

# Computer Networks 1 | CN1

## Summary

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## 1. APPLICATION LAYER (7,6,5)

Combines Layers 7 (Application), 6 (Presentation) and 5 (Session).

### 1.1. COMMON PORTS

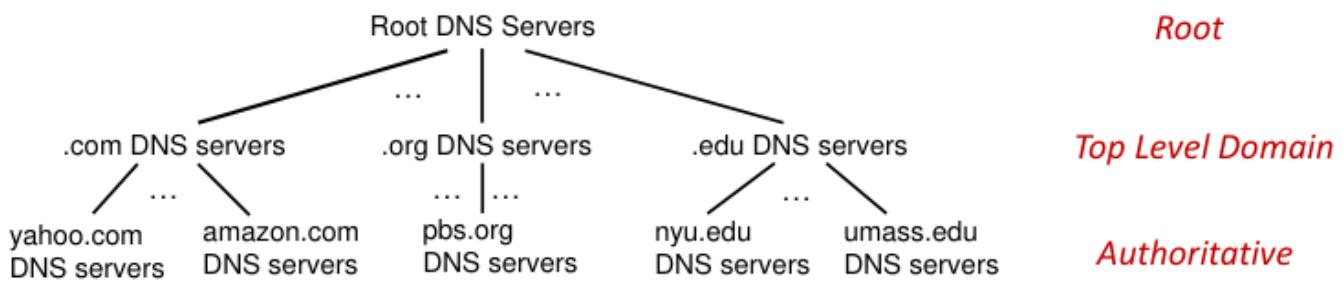
<i>Protocol</i>	<i>Port</i>	<i>Layer 4</i>
DNS	53	UDP, TCP
HTTP	80	TCP
HTTPS	443	TCP
FTP	20, 21	TCP
SMTP	25 (server) 587 (client)	TCP
POP3	110	TCP
DHCP	67 (server) 68 (client)	UDP

## 1.2. HTTP

<b>Feature</b>	<b>HTTP/1.0</b>	<b>HTTP/1.1</b>	<b>HTTP/2</b>	<b>HTTP/3</b>
<b>Connection Management</b>	One request per connection	Persistent connections by default	Multiplexing allows multiple streams	Uses QUIC for multiplexing
<b>Request Methods</b>	Limited (GET, POST, HEAD)	Enhanced (PUT, DELETE, OPTIONS, etc.)	Same as 1.1	Same as 1.1
<b>Caching</b>	Basic caching support	Improved caching with validation	Advanced caching capabilities	Same as 2 but with improved mechanisms
<b>Header Compression</b>	None	None	HPACK (header compression)	QPACK (header compression)
<b>Server Push</b>	Not supported	Not supported	Supported (automatic resource pushing)	Enhanced support for server push
<b>Performance Improvements</b>	None	Minor improvements over 1.0	Significant improvements in performance and latency	Further improvements in speed and efficiency
<b>SSL/TLS Support</b>	Not inherent	Not inherent, but commonly supported	Built-in support with ALPN (Application-Layer Protocol Negotiation)	Uses QUIC, which incorporates TLS 1.3
<b>Transport Protocol</b>	TCP	TCP	TCP	QUIC

### 1.3. DNS

Nameservers resolve domains to IP's through a distributed, hierarchical database.



Term	Definition
Iterated query	Local DNS server iteratively asks one server after the other, descending the domain name hierarchy step after step.
Recursive query	Local DNS server asks root server for domain, which in turn asks the TLD server, which in turn asks the authoritative server etc. until the “call stack” unwinds and returns the fully resolved domain to the query sender.
Caching	

#### 1.3.1. Record types

Term	Definition
A	<b>name:</b> hostname <b>value:</b> IPv4 address
AAAA	<b>name:</b> hostname <b>value:</b> IPv6 address
CNAME	<b>name:</b> alias <b>value:</b> canonical name
NS	<b>name:</b> domain <b>value:</b> hostname of authoritative NS for this domain
MX	<b>name:</b> domain <b>value:</b> name of mailserver

### 1.4. E-MAIL

Term	Definition
ding	
dong	
your	
opinion	
is	
wrong	

## 2. TRANSPORT LAYER (4)

Segment size: 1440-1480b when using IPv4, <=1460b when using IPv6

### 2.1. PRIMARY RESPONSIBILITIES

- Process-to-process delivery (distinguish between multiple applications via ports)
- Ensure reliable transfer (acknowledgments, retransmissions & reordering)
- Flow control (sender does not overwhelm receiver)
- Congestion control (network is not overloaded)

