



Complete Sorting algorithm

- **Bubble Sort**
- **Stable & Unstable sort**
- **Selection sort**
- **Online sorting**



◀ SWIPE ▶



|| follow || for more update

Sorting Algorithm

Sorting →

When we have data and we arrange that data, where we bring the meaning of that data, that called → Sorting

→ Sorting can be easy for human in Short data like

$$[3, 4, 1, 5, 9, 19] \xrightarrow{\text{Sorting}} [1, 3, 4, 5, 9, 19]$$

→ But where Huge data like → in Millions for sorting
We need computer. by this we can solved this problem
Using python programming.



Types of Shorting algorithm →

- Bubble sort algorithm
- Selection sort algorithm
- Insertion sort algorithm



Bubble sort algorithm

inspire by Bubble → when water heated



the same step, follow in the bubble sort algorithm ↴

Algorithm

[5, 3, 1, 2, 4] Apply bubble sort on this list

Step 1 :-  →

[5, 3, 1, 2] → Start moving towards right.

one by one and if we find a number which is bigger, we have to push him / Bubble him to the right.



Step 2 :- [5, 3, 1, 2, 4]

Here when we comparing 5, 3 and i show that 5 is bigger than 3, so we swap 5 to the right. And the list is → [3, 5, 1, 2, 4] → Again go and compare (5, 1)

[3, 1, 5, 2, 4] → Again compare and swap

[3, 1, 2, 5, 4] → compare & swap

[3, 1, 2, 4, 5] → Now Repeat for the Next (3, 1, 2, 4) element.

Rule :-

① `range(5)` → 0, 1, 2, 3, 4

② `range(1, 5)` → 1, 2, 3, 4.

③ `range(5, 0, -1)` → 5, 4, 3, 2, 1 (-ve direction)

⊕

```

: # bubble sorting
def bubbleSort(arr):
    outer loop → for i in range(len(arr)-1, 0, -1):
        inner loop ← for j in range(i):
            if arr[j] > arr[j+1]:
                arr[j], arr[j+1] = arr[j+1], arr[j] → Swaping
    Round
        ① ② ③
        ↘ 3/2/1

```

Let → arr → [4, 2, 3, 1] , len(arr) → 4

here → len(arr)-1 → 4-1 → 3

Means → for i in range(len(arr)-1), here → i=3

and for j in range(i): , here j=0

Sorting → Step: 1 → j=3
Step: 2 → j=2 → [2, 3, 1, 4] → No swapping
↓
[3, 2, 1, 4]

Step 1 → j=3
↓
{4, 2, 3, 1} → Swapping
↓

Step: 3 → j=1 → [2, 1, 3, 4] → No swapping
↓
[1, 2, 3, 4]

→ [2, 4, 3, 1]
↓
[2, 3, 4, 1]
↓
[2, 3, 1, 4]

{Swaping for the 1st element}
done → going 4th outer loop

Bubble sorting done

Q. 2]: # bubble sorting

```

def bubbleSort(arr):
    for i in range(len(arr)-1, 0, -1):
        for j in range(i):
            if (arr[j] > arr[j+1]):
                arr[j], arr[j+1] = arr[j+1], arr[j]

```

outer loop → inner loop → swapping

]: arr = [7, 4, 6, 5, 2]
bubbleSort(arr)
print(arr)

[2, 4, 5, 6, 7]

How this code work internally

Sol →

Step: 1

Step: 1

Step: 2
j=3
[4, 6, 5, 2, 7]
↓
[4, 6, 5, 2, 7]

j=4
[7, 4, 6, 5, 2]
↓
[4, 7, 6, 5, 2]

Step: 3
j=2
[4, 5, 6, 2, 7]
↓
[4, 5, 6, 2, 7]

Step: 3
j=2
[4, 5, 6, 2, 7]
↓
[4, 5, 6, 2, 7] done

Step: 3
j=2
[4, 6, 7, 5, 2]
↓
[4, 6, 5, 7, 2]
↓
[4, 6, 5, 2, 7]

[4, 5, 3, 6, 7]
↓
[4, 2, 5, 6, 7]

[4, 2, 5, 6, 7] done

internal loop done
go to outer loop
and Repeat the
same process
until all value get
Sorted

Step: 4
j=1
[4, 2, 5, 6, 7]
↓
[2, 4, 5, 6, 7]



Bubble Sorting done

Q What is the time complexity of this code.

```
# bubble sorting

def bubbleSort(arr):
    for i in range(len(arr)-1, 0, -1): → outer loop
        for j in range(i):
            if (arr[j] > arr[j+1]):
                arr[j], arr[j+1] = arr[j+1], arr[j]
```

inner loop

- ↳ outer loop runs $\Rightarrow n$ times
- inner loop runs $\Rightarrow n$ times

So that time complexity of this code is

$$T(C) = O(n \times n)$$

$$\boxed{T(C) = O(n^2)}$$

Q In the case of worst case How many times comparison will be happen in this code?

