrix} 1 & -1 \\ -1 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$

$$\begin{aligned} \text{ie} \\ (1) \quad x_1 - x_2 = 0 &\Rightarrow x_1 = x_2 \\ (2) \quad -x_1 + x_2 = 0 &\Rightarrow x_1 = x_2 \end{aligned} \left. \vphantom{\begin{aligned} (1) \\ (2) \end{aligned}} \right\} \text{same equations}$$

2 unknowns  
• 1 indep. constraint  
→ 2 - 1 = 1 free var.

free choice of  $x_2$  (say) } or vice versa  
then determines  $x_1$

$$\underline{N(A)} = \left\{ \begin{pmatrix} x_1 \\ x_2 \end{pmatrix} \mid x_1 = x_2 \right\}$$

Mathematically,  $N(A) = \{ \vec{x} \mid A\vec{x} = \vec{0} \}$

Null space of A

elements (vectors)

set of

such that

condition

Example

$$\underline{S} = \begin{bmatrix} -1 & 1 & 0 & 0 & 0 & 0 \\ 2 & -1 & -2 & 0 & 0 & 0 \\ 0 & 0 & 3 & -3 & 3 & 0 \\ 0 & 0 & 0 & 1 & -1 & 1 \end{bmatrix}$$

Goal:  $N(\underline{S}) = \{ \vec{J} \mid \underline{S}\vec{J} = \vec{0} \}$

we need to solve  $\underline{S}\vec{J} = \vec{0}$

→ usually choose 'free' variables to parameterise the null space with

Example Cont'd

$N=6$

$M=4$

$$\begin{bmatrix} -1 & 1 & 0 & 0 & 0 & 0 \\ 2 & -1 & -2 & 0 & 0 & 0 \\ 0 & 0 & 3 & -3 & 3 & 0 \\ 0 & 0 & 0 & 1 & -1 & 1 \end{bmatrix}$$

row operations

$r_1 \rightarrow r_1$   
 $r_2 \rightarrow r_2 + 2r_1$   
 $r_3 \rightarrow r_3$   
 $r_4 \rightarrow r_4$

to make upper triangular

$$\Rightarrow \begin{bmatrix} -1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & -2 & 0 & 0 & 0 \\ 0 & 0 & 3 & -3 & 3 & 0 \\ 0 & 0 & 0 & 1 & -1 & 1 \end{bmatrix} \vec{J} = \vec{0}$$

$$\begin{cases} -J_1 + J_2 = 0 \\ J_2 - 2J_3 = 0 \\ 3J_3 - 3J_4 + 3J_5 = 0 \\ J_4 - J_5 - J_6 = 0 \end{cases}$$

4 equations  
6 vars.

choose 2: eg  $J_1, J_5$  & write others in terms of these

eg  $J_1=2, J_5=0$

$$\vec{J}^{(1)} = \begin{pmatrix} 2 \\ 2 \\ 1 \\ 0 \\ 1 \\ 0 \end{pmatrix}$$

or  $J_1=0, J_5=1$

$$\vec{J}^{(2)} = \begin{pmatrix} 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 0 \end{pmatrix}$$

&  $N(\underline{S}) = \text{span} \{ \vec{J}^{(1)}, \vec{J}^{(2)} \}$

Span?: all linear combinations of a set of vectors

Here:

$$N(\underline{s}) = \text{span} \left\{ \underbrace{\text{independent vectors}}_{\text{solving } \underline{s} \bar{\underline{J}} = \bar{\underline{0}}} \right\}$$

null space

same number as number free vars

eg  $N(\underline{s}) = \text{span} \left\{ \bar{\underline{J}}^{(1)}, \bar{\underline{J}}^{(2)} \right\}$

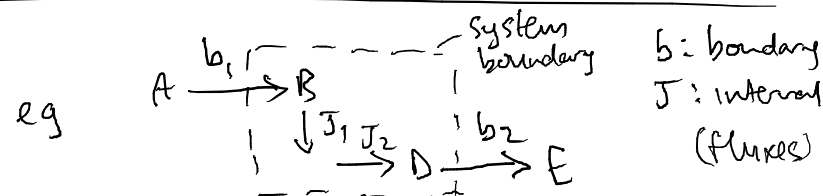
where  $\bar{\underline{J}}^{(1)} = \begin{pmatrix} 2 \\ 2 \\ 1 \\ 0 \\ 1 \end{pmatrix}, \bar{\underline{J}}^{(2)} = \begin{pmatrix} 0 \\ 0 \\ 0 \\ 1 \\ 0 \end{pmatrix}$

Steps: reduce to minimal set of eqns  
choose free vars (eg  $m-n$  of them)  
find set of indep. vectors  
write  $N(\underline{s}) = \text{span} \{ \downarrow \}$

## Boundary vs Internal Fluxes

Sometimes we want to draw

'system boundaries' & call some fluxes 'boundary' fluxes & some 'internal' fluxes.