<i>Term</i>	<i>Definition</i>
Port	<b>16 bit long</b> numbers (0d0-0d65'535) for identifying applications to send packets to. <b>Well-Known:</b> 0d0-0d1'023 for universal TCP/IP applications, managed by the IANA. <b>Registered:</b> 0d1'024-0d49'151 for known applications, also managed by the IANA. <b>Private:</b> 0d49'152-0d65'535 for custom applications, not managed by the IANA.
Socket	Combination of <i>IP:Port</i> .
Multiplexing	Sending data from multiple sockets at sender.
Demultiplexing	Delivering segments to correct socket at receiver.
Checksum	Detect errors (i.e., flipped bits) in transmitted segment.

### 2.2. TCP

Connection-oriented, bidirectional, reliable, managed data flow.

<i>Term</i>	<i>Definition</i>
Handshake	Agreement on <b>starting sequence numbers, maximum segment size</b> and <b>window scaling</b> . 1) SEQ 2) SEQ+ACK 3) ACK
FIN	Termination of a connection. 1) FIN 2) FIN+ACK 3) ACK
Round Trip Time	<b>RTT</b> is the time it takes for a packet to be sent to the receiver and acknowledged back to the sender.
Buffer size	Maximum amount of data (measured in bytes) that can be stored in memory while waiting to be processed or transmitted.
Maximum Segment Size	<b>MSS</b> is the maximum payload size of a TCP packet. In IPv4 networks, typically, the size of the MSS is <b>1460 bytes</b> because it is encapsulated in the data link layer Ethernet frame size of <b>1500 bytes</b> .

#### 2.2.1. Reliability

<i>Term</i>	<i>Definition</i>
Sequence numbers	<b>SEQ</b> ensures that the packets arrive or can be reassembled in order.
Acknowledgement	<b>ACK</b> ensures that the receiver gets all of the packets.

Term	Definition
Retransmission timeout	If an acknowledgment is not received before the timer for a segment expires, a retransmission timeout occurs, and the segment is <b>automatically retransmitted</b> .
Packet loss rate	Measures how many packets of the ones being sent actually arrive.

### 2.2.2. Throughput

Term	Definition
Throughput	Denoted by $T$ , is the amount of data that can be transmitted during a specified time. $T = \frac{W}{R} \leq C_{L3}$
Continuous sending	Sender transmits a stream of data packets in the given window size <b>without waiting for acknowledgments</b> .
Delayed ACK	Receiver waits for a short period to acknowledge <b>multiple segments</b> with a <b>single ACK</b> .
Selective ACK	Instead of asking for a retransmission of all missing segments, <b>SACK</b> (specified by the receiver) allows the sender to send only the lost segments, significantly improving efficiency.