↳ In the worst case every time we will doing the comparison. So

the no. of loop is running $\Rightarrow n^2$

So the in worst case comparison happens $\rightarrow n^2$ time

Q what is the worst case in this situation?

↳ When array already is reversing or sorted in descending order. for example $\{5, 4, 3, 2, 1\} \Rightarrow$ already sorted (worst case)

Stable Sort & Unstable Sort

Stable Sort → if you have two elements with equal keys, and one comes before the other in the input, a stable sort ensures that the element that was initially before the other will be before it in the sorted output.

ex → [Rahul, Mahesh] → Students
 ↓ ↓ ↓ ↓ ↓
 4 5 4 8 3 → heights

Problem

The class teacher said range all the students based on their height

↳ For this problem there could be two possibilities :-

① → {Rahul, Mahesh}
 ↓
 3 4 4 5 8

Both are
sorted

Stable Sorting

before sorting →

Rahul stands before mahesh

After sorting →

Still Rahul Standing before mahesh

↓

This called as → Stable sorting

② → {Mahesh, Rahul}
 ↓
 3 4 4 5 8

UnStable Sorting

Before sorting →

Rahul Stands before mahesh

After sorting →

Rahul Standing after mahesh

This change called as unstable sorting

Real world use case of stable & unstable sorting

fear mongering → spreading fake news.

WhatsApp → spreading fake news that Rohit Sharma Retired



Now people forwarding news.



In all messages → Same Keys → Rohit Sharma Retired



Key

2 → R → RSR

2 → g → RSK

⋮

n → R → RSR

} if j sort in → Unstable sorting
{ order will be Rearrange }
not a good approach

} if j sort → Stable sorting
{ order will be the same }

Right approach

By Using Stable sorting and making data in arranging we will find the person who actually start this news.

① This sorting algorithm follow which one

```
# bubble sorting

def bubbleSort(arr):
    for i in range(len(arr)-1, 0, -1):
        for j in range(i):
            if (arr[j] > arr[j+1]):
                arr[j], arr[j+1] = arr[j+1], arr[j]
```

Ⓐ Stable sorting

Ⓑ Unstable sorting

swapping

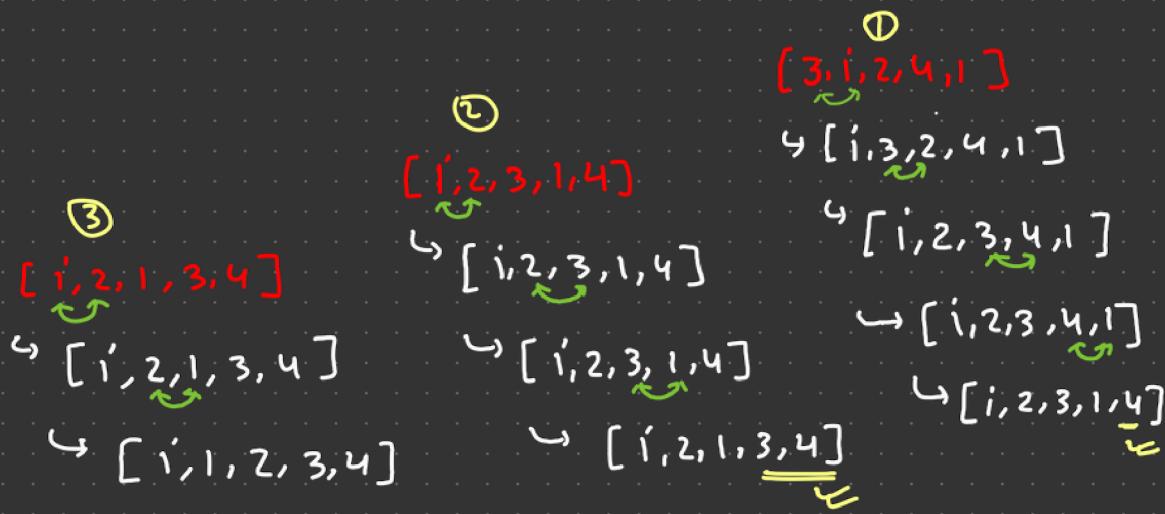
Ans → Swapping happens only when $\text{arr}(j) > \text{arr}[j+1]$

So,

if two value are same then we don't swapping here

So → this algorithm follow → Stable Sorting.

