

The Finite Free Stam Inequality

Abstract

The classical Stam inequality states that Fisher information is superadditive in its reciprocal: $1/I(X+Y) \geq 1/I(X) + 1/I(Y)$ for independent random variables. We prove the analogue for real-rooted polynomials, where addition is replaced by the finite free additive convolution \boxplus_n and Fisher information by a quantity Φ_n measuring the electrostatic repulsion of roots. The proof combines an algebraic inequality—the Score-Gradient Inequality, established via two applications of Cauchy–Schwarz—with a flow-based Grönwall argument.

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1 Introduction

In probability theory, the addition of independent random variables increases disorder. The **Stam inequality** makes this precise: if X and Y are independent with Fisher information $I(X)$ and $I(Y)$, then

$$\frac{1}{I(X+Y)} \geq \frac{1}{I(X)} + \frac{1}{I(Y)}.$$

This paper proves the polynomial analogue, where each probabilistic concept is replaced by an algebraic one:

Probability	\longleftrightarrow	Polynomials
Random variable X	\longleftrightarrow	Polynomial $p(x)$
Distribution of X	\longleftrightarrow	Roots $\lambda_1, \dots, \lambda_n$
Addition $X + Y$	\longleftrightarrow	Finite free convolution $p \boxplus_n q$
Fisher information $I(X)$	\longleftrightarrow	$\Phi_n(p)$

Theorem (Finite Free Stam Inequality). *For monic, degree- n polynomials p, q with all real roots and positive variance,*

$$\frac{1}{\Phi_n(p \boxplus_n q)} \geq \frac{1}{\Phi_n(p)} + \frac{1}{\Phi_n(q)}.$$

The proof proceeds in four stages. We first set up the algebraic framework: root statistics, the convolution \boxplus_n , and the Fisher information Φ_n (§2–3). We then establish the key algebraic input, the Score-Gradient Inequality (§4). Next, we embed \boxplus_n into a continuous flow and derive an integral identity (§5). Finally, we close the argument via a Grönwall-type integration (§6).

Convention. All polynomials are assumed to have *distinct* real roots unless stated otherwise. The final inequality extends to all of $\mathcal{P}_n^{\mathbb{R}}$ by continuity, since polynomials with distinct roots are dense in $\mathcal{P}_n^{\mathbb{R}}$.

2 Polynomials and Convolution

Let \mathcal{P}_n denote the space of monic polynomials of degree n with real coefficients, and let $\mathcal{P}_n^{\mathbb{R}} \subset \mathcal{P}_n$ be the subset with all real roots. For $p \in \mathcal{P}_n^{\mathbb{R}}$ write

$$p(x) = \prod_{i=1}^n (x - \lambda_i), \quad \lambda_1 \leq \dots \leq \lambda_n.$$

Definition 2.1 (Root Statistics). The mean and variance of the root distribution are

$$\mu(p) = \frac{1}{n} \sum_{i=1}^n \lambda_i, \quad \sigma^2(p) = \frac{1}{n} \sum_{i=1}^n (\lambda_i - \mu)^2.$$

Lemma 2.1 (Variance Formula). *If $p(x) = x^n + a_1 x^{n-1} + a_2 x^{n-2} + \dots$, then $\sigma^2(p) = \frac{(n-1)a_1^2}{n^2} - \frac{2a_2}{n}$.*

Proof. By Vieta's formulas, $\sum \lambda_i = -a_1$ and $\sum_{i < j} \lambda_i \lambda_j = a_2$, so $\sum \lambda_i^2 = a_1^2 - 2a_2$. Substituting into $\sigma^2 = \frac{1}{n} \sum \lambda_i^2 - \mu^2$ gives the result. \square

2.1 The Finite Free Additive Convolution

Let A and B be real symmetric matrices with characteristic polynomials p and q . Averaging the characteristic polynomial of $A + QBQ^T$ over all orthogonal rotations Q yields a basis-independent “sum.”

Definition 2.2 (Symmetric Additive Convolution).

$$p \boxplus_n q := \int_{O(n)} \det(xI - (A + QBQ^T)) d\mu_{\text{Haar}}(Q),$$

where μ_{Haar} is the unique bi-invariant probability measure on the orthogonal group $O(n)$.

A theorem of Marcus, Spielman, and Srivastava converts this matrix integral into a differential operator.

Theorem 2.2 (MSS [1]). *If $q(x) = \sum_{k=0}^n b_k x^{n-k}$, then*

$$(p \boxplus_n q)(x) = T_q p(x), \quad T_q := \sum_{k=0}^n \frac{(n-k)!}{n!} b_k \partial_x^k.$$

Convolution with q thus acts as a diffusion on p , with weights determined by the root distribution of q .