We can either include these as usual

$$\begin{matrix} & J_1 & J_2 & b_1 & b_2 \\ \begin{matrix} A \\ B \\ C \\ D \\ E \end{matrix} & \begin{bmatrix} 0 & 0 & -1 & 0 \\ -1 & 0 & 1 & 0 \\ 1 & -1 & 0 & -1 \\ 0 & 0 & 0 & 1 \end{bmatrix} \end{matrix}$$

OR

focus on internal

$$\begin{matrix} & J_1 & J_2 \\ \begin{matrix} B \\ C \\ D \end{matrix} & \begin{bmatrix} -1 & 0 \\ 1 & -1 \\ 0 & 1 \end{bmatrix} \end{matrix}$$

(include b in  $\bar{\underline{J}}$ )

In this later case we can write

$$|\underline{s} \bar{\underline{J}} \geq 0|, \text{ if we choose signs carefully,}$$

$$(\underline{s} \bar{\underline{J}} = \bar{\underline{b}}_{\text{net}} \geq 0)$$

# Geometry: Nullspace vs Feasible region

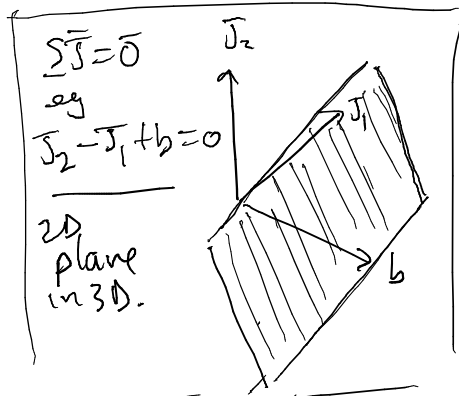
The two forms

$$\underline{\sum \bar{J}} = \bar{0} \quad \text{or} \quad \underline{\sum \bar{J}} \geq 0$$

all fluxes

subset of fluxes eg internal only.

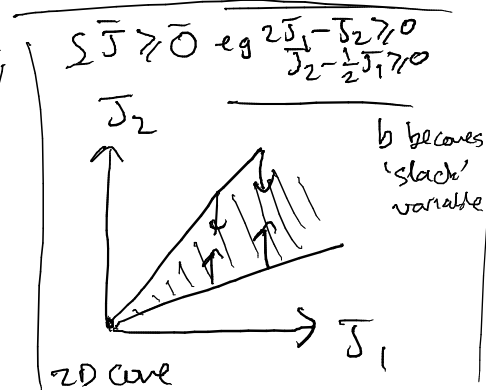
are equivalent but have slightly different geometric pictures



3 vars

1 indep equality constr.

3-1 = 2 dimensional  
(hyper-) plane  
Nullspace



2 vars

2 indep. inequality constraints

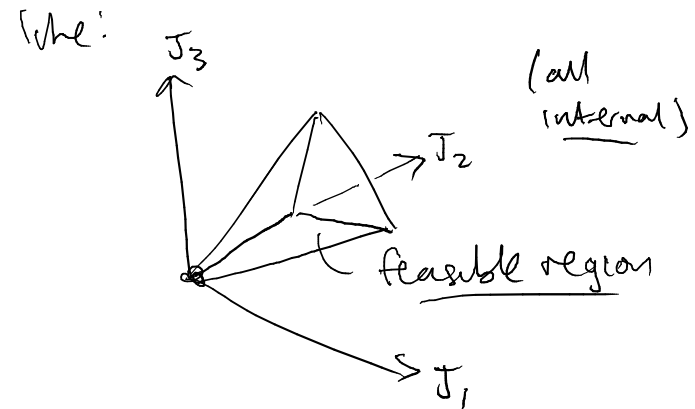
2D cone/polyhedra  
Feasible region

## What? Why?

In general we will put in 'standard form'

$$\underline{\sum \bar{J}} = \bar{0} \quad \& \text{ use } \underline{\text{nullspace}} \quad (\text{include boundary fluxes})$$

But you may see pictures



which come from inequality constraints version.

→ Just be aware

(General principle/trade-off:  
simple in higher dim or complex in lower d)

'Special' solutions: uniqueness?

- Regardless of specific form used, we have linear constraints on fluxes /  $\bar{J}$

(Note: would typically be non-linear in concentrations if using mass action)

The constraints can be thought of as

- Defining a non-trivial null space ( $\underline{S} \bar{J} = \underline{0}$  form)  
or

- A polyhedral feasible region ( $\underline{S} \bar{J} \geq \underline{0}$  form for internal)

$\Rightarrow$  either way we might want to look at ways to 'pick out' particular solutions from these sets of solutions.

## Constraints & Optimality conditions

We can add many types of extra constraints or conditions to narrow down possible solns:

- bounds/signs of fluxes
- thermodynamic feasibility (directional constraints)
- optimality or extreme cases (eg max. energy prod.)