Lets understand with an example:- $\text{arr} \rightarrow [3, 1, 2, 4, 1]$



Before sorting $\text{arr} \rightarrow [3, 1, 2, 4, 1]$

after sorting $\text{arr} \rightarrow [1, 2, 3, 4]$

↳ as we can see that before sorting [1] is before [2] and after the sorting [1] still before [2].

↳ So this is Stable sorting.

{ Because order doesn't change here. }

Interview Q. 2

bubble sorting

```
def bubbleSort(arr):
    for i in range(len(arr)-1, 0, -1):
        for j in range(i):
            if arr[j] > arr[j+1]:
                arr[j], arr[j+1] = arr[j+1], arr[j]
```

n^2 time ↴

Q Apply bubble sorting on this code where $T(c) \neq O(n^2)$
array should be already sorted.
Ans in this code loop runs $\rightarrow n^2$ time
So the $T(c) \Rightarrow n^2$

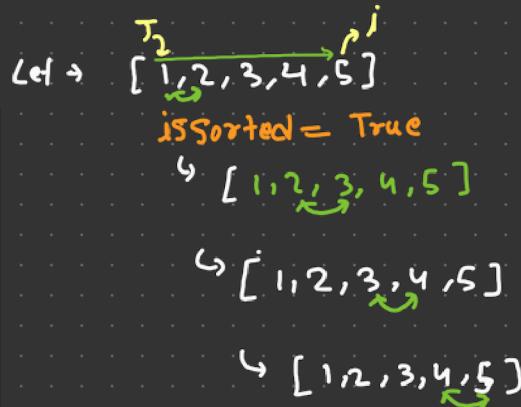
Because bubble sort not that much smart who catch given array is sorted or not \rightarrow He run Both loops,
So,

we need to make some changes in the code in this way where if give sorted array to the code it detect himself the is it sorted or not. and React differently when got sorted array.

Now the optimized code:

```
# bubble sorting

# stable sorting
def bubbleSort_optimized(arr):
    for i in range(len(arr)-1, 0, -1):
        isSorted = True
        for j in range(i):
            if arr[j] > arr[j+1]:
                isSorted = False
                not go here ↴ arr[j], arr[j+1] = arr[j+1], arr[j]
    ↴ Because ↴ if isSorted:
        ↴ already ↴ print('Array is already sorted')
        ↴ sorted ↴ break
```



Now in this case it will not go out from the outer for loop

again did print the statement and Terminate there.
in this way I don't have to run till n^2

So the $T(c) \neq O(n^2)$ where $T(c) = O(n \times 1) \Rightarrow O(n)$

Q What if we are use nearly sorted array?

Let, $[1, 2, 3, 5, 4]$

bubble sorting

```
# stable sorting
def bubbleSort_optimized(arr):
    for i in range(len(arr)-1, 0, -1):
        isSorted = True
        for j in range(i):
            if arr[j] > arr[j+1]:
                isSorted = False
                arr[j], arr[j+1] = arr[j+1], arr[j]

    if isSorted:
        print('Array is already sorted')
        break
```

Ans

(2)

$\overrightarrow{[1, 2, 3, 4, 5]}$

is sorted = True

$\hookleftarrow [\overbrace{1, 2, 3}^{\text{1}}, 4, 5]$

$\hookleftarrow [1, \overbrace{2, 3, 4}^{\text{2}}, 5]$

$\hookleftarrow [1, 2, \overbrace{3, 4}^{\text{3}}, 5]$

$\hookleftarrow [1, 2, 3, \overbrace{4, 5}^{\text{4}}]$

↓

Not sorted Fully

1 loop
Runs only for \Rightarrow 2 times

$\overrightarrow{[1, 2, 3, 4, 5]}$

is sorted = True

$\hookleftarrow [1, 2, 3, 4, 5]$

$\hookleftarrow [1, 2, 3, 4, \overbrace{5}^{\text{5}}]$

$\hookleftarrow [1, 2, 3, 4, 5]$

is sorted = False

$\overline{[1, 2, 3, 4, 5]}$

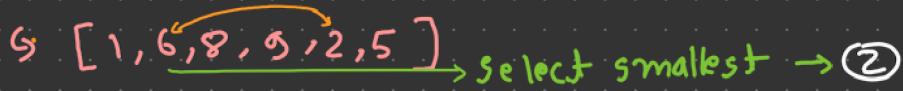
then it came to outer loop

Now then loop Break

Selection sorting

Think of arranging books on a shelf. You find the smallest book and put it first. Then, find the next smallest and put it next. Repeat until all books are in order.