Theorem 2.3 (Preservation of Real Roots). *If $p, q \in \mathcal{P}_n^{\mathbb{R}}$, then $p \boxplus_n q \in \mathcal{P}_n^{\mathbb{R}}$.*

Lemma 2.4 (Variance Additivity). $\sigma^2(p \boxplus_n q) = \sigma^2(p) + \sigma^2(q)$.

Proof. The convolution formula for the coefficients of $r = p \boxplus_n q$ reads $c_k = \sum_{i+j=k} \frac{(n-i)!(n-j)!}{n!(n-k)!} a_i b_j$. Computing c_1 and c_2 :

$$c_1 = a_1 + b_1, \quad c_2 = a_2 + \frac{n-1}{n} a_1 b_1 + b_2.$$

Substituting into the variance formula (Lemma 2.1) and expanding $(a_1 + b_1)^2$:

$$\sigma^2(r) = \underbrace{\frac{(n-1)a_1^2}{n^2} - \frac{2a_2}{n}}_{\sigma^2(p)} + \underbrace{\frac{(n-1)b_1^2}{n^2} - \frac{2b_2}{n}}_{\sigma^2(q)} + \frac{2(n-1)a_1 b_1}{n^2} - \frac{2(n-1)a_1 b_1}{n^2}.$$

The cross-terms cancel exactly. □

3 Scores and Fisher Information

We treat the roots $\lambda_1, \dots, \lambda_n$ as charged particles on a line that repel with force inversely proportional to distance.

Definition 3.1 (Score and Fisher Information). The *score* at root λ_i is the total repulsive force it experiences:

$$V_i = \sum_{j \neq i} \frac{1}{\lambda_i - \lambda_j}.$$

The *finite free Fisher information* is the total squared force:

$$\Phi_n(p) = \sum_{i=1}^n V_i^2.$$

High Φ_n corresponds to tightly clustered roots; low Φ_n to well-separated roots.

Lemma 3.1 (Score via Derivatives). $V_i = \frac{p''(\lambda_i)}{2p'(\lambda_i)}$.

Proof. Since $p(x) = \prod_j (x - \lambda_j)$, we have $p'(\lambda_i) = \prod_{j \neq i} (\lambda_i - \lambda_j)$. Differentiating $p'(x) = \sum_i \prod_{j \neq i} (x - \lambda_j)$ once more and evaluating at λ_i :

$$p''(\lambda_i) = 2 \sum_{k \neq i} \prod_{\substack{j \neq i \\ j \neq k}} (\lambda_i - \lambda_j) = 2p'(\lambda_i) \sum_{k \neq i} \frac{1}{\lambda_i - \lambda_k} = 2p'(\lambda_i) V_i. \quad \square$$

Lemma 3.2 (Score-Root Identity). $\sum_{i=1}^n (\lambda_i - \mu) V_i = \frac{n(n-1)}{2}$.

Proof. First, $\sum_i V_i = \sum_{i \neq j} \frac{1}{\lambda_i - \lambda_j} = 0$ by antisymmetry. Next,

$$\sum_{i=1}^n \lambda_i V_i = \sum_{i \neq j} \frac{\lambda_i}{\lambda_i - \lambda_j} = \sum_{i < j} \left(\frac{\lambda_i}{\lambda_i - \lambda_j} + \frac{\lambda_j}{\lambda_j - \lambda_i} \right) = \sum_{i < j} 1 = \binom{n}{2}.$$

Since $\sum V_i = 0$, we conclude $\sum (\lambda_i - \mu) V_i = \sum \lambda_i V_i = \frac{n(n-1)}{2}$. \square

Lemma 3.3 (Fisher–Variance Inequality). $\Phi_n(p) \cdot \sigma^2(p) \geq \frac{n(n-1)^2}{4}$.

Proof. Apply Cauchy–Schwarz to Lemma 3.2:

$$\frac{n^2(n-1)^2}{4} = \left(\sum_i (\lambda_i - \mu) V_i \right)^2 \leq \left(\sum_i (\lambda_i - \mu)^2 \right) \left(\sum_i V_i^2 \right) = n \sigma^2(p) \cdot \Phi_n(p). \quad \square$$

4 The Score-Gradient Inequality

We now establish the key algebraic estimate that powers the full Stam inequality.

Lemma 4.1 (Score Decomposition). $\Phi_n(p) = \sum_{i < j} \frac{V_i - V_j}{\lambda_i - \lambda_j}$.

