

# Algorithm task

## Dominator

---

### 1) First approach in c (for loop approach)

```
#include <stdio.h>
#include <stdlib.h>

int main()
{
    int n; scanf("%d",&n);
    int arr[n];
    for(int i=0;i<n;i++){
        scanf("%d",&arr[i]);
    }
    int count1 = 0 , count2 = 0, dom = arr[0] ;

    for(int i=0;i<n;i++){
        int e = arr[i];
        count1=0;
        for(int j=0;j<n;j++){
            if(arr[j]==e) count1++;
        }
        if(count2<count1){
            count2 = count1;
            dom = arr[i];
        }
    }
    // printf("%d\n",count2);
    // printf("%d\n",dom);
    if(count2 <= n/2) printf("-1");
    else{
        for(int i=0;i<n;i++){
            if(arr[i]==dom) printf("%d ",i);
        }
    }
    return 0;
}
```

Note : this approach is the worst in terms of time complexity

As it has a time complexity of  $O(n^2)$

$$T(n) \approx n^2 + 3n + 4 + c$$

---

Pseudocode :

```
For( i=0 to n-1 ){
```

```
  Elem ← arr[i];    //variables/array scanned from user :
```

```
  Count_in ←0;      // n, arr[];
```

```
    For( j=0 to n-1){    //variables initialized :
```

```
      If( arr[j] = Elem)    //count_in ←0,count_fin←0
```

```
        Count_in +=1; //Dominator ←arr[0]
```

```
    }
```

```
  If( count_fin<count_in ){
```

```
    Count_fin ←count_in;
```

```
    Dominator ←arr[i]
```

```
  }
```

```
}
```

---

If( count\_fin <= (n/2) )

Print( -1 );

Else{

For( i=0 to n-1 ){

If( arr[i] = Dominator )

Print( i );

}

}

---

## 2) Second approach in c( frequency arrays)

```
#include <stdio.h>
#include <stdlib.h>
int main()
{
    int x; scanf("%d",&x);
    int arr[x];
    int freq[100000]={0};
    for(int i=0;i<x;i++){
        scanf("%d",&arr[i]);
        freq[arr[i]]++;
    }
    int maxe = arr[0] , maxx = 0;

    for(int i=0;i<x;i++){
        if(maxx<freq[arr[i]]){
            maxx = freq[arr[i]];
            maxe = arr[i];
        }
    }
    if(maxx<=(x/2))
        printf("-1");
}
```

```

else{
for(int i=0;i<x;i++){
if(arr[i]==maxe){
printf("%d ",i);
}
}
}
return 0;
}

```

Note : this approach is less stable when dealing with numbers more than 100000 and negative numbers but has less time complexity than the previous one  $O(n)$

$$T(n) \approx 3n + 3 + c$$

Pseudocode :

//after scanning number of elements and array elements from the user , we initialized an array called freq[] and initialized its elements with zeros.

//this algorithm uses elements of the original array as an index of the freq[] array

//we initialized elements :  $\text{maxE} \leftarrow \text{arr}[0]$  and

$\text{max} \leftarrow 0;$

---

```

For( i=0 to n-1 ){
Freq[arr[i]]++;
}

For( i=0 to n-1 ){
If( max<freq[ arr[i] ] ){
max ←freq[arr[i]];
maxe ←arr[i];
}
}

If( max <=(n/2))
    Printf( -1 );

Else{
For( i=0 to n-1 ){
If(arr[i] = maxe){
print( i );
}
}

```

---

}

}

---

### 3) Third approach in c++ (most optimal)

```
#include <iostream>
#include <bits/stdc++.h>

using namespace std;

int main()
{
    unordered_map<int,int> mp;
    int x; cin>>x;
    int arr[x];
    for(int i=0;i<x;i++){
        cin>>arr[i];
        mp[arr[i]]++;
    }
    int mx= mp[arr[0]];
    int maxe = arr[0];
    for(int i=0;i<x;i++){
        if(mx<mp[arr[i]]){
            mx = mp[arr[i]];
            maxe = arr[i];
        }
    }
    // cout<<mx<<"\n"<<maxe<<"\n";
    if(mx <=(x/2)){
        cout<<-1;
    }else{
        for(int i=0;i<x;i++){
            if(arr[i]==maxe) cout<<i<<" ";
        }
    }

    return 0;
}
```

This approach we used unordered map and this code has  $O(n)$  also

And a  $T(n) \approx 3n + 3 + c$

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But this code is able to handle larger input values with better time complexity .

