XMonad KeyMap QuasiQuoter

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Abstract

Play around with the idea of using a QuasiQuoter for key bindings in xmonad configurations. Currently, it is extremely limited (it can only bind keys to spawn actions), has some very overcomplicated nonsense, and is not tested well.

Source files are literate Haskell primarily so I can keep notes of what all this junk is actually supposed to be doing and why it's doing it that way.

Not actually intended for practical use. Maybe after lots of expansion, simplification, and testing, but certainly not now.

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Part I

Internals

1 Types

Custom data types and structures used throughout. TODO: can these be split into two seperate modules?

```
\begin{tabular}{ll} \textbf{module } \textit{Types where} \\ \textbf{import } \textit{Data.Bifoldable} \\ \textbf{import } \textit{Data.List.NonEmpty } (\textit{NonEmpty } (..), (< |), \textit{nonEmpty}) \\ \textbf{import } \textit{Data.Map} & (\textit{Map, elemAt, foldMapWithKey, mapKeysMonotonic, singleton, unionWithMapWithKey, mapKeysMonotonic, singleton, unionWithMapWithKey, mapKeysMonotonic, singleton, unionWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWithMapWit
```

1.1 Key actions

Data type representing actions so that they can be lifted to Template Haskell experessions later.

```
data KeyAction = ActionSpawn String
deriving (Eq, Show)
```

1.2 Submaps

A nice feature of XMonad.Util.EZConfig.additionalKeysP is that it allows for the easy binding of sequences of keys, via some "sugar" on top of XMonad.Actions.Submap.submap. I want to recreate this feature here.

The current solution is way too complicated and needs to be simplified.

A key may be bound directly to an action or be bound to a submap. This choice is represented by the ESM type. SM wraps a plain Map's values with ESM.

1.2.1 *ESM*

Either a submap (LSM) or just a value (RSM). TODO: could a recursive type be used instead of this and a **newtype**?

```
data ESM \ k \ v = LSM \ (SM \ k \ v) \mid RSM \ v \ deriving \ (Eq. Show)
```

1.2.2 *SM*

The SM type wraps the keys of a Map with ESM.

```
newtype SM \ k \ v = SM \ \{unSM :: Map \ k \ (ESM \ k \ v)\} deriving (Eq, Show)
```

1.2.3 Instances

Semigroup and monoid instances inherit an Ord constraint from Map.

```
instance Ord \ k \Rightarrow Semigroup \ (ESM \ k \ v) where LSM \ x \diamondsuit LSM \ y = LSM \ (x \diamondsuit y) instance Ord \ k \Rightarrow Semigroup \ (SM \ k \ v) where SM \ m \diamondsuit SM \ n = SM \ \$ \ unionWith \ (\diamondsuit) \ m \ n instance Ord \ k \Rightarrow Monoid \ (SM \ k \ v) where mempty = SM \ mempty
```

Folding, mapping, and traversing ESMs recurses on left values and uses rights directly. SM instances just have to apply functions recursively to their internal Maps and ESMs.

```
instance Foldable
                     (ESM \ k) where foldMap \ f \ (RSM \ x) =
                                                                       f x
                                     foldMap \ f \ (LSM \ sm) =
                                                                       foldMap \ f \ sm
                                             f(RSM x) = RSM \$ f x
                     (ESM \ k) where fmap
instance Functor
                                     fmap
                                              f(LSM \ sm) = LSM \ \$ \ fmap \ f \ sm
instance Traversable (ESM k) where traverse f(RSM x) = RSM \ll f x
                                     traverse\ f\ (LSM\ sm) = LSM\ <\$>\ traverse\ f\ sm
instance Foldable
                     (SM \ k) where foldMap \ f \ (SM \ m) =
                                                                 foldMap \ (foldMap \ f) \ m
                                          f(SM m) = SM \$ fmap
instance Functor
                     (SM \ k) where fmap
                                                                          (fmap)
instance Traversable (SM k) where traverse f (SM m) = SM <$> traverse (traverse f) m
```

1.2.4 Probably superfluous instances

I don't think any of these are used, will be ever used, or should ever be used, but they are defined here anyway for... fun.

