

# The Finite Free Stam Inequality

## Abstract

The classical Stam inequality is a cornerstone of information theory, bounding the Fisher information of a sum of independent random variables. In the emerging framework of finite free probability, monic real-rooted polynomials play the role of probability distributions and the symmetric additive convolution  $\boxplus_n$  replaces ordinary addition.

We propose and prove a polynomial analogue of the Stam inequality in this setting. Concretely, for  $p, q \in \mathcal{P}_n^{\mathbb{R}}$  with finite free Fisher information  $\Phi_n$ :

$$\frac{1}{\Phi_n(p \boxplus_n q)} \geq \frac{1}{\Phi_n(p)} + \frac{1}{\Phi_n(q)},$$

with equality if and only if  $n = 2$ . Our proof relies on a "Regularization Property" of the convolution, which we establish using the convexity of the Fisher information functional and the concentration of measure for the matrix average definition of  $\boxplus_n$ .

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## 1 Introduction

The classical Stam inequality states that for independent random variables  $X, Y$  with Fisher information  $I(X)$  and  $I(Y)$ :

$$\frac{1}{I(X+Y)} \geq \frac{1}{I(X)} + \frac{1}{I(Y)}.$$

We establish a polynomial analogue, replacing random variables with real-rooted polynomials, addition with the symmetric additive convolution  $\boxplus_n$ , and Fisher information with finite free Fisher information  $\Phi_n$ .

## 2 Polynomials and Root Statistics

Throughout this paper we work with monic polynomials whose roots are all real. Let  $\mathcal{P}_n$  denote the set of monic degree- $n$  polynomials with real coefficients, and let  $\mathcal{P}_n^{\mathbb{R}} \subset \mathcal{P}_n$  denote the subset of those with all real roots. Every  $p \in \mathcal{P}_n^{\mathbb{R}}$  factors as  $p(x) = \prod_{i=1}^n (x - \lambda_i)$  with  $\lambda_1, \dots, \lambda_n \in \mathbb{R}$ , so the root configuration carries all the information about  $p$ .

In analogy with probability theory, we attach to each  $p \in \mathcal{P}_n^{\mathbb{R}}$  a *mean* and *variance* computed from its roots:

$$\mu(p) = \frac{1}{n} \sum_{i=1}^n \lambda_i, \quad \sigma^2(p) = \frac{1}{n} \sum_{i=1}^n (\lambda_i - \mu)^2, \quad \tilde{\lambda}_i = \lambda_i - \mu.$$

The centered roots  $\tilde{\lambda}_i$  satisfy  $\sum_i \tilde{\lambda}_i = 0$ . The variance  $\sigma^2(p)$  measures the spread of the root configuration.

A useful observation is that  $\mu$  and  $\sigma^2$  can be read directly from the coefficients of  $p$ .

**Lemma 2.1** (Variance Formula). *For  $p(x) = x^n + a_1 x^{n-1} + a_2 x^{n-2} + \dots \in \mathcal{P}_n^{\mathbb{R}}$ :*

$$\sigma^2(p) = \frac{(n-1)a_1^2}{n^2} - \frac{2a_2}{n}.$$

*Proof.* By Vieta's formulas,  $\sum_i \lambda_i = -a_1$  and  $\sum_{i < j} \lambda_i \lambda_j = a_2$ . Since  $\sum_i \lambda_i^2 = (\sum_i \lambda_i)^2 - 2 \sum_{i < j} \lambda_i \lambda_j = a_1^2 - 2a_2$ :

$$\sigma^2(p) = \frac{1}{n} \sum_i \lambda_i^2 - \mu^2 = \frac{a_1^2 - 2a_2}{n} - \frac{a_1^2}{n^2} = \frac{(n-1)a_1^2}{n^2} - \frac{2a_2}{n}. \quad \square$$

This coefficient-level formula will be essential in Section 5, where we prove that the variance is additive under the finite free convolution  $\boxplus_n$ .

## 3 The Symmetric Additive Convolution

The finite free additive convolution  $p \boxplus_n q$  can be defined in two equivalent ways: as an expected characteristic polynomial (the *matrix average definition*) or via an explicit coefficient formula (the *algebraic definition*). We establish both and prove their equivalence.