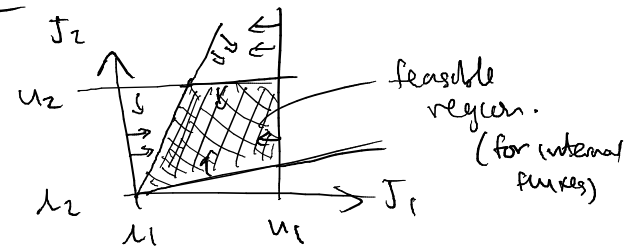
The most natural to include first in general are lower & upper flux bounds, ie

Capacity constraints:

$$| \underline{l}_i \leq \bar{J}_i \leq \underline{u}_i |$$

(note: without these only know ratios)

Gives



For irreversible, use  $| \underline{J}_i \geq 0 |$  (why?)

## 'Optimal' solutions

A useful way to pick out 'special' solutions is to use optimality (max/min) conditions

— these tell us about limits on what is possible

eg 'maximal rate of ATP production is...'

Note :

— a real system may or may not reach these limits!

↳ many competing goals so may not be optimal for any one

↳ still a good way to predict/understand

↳ many sensible constraints can be rewritten as max/min conditions.

## Example : cont'd

Recall our example where we wrote our 2D null space in terms of  $J_1$  &  $J_5$

$\begin{aligned} J_2 &= J_1 \\ J_3 &= J_1/2 \\ J_4 &= J_1/2 + J_5 \\ J_6 &= J_1/2 \end{aligned}$	<p>ie <math>J_2, J_3, J_4, J_6 = f(J_1, J_5)</math> 6 vars, 2 <u>free</u> 4 det.</p>
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can write as  $\underline{J} = \underline{0}$

To find a particular sol<sup>n</sup>.

• suppose  $2 \leq J_1, J_5 \leq 10$

•  $J_i \geq 0$  for all  $i$  bounds apply to both

•  $J_4$  is ATP production &  $J_6$  is lactate prod. } for example

Goals Case 1. Max ATP prod.

Case 2. Max ATP prod. while minimising lactate prod.

## Case 1.

$$\left. \begin{array}{l} \max J_4 \\ \text{subject to} \\ \bullet \sum \bar{J} = 0 \\ \bullet 2 \leq J_1 \leq 10 \\ \bullet 2 \leq J_5 \leq 10 \end{array} \right\} \Leftrightarrow \left\{ \begin{array}{l} J_2 = J_1 \\ J_3 = J_1/2 \\ J_4 = J_1/2 + J_5 \\ J_6 = J_1/2 \end{array} \right.$$

we re-wrote as

Sol<sup>n</sup>: want  $J_4 = \frac{J_1}{2} + J_5$  max.  
 $\Rightarrow$  set  $J_1 = J_5 = 10$  (can vary indep.)  
 $\Rightarrow J_4 = 15$

## Case 2

$$\left. \begin{array}{l} \max aJ_4 - bJ_6 \\ a, b > 0 \text{ weights} \\ \text{('values' or 'costs')} \\ \text{s.t. same constraints} \end{array} \right\} \begin{array}{l} \max -J \\ \equiv \\ \min J \end{array}$$

$$\max aJ_4 - bJ_6 = \frac{(a-b)J_1 + aJ_5}{2} \left\{ \begin{array}{l} \text{if } a > b \\ \text{same: } J_1 = J_5 = 10 \end{array} \right.$$

$$\left\{ \begin{array}{l} \text{if } a < b \\ J_1 = 2, J_5 = 10 \end{array} \right.$$

not enough  
info  
for unique

$$\left\{ \begin{array}{l} \text{if } a = b \\ J_1 = ?, J_5 = 10 \end{array} \right.$$

## General Optimisation Framework

$$\min z = \bar{c}^T \bar{J} \quad \left. \vphantom{\min} \right\} \text{objective function}$$

$$\text{s.t. } \left. \begin{array}{l} \sum \bar{J} = 0 \\ \lambda_i \leq \bar{J}_i \leq u_i \end{array} \right\} \text{constraints}$$

$z$ : scalar (number)

$\bar{c}$ : vector of weights  
(here, 'costs')

eg

$$\begin{aligned} \bar{c}^T \bar{J} &= (1 \ 2 \ -1) \begin{pmatrix} J_1 \\ J_2 \\ J_3 \end{pmatrix} \\ &= J_1 + 2J_2 - J_3 = z \end{aligned}$$

Linear objective function +  
Linear constraints } "Linear Programming"

What should I be able to do?

- Given network find  $\underline{S}$
- Find nullspace for simple  $\underline{S}$ 
  - ↳ reduce to indep. eq's
  - ↳ choose free fluxes
  - ↳ find indep. vectors
  - ↳ write as  $\text{span}\{ \quad \}$
- Describe typical/useful additional constraints
- Write down optimisation problem for given description
- Solve very simple optimisation problems