ESM can be lifted to an applicative. When combining two maps, each of the right values of the "argument" map are passed to each of the right values of the "function" map. A similar method is used to lift ESM to a monad.

```
instance Applicative (ESM k) where

pure = RSM
RSM f <*>RSM x = RSM (f x)
RSM f <*>LSM sm = LSM (fmap f sm)
LSM sm <*>RSM x = LSM (fmap ($x) sm)
fs <*>LSM (SM m) = LSM (SM (fmap (fs<*>) m))
instance Monad (ESM k) where
RSM x \gg f = f x
LSM (SM m) \gg f = LSM (SM (fmap ($\approx f) m))
```

 $Data.Map.foldMap\,WithKeys$ can easily be reworked into bifoldMap and a Bifoldable instance.

```
instance Bifoldable\ SM where bifoldMap\ f\ g\ (SM\ m) = foldMap\ WithKey\ (\lambda k\ v \to f\ k \diamondsuit bifoldMap\ f\ g\ v)\ m instance Bifoldable\ ESM where bifoldMap\ g\ (RSM\ x)\ = g\ x bifoldMap\ f\ g\ (LSM\ sm)\ = bifoldMap\ f\ g\ sm
```

In theory, a combination of mapKeys and plain fmap theoretically can be composed into bimap. However, mapKeys has an Ord constraint and will not work with Bifunctor.

Bitraversing is more insteresting. There is no function provided by *Data.Map* that allows for traversing over keys. Defining one, however, would not be too complicated. Two possibilities are:

Reconstructing the *Map* within the applicative. A function declared like this would work similarly to *mapKeysMonotonic*—and have the same restriction that function being applied must be monotonic, so its no good.

```
\begin{array}{ll} \textbf{import } \textit{Data.Map.Internal } (\textit{Map } (..)) \\ \textit{bitraverseMonotonic} :: \textit{Applicative } f \Rightarrow (k \rightarrow f \ g) \rightarrow (v \rightarrow f \ w) \rightarrow \textit{Map } k \ v \rightarrow f \ (\textit{Map } g \ w) \\ \textit{bitraverseMonotonic} \ \_\_\textit{Tip} &= \textit{pure } \textit{Tip} \\ \textit{bitraverseMonotonic } f \ g \ (\textit{Bin } s \ k \ v \ l \ r) = \textit{Bin} < pure \ s < f \ k < g \ v \\ < *> \textit{bitraverseMonotonic } f \ g \ l \\ < *> \textit{bitraverseMonotonic } f \ g \ r \\ \end{array}
```

Mapping each element into a singleton and folding results. Using foldMapWithKey and the Ap data type (from Data.Monoid), this task is trivial. Sequentially applying singleton to the new lifted keys, then wrapping this new

Map in Ap produces a monoid that is easily folded into a single Map, as long as Map itself is a monoid—which it only is if its keys are Orderable, thus having the same constraint as the previous method and being equally useless for this purpose.

```
\begin{array}{l} \textbf{import } \textit{Data.Monoid } (\textit{Ap } (..)) \\ \textit{traverseKeysMonoid } :: (\textit{Applicative } f, \textit{Ord } g) \Rightarrow (k \rightarrow f \ g) \rightarrow \textit{Map } k \ (\textit{ESM } k \ a) \rightarrow f \ (\textit{Map } g \ (\textit{ESM } k \ a)) \\ \textit{traverseKeysMonoid } f = \textit{getAp} \circ \textit{foldMapWithKey } \textit{go} \\ \textbf{where } \textit{go } k \ v = \textit{Ap } (\textit{singleton } <\$>f \ k <*>\textit{recurse } v) \\ \textit{recurse } (\textit{RSM } x) = \textit{RSM } <\$>x \\ \textit{recurse } (\textit{LSM } \textit{sm}) = \textit{LSM } <\$>\textit{traverseKeysMonoid } f \ \textit{sm} \\ \end{array}
```

1.2.5 Other Map functions for SM

mapKeysMonotonic; an unsafe function that applies a given function over the keys of a submap. It does not check if the ordering of the keys has changed, so if they do change, the Map will be broken, hence it being unsafe.

```
smMapKeysMonotonic :: (k 	o g) 	o SM \ k \ a 	o SM \ g \ a smMapKeysMonotonic \ f \ (SM \ m) = SM \ \$ \ fmap \ recurse \ \$ \ mapKeysMonotonic \ f \ m where recurse \ (RSM \ x) = RSM \ x recurse \ (LSM \ x) = LSM \ \$ \ smMapKeysMonotonic \ f \ x
```

1.2.6 Creating and eliminating submaps

A "singleton" submap can be created by providing a "path" and its value, the path being a non-empty list of keys to the value.

```
toSM :: NonEmpty \ k \to t \to SM \ k \ t

toSM \ (k : | ks) = SM \circ singleton \ k \circ maybe \ pure \ (\lambda p \to LSM \circ toSM \ p) \ (nonEmpty \ ks)
```

