

AUTOMATA THEORY AND COMPUTABILITY

V Sem CSE

Module-2 REGULAR EXPRESSIONS

Regular Expressions (RE) is a string that can be framed according to the following rules.

1. \emptyset is a regular expression
2. ϵ " "
3. Every element in Σ is a RE
4. Given 2 REs α and β , $\alpha\beta$ is RE
5. " " " $\alpha \cup \beta$ is RE
6. Given a RE α , α^* is a RE
7. " " " α^+ is a RE
8. " " " (α) is a RE

Every regular expression contains 2 kinds of symbols

1. an alphabet Σ
2. Special symbol $\emptyset, \cup, \epsilon, (), *, +$

Example:

Analyze simple RE

$$\begin{aligned} L((a \cup b)^* b) &= L((a \cup b)^*) L(b) \\ &= (L(a \cup b))^* L(b) \\ &= (L(a) \cup L(b))^* L(b) \\ &= (\{a\} \cup \{b\})^* \{b\} \\ &= \{a, b\}^* \{b\} \end{aligned}$$

The meaning of RE $(a \cup b)^*b$ is the set of all strings over the alphabet $\{a, b\}$ that end in b .

Given a language, find a regular expression.

Ex. 1. $L = \{w \in \{a, b\}^* : |w| \text{ is even}\}$

$$((a \cup b)(a \cup b))^* \text{ or } ((a+b)(a+b))^*$$

Ex. 2. $L = \{w \in \{a, b\}^* : |w| \leq 2\}$

i.e. $|w|$ can be 0, 1 or 2

ϵ for 0
 $\left. \begin{matrix} a \\ b \end{matrix} \right\}$ for 1

i.e. $(\epsilon + a + b)$

$$\begin{aligned} \text{for } |w| = 2 & (\epsilon + a + b)(\epsilon + a + b) \\ &= (\epsilon + a + b)^2 \text{ or } (\epsilon \cup a \cup b)^2 \end{aligned}$$

Ex. 3. $L = \{w \in \{a, b\}^* : |w| \text{ is odd}\}$

$$(a \cup b)((a \cup b)(a \cup b))^*$$

or

$$(a + b)((a + b)(a + b))^*$$

Ex. 4. $L = \{w \in \{a, b\}^* : w \text{ contains alternate } a\text{'s and } b\text{'s}\}$

$(ab)^*$ - but this always begins with a and ends with b .

Hence add $(\epsilon + b)$ at beginning &
 $(\epsilon + a)$ at end.

$$\text{i.e. } (\epsilon + b)(ab)^*(\epsilon + a).$$

Ex. 5. $L = \{w \in \{0,1\}^* : w \text{ contains at most one pair of consecutive 0's}\}$

- No zero's 1^*
- with one zero 01 or 10
- with a single pair of 00 .

$$(1+01)^*0(1+01)^* \quad \text{with } 01$$

$$(1+10)^*0(1+10)^* \quad \text{with } 10.$$

Hence $RE = 1^* + (1+01)^*0(1+01)^* + (1+10)^*0(1+10)^*$

Ex. 6. $L = \{w \in \{a,b,c\}^* : w \text{ contains at least one } a \text{ and at least one } b\}$

minimum $w = ab$ or ba

It can begin with anything, end with anything and in between also can be anything.

$$(a+b)^*a(a+b)^*b(a+b)^* \quad \text{for } \underline{ab}$$

$$(a+b)^*b(a+b)^*a(a+b)^* \quad \text{for } \underline{ba}$$

$$\text{hence } RE = [(a+b)^*a(a+b)^*b(a+b)^*] + [(a+b)^*b(a+b)^*a(a+b)^*]$$

Ex. 7

$L = \{w \in \{a, b\}^* : w \text{ ends with } b \text{ and has no substring } aa\}$

end with b but no aa .

is either \underline{b} or $a\underline{b}$.

$(b+ab)$ \underline{b} or $a\underline{b}$

but b or ab can come any no. of times

$$(b+ab)(b+ab)^* = (b+ab)^+$$

Ex. 8 $L = \{w \in \{a, b\}^* : w \text{ begins with } a \text{ and ends with } b\}$

begins with a and ends with $b = ab$
in between it could be anything.

$$RE = a(a+b)^*b.$$

Ex. 9 $L = \{w \in \{a, b\}^* : \text{second symbol from right end is } a\}$

Initially anything. $(a+b)^*$.

second sym from right is a $(a+b)^*a(a+b)$

last sym can be \swarrow
either a or b .

$$RE = (a+b)^*a(a+b)$$

Ex. 10.

