

Homework 2. Naive parsing of context free grammars

Theoretical background

A *derivation* is a rule list that describes how to derive a phrase from a nonterminal symbol. For example, suppose we have the following grammar with start symbol Expr:

Expr \rightarrow Term Binop Expr

Expr \rightarrow Term

Term \rightarrow Num

Term \rightarrow Lvalue

Term \rightarrow Incrop Lvalue

Term \rightarrow Lvalue Incrop

Term \rightarrow "(" Expr ")"

Lvalue \rightarrow \$ Expr

Incrop \rightarrow "++"

Incrop \rightarrow "--"

Binop \rightarrow "+"

Binop \rightarrow "-"

Num \rightarrow "0"

Num \rightarrow "1"

Num \rightarrow "2"

Num \rightarrow "3"

Num \rightarrow "4"

Num \rightarrow "5"

Num \rightarrow "6"

Num \rightarrow "7"

Num \rightarrow "8"

Num \rightarrow "9"

Then here is a derivation for the phrase "3" "+" "4" from the nonterminal Expr. After each rule is applied, the resulting list of terminals and nonterminals is given.

rule	after rule is applied
(at start)	Expr
Expr \rightarrow Term Binop Expr	Term Binop Expr
Term \rightarrow Num	Num Binop Expr
Num \rightarrow "3"	"3" Binop Expr
Binop \rightarrow "+"	"3" "+" Expr
Expr \rightarrow Term	"3" "+" Term
Term \rightarrow Num	"3" "+" Num

Num \rightarrow "4"

"3" "+" "4"

In a *leftmost derivation*, the leftmost nonterminal is always the one that is expanded next. The above example is a leftmost derivation.

Motivation

You'd like to test grammars that are being proposed as test cases for CS 132 projects. One way is to test it on actual CS 132 projects, but those projects aren't done yet and anyway you'd like a second opinion in case the student projects are incorrect. So you decide to write a simple parser generator. Given a grammar in the style of Homework 1, your program will generate a function that is a parser. When this parser is given a program to parse, it produces a derivation for that program, or an error indication if the program contains a syntax error and cannot be parsed.

The key notion of this assignment is that of a *matcher*. A *matcher* is a function that inspects a given string of terminals to find a match for a prefix that corresponds to a nonterminal symbol of a grammar, and then checks whether the match is acceptable by testing whether a given acceptor succeeds on the corresponding derivation and suffix. For example, a matcher for `awkish_grammar` below might inspect the string `["3"; "+"; "4"; "-"]` and find two possible prefixes that match, namely `["3"; "+"; "4"]` and `["3"]`. The matcher will first apply the acceptor to a derivation for the first prefix `["3"; "+"; "4"]`, along with the corresponding suffix `["-"]`. If this is accepted, the matcher will return whatever the acceptor returns. Otherwise, the matcher will apply the acceptor to a derivation for the second prefix `["3"]`, along with the corresponding suffix `["+"; "4"; "-"]`, and will return whatever the acceptor returns. If a matcher finds no matching prefixes, it returns the special value `None`.

As you can see by mentally executing the example, matchers sometimes need to try multiple alternatives and to backtrack to a later alternative if an earlier one is a blind alley.

An *acceptor* is a function that accepts a rule list and a suffix by returning some value wrapped inside the [Some constructor](#). The acceptor rejects the rule list and suffix by returning `None`. For example, the acceptor `(fun d -> function | "+"::t -> Some (d, "+"::t) | _ -> None)` accepts any rule list but accepts only suffixes beginning with `"+"`. Such an acceptor would cause the example matcher to fail on the prefix `["3"; "+"; "4"]` (since the corresponding suffix begins with `"-"`, not `"+"`) but it would succeed on the prefix `["3"]`.

By convention, an acceptor that is successful returns `Some (d, s)`, where `d` is a rule list that typically contains the acceptor's input rule list as a sublist (because the acceptor may do further parsing, and therefore has applied more rules than before), and `s` is a tail of the input suffix (again, because the acceptor may have parsed more of the input, and has therefore consumed some of the suffix). This allows the matcher's caller to retrieve the derivation for the matched prefix, along with an indication where the matched prefix ends (since it ends just before the suffix starts). Although this behavior is crucial for the internal acceptors used by your code, it is not required for top-level acceptors supplied by test programs: a top-level acceptor needs only to return a `Some x` value to succeed.

