

# AUTOMATA THEORY AND COMPUTABILITY

V Sem CSE

## Module-2 REGULAR EXPRESSIONS

Regular Expressions (RE) is a string that can be formed according to the following rules.

1.  $\emptyset$  is a regular expression
2.  $a \in \Sigma$
3. Every element in  $\Sigma$  is a RE
4. Given 2 REs  $\alpha$  and  $\beta$ ,  $\alpha\beta$  is RE
5.  $\alpha^*$  is RE
6. Given a RE  $\alpha$ ,  $\alpha^*$  is a RE
7.  $\alpha^+$  is a RE
8.  $(\alpha)$  is a RE

Every regular expression contains 2 kinds of symbols

1. an alphabet  $\Sigma$
2. Special symbol  $\emptyset$ ,  $\cup$ ,  $\in$ ,  $\subset$ ,  $*$ ,  $+$

### Example:

Analyze simple RE

$$\begin{aligned} L((a \cup b)^* b) &= L((a \cup b)^*) L(b) \\ &= (L(a \cup b))^* L(b) \\ &= (L(a) \cup L(b))^* L(b) \\ &= (\{a\} \cup \{b\})^* \{b\} \\ &= \{a, b\}^* \{b\} \end{aligned}$$

The meaning of RE  $(a \cup b)^* b$  is the set of all strings over the alphabet  $\{a, b\}$  that end in  $b$ .

Given a language, find a regular expression

Ex.  $L = \{w \in \{a, b\}^*: |w| \text{ is even}\}$

$$((a \cup b)(a \cup b))^* \text{ or } ((a + b)(a + b))^*$$

Ex2.  $L = \{w \in \{a, b\}^*: |w| \leq 2\}$

i.e.  $|w|$  can be 0, 1 or 2

$$\begin{cases} \epsilon & \text{for 0} \\ a \\ b \end{cases} \begin{cases} \epsilon + a + b & \text{i.e. } (\epsilon + a + b) \\ \} & \text{for 1} \end{cases}$$

$$\begin{aligned} \text{for } |w|=2 & (\epsilon + a + b)(\epsilon + a + b) \\ & = (\epsilon + a + b)^2 \text{ or } (\epsilon + a \cup b)^2 \end{aligned}$$

Ex.3.  $L = \{w \in \{a, b\}^*: |w| \text{ is odd}\}$

$$(a \cup b) - ((a \cup b)(a \cup b))^*$$

or

$$(a + b)((a + b)(a + b))^*$$

Ex.4.  $L = \{w \in \{a, b\}^*: w \text{ contains alternate a's and b's}\}$

$(ab)^*$  - but this always begins with a and ends never b

Hence add  $(\epsilon + b)$  at beginning &  
 $(\epsilon + a)$  at end.

$$\text{i.e. } (\epsilon + b)(ab)^*(\epsilon + a).$$

Ex 5.  $L = \{w \in \{0,1\}^*: w \text{ contains at most one pair of consecutive } 0's\}$

- No zero's  $1^*$
- with one zero  $01$  or  $10$
- with a single pair of  $00$ .

$$(1+01)^* 0 (1+01)^* \quad \text{with } 01$$

$$(1+10)^* 0 (1+10)^* \quad \text{with } 10.$$

Hence

$$RE = 1^* + (1+01)^* 0 (1+01)^* + (1+10)^* 0 (1+10)^*$$

Ex. 6.  $L = \{w \in \{a,b,c\}^*: w \text{ contains at least one } a \text{ and at least one } b\}$

minimum  $w = ab$  or  $ba$

It can begin with anything, end with anything and in between also can be anything.

$$(a+b)^* a (a+b)^* b (a+b)^* \text{ for } ab$$

$$(a+b)^* b (a+b)^* a (a+b)^* \text{ for } ba$$

$$\text{hence } RE = [(a+b)^* a (a+b)^* b (a+b)^*] + [(a+b)^* b (a+b)^* a (a+b)^*]$$

Ex. 7  $L = \{w \in \{a, b\}^*: w \text{ ends with } b \text{ and has no substring aa}\}$

end with  $b$  but no  $aa$ .

is either  $\underline{b}$  or  $\underline{ab}$ .

$(b+ab) \quad \underline{b} \text{ or } ab$

but  $b$  or  $ab$  can come any no. of times

$$(b+ab)(b+ab)^* = (b+ab)^+$$

Ex. 8  $L = \{w \in \{a, b\}^*: w \text{ begins with } a \text{ and ends with } b\}$

begins with  $a$  and ends with  $b = ab$   
in between it could be anything.

$$RE = a(a+b)^*b.$$

Ex. 9  $L = \{w \in \{a, b\}^*: \text{second symbol from right end is } a\}$

Initially anything.  $(a+b)^*$

Second sym from right is  $a$

$$(a+b)^*a(a+b)$$

last sym can be  
either  $a$  or  $b$ .

$$RE = (a+b)^*a(a+b)$$

Ex. 10.