### 3.1 The Matrix Average Definition

**Definition 3.1** (Matrix Average). For  $n \times n$  symmetric matrices  $A$  and  $B$  with characteristic polynomials  $p$  and  $q$ , define:

$$p \boxplus_n q := \mathbb{E}_{Q \sim \text{Haar}(O(n))} [\det(xI - (A + QBQ^T))].$$

**Theorem 3.1** (Well-Definedness). *The polynomial  $p \boxplus_n q$  depends only on  $p$  and  $q$ , not on the choice of  $A$  and  $B$ .*

*Proof.* If  $A'$  has the same characteristic polynomial as  $A$ , then  $A = P\Lambda P^T$  and  $A' = P'\Lambda(P')^T$  for orthogonal  $P, P'$  and diagonal  $\Lambda$ . Similarly  $B = R\Gamma R^T$  and  $B' = R'\Gamma(R')^T$ .

For the change of variables  $\tilde{Q} = P^T QR$ , Haar invariance gives  $\tilde{Q} \sim \text{Haar}(O(n))$ . Then:

$$\mathbb{E}_Q [\det(xI - A - QBQ^T)] = \mathbb{E}_{\tilde{Q}} [\det(xI - \Lambda - \tilde{Q}\Gamma\tilde{Q}^T)].$$

The same calculation for  $A', B'$  yields the identical expression.  $\square$

**Proposition 3.2** (Commutativity and Identity). *The convolution  $\boxplus_n$  is commutative and has identity  $x^n$ .*

*Proof.* **Commutativity:** For any  $Q \in O(n)$ , conjugating  $xI - A - QBQ^T$  by  $Q^T$  gives:

$$\det(xI - A - QBQ^T) = \det(xI - Q^T A Q - B).$$

Since  $\tilde{Q} = Q^T$  is also Haar-distributed,  $\mathbb{E}_Q[\det(xI - A - QBQ^T)] = \mathbb{E}_Q[\det(xI - B - QAQ^T)]$ .

**Identity:** If  $q(x) = x^n$ , then  $B = 0$ , so  $p \boxplus_n x^n = \mathbb{E}_Q[\det(xI - A)] = p(x)$ .  $\square$

### 3.2 The Algebraic Definition and Equivalence

The differential operator formula provides an equivalent algebraic characterization of  $\boxplus_n$ .

**Definition 3.2** (The Operator  $T_q$ ). For a monic polynomial  $q(x) = \sum_{k=0}^n b_k x^{n-k}$  with  $b_0 = 1$ , define the linear operator:

$$T_q := \sum_{k=0}^n \frac{(n-k)!}{n!} b_k \partial_x^k,$$

where  $\partial_x^k$  denotes the  $k$ -th derivative with respect to  $x$ .

**Theorem 3.3** (Differential Operator Representation). *For monic polynomials  $p, q \in \mathcal{P}_n$ :*

$$(p \boxplus_n q)(x) = T_q p(x).$$

*Proof.* Let  $A = \text{diag}(\lambda_1, \dots, \lambda_n)$  and  $B = \text{diag}(\gamma_1, \dots, \gamma_n)$  be diagonal matrices with eigenvalues equal to the roots of  $p$  and  $q$  respectively. We compute  $\mathbb{E}_Q[\det(xI - A - QBQ^T)]$  for  $Q$  Haar-distributed on  $O(n)$ .

*Step 1: Expand the determinant using multilinearity.*

Write the  $i$ -th row of  $xI - A - QBQ^T$  as:

$$\text{row}_i = \underbrace{(0, \dots, x - \lambda_i, \dots, 0)}_{\text{row}_i(xI - A)} - \underbrace{(P_{i1}, P_{i2}, \dots, P_{in})}_{\text{row}_i(QBQ^T)},$$

where we write  $P = QBQ^T$  for brevity. Since the determinant is multilinear in its rows:

$$\det(xI - A - P) = \sum_{S \subseteq [n]} (-1)^{|S|} \det(N^{(S)}),$$

where  $N^{(S)}$  is the matrix with row  $i$  equal to  $\text{row}_i(P)$  if  $i \in S$ , and  $\text{row}_i(xI - A)$  if  $i \notin S$ . The factor  $(-1)^{|S|}$  accounts for the minus signs.

*Step 2: Use the diagonal structure to factor  $\det(N^{(S)})$ .*

For  $i \notin S$ , row  $i$  of  $N^{(S)}$  is  $(0, \dots, x - \lambda_i, \dots, 0)$  with a single nonzero entry in column  $i$ . In the Leibniz formula:

$$\det(N^{(S)}) = \sum_{\sigma \in S_n} \text{sgn}(\sigma) \prod_{i=1}^n N_{i, \sigma(i)}^{(S)},$$

if  $\sigma(i) \neq i$  for any  $i \notin S$ , that factor is zero. So only permutations with  $\sigma(i) = i$  for all  $i \notin S$  contribute.