Whenever there are several rules to try for a nonterminal, you should always try them left-to-right. For example, `awkish_grammar` below contains this:

```
| Expr ->
  [[N Term; N Binop; N Expr];
  [N Term]]
```

and therefore, your matcher should attempt to use the rule `"Expr \rightarrow Term Binop Expr"` before attempting to use the simpler rule `"Expr \rightarrow Term"`.

Definitions

symbol, right hand side, rule

same as in Homework 1.

alternative list

A list of right hand sides. It corresponds to all of a grammar's rules for a given nonterminal symbol. By convention, an empty alternative list `[]` is treated as if it were a singleton list `[[]]` containing the empty symbol string.

production function

A function whose argument is a nonterminal value. It returns a grammar's alternative list for that nonterminal.

grammar

A pair, consisting of a start symbol and a production function. The start symbol is a nonterminal value.

derivation

a list of rules used to derive a phrase from a nonterminal. For example, the OCaml representation of the example derivation shown above is as follows:

```
[Expr, [N Term; N Binop; N Expr];
Term, [N Num];
Num, [T "3"];
Binop, [T "+"];
Expr, [N Term];
Term, [N Num];
Num, [T "4"]]
```

fragment

a list of terminal symbols, e.g., `["3"; "+"; "4"; "xyzzz"]`.

acceptor

a curried function with two arguments: a derivation *d* and a fragment *frag*. If the fragment is not acceptable, it returns `None`; otherwise it returns `Some x` for some value *x*.

matcher

a curried function with two arguments: an acceptor *accept* and a fragment *frag*. A matcher matches a prefix *p* of *frag* such that *accept* (when passed a derivation and the corresponding suffix) accepts the corresponding suffix (i.e., the suffix of *frag* that remains after *p* is removed). If there is such a match, the matcher returns whatever *accept* returns; otherwise it returns `None`.

Assignment

1. To warm up, notice that the format of grammars is different in this assignment, versus Homework 1. Write a function `convert_grammar gram1` that returns a Homework 2-style grammar, which is converted from the Homework 1-style grammar *gram1*. Test your implementation of `convert_grammar` on the test grammars given in Homework 1. For example, the top-level definition `let awksub_grammar_2 = convert_grammar awksub_grammar` should bind `awksub_grammar_2` to a Homework 2-style grammar that is equivalent to the Homework 1-style grammar `awksub_grammar`.
2. Write a function `parse_prefix gram` that returns a matcher for the grammar *gram*. When applied to an acceptor *accept* and a fragment *frag*, the matcher must return the first acceptable match of a prefix of *frag*, by trying the grammar rules in order; this is not necessarily the shortest nor the longest acceptable match. A match is considered to be acceptable if *accept* succeeds when given a derivation and the suffix fragment that immediately follows the matching prefix. When this happens, the matcher returns whatever the acceptor returned. If no acceptable match is found, the matcher returns `None`.
3. Write two good, nontrivial test cases for your `parse_prefix` function. These test cases should all be in the style of the test cases given below, but should cover different problem areas. Your test cases should be named `test_1` and `test_2` (note the underscores; this distinguishes your test cases from the standard ones given below). Your test cases should test at least one grammar of your own. You may reuse your test cases for Homework 1 as part of `test_1`, but `test_2` should be new.
4. Assess your work by writing an after-action report that summarizes why you solved the problem the way you did, other approaches that you considered and rejected (and why you rejected them), and any

weaknesses in your solution in the context of its intended application. If possible, illustrate weaknesses by test cases that fail with your implementation. This report should be a simple [ASCII plain text](#) file that consumes a page or so (at most 100 lines and 80 columns per line, and at least 50 lines, please). See [Resources for oral presentations and written reports](#) for advice on how to write assessments; admittedly much of the advice there is overkill for the simple kind of report we're looking for here.

Unlike Homework 1, we are expecting some weaknesses here, so your assessment should talk about them. For example, we don't expect that your implementation will work with all possible grammars, but we would like to know which sort of grammars it will have trouble with.

As with Homework 1, your code may use the [Pervasives](#) and [List](#) modules, but it should use no other modules. Your code should be free of [side effects](#). Simplicity is more important than efficiency, but your code should avoid using unnecessary time and space when it is easy to do so.

Submit

We will test your program on the SEASnet Linux servers as before, so make sure that `/usr/local/cs/bin` is at the start of your path, using the same technique as in Homework 1.

Submit three files:

- `hw2.ml` should define `parse_prefix` along with any auxiliary types and functions needed to define `parse_prefix`.
- `hw2test.ml` should contain your test cases.
- `hw2.txt` should hold your assessment.

Please do not put your name, student ID, or other personally identifying information in your files.