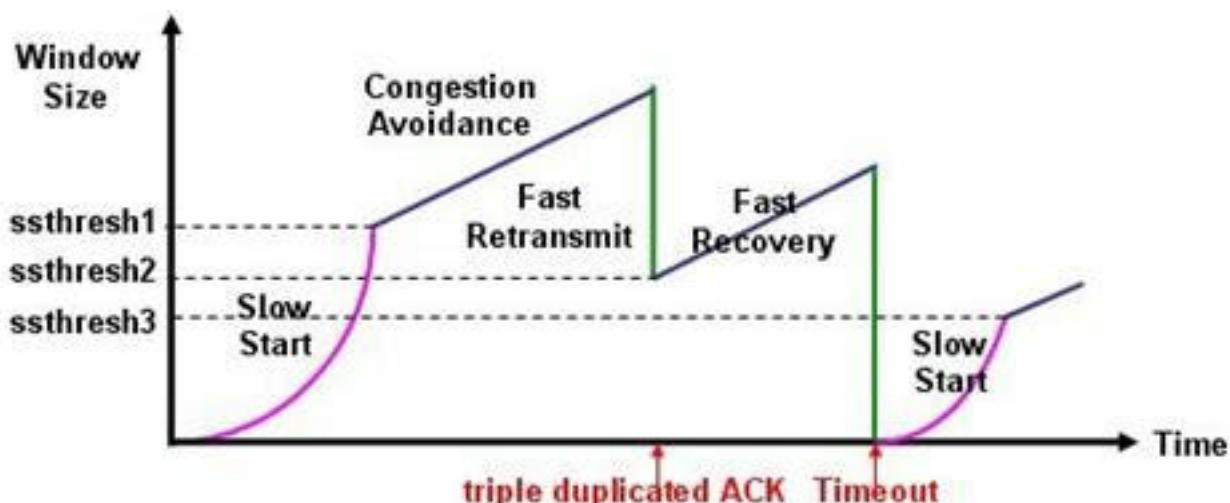
### 2.2.3. Flow control

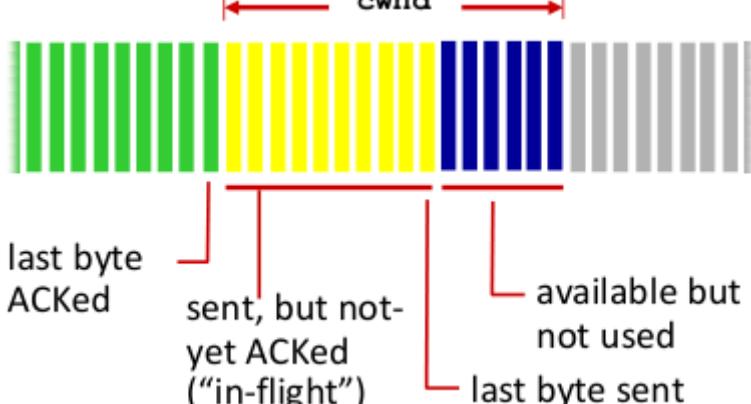
So that the sender does not overwhelm the receiver.

Term	Definition
Window Size	Denoted by $W$ , is a <b>16 bit</b> number sent with each packet by the receiver inside of the <b>rwnd</b> header field, indicating the amount of data he still has space for.
Window scale	Used when the TCP window size needs to be increased beyond the traditional maximum of 65,535 bytes due to the demands of high-speed networks. If the handshake header includes the <b>window scale option</b> and the packet header includes the <b>scaling factor</b> then the effective window size is calculated as such: <b>window size * scaling factor</b>

### 2.2.4. Congestion control

To prevent network congestion.



Term	Definition
Congestion window	 <p>Diagram illustrating the Congestion Window (cwnd) over 16 slots. The first 10 slots are green ('ACKed'), the next 5 are yellow ('sent, but not-yet ACKed ("in-flight")'), and the last 1 slot is blue ('available but not used'). A red double-headed arrow at the top spans the first 14 slots and is labeled 'cwnd'.</p>
Sliding Window	Describes the process of the congestion window sliding to the right after receiving ACKs.
Slow start	<p>Gradual growth (doubling <b>cwnd</b> every <b>RTT</b>) within the congestion window size at the start of a connection or after a period of state of no activity.</p> <p><b>Purpose:</b> Allows the sender to probe the available bandwidth in a controlled way.</p>
Congestion avoidance	<p>Transition from sluggish start to congestion avoidance segment after accomplishing a threshold.</p> <p><b>Purpose:</b> Maintains a truthful share of the community bandwidth even as heading off excessive congestion.</p>
Fast Retransmit	<p>Detects packet loss through duplicate acknowledgments and triggers speedy retransmission without waiting for the <b>retransmission timeout</b>.</p> <p><b>Purpose:</b> Speeds up the recuperation method with the aid of retransmitting lost packets without looking ahead to a timeout.</p>
Fast Recovery	<p>Enters a quick healing state after detecting packet loss, lowering congestion window and transitioning to congestion avoidance.</p> <p><b>Purpose:</b> Accelerates healing from congestion by way of avoiding a complete go back to slow begin after packet loss.</p>
AIMD	<p>Adjusts the congestion window size based on network situations following the <b>Additive Increase, Multiplicative Decrease</b> principle.</p> <p><b>Purpose:</b> Provides a balanced approach by way of linearly growing the window all through congestion avoidance and halving it on packet loss.</p>

## 2.3. UDP

## 2.4. QUIC

Actually a layer 7 Protocol, running on top of UDP

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## 3. NETWORK LAYER (3)

Packet size: **1500b**

### 3.1. SUBNETTING

Dividing a **/X** network into **n** amount of **/Y** subnets:  $2^{Y-X} = n$ .