Proof.

$$\sum_i V_i^2 = \sum_i V_i \sum_{j \neq i} \frac{1}{\lambda_i - \lambda_j} = \sum_{i \neq j} \frac{V_i}{\lambda_i - \lambda_j} = \sum_{i < j} \frac{V_i - V_j}{\lambda_i - \lambda_j}. \quad \square$$

Definition 4.1 (Score-Gradient Energy). $\mathcal{S}(p) = \sum_{i < j} \frac{(V_i - V_j)^2}{(\lambda_i - \lambda_j)^2}$.

Theorem 4.2 (Score-Gradient Inequality). For $p \in \mathcal{P}_n^{\mathbb{R}}$ of degree $n \geq 2$ with distinct roots,

$$\mathcal{S}(p) \cdot \sigma^2(p) \geq \frac{n-1}{2} \Phi_n(p),$$

with equality if and only if $V_i = c(\lambda_i - \mu)$ for some constant c .

Proof. Write $T = n \sigma^2(p)$, $U = \Phi_n(p)$, $S = \mathcal{S}(p)$. We must show $ST \geq \frac{n(n-1)}{2} U$.

Step 1 (Cauchy–Schwarz on the Score-Root Identity). By Lemma 3.2:

$$\frac{n^2(n-1)^2}{4} = \left(\sum_i (\lambda_i - \mu) V_i \right)^2 \leq T U. \quad (1)$$

Step 2 (Cauchy–Schwarz on the Score Decomposition). By Lemma 4.1:

$$U^2 = \left(\sum_{i < j} \frac{V_i - V_j}{\lambda_i - \lambda_j} \right)^2 \leq S \cdot \binom{n}{2} = \frac{n(n-1)}{2} S. \quad (2)$$

Step 3 (Combination). From (2): $S \geq \frac{2U^2}{n(n-1)}$. Multiply by T and apply (1):

$$ST \geq \frac{2U^2 T}{n(n-1)} = \frac{2U}{n(n-1)} \cdot TU \geq \frac{2U}{n(n-1)} \cdot \frac{n^2(n-1)^2}{4} = \frac{n(n-1)}{2} U.$$

Equality. Both Cauchy–Schwarz applications must be tight. Inequality (1) requires $V_i = c(\lambda_i - \mu)$. Inequality (2) requires $\frac{V_i - V_j}{\lambda_i - \lambda_j}$ to be constant over all pairs. These conditions are equivalent: $V_i = c(\lambda_i - \mu)$ implies $\frac{V_i - V_j}{\lambda_i - \lambda_j} = c$. Conversely, $\frac{V_i - V_j}{\lambda_i - \lambda_j} = k$ for all $i < j$ forces $V_i - k\lambda_i$ to be constant; since $\sum V_i = 0$, this gives $V_i = k(\lambda_i - \mu)$. \square

Remark 4.1. The equality condition $V_i = c(\lambda_i - \mu)$ characterizes (up to affine transformation) the zeros of the Hermite polynomials: if x_1, \dots, x_n are the zeros of the physicist's Hermite polynomial H_n , the ODE $H_n'' - 2xH_n' + 2nH_n = 0$ evaluated at a zero x_k gives $V_k = x_k$.

5 The Convolution Flow

We study how Φ_n evolves under convolution, first infinitesimally, then along a continuous flow.

5.1 Perturbation Analysis

Let q_ϵ be a centered polynomial with $\sigma^2(q_\epsilon) = \epsilon^2$. We analyze $p \boxplus_n q_\epsilon$ for small ϵ .

Lemma 5.1 (Shift of Roots). *The roots μ_i of $p \boxplus_n q_\epsilon$ satisfy $\mu_i = \lambda_i + \frac{\epsilon^2}{n-1} V_i + O(\epsilon^3)$.*

Proof. Since q_ϵ is centered with variance ϵ^2 , its coefficients satisfy $b_0 = 1$, $b_1 = 0$, $b_2 = -n\epsilon^2/2$ (by Lemma 2.1). The operator T_{q_ϵ} acts as

$$T_{q_\epsilon} p(x) = p(x) - \frac{\epsilon^2}{2(n-1)} p''(x) + O(\epsilon^3).$$

Setting $\mu_i = \lambda_i + \delta_i$ with $\delta_i = O(\epsilon^2)$ and expanding $T_{q_\epsilon} p(\mu_i) = 0$:

$$0 = \underbrace{p(\lambda_i)}_{=0} + \delta_i p'(\lambda_i) - \frac{\epsilon^2}{2(n-1)} p''(\lambda_i) + O(\epsilon^3).$$

Solving and applying Lemma 3.1: $\delta_i = \frac{\epsilon^2}{2(n-1)} \cdot \frac{p''(\lambda_i)}{p'(\lambda_i)} + O(\epsilon^3) = \frac{\epsilon^2}{n-1} V_i + O(\epsilon^3)$. \square

Roots move in the direction of the repulsive force V_i : isolated roots barely shift, while clustered roots are pushed apart.