Let,



in this example we find → smallest

we also find the largest value as value with selection sort technique

Ex. 2

$\left[\begin{array}{c} 3, 1, 5, 4, 2 \end{array} \right]$ *replace* \rightarrow Select largest value $\rightarrow 5$

$\hookrightarrow \left[\begin{array}{c} 3, 1, 2, 4, 5 \end{array} \right] \rightarrow$ Select largest $\rightarrow 4$

$\hookrightarrow \left[\begin{array}{c} 3, 1, 2, 4, 5 \end{array} \right]$ *replace* \rightarrow Select largest $\rightarrow 3$

$\hookrightarrow \left[\begin{array}{c} 2, 1, 3, 4, 5 \end{array} \right] \rightarrow$ Select largest $\rightarrow 2$

$\hookrightarrow \left[\begin{array}{c} 1, 2, 3, 4, 5 \end{array} \right] \rightarrow$ sorted



Selection Sort which technique follow
stable OR Unstable sort

$\rightarrow \left[\begin{array}{c} 3, 5, 2, 5', 1 \end{array} \right] \rightarrow$ Select largest $\rightarrow 5$

$\hookrightarrow \left[\begin{array}{c} 3, 1, 2, 5', 5 \end{array} \right] \rightarrow$ " $\rightarrow 5'$

$\hookrightarrow \left[\begin{array}{c} 3, 1, 2, 5', 5 \end{array} \right] \rightarrow$ " $\rightarrow 3$

$\hookrightarrow \left[\begin{array}{c} 2, 1, 3, 5', 5 \end{array} \right] \rightarrow$ " $\rightarrow 2$

$\hookrightarrow \left[\begin{array}{c} 1, 2, 3, 5', 5 \end{array} \right] \rightarrow$ Before sorting $\rightarrow 5$ Before $5'$
After sorting $\rightarrow 5'$ Before 5

Selection sorting follow Unstable sort

Change happen

Code for the selection sort

```
# Selection sorting

def selection_sort(arr):
    for i in range(len(arr)):
        min = i
        for j in range(i+1, len(arr)):
            if arr[j] < arr[min]:
                min = j

        # find the smallest value
        arr[i], arr[min] = arr[min], arr[i]
```

Selection sort Vs Bubble sort

```
# Selection sorting

def selection_sort(arr):
    for i in range(len(arr)):      ## iterating through range
        min = i
        for j in range(i+1, len(arr)):  ## finding the min values
            if arr[j] < arr[min]:
                min = j

        # find the smallest value
        arr[i], arr[min] = arr[min], arr[i]
```

V/S

```
: # bubble sorting

def bubbleSort(arr):
    for i in range(len(arr)-1, 0, -1):
        for j in range(i):
            if (arr[j]>arr[j+1]):
                arr[j], arr[j+1] = arr[j+1], arr[j]
```

1 $T(c) \Rightarrow O(n^2)$

$T(c) = O(n^2)$

2 No. of Comparision $\rightarrow \Delta n^2$

Comparision $\hookrightarrow \Delta n^2$

3 No. of swapping
 $\rightarrow n$ time

Swapping $\hookrightarrow \Delta n^2$

4 Stability \rightarrow Unstable

Stability \rightarrow Stable

selection sort better Performance than bubble sort because,
selection sort having less swapping Compare to
bubble sort.

↪ In real world, (Practically) ↪

Bubble, Selection Sort → ① Very expensive
Selection → ④ Slow

Use → Fast → like → ① Merge sort
② Quick sort

Insertion Sorting/ Online Sorting

Whenever you have incoming stream of data and regularly data coming in and you want to keep that data in sorted order. Here **Insertion sort** becomes very handy.

For ex,

zoom chat ↗

arrange the chat → Based on students name ↘

Algorithm that being used
insertion sort ↙

definition:- Insertion Sort builds the sorted part of the list by repeatedly picking elements from the unsorted part and inserting them in the correct position within the sorted part.