The inverse of the above, a path and its value can be extracted from a non-empty submap.

```
\begin{array}{lll} fromSM :: SM \ k \ t \rightarrow Maybe \ (NonEmpty \ k, t) \\ fromSM \ (SM \ m) \ | \ null \ m &= Nothing \\ & \mid (k, RSM \ x) &\leftarrow head' \ m &= Just \ (k : \mid [\,], x) \\ & \mid (k, LSM \ sm) \leftarrow head' \ m, \\ & \quad Just \ (ks, x) &\leftarrow fromSM \ sm = Just \ (k < \mid ks, x) \\ & \mid otherwise &= Nothing \\ & \mathbf{where} \ head' = elemAt \ 0 \end{array}
```

2 Parsing

Parsers used for producing data structres from input to the QuasiQuoter. The syntax here is more similar to the ones found in vis or If than additionalKeysP. TODO: better error messages for just about everything here.

```
module Parsers where
import Control.Monad
{\bf import}\ Control. Applicative
import Text.ParserCombinators.Parsec hiding ((<math><|>), many, optional)
                Data.List
import
import
                Data.List.NonEmpty (NonEmpty, some1)
import qualified Data.List.NonEmpty as NE
import
                Data.Set
                                     (Set)
import qualified Data.Set
                                     as S
import XMonad
import Types
```

2.1 Utilities

```
\begin{array}{ll} someOf &= some & \circ \ oneOf \\ tryMany &= many & \circ \ try \\ trySome &= some1 \circ try \\ testAhead &= \lambda p \rightarrow lookAhead \ (\textit{True} < p) < |> pure \ \textit{False} \end{array}
```

2.2 Parsing regular keys

Some keys are allowed to specified literally, i.e., by the actual character instead of the symbolic name.

```
allowedLiterally = [,!,.,,,] \setminus "<->"
```

Parser for a single, non-whitespace, ASCII character.

```
 \begin{array}{l} asciiKeysym :: CharParser \ st \ KeySym \\ asciiKeysym = choice \ xs <?> \ "an \ ASCII \ character" \\ \mathbf{where} \ xs = zipWith \ (\lambda c \ s \rightarrow s <\$ \ char \ c) \ [ '!' . . , `~' ] \ [xK\_exclam . . xK\_asciitilde] \\ \end{array}
```

Parser for known symbolic key name.

```
keysymName :: CharParser st KeySym
keysymName = \mathbf{do}
str \leftarrow someOf \ allowedLiterally
\mathbf{let} \ sym = stringToKeysym \ str
```

```
if sym \equiv noSymbol then fail \$ "Invalid key symbol name \"" + str ++ "\"" else return \ sym
```

Parser that tries to read a symbolic name, or ASCII character if that fails.

```
keysym' :: CharParser st KeySym
keysym' = try keysymName <|> asciiKeysym
```

Same as keysym' except the names are surrounded with angled brackets.

TODO: this is Special pattern is ugly.

```
keysym :: CharParser st KeySym
keysym = do
isSpecial ← testAhead (char '<')
if isSpecial then char '<' *> keysymName <* char '>'
else asciiKeysym
```

2.3 Parsing modifier keys

Parser that reads an abbreviation for a modifier key and returns its mask.

```
\begin{aligned} keymask' &:: CharParser\ st\ KeyMask\\ keymask' &= choice\ xs\\ &\text{ where } xs = zip\ With\ (\lambda s\ m \to m < \ try\ (string\ s))\\ & \left[\text{"S", "C", "M1", "M2", "M4", "M5"}\right]\\ & \left[shiftMask,\ controlMask,\ mod1Mask,\ mod2Mask,\ mod3Mask,\ mod4Mask,\ mod5Mask]\right] \end{aligned}
```

Same as keymask' but also accepts "M" for modMask, represented as Nothing.

```
keymask :: CharParser st (Maybe KeyMask) \\ keymask = Just <\$> keymask' <|> Nothing <\$ char 'M'
```

2.4 Parsing key combinations

Parser that either reads a single asciiKeysym, or many (or no) hypen-seperated keymasks and a keysym' between angled brackets.

```
keyCombo :: CharParser \ st \ (Set \ (Maybe \ KeyMask), KeySym)
keyCombo = \mathbf{do}
isSpecial \leftarrow testAhead \ (char \ '<')
\mathbf{if} \ isSpecial \ \mathbf{then} \ char \ '<' \ *> pure \ withMasks <*> tryMany \ (keymask <* char \ '-') <*> keysym' <* char \ 'else \ noMasks <$> asciiKeysym
\mathbf{where} \ noMasks = \lambda s \ \rightarrow (mempty, \ s)
withMasks = \lambda m \ s \rightarrow (S.fromList \ m, s)
```