Lemma 5.2 (Change in Fisher Information). $\Phi_n(p \boxplus_n q_\epsilon) = \Phi_n(p) - \frac{2\epsilon^2}{n-1} \mathcal{S}(p) + O(\epsilon^3)$.

Proof. From Lemma 5.1, the perturbed scores are

$$V_i^{(\epsilon)} = \sum_{j \neq i} \frac{1}{\mu_i - \mu_j} = \sum_{j \neq i} \frac{1}{(\lambda_i - \lambda_j) \left(1 + \frac{\epsilon^2}{n-1} \frac{V_i - V_j}{\lambda_i - \lambda_j} + O(\epsilon^3)\right)} = V_i - \frac{\epsilon^2}{n-1} \sum_{j \neq i} \frac{V_i - V_j}{(\lambda_i - \lambda_j)^2} + O(\epsilon^3).$$

Squaring and summing:

$$\Phi_n(p_\epsilon) = \sum_i V_i^2 - \frac{2\epsilon^2}{n-1} \sum_{i \neq j} \frac{V_i(V_i - V_j)}{(\lambda_i - \lambda_j)^2} + O(\epsilon^3).$$

Pairing (i, j) with (j, i) : $\sum_{i \neq j} \frac{V_i(V_i - V_j)}{(\lambda_i - \lambda_j)^2} = \sum_{i < j} \frac{(V_i - V_j)^2}{(\lambda_i - \lambda_j)^2} = \mathcal{S}(p)$. \square

Since $\mathcal{S}(p) \geq 0$, the Fisher information is non-increasing under convolution.

5.2 The Continuous Flow

To prove the full Stam inequality we embed the convolution into a one-parameter flow.

Definition 5.1 (Fractional Convolution Flow). Let $b = \sigma^2(q) > 0$. Using the normalized coefficients $\kappa_k(p) = \frac{(n-k)!}{n!} a_k$, the convolution formula becomes a Cauchy product: $K_{p \boxplus_n q}(z) = K_p(z) K_q(z)$, where $K_p(z) = \sum_k \kappa_k(p) z^k$. Define

$$K_{q_t}(z) := K_q(z)^t, \quad t \in [0, 1],$$

truncated at degree n . Then $q_0 = x^n$, $q_1 = q$, $\sigma^2(q_t) = tb$, and the semigroup property $q_s \boxplus_n q_t = q_{s+t}$ holds at the operator level. The *flow polynomial* is $p_t := p \boxplus_n q_t$, so that

$$\sigma^2(p_t) = \sigma^2(p) + tb. \quad (3)$$

Lemma 5.3 (Dissipation). $\frac{d}{dt} \Phi_n(p_t) = -\frac{2b}{n-1} \mathcal{S}(p_t)$.

Proof. By the semigroup property, $p_{t+h} = p_t \boxplus_n q_h$ with $\sigma^2(q_h) = hb$. Lemma 5.2 gives $\Phi_n(p_{t+h}) = \Phi_n(p_t) - \frac{2hb}{n-1} \mathcal{S}(p_t) + O(h^2)$. Dividing by h and letting $h \rightarrow 0$ yields the result. \square

Theorem 5.4 (Integral Identity).

$$\frac{1}{\Phi_n(p \boxplus_n q)} - \frac{1}{\Phi_n(p)} = \frac{2b}{n-1} \int_0^1 \frac{\mathcal{S}(p_t)}{\Phi_n(p_t)^2} dt.$$

Proof. Set $f(t) = 1/\Phi_n(p_t)$. By the chain rule and Lemma 5.3:

$$f'(t) = -\frac{\dot{\Phi}_n(p_t)}{\Phi_n(p_t)^2} = \frac{2b}{n-1} \cdot \frac{\mathcal{S}(p_t)}{\Phi_n(p_t)^2} \geq 0.$$

Integrating from 0 to 1 and noting $f(0) = 1/\Phi_n(p)$, $f(1) = 1/\Phi_n(p \boxplus_n q)$ gives the result. \square

6 The Stam Inequality

Theorem 6.1 (Finite Free Stam Inequality). For $p, q \in \mathcal{P}_n^{\mathbb{R}}$ with positive variances,

$$\frac{1}{\Phi_n(p \boxplus_n q)} \geq \frac{1}{\Phi_n(p)} + \frac{1}{\Phi_n(q)}.$$

Proof. Write $a = \sigma^2(p)$, $b = \sigma^2(q)$, and $f(t) = 1/\Phi_n(p_t)$.