It is much less efficient on large list than more advanced algorithm such as quicksort, heapsort, mergesort. However it provide several advantages like :-

①. Simple implementation

②. Efficient for small data sets or nearly sorted.

ex,

$$\{6, 7, 2, 6, 9\}$$

Sorted Unsorted

→ Try to find any No in the left who is bigger than $\boxed{7}$

$$\{6, 7, \cancel{2}, 6, 9\}$$

\cancel{S} Us

Remain same \leftarrow NO \leftarrow in our cond:

$$\{2, \cancel{6}, 7, \cancel{6}, 9\}$$

\cancel{S} Us

$$\{6, 2, 7, 6, 9\}$$

Again is there any bigger than

$\cancel{2}$ in sorted. → Yes

Check who is bigger than $\cancel{6}$ in sorted

$$\{2, \cancel{6}, 6, \cancel{7}, 9\}$$

\cancel{S} Us

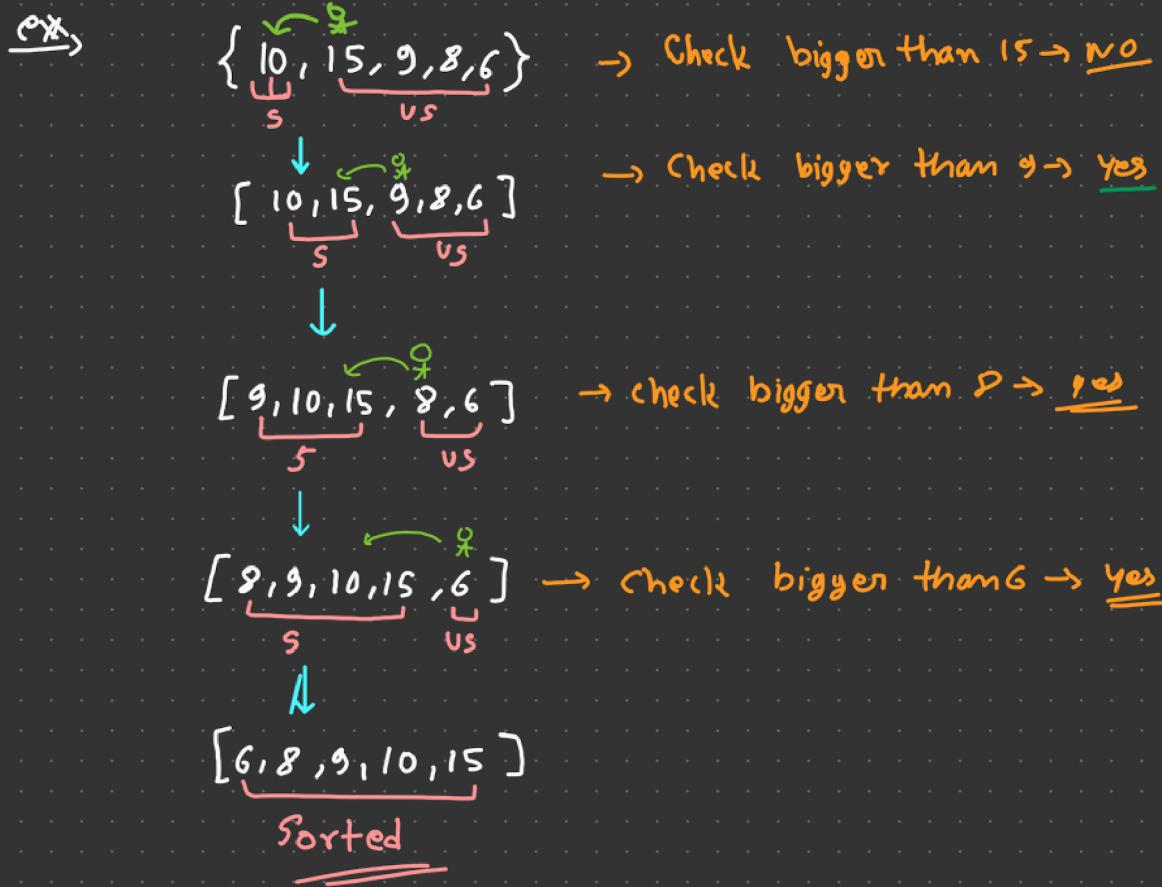
Check who is bigger than $\cancel{9}$ in sorted → No

$$\{2, 6, 6, 7, 9\}$$

Sorted

Here we keep trying to inserting the element so that the array get sorted in specific order

↪ Insertion / online sorting follow \Rightarrow Stable sorting



Code for insertion sort :

3. Online Sorting / insertion sort

```
def insertion_sort(arr):
    for i in range(1, len(arr)):
        v = arr[i]
        j = i
        while (j >= 1 and arr[j-1] > v):
            arr[j] = arr[j-1]
            j -= 1
        arr[j] = v
```