Sample test cases

```
let accept_all derivation string = Some (derivation, string)
let accept_empty_suffix derivation = function
  | [] -> Some (derivation, [])
  | _ -> None

(* An example grammar for a small subset of Awk, derived from but not
   identical to the grammar in
   <http://web.cs.ucla.edu/classes/winter06/cs132/hw/hw1.html>.
   Note that this grammar is not the same as Homework 1; it is
   instead the same as the grammar under "Theoretical background"
   above. *)

type awksub_nonterminals =
  | Expr | Term | Lvalue | Incrop | Binop | Num

let awkish_grammar =
  (Expr,
   function
     | Expr ->
       [[N Term; N Binop; N Expr];
        [N Term]]
     | Term ->
       [[N Num];
        [N Lvalue];
        [N Incrop; N Lvalue];
        [N Lvalue; N Incrop];
        [T "("; N Expr; T ")"]]
     | Lvalue ->
```

```

    [[T"$"; N Expr]]
| Incrop ->
    [[T"++"];
    [T"--"]]
| Binop ->
    [[T"+"];
    [T"-"]]
| Num ->
    [[T"0"]; [T"1"]; [T"2"]; [T"3"]; [T"4"];
    [T"5"]; [T"6"]; [T"7"]; [T"8"]; [T"9"]]

let test0 =
  ((parse_prefix awkish_grammar accept_all ["ouch"]) = None)

let test1 =
  ((parse_prefix awkish_grammar accept_all ["9"])
   = Some ([ (Expr, [N Term]); (Term, [N Num]); (Num, [T "9"])] , []))

let test2 =
  ((parse_prefix awkish_grammar accept_all ["9"; "+"; "$"; "1"; "+"])
   = Some
    ([ (Expr, [N Term; N Binop; N Expr]); (Term, [N Num]); (Num, [T "9"]);
      (Binop, [T "+"]); (Expr, [N Term]); (Term, [N Lvalue]);
      (Lvalue, [T "$"; N Expr]); (Expr, [N Term]); (Term, [N Num]);
      (Num, [T "1"])] ,
      [ "+"] ))

let test3 =
  ((parse_prefix awkish_grammar accept_empty_suffix ["9"; "+"; "$"; "1"; "+"])
   = None)

(* This one might take a bit longer.... *)
let test4 =
  ((parse_prefix awkish_grammar accept_all
    ["("; "$"; "8"; ")"; "-"; "$"; "++"; "$"; "--"; "$"; "9"; "+";
    "("; "$"; "++"; "$"; "2"; "+"; "("; "8"; ")"; "-"; "9"; ")";
    "-"; "("; "$"; "$"; "$"; "$"; "$"; "++"; "$"; "$"; "5"; "++";
    "++"; "--"; ")"; "-"; "++"; "$"; "$"; "("; "$"; "8"; "++"; ")";
    "++"; "+"; "0"])
   = Some
    ([ (Expr, [N Term; N Binop; N Expr]); (Term, [T "("; N Expr; T ")"]);
      (Expr, [N Term]); (Term, [N Lvalue]); (Lvalue, [T "$"; N Expr]);
      (Expr, [N Term]); (Term, [N Num]); (Num, [T "8"]); (Binop, [T "-"]);
      (Expr, [N Term; N Binop; N Expr]); (Term, [N Lvalue]);
      (Lvalue, [T "$"; N Expr]); (Expr, [N Term; N Binop; N Expr]);
      (Term, [N Incrop; N Lvalue]); (Incrop, [T "++"]);
      (Lvalue, [T "$"; N Expr]); (Expr, [N Term; N Binop; N Expr]);
      (Term, [N Incrop; N Lvalue]); (Incrop, [T "--"]);
      (Lvalue, [T "$"; N Expr]); (Expr, [N Term; N Binop; N Expr]);
      (Term, [N Num]); (Num, [T "9"]); (Binop, [T "+"]); (Expr, [N Term]);
      (Term, [T "("; N Expr; T ")"]); (Expr, [N Term; N Binop; N Expr]);
      (Term, [N Lvalue]); (Lvalue, [T "$"; N Expr]);
      (Expr, [N Term; N Binop; N Expr]); (Term, [N Incrop; N Lvalue]);
      (Incrop, [T "++"]); (Lvalue, [T "$"; N Expr]); (Expr, [N Term]);
      (Term, [N Num]); (Num, [T "2"]); (Binop, [T "+"]); (Expr, [N Term]);
      (Term, [T "("; N Expr; T ")"]); (Expr, [N Term]); (Term, [N Num]);
      (Num, [T "8"]); (Binop, [T "-"]); (Expr, [N Term]); (Term, [N Num]);
      (Num, [T "9"]); (Binop, [T "-"]); (Expr, [N Term]);
      (Term, [T "("; N Expr; T ")"]); (Expr, [N Term]); (Term, [N Lvalue]);
      (Lvalue, [T "$"; N Expr]); (Expr, [N Term]); (Term, [N Lvalue]);
      (Lvalue, [T "$"; N Expr]); (Expr, [N Term]); (Term, [N Lvalue]);
      (Lvalue, [T "$"; N Expr]); (Expr, [N Term]); (Term, [N Lvalue; N Incrop]);
      (Lvalue, [T "$"; N Expr]); (Expr, [N Term]); (Term, [N Lvalue; N Incrop]);
      (Lvalue, [T "$"; N Expr]); (Expr, [N Term]); (Term, [N Incrop; N Lvalue]);
      (Incrop, [T "++"]); (Lvalue, [T "$"; N Expr]); (Expr, [N Term]);
```

```
(Term, [N Lvalue; N Incrop]); (Lvalue, [T "$"; N Expr]); (Expr, [N Term]);
(Term, [N Num]); (Num, [T "5"]); (Incrop, [T "++"]); (Incrop, [T "++"]);
(Incrop, [T "--"]); (Binop, [T "-"]); (Expr, [N Term]);
(Term, [N Incrop; N Lvalue]); (Incrop, [T "++"]);
(Lvalue, [T "$"; N Expr]); (Expr, [N Term]); (Term, [N Lvalue; N Incrop]);
(Lvalue, [T "$"; N Expr]); (Expr, [N Term]);
(Term, [T "("; N Expr; T ")"]); (Expr, [N Term]);
(Term, [N Lvalue; N Incrop]); (Lvalue, [T "$"; N Expr]); (Expr, [N Term]);
(Term, [N Num]); (Num, [T "8"]); (Incrop, [T "++"]); (Incrop, [T "++"]);
(Binop, [T "+"]); (Expr, [N Term]); (Term, [N Num]); (Num, [T "0"]),
[]))
```

```
let rec contains_lvalue = function
| [] -> false
| (Lvalue,_)::_ -> true
| _::rules -> contains_lvalue rules
```

```
let accept_only_non_lvalues rules frag =
  if contains_lvalue rules
  then None
  else Some (rules, frag)
```

```
let test5 =
  ((parse_prefix awkish_grammar accept_only_non_lvalues
    ["3"; "-"; "4"; "+"; "$"; "5"; "-"; "6"])
  = Some
    ([ (Expr, [N Term; N Binop; N Expr]); (Term, [N Num]); (Num, [T "3"]);
      (Binop, [T "-"]); (Expr, [N Term]); (Term, [N Num]); (Num, [T "4"]) ],
      ["+"; "$"; "5"; "-"; "6"])))
```