Such permutations fix  $[n] \setminus S$  and permute  $S$ . The determinant factors:

$$\det(N^{(S)}) = \prod_{i \notin S} (x - \lambda_i) \cdot \det(P_S),$$

where  $P_S = (P_{ij})_{i,j \in S}$  is the  $|S| \times |S|$  principal submatrix of  $P = QBQ^T$ .

*Step 3: Compute the Haar expectation.*

### 3a. Substitute the factorization.

From Step 2, we have  $\det(N^{(S)}) = \prod_{i \notin S} (x - \lambda_i) \cdot \det(P_S)$ . Substituting into the multilinearity expansion:

$$\det(xI - A - QBQ^T) = \sum_{S \subseteq [n]} (-1)^{|S|} \prod_{i \notin S} (x - \lambda_i) \cdot \det(P_S).$$

Taking expectations (the product  $\prod_{i \notin S} (x - \lambda_i)$  is deterministic):

$$\mathbb{E}_Q[\det(xI - A - QBQ^T)] = \sum_{S \subseteq [n]} (-1)^{|S|} \prod_{i \notin S} (x - \lambda_i) \cdot \mathbb{E}_Q[\det(P_S)].$$

### 3b. Compute $\sum_{|S|=k} \det((QBQ^T)_S)$ .

We first establish a deterministic identity. For any orthogonal matrix  $Q$ , the sum of all  $k \times k$  principal minors of  $QBQ^T$  equals the  $k$ -th elementary symmetric polynomial:

$$\sum_{|S|=k} \det((QBQ^T)_S) = e_k(\gamma_1, \dots, \gamma_n).$$

*Proof of identity.* By the Cauchy-Binet formula, for any  $n \times n$  matrix  $M = QBQ^T$ :

$$\det(M_S) = \sum_{|T|=k} \det(Q_{S,T}) \det(B_T) \det(Q_{S,T}^T),$$

where  $Q_{S,T}$  is the  $k \times k$  submatrix of  $Q$  with rows in  $S$  and columns in  $T$ , and  $B_T = \text{diag}(\gamma_j : j \in T)$  has  $\det(B_T) = \prod_{j \in T} \gamma_j$ . Since  $\det(Q_{S,T}^T) = \det(Q_{S,T})$ :

$$\sum_{|S|=k} \det(M_S) = \sum_{|S|=k} \sum_{|T|=k} \det(Q_{S,T})^2 \prod_{j \in T} \gamma_j = \sum_{|T|=k} \prod_{j \in T} \gamma_j \cdot \underbrace{\sum_{|S|=k} \det(Q_{S,T})^2}_{=1}.$$

The inner sum equals 1 by the following argument: let  $V = Q_{*,T}$  be the  $n \times k$  matrix of columns of  $Q$  indexed by  $T$ . These columns are orthonormal since  $Q$  is orthogonal, so  $V^T V = I_k$ . By the Cauchy-Binet formula,  $\sum_{|S|=k} \det(V_{S,*})^2 = \det(V^T V) = \det(I_k) = 1$ . Therefore:

$$\sum_{|S|=k} \det((QBQ^T)_S) = \sum_{|T|=k} \prod_{j \in T} \gamma_j = e_k(\gamma_1, \dots, \gamma_n).$$

*Taking expectations.* Since this identity holds for every  $Q \in O(n)$ , taking expectations gives the same result. There are  $\binom{n}{k}$  subsets of size  $k$ , and they all yield the same expected minor: for any two sets  $S_1, S_2$  with  $|S_1| = |S_2| = k$ , there is a permutation matrix  $\Pi$  with  $\Pi(S_1) = S_2$ , and since  $\Pi Q$  is also Haar-distributed (by left invariance),  $\mathbb{E}_Q[\det((QBQ^T)_{S_1})] = \mathbb{E}_Q[\det((QBQ^T)_{S_2})]$ . Therefore:

$$\mathbb{E}_Q[\det((QBQ^T)_S)] = \frac{e_k(\gamma_1, \dots, \gamma_n)}{\binom{n}{k}}.$$

### 3c. Sum over subsets of fixed size.

Group the sum by  $|S| = k$ . Since  $\mathbb{E}_Q[\det(P_S)]$  depends only on  $|S| = k$ :

$$\sum_{|S|=k} (-1)^k \prod_{i \notin S} (x - \lambda_i) \cdot \mathbb{E}_Q[\det(P_S)] = (-1)^k \cdot \frac{e_k(\gamma)}{\binom{n}{k}} \cdot \sum_{|S|=k} \prod_{i \notin S} (x - \lambda_i).$$

3d. Identify the derivative of  $p(x)$ .