$L = \{ w \in \{a, b\}^*: w \text{ begins \& ends with same symbol} \}$

begins ~~and~~ & ends with  $a = a(a+b)^*a$   
 " " " "  $b = b(a+b)^*b$

$$\therefore RE = a(a+b)^*a + b(a+b)^*b.$$

Ex. 11.

$L = \{ w \in \{a, b\}^*: |w| \text{ is even or multiple of } 3 \}$

length is even  $[(a+b)(a+b)]^*$

length multiple of three  $((a+b)(a+b)(a+b))^*$ .

$$RE = ((a+b)(a+b))^* + ((a+b)(a+b)(a+b))^*.$$

Ex. 12.

$L = \{ w \in \{a, b\}^*: w = a^m b^n \text{ m is even \& n is odd} \}$

i.e even no. of a's followed by odd no. of b's

even no. of a's  $(aa)^*$

odd " " b's  $b(bb)^*$

$$RE = (aa)^* b(bb)^*.$$

Ex. 13.  $L = \{ w \in \{a, b\}^*: w = a^m b^n \text{ m+n is even} \}$

we know that even + even = even  
 odd + odd = even

even no of a's & even no of b's  $-(aa)^*(bb)^*$   
 odd " " & odd " "  $= a(aa)^* b(bb)^*$

$$RE = (aa)^*(bb)^* + a(aa)^* b(bb)^*$$

Ex.14.  $L = \{a^n b^m \mid n \geq 1, m \geq 1, nm \geq 3\}$

we know that

$$nm \geq 3 \text{ when}$$

$$1) n=1 \& m \geq 3 \quad \text{or}$$

$$2) n \geq 3 \& m=1 \quad \text{or}$$

$$3) n \geq 2 \& m \geq 2$$

$$1) \underline{abb} b^* \quad |a|=1 \quad |b| \geq 3$$

$$2) aaa a^* b \quad |a| \geq 3 \quad |b|=1$$

$$3) aaa^* bbb^* \quad |a| \geq 2 \quad |b| \geq 2$$

$$RE = abbbb^* + aaa a^* b + aaa^* bbb^*.$$

Ex.15.  $L = \{w \in \{a,b\}^* : |w| \bmod 3 = 0\}$

$$(a+b)(a+b)(a+b)^*$$

Ex.16.  $L = \{w \in \{a,b\}^* : \#_a(w) \bmod 3 = 0\}$

$\epsilon, aaa, abaa, aaabababba,$

$$b^* (a (b^* a b^* a b^*)^*)$$

## Building an FSM from a regular expression (RE).

Theorem :- For any every RE there is an equivalent FSM.

Proof :- The proof is by construction of FSM.

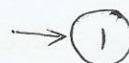
We will show that, given a RE  $\alpha$ , we can construct an FSM  $M$ . Sketch that  $L(\alpha) = L(M)$ .

There exists an FSM that corresponds to each primitive RE.

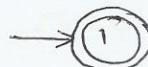
If  $\alpha$  is any  $c \in \Sigma$ .