2.5 Parsing actions

Parser that reads a spawn action.

```
keyActionSpawn :: CharParser \ st \ KeyAction \\ keyActionSpawn = char \ `\$' \ *> \\ (ActionSpawn <\$> ((:) <\$> (noneOf "|" <?> "a command") \\ <*> (manyTill \ anyChar \$ \ lookAhead \$ \ void \ (char \ '|') <|> \ eof)))
```

Parser that reads any action.

```
keyAction :: CharParser st KeyAction \\ keyAction = keyActionSpawn
```

2.6 Parsing key bindings

Parser that reads a sequence space-sperated of key combinations followed by an action.

```
keybind :: CharParser \ st \ (SM \ (Set \ (Maybe \ KeyMask), KeySym) \ KeyAction) keybind = toSM < sym(keyCombo < some \ space) < keyAction
```

Parser that reads a list of *keybinds*, separated by pipes (i.e., |).

```
keybinds :: CharParser st (SM (Set (Maybe KeyMask), KeySym) KeyAction)
keybinds = spaces *> pure mconcat <*> keybind 'sepBy1' (char '| '*> spaces) <* spaces
```

2.7 Putting it all together

Wraps *keybinds* with *parse*. Intended to the be main export to other modules (so they do not also need to import Parsec).

```
parseKeymap :: String \rightarrow Either \ ParseError \ (SM \ (Set \ (Maybe \ KeyMask), KeySym) \ KeyAction)
parseKeymap = parse \ keybinds ""
```

3 Lifting

This module defines functions for lifting values into Template Haskell typed expressions.

```
{-# LANGUAGE TemplateHaskell #-} module Lift where
```

The "goal" function is one that works on the output of the parser and returns an expression representing a lambda taking an XConfig and returning the key map; with a type signature

```
goal :: SM \quad (Set (Maybe KeyMask), KeySym) KeyAction 
 <math>\rightarrow TExpQ (XConfig Layout \rightarrow Map (KeyMask, KeySym) (X ()))
```

This function will be broken up into smaller "steps":

- 1. Lift the key actions to expressions representing X actions.
- 2. Lift the key combinations to expressions representing pairs of masks and symbols.
- 3. Lift the ESMs to expressions representing single X actions, taking right values directly and merging left ones with submap.
- 4. Lift the SMs to expressions representing plain Maps.
- 5. Finally, splice resulting expression into one with the XConfig lambda.

3.1 Imports

For the Map and Set data structures:

```
importData.Map(Map)importqualifiedData.MapasimportData.Set(Set)importqualifiedData.Setas
```

For Template Haskell types and functions:

```
import Language. Haskell. TH. Lib
import Language. Haskell. TH. Syntax
```

For working with KeyMasks and KeySyms:

```
import XMonad (KeyMask, KeySym, (.|.))
import Foreign. C. Types (CUInt)
```

For working with XConfig:

```
import XMonad (Layout, XConfig (XConfig, modMask))
```

For working with KeyActions:

```
import XMonad (X, spawn) import XMonad.Actions.Submap (submap)
```

Finally, custom types:

```
\mathbf{import}\ \mathit{Types}
```

3.2 Lifting key actions

```
liftKeyAction :: KeyAction \rightarrow TExpQ(X())
liftKeyAction (ActionSpawn x) = [\lor spawn x \lor]
```

3.3 Lifting key combinations

Before key masks can be lifted, a Lift instance for its "root" type, CUInt , must be declared:

```
instance Lift CUInt where lift = litE \circ integerL \circ toInteger
```

Compared to the rest, this function takes an extra argument, varName :: Name. This represents the value of modMask.

```
\begin{array}{l} \textit{liftKeyCombo} :: Name \rightarrow (Set \ (Maybe \ KeyMask), KeySym) \rightarrow TExpQ \ (KeyMask, KeySym) \\ \textit{liftKeyCombo} \ varName \ (set, sym) = \\ \textbf{let} \ knownMasks \ = foldr \ (\lambda mask \ rest \rightarrow maybe \ rest \ (.|.rest) \ mask) \ 0 \ set \\ varExpression = TExp < > varE \ varName \\ \textbf{in} \ \ \textbf{if} \ Nothing \in set \\ \textbf{then} \ [\lor (knownMasks \ .|. \$\$(varExpression), sym) \lor] \\ \textbf{else} \ [\lor (knownMasks, sym) \ \lor] \\ \end{array}
```

3.4 Lifting ESMs

Right values can be returned as-is, but left values require more work:

