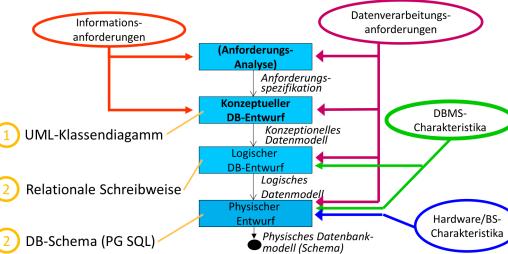


Georgiy Shevoroshkin

- uml diagramm (konzeptionell)
 - assoziationen, bedingungen
- transaktionen
 - fuzzy read, deadlock, dirty read, write skew, phantom read, serializable snapshot isolation, cascading rollbacks
 - schedule analysieren + Serialisierbarkeitsgraph
- begriffe (physisches schema, DBMS)
- bbaum indexe einfügen



DataBase System

Besteht aus Datenbankmanagementsystem und Datenbasen

DataBase Management System

- Redundanzfreiheit
- Datenintegrität
- Kapselung

ANSI Modell Logische Ebene:

Interne Ebene:

Externe Ebene:

Mapping:

Unified Modeling Language

- Assoziation
- ◆ Komposition
- Aggregation
- Vererbung

Complete: Alle Subklassen sind definiert

Incomplete: Zusätzliche Subklassen sind erlaubt

Disjoint: Ist Instanz von genau einer Unterklassie

Overlapping: Kann Instanz von mehreren überlappenden Unterklassen sein

Normalisierung

1NF: Atomare Attributwerte

2NF: Nichtschlüsselattr. voll vom Schlüssel abhängig

3NF: Keine transitiven Abhängigkeiten

BCNF: Nur abhängigkeiten vom Schlüssel

(Voll-)funktionale Abhängigkeit:

Transitive Abhängigkeit:

Einfügeanomalie, Löschanomalie, Änderungsanomalie

Vererbung

Tabelle pro Sub- und Superklasse:

-- TODO: check if correct

```
CREATE TABLE sup (id SERIAL PRIMARY KEY, -- 3.a
  name TEXT UNIQUE);
```

```
CREATE TABLE sub1 (id SERIAL PRIMARY KEY, age INT);
CREATE TABLE sub2 (id SERIAL PRIMARY KEY);
ALTER TABLE sub1 ADD CONSTRAINT id FOREIGN KEY
REFERENCES sup (id); -- Auch für sub2
```

Tabelle pro Subklasse: Enthält jeweil. Subklassattribute

```
CREATE TABLE sub1 (id SERIAL PRIMARY KEY, -- 3.b
  name TEXT UNIQUE, age INT);
CREATE TABLE sub2 (id SERIAL PRIMARY KEY,
  name TEXT UNIQUE);
```

Einige Tabelle für Superklasse: Enthält alle Attribute

```
CREATE TABLE sup (id SERIAL PRIMARY KEY, -- 3.c
  name TEXT UNIQUE, age INT);
```

Data Definition Language

```
CREATE SCHEMA s;
CREATE TABLE t (id SERIAL PRIMARY KEY,
  name TEXT UNIQUE,
  grade DECIMAL(2,1) NOT NULL,
  fk INT FOREIGN KEY REFERENCES t2.id ON DELETE CASCADE,
  added TIMESTAMP DEFAULT CURRENT_TIMESTAMP,
  u VARCHAR(9) DEFAULT CURRENT_USER,
  CHECK (grade between 1 and 6));
ALTER TABLE t2 ADD CONSTRAINT c PRIMARY KEY (a, b);
TRUNCATE/DROP TABLE t;
```

Data Manipulation Language

```
INSERT INTO t (id, grade) VALUES (1, 1) RETURNING id;
```

Views

```
CREATE VIEW v (id, grade, u) AS SELECT id, grade, u
FROM t;
```

Data Control Language

```
CREATE ROLE r WITH LOGIN PASSWORD '';
GRANT INSERT ON TABLE t TO r;
REVOKE CREATE ON SCHEMA s FROM r;
ALTER ROLE r CREATEROLE, CREATEDB, INHERIT;
GRANT r TO user_name;
-- read all future created tables
ALTER DEFAULT PRIVILEGES IN SCHEMA s GRANT SELECT ON
TABLES TO readonlyuser;
CREATE POLICY p ON t FOR ALL TO PUBLIC USING (u =
current_user);
ALTER TABLE t ENABLE ROW LEVEL SECURITY;
```

Common Table Expressions

```
WITH RECURSIVE q AS (SELECT * FROM t WHERE grade > 1
UNION ALL SELECT * FROM t INNER JOIN q ON q.u =
t.name) SELECT id as 'ID' FROM q;
```

Window Functions

```
SELECT id, RANK() OVER
  (ORDER BY grade DESC) as r FROM t;
SELECT id, u, LAG(name, 1) OVER
  (PARTITION BY fk ORDER BY id DESC) FROM t;
-- PERCENT/DENSE_RANK(), FIRST_VALUE(v), LAST_VALUE(n)
-- NTH_VALUE(v,n), NTILE(n), LEAD(v,o), ROW_NUMBER()
```

Subqueries

```
SELECT * FROM t WHERE grade > ANY (SELECT g FROM t2);
SELECT * FROM t WHERE EXISTS (SELECT g FROM t2);
-- ALL, IN, =
```

JOIN

```
SELECT a.* , b.* FROM a INNER JOIN b ON a.id = b.id;
SELECT y.* , x.* FROM t AS y JOIN LATERAL
  (SELECT * FROM t2 WHERE t2.id = y.id) AS x;
```

GROUP BY

```
SELECT id, COUNT(*) FROM t
  GROUP BY grade, id HAVING COUNT(*) > 2;
```

WHERE

```
BETWEEN 1 AND 5; LIKE '___%'; AND; IS (NOT) NULL
IN (1, 5) ; LIKE '%asd'; OR ;
```

INDEX

```
CREATE INDEX i ON t /*USING BTREE*/ (grade, UPPER(u))
INCLUDE added;
DROP INDEX i;
```

Transaktionen

```
BEGIN; SAVEPOINT s;
COMMIT; ROLLBACK /*TO SAVEPOINT s*/;
```

Isolation

```
SET TRANSACTION ISOLATION LEVEL ...;
SET SESSION CHARACTERISTICS AS TRANSACTION ISOLATION
LEVEL ...;
```

	Read Un-committed	Read Committed	Repeatable Read	Serializable
Dirty Write	*	*	*	X
Dirty Read	✓	X	X	X
Lost Update	✓	✓	X	X
Fuzzy Read	✓	✓	X	X
Phantom Read	✓	✓	✓	X
Read Skew	✓	✓	X	X
Write Skew	✓	✓	✓	*
Deadlock				X
Cascading				X
Rollback				

Dirty Read: Lese Daten von nicht committed T's

Fuzzy Read: Versch. Werte beim mehrmaligen Lesen gleicher Daten (da durch andere T geändert)

Phantom Read: Neue/Gelöschte Rows einer anderen T

Relationale Algebra

```
πR1,R4(R) SELECT R1,R4 FROM R;
σR1>30(R) SELECT * FROM R WHERE R1 > 30;
ρa←R SELECT * FROM R AS a;
R × S SELECT * FROM R,S;
```

```
R ⋈A=B S SELECT * FROM R JOIN S ON R.A=S.B;
```

Serialisierbarkeit

Backup

B-Baum