Sample use of test cases

If you put the sample test cases into a file `hw2sample.ml`, you should be able to use it as follows to test your `hw2.ml` solution on the SEASnet implementation of OCaml. Similarly, the command `#use "hw2test.ml";;` should run your own test cases on your solution.

```
$ ocaml
      OCaml version 4.03.0

# #use "hw2.ml";;
...
val parse_prefix :
  'a * ('a -> ('a, 'b) symbol list list) ->
  (('a * ('a, 'b) symbol list) list -> 'b list -> ('c list * 'd) option) ->
  'b list -> ('c list * 'd) option = <fun>
...
# #use "hw2sample.ml";;
val accept_all : 'a -> 'b -> ('a * 'b) option = <fun>
val accept_empty_suffix : 'a -> 'b list -> ('a * 'c list) option = <fun>
type awksub_nonterminals = ...
val awkish_grammar :
  awksub_nonterminals *
  (awksub_nonterminals -> (awksub_nonterminals, string) symbol list list) =
  (Expr, <fun>)
val test0 : bool = true
val test1 : bool = true
val test2 : bool = true
val test3 : bool = true
val test4 : bool = true
val contains_lvalue : (awksub_nonterminals * 'a) list -> bool = <fun>
val accept_only_non_lvalues :
  (awksub_nonterminals * 'a) list ->
  'b -> ((awksub_nonterminals * 'a) list * 'b) option = <fun>
```

```
val test5 : bool = true  
#
```

Hint

You can use [a previous Homework 2](#) as a hint. It is a tough homework and is not the same problem but there are some common ideas. Look for the sample solution at the end.

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