- 1. Recurse on child maps.
- 2. Convert the map to a (sorted) list of pairs.
- 3. Lift each pair in the list.
- 4. Lift the whole list.
- 5. Splice the list's expression into one turning it back into a map.
- 6. Splice the map's expression into one calling *submap* on it.

```
\begin{array}{lll} \mathit{liftESM} :: \mathit{ESM} \; (\mathit{TExpQ} \; (\mathit{KeyMask}, \mathit{KeySym})) \; (\mathit{TExpQ} \; (X \; ())) \to \mathit{TExpQ} \; (X \; ()) \\ \mathit{liftESM} \; (\mathit{RSM} \; x) & = x \\ \mathit{liftESM} \; (\mathit{LSM} \; (\mathit{SM} \; \mathit{rawMap})) = \\ \mathbf{let} \; \mathit{rightMap} & = \; fmap \; \mathit{liftESM} \; \mathit{rawMap} \\ \mathit{rightList} & = \; M.toAscList \; \mathit{rightMap} \\ \mathit{rightList}' & = \; fmap \; tupToExp \; \mathit{rightList} \\ \mathbf{where} \; \mathit{tupToExp} \; (\mathit{comboExp}, \mathit{actionExp}) = \\ \end{array}
```

```
[\lor (\$\$(comboExp), \$\$(actionExp)) \lor] rightExpression = fmap \ TExp \ \$ \ listE \ \$ \ map \ (fmap \ unType) \ rightList' mapExpression = [\lor M.fromDistinctAscList \ \$\$ \ (rightExpression) \lor] \textbf{in} \ [\lor submap \ \$\$ \ (mapExpression) \lor]
```

3.5 Lifting SMs

- 1. Call liftESM on child maps.
- 2. Convert the map to a (sorted) list of pairs.
- 3. Lift each pair in the list.
- 4. Lift the whole list.
- 5. Splice the list's expression into one turning it back into a map.

TODO: switch name with liftSM? feels like I've been using "submap" to refer more general functions, unlike this.

```
liftSubmaps :: SM
                        (TExpQ\ (KeyMask, KeySym))\ (TExpQ\ (X\ ()))
            \rightarrow TExpQ (Map
                                 (KeyMask, KeySym) (X ()))
liftSubmaps\ (SM\ rawMap) =
  let flatMap
                      = fmap \ liftESM \ rawMap
     flatList
                      = M.toAscList flatMap
     flatList'
                      = fmap \quad tup ToExp \ flatList
                         where tup ToExp (comboExp, actionExp) =
                                   [\lor (\$\$(comboExp), \$\$(actionExp)) \lor]
     listExpression = fmap \ TExp \ \ listE \ \ \ map \ (fmap \ unType) \ flatList'
     mapExpression = [\lor M.fromDistinctAscList \$\$ (listExpression) \lor]
 in mapExpression
```

3.6 Putting it all together

Primary export of this module.

```
liftKeymap :: SM \ (Set \ (Maybe \ KeyMask), KeySym) \ KeyAction \rightarrow ExpQ \ liftKeymap = fmap \ unType \circ liftSM
```

Part II

Exports

4 QuasiQuoters

Module that defines the QuasiQuoter.

```
module QQ where import Language.Haskell.TH.Quote import Parsers import Lift
```

Quotes a keymap from a string or throws an error if it can't.

```
quoteKeymap = either (fail \circ show) \ liftKeymap \circ parseKeymap
```

The QuasiQuoter. It does not support quoting anything other than expressions.

```
\label{eq:keymap} keymap :: QuasiQuoter \\ keymap = QuasiQuoter \\ \{quoteExp = quoteKeymap \\ , quotePat = fail \text{"keymap can only quote expressions."} \\ , quoteDec = fail \text{"keymap can only quote expressions."} \\ , quoteType = fail \text{"keymap can only quote expressions."} \\ \}
```

5 Utilities

Utility functions to be exported by the main module alongside the quasiquoter.

```
module Utilities where
import Data.Map (Map)
import XMonad
```

Roughly equivalent to XMonad. Util. EZConfig. additional Keys except that the argument here is the same type as keys from XConfig.

```
 \begin{array}{ll} additional Keys :: & XConfig \ l \\ & \rightarrow (XConfig \ Layout \rightarrow Map \ (KeyMask, KeySym) \ (X \ ())) \\ & \rightarrow & XConfig \ l \\ additional Keys \ xc \ newKeys = xc \ \{keys = newKeys < + > keys \ xc \} \end{array}
```

The same function, with the arguments flipped. Fits better with functions like with Urgency Hook, etc.

```
withAdditionalKeys = flip \ additionalKeys
```