Eg: Dividing a **/16** network into **/24** subnets will yield **256** subnets, because  $2^{24-16} = 2^8 = 256$

### 3.2. IPV6

#### 3.2.1. Glossary

Term	Definition
Extension Header	Additional headers used in IPv6 to provide optional information. These can define aspects like payload size, routing, or fragmentation.
DHCPv6	Dynamic Host Configuration Protocol for IPv6; this allows servers to assign IPv6 addresses dynamically from a pool, similar to DHCP for IPv4.
NAT64	Network Address Translation from IPv6 to IPv4 and vice versa; it facilitates communication between IPv6 and IPv4 networks.
Neighbor Discovery Protocol ( <b>NDP</b> )	A protocol in IPv6 for discovering other network nodes, determining their link-layer addresses, and ensuring that addresses are valid and reachable.
Internet Control Message Protocol ( <b>ICMPv6</b> )	A crucial part of IPv6 that handles error messages and operational queries, with an expanded role compared to ICMP in IPv4.
MTU	Maximum Transmission Unit; the size of the largest packet that can be sent in a single frame over a network medium. IPv6 can handle larger MTUs compared to IPv4.
Multicast Listener Discovery ( <b>MLD</b> )	IPv6 multicast routers can use MLD to discover multicast listeners on a directly attached link.
Path MTU Discovery ( <b>PMTUD</b> )	Protocol for determining the Maximum Transmission Unit (MTU) size on the network path between two hosts, usually with the goal of avoiding IP fragmentation.

#### 3.2.2. Special addresses

Term	Definition
Link-local Address	<b>FE80::/10</b> Used for local communication between devices on the same network segment.
Global Unicast Address	<b>2000::/3</b> A globally routable address, these addresses are equivalent to public IPv4 addresses and can be reached over the internet.
Unique Local Address ( <b>ULA</b> )	<b>FC00::/7</b> An address for local communication that is not routable on the global internet, similar to private addresses in IPv4.
Multicast Address	<b>FF00::/8</b> An address that enables a single packet to be sent to multiple destinations simultaneously.
Anycast Address	An address assigned to multiple interfaces, where a packet sent to an anycast address is routed to the nearest (in terms of routing distance) interface.
Reserved Address	Certain ranges in IPv6 are reserved for future use or specific functions. For example, addresses starting with <b>::/128</b> are reserved for unspecified addresses.
Documentation Address	<b>2001:DB8::/32</b> Designated specifically for use in documentation and examples, ensuring it does not conflict with real-world addresses.

<b>Term</b>	<b>Definition</b>
Link-local Multicast Address	<b>FF02::/16</b> Part of the link-local address range; it enables devices to communicate within a local network without requiring an external routing address.

<b>Addresses</b>	<b>Range</b>	<b>Scope</b>
Unspecified	::/128	n/a
Loopback	::1	Host
IPv4-Embedded	64:ff9b::/96	n/a
Discard-Only	100::/64	n/a
Link-Local	fe80::/10	Link
Global Unicast	2000::/3	Global
Unique Local (ULA)	fc00::/7	Global
Multicast	ff00::/8	Variable

### 3.2.2.1. Multicast

<b>Term</b>	<b>Definition</b>
ff02::1	All nodes, within scope 2 (link-local).
ff02::2	All routers, within scope 2 (link-local).
ff02::1:ffff:xxxx	The IPv6 node joins a solicited multicast address group from all the interfaces where unicast and anycast addresses are configured. Its scope is the link-local.

### 3.2.2.2. DHCPv6

<b>Term</b>	<b>Definition</b>
ff02::1:2	A link-scoped multicast address used by a client to communicate with neighboring (i.e., on-link) relay agents and servers. All servers and relay agents are members of this multicast group.
ff05::1:3	A site-scoped multicast address used by a relay agent to communicate with servers, either because the relay agent wants to send messages to all servers or because it does not know the unicast address of the servers.

### 3.2.3. Header

<b>Term</b>	<b>Definition</b>
Version	Always 6 with IPv6. IPv4 would be 4.
Flow Label	For identifying packets that require special handling, like real-time streaming.
Traffic Class	Priority or type of traffic.
Payload Length	Size of the payload in bytes.
Next Header	Type of optional header following the IPv6 header.
Hop Limit	Maximum number of hops a packet can take before being discarded.