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Pseudocode :

```
// initialize an unordered map from the STL in c++
```

```
//unordered_map<int,int> mp;
```

```
//we will use the same method as the frequency array  
but inserting and counting will be more efficient as we  
can handle negative and positive numbers
```

```
[-2147483648 , 2147483647]
```

```
For(i=0 to n-1 )
```

```
{
```

```
mp[arr[i]]++;
```

```
}
```

---

```
For(i=0 to n-1){  
  If(max < mp[arr[i]]){  
    max = mp[arr[i]];  
    maxe = arr[i];  
  }  
}  
  
If(max<=(n/2)){  
  Print( -1 );  
}  
else{  
  For(i=0 to n-1){  
    If(arr[i] = maxe)  
      Print (i);  
  }  
}
```

---



#### 4)fourth and final approach (recursive)

```
#include <bits/stdc++.h>

using namespace std;

int find_candidate(int A[], int size)
{
    int candidate = A[0];
    int count = 1;

    for (int i = 1; i < size; i++)
    {
        if (A[i] == candidate)
        {
            count++;
        }
        else
        {
            count--;
            if (count == 0)
            {
                candidate = A[i];
                count = 1;
            }
        }
    }

    return candidate;
}

int count_occurrences(int A[], int size, int candidate)
{
    int count = 0;
    for (int i = 0; i < size; i++)
    {
        if (A[i] == candidate)
        {
            count++;
        }
    }

    return count;
}
```

```

}

int find_dominator_index(int A[], int start, int end)
{
    if (start == end)
    {
        return start;
    }

    int mid = start + (end - start) / 2;

    int left_dominator = find_dominator_index(A, start, mid);
    int right_dominator = find_dominator_index(A, mid + 1, end);

    if (left_dominator == right_dominator)
    {
        return left_dominator;
    }

    int left_candidate = A[left_dominator];
    int right_candidate = A[right_dominator];

    int left_count = count_occurrences(A, end - start + 1, left_candidate);
    int right_count = count_occurrences(A, end - start + 1, right_candidate);

    if (left_count > (mid - start + 1) / 2)
    {
        return left_dominator;
    }
    else if (right_count > (end - mid) / 2)
    {
        return right_dominator;
    }
    else
    {
        return -1;
    }
}

int main()
{
    int A[] = {3, 4, 3, 2, 3, -1, 3, 3};
    int size = sizeof(A) / sizeof(A[0]);
    int result = find_dominator_index(A, 0, size - 1);
    if (result != -1)

```

```

{
    cout << "Dominator index: " << result << endl;
}
else
{
    cout << "No dominator found." << endl;
}
return 0;
}

```

In this approach we used a recursive divide and conquer algorithm (binary search like) to find the index of the dominator

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This algorithm has a time complexity of  $O(n \log n)$

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Pseudocode :

Find\_candidate(int Arr[],size){

    Candidate  $\leftarrow$  arr[0]

    Count  $\leftarrow$  1

For(i=1 to size){

    If(arr[i]=candidate) count +=1;

    Else{

        Count-=1;

```
If(count =0)
```

```
{
```

```
  Candidate = Arr[i];
```

```
  Count = 1;
```

```
}
```

```
}
```

```
}
```

```
Return candidate;
```

```
}
```

```
int count_occurrences(A,size, candidate) {
```

```
  count  $\leftarrow$  0
```

```
  for (i = 0 to size – 1) {
```

```
    if A[i] = candidate {
```

```
      count  $\leftarrow$  count + 1
```

```
    }
```

```
  }
```

```
    return count
}

int find_dominator_index(A, start, end) {
    if start == end {
        return start
    }

    mid <-- start + (end - start) / 2

    left_dominator <-- find_dominator_index(A, start, mid)
    right_dominator <-- find_dominator_index(A, mid + 1,
end)

    if left_dominator == right_dominator {
        return left_dominator
    }
```

```
left_candidate <-- A[left_dominator]
```

```
right_candidate <-- A[right_dominator]
```

```
left_count <-- count_occurrences(A, left_candidate)
```

```
right_count <-- count_occurrences(A, right_candidate)
```

```
if left_count > (mid - start + 1) / 2 {
```

```
    return left_dominator
```

```
} else if right_count > (end - mid) / 2 {
```

```
    return right_dominator
```

```
} else {
```

```
    return -1
```

```
}}
```

```
//main function with array declaration and function  
calling
```

Points of comparison	Algorithm 1 in c	Algorithm 2 in c	Algorithm 3 in c++	Algorithm 4 in C++
Time complexity	$O(n^2)$	$O(n)$	$O(n)$	$O(n \log n)$
Accuracy	Pretty accurate except for the time complexity part	Not accurate with edge cases	Most accurate one	Accurate enough but only returns 1 index
Recurrence relation	$T(n) \approx n^2 + 3n + 4 + c$	$T(n) \approx 3n + 3 + c$	$T(n) \approx 3n + 3 + c$	$T(n) \approx 2(t/2) + n + c$