$L = \{w \in \{a, b\}^* : w \text{ begins \& ends with same symbol}\}$

begins ~~with~~ & ends with a = $a(a+b)^*a$

" " " b = $b(a+b)^*b$

$$\therefore RE = a(a+b)^*a + b(a+b)^*b.$$

Ex. 11.

$L = \{w \in \{a, b\}^* : |w| \text{ is even or multiple of } 3\}$

length is even $[(a+b)(a+b)]^*$

length multiple of three $((a+b)(a+b)(a+b))^*$

$$RE = ((a+b)(a+b))^* + ((a+b)(a+b)(a+b))^*.$$

Ex. 12.

$L = \{w \in \{a, b\}^* : w = a^m b^n \text{ } m \text{ is even \& } n \text{ is odd}\}$

i.e. even no. of a's followed by odd no. of b's

even no. of a's $(aa)^*$
odd " " b's $b(bb)^*$

$$RE = (aa)^* b (bb)^*.$$

Ex. 13. $L = \{w \in \{a, b\}^* : w = a^m b^n \text{ } m+n \text{ is even}\}$

we know that even + even = even
odd + odd = even

even No of a's & even No of b's $= (aa)^*(bb)^*$
 odd " " & odd " " $= a(aa)^*b(bb)^*$

$$RE = (aa)^*(bb)^* + a(aa)^*b(bb)^*$$

Ex. 14. $L = \{a^n b^m \mid m \geq 1, n \geq 1, mn \geq 3\}$

we know that

$mn \geq 3$ when

- 1) $n=1$ & $m \geq 3$ or
- 2) $n \geq 3$ & $m=1$ or
- 3) $n \geq 2$ & $m \geq 2$

- 1) $a b b b b^*$ $|a|=1$ $|b| \geq 3$
- 2) $a a a a^* b$ $|a| \geq 3$ $|b|=1$
- 3) $a a a^* b b b^*$ $|a| \geq 2$ $|b| \geq 2$

$$RE = a b b b b^* + a a a a^* b + a a a^* b b b^*.$$

Ex. 15. $L = \{w \in \{a, b\}^* : |w| \bmod 3 = 0\}$

$$((a+b)(a+b)(a+b))^*$$

Ex. 16. $L = \{w \in \{a, b\}^* : \#_a(w) \bmod 3 = 0\}$

$\epsilon, aaa, abaa, aaabababba,$

$$b^*(a(b^*ab^*ab^*ab^*))^*$$

Building an FSM from a regular expression (RE)

Theorem :- For any every RE there is an equivalent FSM.

proof :- The proof is by construction of FSM.

We will show that, given a RE α , we can construct an FSM M . Such that $L(\alpha) = L(M)$.

There exists an FSM that corresponds to each primitive RE.