### 3.2.3.1. IPv6 Extension Headers currently defined

<i>Term</i>	<i>Definition</i>
Routing	Extended routing, like IPv4 loose source route. The Routing header is used by an IPv6 source to list one or more intermediate nodes to be “visited” on the way to a packet’s destination. There are different types of routing headers defined for different uses.
Fragmentation	Fragmentation and reassembly. The Fragment header is used by an IPv6 source to send a packet larger than would fit in the path MTU to its destination.
Authentication	Integrity and authentication, security. The Authentication Header (AH) is used by IPsec to provide security services like integrity and data origin authentication to IPv6 traffic.
Encapsulating Security Payload	Confidentiality. Encapsulating Security Payload (ESP) Extension Header [RFC2406(opens in a new tab)] is used by IPsec to provide security services like confidentiality and/or integrity to IPv6 packets. The ESP Extension Header can be followed by an additional Destination Options Extension Header and the upper layer datagram.
Hop-by-Hop Option	Special options that require hop-by-hop processing. The Hop-by-Hop Options header is used to carry optional information that may be examined and processed by every node along a packet’s delivery path. The information is included in the form of one or more options using a TLV (Type-Length-Value) format.
Destination Options	Optional information to be examined by the destination node. The Destination Options header is used to carry optional information that is meant to be examined only by a packet’s destination node(s). The information is included in the form of one or more options using a TLV (Type-Length-Value) format.

### 3.2.4. Neighbor Discovery protocol (ND)

#### 3.2.4.1. Host - Router Discovery Functions

<i>Term</i>	<i>Definition</i>
Router discovery	Hosts can locate routers residing on attached links.
Prefix discovery	Hosts can discover address prefixes that are on-link for attached links.
Parameter discovery	Hosts can find parameters (e.g., MTU).
Address autoconfiguration	Stateless configuration of addresses of network interfaces.
Redirect	Provide a better next-hop route for certain destinations.

#### 3.2.4.2. Host - Host Communication Functions

<i>Term</i>	<i>Definition</i>
Address resolution	Mapping between IP addresses and link-layer addresses. This is equivalent to ARP for IPv4. This function allows to resolve the link-layer address of another node in the link when only the IPv6 address of that node is known.

<b>Term</b>	<b>Definition</b>
Next-hop determination	Hosts can find next-hop routers for a destination.
Neighbor unreachability detection (NUD)	Determine that a neighbor is no longer reachable on the link.
Duplicate address detection (DAD)	Nodes can check whether an address is already in use.

### 3.2.4.3. Packet types

<b>Name</b>	<b>Type</b>	<b>Description</b>
Router Solicitation ( <b>RS</b> )	133	To locate routers on an attached link.
Router Advertisement ( <b>RA</b> )	134	Used by routers to advertise their presence periodically or in response to a RS message.
Neighbor Solicitation ( <b>NS</b> )	135	To find the MAC-address of the neighbor or to check if the neighbor is still reachable.
Neighbor Advertisement ( <b>NA</b> )	136	To respond to a Neighbor Solicitation message.
Redirect	137	To point the host to a better first hop router for a destination.

### 3.2.5. Stateless Address Autoconfiguration (SLAAC)

A method for automatically configuring IPv6 addresses without a DHCP server, relying on local network information.

#### 3.2.5.1. Autoconfigure link-local address

Mac address: **70:07:12:34:56:78**

- 1) Flip **7th** bit: **72:07:12:34:56:78**. If it is “0”, the address is locally administered and if it is “1”, the address is globally unique.
- 2) Insert **FFEE** in the middle: **7207:12FF:EE34:5678**
- 3) Combine with link-local prefix: **FE80::7207:12FF:EE34:5678**

New address: **FE80::7207:12FF:EE34:5678**

#### 3.2.5.2. Perform Duplicate Address Detection (DAD)

To make sure that the address is actually unique in the local segment.

Upon configuring an IPv6 address, every node joins a **multicast group** identified by the address **FF02::1:FFxx:xxxx** where xx:xxxx are the **last 6 hexadecimal values** in the IPv6 unicast address, e.g. FF02::1:FF34:5678

- 1) The host sends a Neighbor Solicitation message from the Unspecified Address (::) to the Solicited Node multicast address.
- 2) If the generated address is in use, the host using that address sends a Neighbor Advertisement back. The sending host then knows the tentative address can not be used.
- 3) The host then proceeds to generate a new address and sends a new Neighbor Solicitation message to the link.
- 4) If there is no reply after some time, the host informs all the other hosts that it uses this address and it sends a Neighbor Advertisement message to the All Nodes address.
- 5) The host assigns the address to the interface and now has an active IPv6 link. This is the so-called Link-local Address Assignment.

#### 3.2.5.3. Router search

- 1) Router solicitation
- 2) Router advertisement

### 3.2.5.4. Generating global unicast address

Based on the information from the Router Advertisement, the host generates a global unicast address and wants to know if it is available to use, so it does the DAD process again. If it is not a duplicate, the host will use it.

## 3.2.6. DHCPv6

### 3.2.6.1. Flags

Term	Definition
A	Host can perform SLAAC to generate its own IPv6 address based on the prefix(es) contained in the RA message.
O	Host can fetch additional options from the DHCPv6. The DHCPv6 does not provide IPv6 addresses in this case.
M	Host will get its IP address and additional options from a DHCPv6 server.
L	The prefix shared in the RA is reachable on the link.

