

- Georgiy Shevoroshkin
- uml diagramm (konzeptionell)
 - assoziationen, bedingungen
- transaktionen
 - isolation levels erklären, welche fehler sie beheben
 - fuzzy read, deadlock, dirty read, write skew, phantom read, serializable snapshot isolation, cascading rollbacks
 - schedule analysieren + Serialisierbarkeitsgraph
- begriffe (physisches schema, DBMS)
- b-baum indexe einfügen

DataBase System

Besteht aus Datenbankmanagementsystem und Datenbasen

DataBase Management System

- Redundanzfreiheit
- Datenintegrität
- Kapselung

ANSI Modell

Logische Ebene:

Interne Ebene:

Externe Ebene:

Mapping:

Unified Modeling Language

- Assoziation ◆ Komposition
- ◊ Aggregation ▷ Vererbung

Complete:

Incomplete:

Disjoint:

Overlapping:

Normalisierung

1NF: Atomare Attributwerte

2NF: Nichtschlüsselattr. voll vom Schlüssel abhängig

3NF: Keine transitiven Abhängigkeiten

BCNF: Nur abhängigkeiten vom Schlüssel

(Voll-)funktionale Abhängigkeit:

Transitive Abhängigkeit:

Einfügeanomalie, Löschanomalie, Änderungsanomalie

Vererbung

TODO: sql beispiele Einzige Tabelle für Superklasse:

Tabelle pro Subklasse:

Tabelle pro Sub- und Superklasse:

Data Definition Language

```
CREATE SCHEMA s;
CREATE TABLE t (
  id SERIAL PRIMARY KEY,
  name TEXT UNIQUE,
  grade DECIMAL(2,1) NOT NULL,
  added TIMESTAMP DEFAULT CURRENT_TIMESTAMP,
  u VARCHAR(9) DEFAULT CURRENT_USER,
  fk INT FOREIGN KEY REFERENCES t2.id ON DELETE
  CASCADE,
  CHECK (grade between 1 and 6)
);
```

```
ALTER TABLE t2 ADD CONSTRAINT c PRIMARY KEY (a, b);
TRUNCATE/DROP TABLE t;
```

Data Manipulation Language

```
INSERT INTO t (added, grade) VALUES ('2002-10-10', 1)
RETURNING id;
```

Views

```
CREATE VIEW v (id, grade, u) AS SELECT id, grade, u
FROM t;
```

Data Control Language

```
CREATE ROLE r WITH LOGIN PASSWORD '';
GRANT INSERT ON TABLE t TO r;
REVOKE CREATE ON SCHEMA s FROM r;
ALTER ROLE r CREATOROLE, CREATEDB, INHERIT;
GRANT r TO user_name;
-- read all future created tables
ALTER DEFAULT PRIVILEGES IN SCHEMA s GRANT SELECT ON
TABLES TO readonlyuser;
CREATE POLICY p ON t FOR ALL TO PUBLIC USING (u =
current_user);
ALTER TABLE t ENABLE ROW LEVEL SECURITY;
```

Common Table Expressions

```
WITH RECURSIVE q AS (SELECT * FROM t WHERE grade>1
UNION ALL SELECT * FROM t INNER JOIN q ON q.u =
t.name) SELECT id as "ID" FROM q;
```

Window Functions

```
SELECT id, RANK() OVER (ORDER BY grade DESC) as r FROM
t;
```

Subqueries

```
SELECT * FROM t WHERE grade > ANY (SELECT g FROM t2);
SELECT * FROM t WHERE EXISTS (SELECT g FROM t2);
```

JOIN

```
SELECT y.*, x.* FROM t AS y, JOIN LATERAL (SELECT *
FROM t2 WHERE t2.id = y.id) AS x;
```

GROUP BY

```
SELECT id, COUNT(*) FROM t GROUP BY grade, id HAVING
COUNT(*) > 2;
```

WHERE

```
BETWEEN 1 AND 5; LIKE '___%'
IN (1, 5) ; LIKE '%asd'
```

INDEX

```
CREATE INDEX i ON t /*USING BTREE*/ (grade, UPPER(u))
INCLUDE (added);
CREATE INDEX i ON t (grade) WHERE grade > 4;
DROP INDEX i;
```

Transaktionen

```
BEGIN; SAVEPOINT s;
COMMIT; ROLLBACK /*TO SAVEPOINT s*/;
```

Isolation

```
SET TRANSACTION ISOLATION LEVEL ...;
SET SESSION CHARACTERISTICS AS TRANSACTION ISOLATION
LEVEL ...;
```

TODO: check

	Read mitted	Uncom- mitted	Read ted	Commit- ted	Repeatable Read
Dirty Write	*		*		*
Dirty Read	✓		X		X
Lost Update	✓		✓		X
Fuzzy Read	✓		✓		X
Phantom Read	✓		✓		✓
Read Skew	✓		✓		X
Write Skew	✓		✓		✓

Dirty Read: Lese Daten von nicht committed T's

Fuzzy Read: Versch. Werte beim mehrmaligen Lesen gleicher Daten (da durch andere T geändert)

Phantom Read: Neue/Gelöschte Rows einer anderen T

Deadlock:

Cascading Rollback:

Relationale Algebra

```
πR1,R4(R) SELECT R1,R4 FROM R;
σR1>30(R) SELECT * FROM R WHERE R1 > 30;
ρa ← R SELECT * FROM R AS a;
R ⋈ S SELECT * FROM R,S;
R ⋈A=B S SELECT * FROM R JOIN S ON R.A=S.B;
```

Serialisierbarkeit

Backup

Write-Ahead Log

TODO: LSN, TaID, PageID, Redo, Undo, PrevLSN

Dreiwertige Logik (cursed)

```
SELECT NULL IS NULL; -- true
SELECT NULL = NULL; -- [null]
```