If α is any $c \in \Sigma$



If α is \emptyset

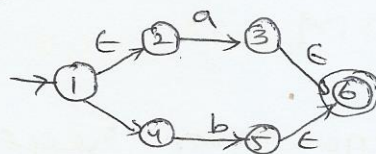


If α is ϵ



If α is RE = $\beta \cup \gamma$

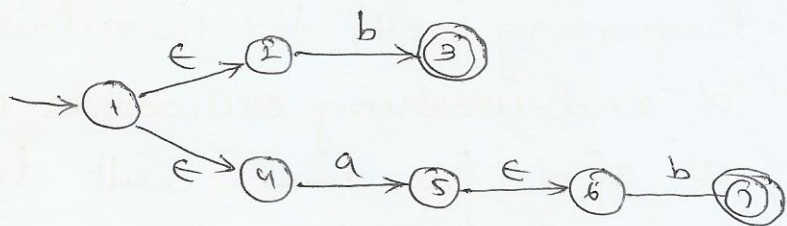
$\alpha = a \cup b$



$\alpha = \beta \cup \gamma$

$\beta = b$ $\gamma = ab$

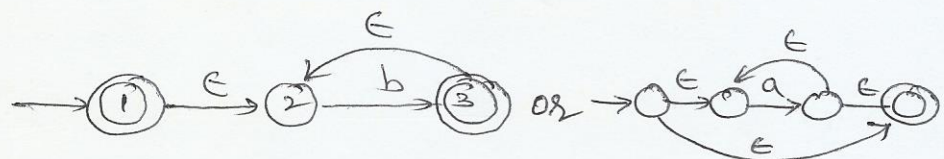
ie $\alpha = b \cup ab$



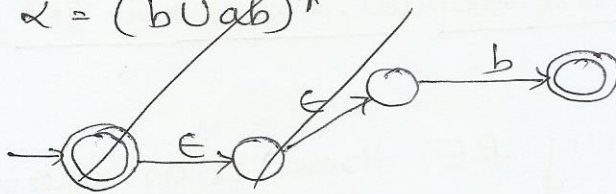
If $\alpha = ab$



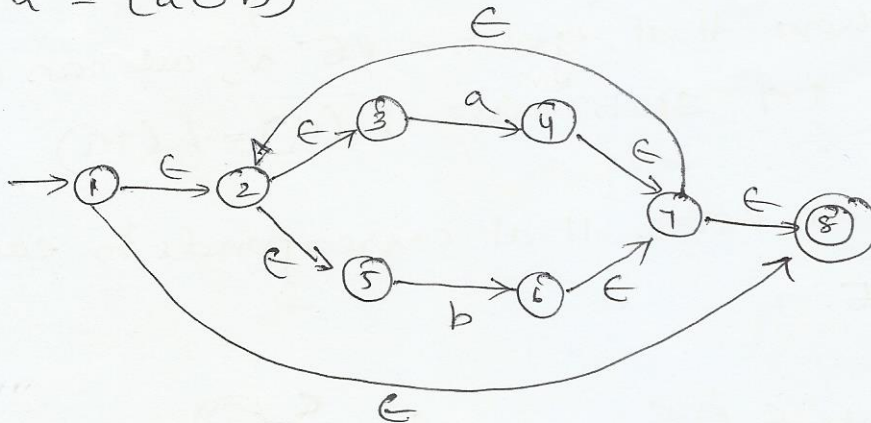
If $\alpha = b^*$



If $\alpha = (b \cup ab)^*$



If $\alpha = (a \cup b)^*$



Based on the construction, we define algorithm to construct, given a RE α , a corresponding PSTN.

Algorithm regextofsm ($\alpha: RE$) =

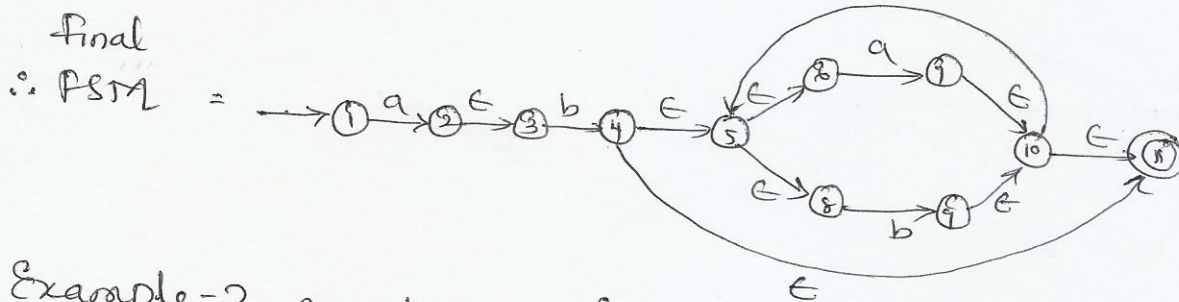
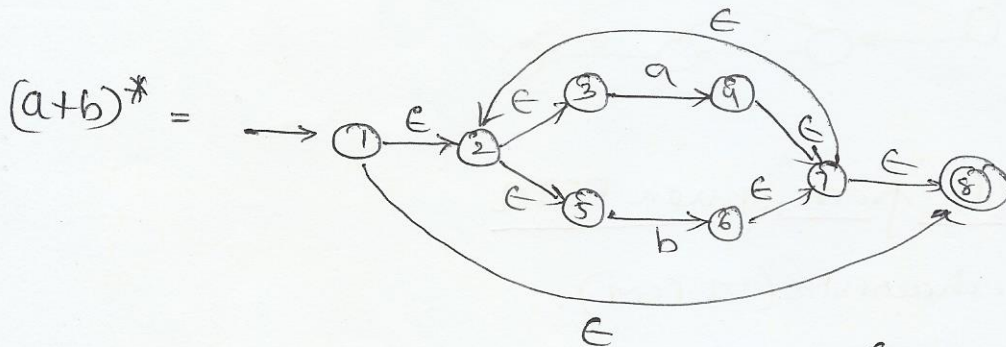
Beginning with the primitive subexpression of α and working outwards until an PSTN for all of α has been built do:

Construct an PSTN as described above.

Proof:

Example 1: Build FSM from a given RE

$$RE = ab(a+b)^* \quad \text{or} \quad ab(a \cup b)^*$$



Example-2. Build FSM from a given RE

$$RE = a^* + b^* + c^*$$