## 3.3. IPV4

### 3.3.1. Network classes (private nets)

Term	Definition
A	<b>10.0.0.0</b> - 10.255.255.255 (10/ <b>8</b> prefix)
B	<b>172.16.0.0</b> - 172.31.255.255 (172.16/ <b>12</b> prefix)
C	<b>192.168.0.0</b> - 192.168.255.255 (192.168/ <b>16</b> prefix)

### 3.3.2. Subnetting

#### 3.3.2.1. Calculating subnet mask

$$/24 = \underline{\text{1111}} \underline{\text{1111}} . \underline{\text{1111}} \underline{\text{1111}} . \underline{\text{1111}} \underline{\text{1111}} . \underline{\text{0000}} \underline{\text{0000}} = 255.255.255.0$$

$$/10 = \underline{\text{1111}} \underline{\text{1111}} . \underline{\text{1100}} \underline{\text{0000}} . \underline{\text{0000}} \underline{\text{0000}} . \underline{\text{0000}} \underline{\text{0000}} = 255.192.0.0$$

#### 3.3.2.2. Calculating increment

$$\text{Address increment} = \frac{\text{amount of addresses}}{256} \text{ or } 2^{8-(\text{mask mod 8})}$$

Let there be 4 subsequent networks starting with 10.0.0.0, each being /20

$$\text{Amount of addresses} = 2^{32-20} = 2^{12} = 4096. \text{ Increment} = \frac{4096}{256} = 16$$

$$\text{Alternatively: } 2^{8-(20 \text{ mod } 8)} = 2^{8-4} = 2^4 = 16$$

$$\text{Networks} = 10.0.0.0/20, 10.0.\textbf{16}.0/20, 10.0.\textbf{32}.0/20, 10.0.\textbf{48}.0/20$$

## 3.4. ROUTING

- routing table

### 3.4.1. Data plane

- local, per-router function
- determines how datagram arriving on router input port is forwarded to router output port

### 3.4.2. Control plane

- network-wide logic
- determines how datagram is routed among routers along end-to-end path from source to destination host

### 3.4.3. Static

### 3.4.4. Dynamic

#### 3.4.4.1. Algorithms

##### 3.4.4.1.1. Dijkstra's algorithm (Link State)

##### 3.4.4.1.2. ? algorithm (Distance vector)

#### 3.4.4.2. Protocols

##### 3.4.4.2.1. OSPF (Open Shortest Path First)

##### 3.4.4.2.2. BGP (Border Gateway Protocol)

### 3.4.5. Fragmentation

- offset = transferred bytes / 8
- 

## 4. DATA LINK LAYER (2)

Functions:

- error detection
- flow control
- addressing
- layer 2 packets change after each intermediary node (switch/router)
- ethernet
  - full duplex / half duplex
  - collision domain
  - CSMA/CD
- switching
  - switch table
  - flooding
  - learning
- vlans
  - forwarding between vlans is done via routing
  - trunk port
    - carries frames between vlans defined over multiple physical switches
  - frame format
  - spanning tree
    - so that broadcast packets will not continuously loop
    - a switch is selected as root
    - a tree-like loop-free topology is established
    - bridge protocol data units
      - Hello BPDU
      - Topology Change Notification (TCN) BPDU
      - comparison algorithm

## 4.1. ETHERNET FRAMES

### 4.1.1. Ethernet II Frame

Size: 64B (1518 Bytes)						
DA 6B	SA 6B	Type 2B	DATA (MAC SDU) 0 (+64 padding) ... 1500B			FCS 4B

Most common type in use today. Also called the DIX frame.

MAC PDU must be at least 64B to guarantee that all collisions can be detected. If it's smaller, the frame must be filled with Padding Bytes.

### 4.1.2. IEEE 802.3 Frame

Size: 64B (1518 Bytes)					
DA 6B	SA 6B	Length 2B	LLC 802.2	LLC SDU	FCS 4B

## 4.2. MAC-ADDRESS

Size: 6B (48bit)	
Organizationally Unique Identifier (OUI) 3B	NIC specific 3B

Used for identifying interfaces.

7th bit: Globally unique (0) or locally administered (1)

8th bit: Unicast (0) or multicast (1)

## 4.3. ADDRESS RESOLUTION PROTOCOL (ARP)

Maps network addresses to data link layer addresses. Resolves IPv4 addresses to MAC addresses.