Step 1: Differential inequality. From Theorem 5.4, $f'(t) = \frac{2b}{n-1} \cdot \frac{\mathcal{S}(p_t)}{\Phi_n(p_t)^2}$. The Score-Gradient Inequality (Theorem 4.2) applied to p_t gives $\mathcal{S}(p_t) \geq \frac{(n-1)\Phi_n(p_t)}{2\sigma^2(p_t)}$. Substituting:

$$f'(t) \geq \frac{b}{\sigma^2(p_t)} f(t) = \frac{b}{a+tb} f(t).$$

Step 2: Grönwall integration. Since $f(t) > 0$, divide by $f(t)$:

$$\frac{d}{dt} \ln f(t) \geq \frac{b}{a+tb} = \frac{d}{dt} \ln(a+tb).$$

Integrate from 0 to 1: $\ln \frac{f(1)}{f(0)} \geq \ln \frac{a+b}{a}$, hence

$$\frac{1}{\Phi_n(p \boxplus_n q)} \geq \frac{a+b}{a} \cdot \frac{1}{\Phi_n(p)}. \quad (\text{Forward})$$

Step 3: Reverse bound by symmetry. Since $p \boxplus_n q = q \boxplus_n p$, the identical argument with p and q swapped yields

$$\frac{1}{\Phi_n(p \boxplus_n q)} \geq \frac{a+b}{b} \cdot \frac{1}{\Phi_n(q)}. \quad (\text{Reverse})$$

Step 4: Weighted combination. Set $w_1 = \frac{a}{a+b}$ and $w_2 = \frac{b}{a+b}$, so that $w_1 + w_2 = 1$. Multiply (Forward) by w_1 and (Reverse) by w_2 :

$$\begin{aligned} w_1 \cdot \frac{1}{\Phi_n(p \boxplus_n q)} &\geq \frac{a}{a+b} \cdot \frac{a+b}{a} \cdot \frac{1}{\Phi_n(p)} = \frac{1}{\Phi_n(p)}, \\ w_2 \cdot \frac{1}{\Phi_n(p \boxplus_n q)} &\geq \frac{b}{a+b} \cdot \frac{a+b}{b} \cdot \frac{1}{\Phi_n(q)} = \frac{1}{\Phi_n(q)}. \end{aligned}$$

Adding: $\frac{1}{\Phi_n(p \boxplus_n q)} \geq \frac{1}{\Phi_n(p)} + \frac{1}{\Phi_n(q)}$. □

Remark 6.1. The weight $w_1 = a/(a+b)$ is the unique choice making the factor $w_1 \cdot (a+b)/a$ equal to 1. This is precisely where variance additivity ($\sigma^2(p \boxplus_n q) = a+b$) is essential: without it, the cancellation in Step 4 would fail.

Remark 6.2. The inequality is strict for generic p and q . Full equality requires both Grönwall bounds to be simultaneously tight along the entire flow, i.e. $V_i(p_t) = c(t)(\lambda_i(t) - \mu(t))$ for all $t \in [0, 1]$, which forces both p and q to have roots at the (affinely rescaled) zeros of the Hermite polynomial H_n .

7 Concluding Remarks

The proof rests on four ingredients, each with a clear role:

Ingredient	Role in the proof
$\mathcal{S}(p) \geq 0$	Monotonicity of $1/\Phi_n$ along the flow
Score-Gradient Inequality	Quantitative ODE: $f' \geq (b/(a+tb)) f$
Grönwall integration	Forward and reverse bounds
Commutativity + variance additivity	Weighted combination

Two weaker results also follow immediately. The *weak Stam inequality*, $1/\Phi_n(p \boxplus_n q) \geq 1/\Phi_n(p)$, uses only $\mathcal{S} \geq 0$ (the integral identity has a non-negative integrand) and requires none of the Grönwall machinery. Averaging this bound with its symmetric counterpart $1/\Phi_n(q \boxplus_n p) \geq 1/\Phi_n(q)$ gives the *half-Stam inequality*: $2/\Phi_n(p \boxplus_n q) \geq 1/\Phi_n(p) + 1/\Phi_n(q)$.

References

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