If  $\alpha = \emptyset$

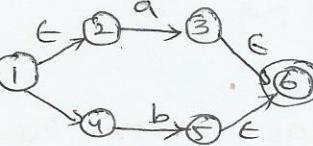


If  $\alpha \in \epsilon$



If  $\alpha$  is RE =  $B \cup r$

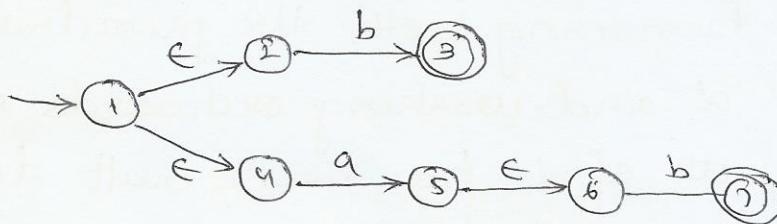
$$L = a \cup b$$



$\alpha = B \cup r$

$$B = b \quad r = ab$$

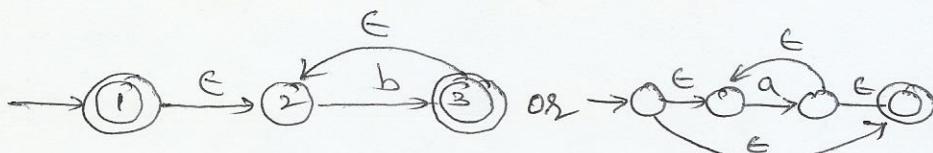
$$\text{i.e } \alpha = b \cup ab$$

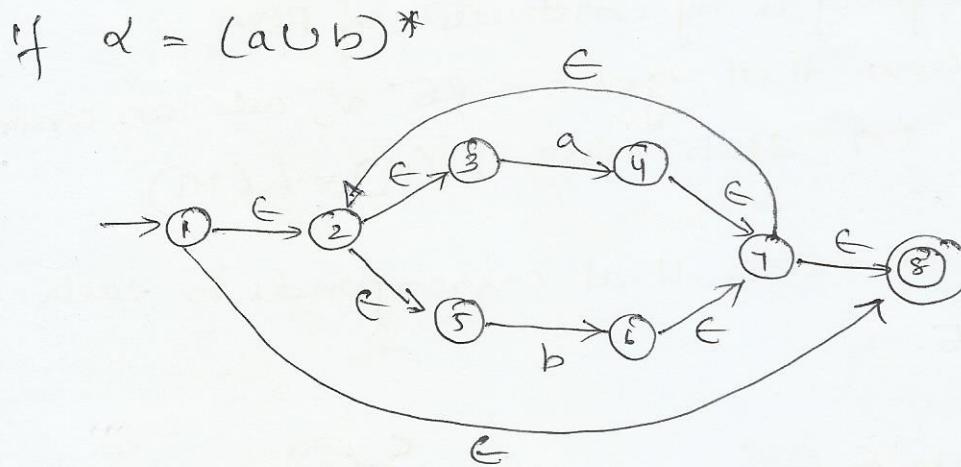
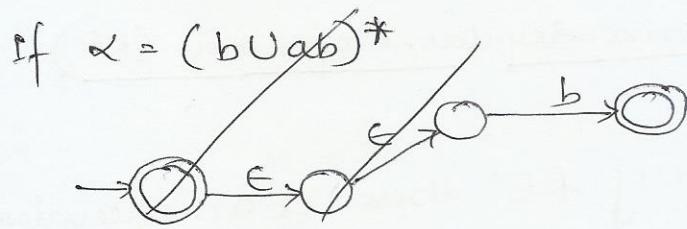


If  $\alpha = ab$



If  $\alpha = b^*$





Based on the construction, we define algorithm to construct, given a RE  $\alpha$ , a corresponding PSM.

Algorithm to construct PSM for  $\alpha$  (RE) =

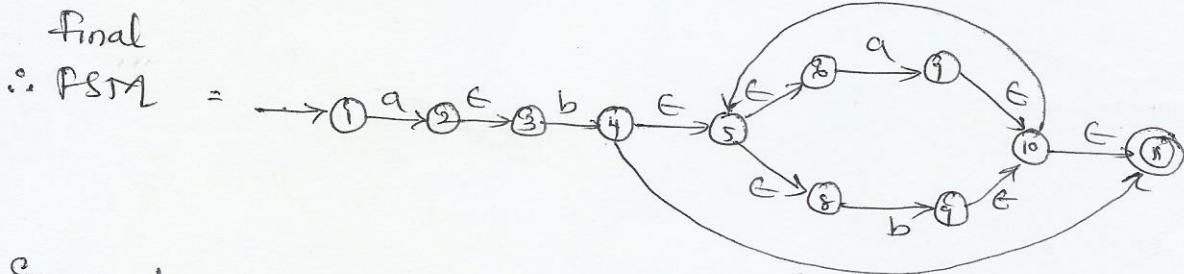
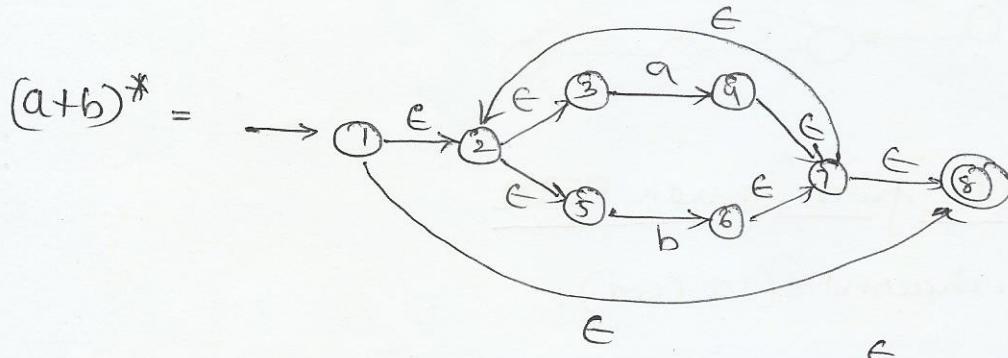
Beginning with the primitive subexpression of  $\alpha$  and working outwards until an PSM for all of  $\alpha$  has been built do:

    Construct an PSM as described above.

Point:

Example 1: Build FSM from a given RE

$$RE = ab(a+b)^* \quad \text{or} \quad ab(a \cup b)^*$$



Example - 2: Build FSM from a given RE

$$RE = a^* + b^* + c^*$$