IPv6 does not need ARP because it uses the Neighbor Discovery Protocol (NDP).

ARP table holds mappings from IPv4 to MAC addresses. Entries are added by monitoring the traffic and adding source IP and MAC addresses of the incoming packets to the table. If no entry is found inside of the ARP table, then the node launches an ARP discovery process. This is done by sending an ARP broadcast request and receiving an ARP reply from the requested MAC addresses' host. When a node receives a packet with a destination IP address where no cached entry for the MAC address can be found, the encapsulation of the IPv4 packet fails and the packet gets dropped.

ARP has no validation if the sender of a frame is correct. ARP spoofing, also called ARP poisoning, refers to the method of inserting the wrong MAC address into ARP requests and responses by the node. An attacker can lead sent frames to the wrong destination and has the ability to read the traffic (MITM attack). Configuring static ARP entries is one way to prevent ARP spoofing.

## 4.4. SWITCH

- All devices connected to the switch ports form a *broadcast domain*

## 4.5. VLAN

LAN: all devices in the same broadcast domain

VLAN: Virtual separation of LAN on a switch

Reasons for using VLANs:

- To reduce CPU overhead on each device by reducing the number of devices that receive each broadcast frame

- To reduce security risks by reducing the number of hosts that receive copies of frames that the switches flood (broadcasts, multicasts, and unknown unicasts)
- To improve security for hosts that send sensitive data by keeping those hosts on a separate VLAN
- To create more flexible designs that group users by department, or by groups that work together, instead of by physical location
- To solve problems more quickly, because the failure domain for many problems is the same set of devices as those in the same broadcast domain
- To reduce the workload for the Spanning Tree Protocol (STP) by limiting a VLAN to a single access switch

#### 4.5.1. Trunking

With **trunking**, only a single cable is needed to carry traffic for all VLANs. VLAN trunking works by applying **VLAN tagging**, where the sending switch adds an extra header to each frame before sending it across the trunk link. This trunking header contains a VLAN Identifier (VLAN ID), allowing the receiving switch to determine the VLAN to which each frame belongs. Switch ports that are assigned to a single VLAN and carry traffic for only that VLAN are referred to as **access ports**. Ports that carry traffic for multiple VLANs using VLAN tagging are called **trunk ports**.

#### 4.5.2. 802.1Q

The standard of how to tag an ethernet frame in a trunk is defined in IEEE 802.1Q. 802.1Q inserts an extra 4 byte 802.1Q VLAN header into the original frame's Ethernet header.

#### 4.5.3. Inter-VLAN Routing

##### 4.5.3.1. Attaching a Router

A router can be added to a switch using multiple VLANs. The cable from the switch to the router gets configured as a trunk. The router then can simply perform its usual routing logic between the subnets. This concept is called **Router-on-a-Stick**.

##### 4.5.3.2. Using a Layer 3 Switch

With the use of a switch with layer 3 capabilities, the need for a separate router is omitted, as the switch brings the ability for routing by itself. Routing can be turned on that switch and packets between the VLANs get routed.

#### 4.5.4. Link Aggregation Group (LAG)

Combine a number of physical ports together to one logical port.

#### 4.5.5. Link Aggregation Control Protocol (LACP)

IEEE specification (802.3ad) that also enables several physical ports to be bundled together to form a LAG. LACP enables a switch to negotiate an automatic bundle by sending LACP packets to the peer.

### 4.6. ERROR DETECTION

EDC

#### 4.6.1. Cyclic Redundancy Check (CRC)

---

## 5. CISCO

### 5.1. ROUTER SETUP

```
Router> enable  
Router# configure terminal  
Router(config)#
```

## 5.2. INTERFACES

### 5.2.1. Static IP Assignment

```
Router(config)# interface GigabitEthernet 0/0/1
Router(config-if)# ip address 172.16.0.0 255.255.255.252
Router(config-if)# no shutdown
Router(config-if)# exit
```

### 5.2.2. DHCP Assignment

```
Router(config)# interface GigabitEthernet 0/1/1
Router(config-if)# ip address dhcp
Router(config-if)# no shutdown
Router(config-if)# exit
```

### 5.2.3. Show

```
Router(config)# do show ip interface brief
Router# show ip interface brief
Router# show ip interface GigabitEthernet 0/0/1
```

## 5.3. VLAN

```
Switch(config)# vlan 120
Switch(config-if)# name vlan-server
Switch(config-if)# exit
```

### 5.3.1. Assign IP

```
Switch(config)# interface vlan 120
Switch(config-if)# ip address 10.120.0.10 255.255.255.0
```

### 5.3.2. Access Port

```
Switch(config)# interface GigabitEthernet 0/0/1
Switch(config-if)# switchport mode access
Switch(config-if)# switchport access vlan 120
Switch(config-if)# exit
Switch(config)# interface GigabitEthernet 0/0/1-5
Switch(config-if)# switchport mode access
Switch(config-if)# switchport access vlan 120
Switch(config-if)# exit
```

### 5.3.3. Access Port

```
Switch(config-if)# switchport mode trunk
```

### 5.3.4. VTP (Virtual Trunk Protocol)

#### 5.3.4.1. Server

```
Switch(config)# vtp domain ins
Switch(config)# vtp mode server
```

#### 5.3.4.2. Client

```
Switch(config)# vtp domain ins
Switch(config)# vtp mode client
```

### 5.3.5. LACP (Link Aggregation Control Protocol)

```
Switch(config-if)# channel-group 5 mode active
Switch(config-if)# channel-group 5 mode passive
```

### 5.3.6. Load Balancing

```
Switch(config)# port-channel load-balance <strategy>
```

### 5.3.7. STP (Spanning Tree Protocol)

#### 5.3.7.1. Bridge priority

```
Switch(config)# spanning-tree vlan 1 priority <priority>
```

#### 5.3.7.2. Interface costs

```
Switch(config-if)# spanning-tree cost 100
```

#### 5.3.7.3. PortFast mode

```
Switch(config-if)# spanning-tree portfast
```

#### 5.3.7.4. Show

```
Switch# show spanning-tree
```

```
Switch# show spanning-tree root
```

## 5.4. ROUTING

### 5.4.1. Static

#### 5.4.1.1. IPv4

```
Router(config-if)# ip route <destination_network_id> <subnet_mask> <next_hop_router>
<administrative_distance>?
```

```
Router(config-if)# ip route 10.0.0.0 255.0.0.0 192.168.1.1
```

#### 5.4.1.2. IPv6

```
Router(config-if)# ip route <ipv6_prefix> <outgoing_interface> <next-hop>
<administrative_distance>?
```

```
Router(config-if)# ipv6 route 2001:db8:2103:a::/64 GigabitEthernet1/0/1
fe80::ba27:ebff:fea8:3e50
```

### 5.4.2. OSPF

#### 5.4.2.1. IPv4

```
Router(config)# router ospf <process-id>
```

```
Router(config-if)# ip ospf <process-id> area <area-nr>
```

#### 5.4.2.2. IPv6

```
Router(config)# ipv6 router ospf <process-id>
```

```
Router(config-if)# ipv6 ospf <process-id> area <area-nr>
```

#### 5.4.3. Show

```
Router# show ip route
```

```
Router# show ip ospf route
```

## 5.5. DHCP

### 5.5.1. Create Pool

```
Router#(config-if) ip dhcp pool DEV
network 192.168.1.0 255.255.255.0
default-router 192.168.1.1
dns-server 1.1.1.1 8.8.8.8
lease 5
domain-name enterprise.com
```

### 5.5.2. Relay Agent

```
Router(config-if)# ip helper-address 176.16.12.10
```

## 5.6. NAT

### 5.6.1. IF Inside

```
Router(config-if)# ip nat inside
```

### 5.6.2. IF Outside

```
Router(config-if)# ip nat outside
```

### 5.6.3. ACL (Access Control List)

```
Router(config-if)# access-list 1 permit 192.168.1.0 0.0.0.255
```

### 5.6.4. PAT (Nat overload)

```
Router(config-if)# ip nat inside source list 1 interface GigabitEthernet0/1 overload
```

## 5.7. IPV6

```
Router(config-if)# ipv6 enable
```

### 5.7.1. DHCPv6

```
Router(config-if)# ipv6 dhcp client pd MY_PREFIX
```

```
Router(config-if)# ipv6 address autoconfig default
```

## 5.8. TROUBLESHOOTING

### 5.8.1. Ping

```
Router# ping <destination-ip> source <interface-name>
```

### 5.8.2. Traceroute

```
Router# traceroute <destination-ip> source <interface-name> numeric
```

---

## 6. BINARY, DECIMAL, HEX

**0xA4 6A = 0b1010 0100 0110 1010 = 0d42'090**

**0x04 B0 = 0b0100 1011 0000 = 0d1'200**

**0x01 D4 C0 = 0b0001 1101 0100 1100 0000 = 0d120'000**