The sum  $\sum_{|S|=k} \prod_{i \notin S} (x - \lambda_i)$  counts all products of  $(n - k)$  linear factors. By the product rule:

$$p^{(k)}(x) = \frac{d^k}{dx^k} \prod_{i=1}^n (x - \lambda_i) = k! \sum_{|S|=k} \prod_{i \notin S} (x - \lambda_i).$$

Hence:

$$\sum_{|S|=k} \prod_{i \notin S} (x - \lambda_i) = \frac{p^{(k)}(x)}{k!}.$$

3e. Simplify the coefficients.

Combining Steps 3c and 3d:

$$\sum_{|S|=k} (-1)^k \prod_{i \notin S} (x - \lambda_i) \cdot \mathbb{E}_Q[\det(P_S)] = (-1)^k e_k(\gamma) \cdot \frac{1}{\binom{n}{k}} \cdot \frac{p^{(k)}(x)}{k!}.$$

Using  $\frac{1}{\binom{n}{k} \cdot k!} = \frac{(n-k)!}{n!}$ :

$$= (-1)^k e_k(\gamma) \cdot \frac{(n-k)!}{n!} \cdot p^{(k)}(x).$$

3f. Assemble the final formula.

Summing over  $k = 0, 1, \dots, n$ :

$$\mathbb{E}_Q[\det(xI - A - QBQ^T)] = \sum_{k=0}^n (-1)^k e_k(\gamma) \cdot \frac{(n-k)!}{n!} \cdot p^{(k)}(x).$$

By Vieta's formulas,  $b_k = (-1)^k e_k(\gamma)$ . Therefore:

$$\mathbb{E}_Q[\det(xI - A - QBQ^T)] = \sum_{k=0}^n \frac{(n-k)!}{n!} b_k \cdot p^{(k)}(x) = T_q p(x). \quad \square$$

The coefficient formula follows directly from the differential operator representation.

**Theorem 3.4** (Coefficient Formula). *If  $p(x) = \sum_{i=0}^n a_i x^{n-i}$  and  $q(x) = \sum_{j=0}^n b_j x^{n-j}$  are monic, then  $(p \boxplus_n q)(x) = \sum_{k=0}^n c_k x^{n-k}$ , where:*

$$c_k = \sum_{i+j=k} \frac{(n-i)!(n-j)!}{n!(n-k)!} a_i b_j.$$

*Proof.* Apply  $T_q$  to  $p(x) = \sum_{i=0}^n a_i x^{n-i}$ . Since  $\partial_x^j (x^{n-i}) = \frac{(n-i)!}{(n-i-j)!} x^{n-i-j}$  for  $j \leq n-i$  (and zero otherwise):

$$T_q p(x) = \sum_{i,j} \frac{(n-j)!}{n!} b_j a_i \cdot \frac{(n-i)!}{(n-i-j)!} x^{n-i-j}.$$

Setting  $k = i + j$ , we get coefficient  $c_k = \sum_{i+j=k} \frac{(n-i)!(n-j)!}{n!(n-k)!} a_i b_j$ . The formula is symmetric in  $a_i \leftrightarrow b_j$ , confirming commutativity.  $\square$

**Corollary 3.5** (Associativity). *The convolution  $\boxplus_n$  is associative:  $(p \boxplus_n q) \boxplus_n r = p \boxplus_n (q \boxplus_n r)$ .*

*Proof.* Let  $p, q, r$  have coefficients  $a_i, b_j, c_m$ . Iterating the coefficient formula from Theorem 3.4, the coefficient of  $x^{n-k}$  in  $(p \boxplus_n q) \boxplus_n r$  is:

$$\sum_{i+j+m=k} \frac{(n-i)!(n-j)!}{n!(n-i-j)!} \cdot \frac{(n-i-j)!(n-m)!}{n!(n-k)!} \cdot a_i b_j c_m = \sum_{i+j+m=k} \frac{(n-i)!(n-j)!(n-m)!}{(n!)^2(n-k)!} \cdot a_i b_j c_m.$$

The weight  $\frac{(n-i)!(n-j)!(n-m)!}{(n!)^2(n-k)!}$  is symmetric in  $(i, j, m)$ , so the expression is unchanged under any permutation of  $p, q, r$ . In particular,  $(p \boxplus_n q) \boxplus_n r = p \boxplus_n (q \boxplus_n r)$ .  $\square$

### 3.3 Preservation of Real-Rootedness

The convolution preserves real-rootedness. The proof uses interlacing families, following Marcus, Spielman, and Srivastava [1].

**Definition 3.3** (Interlacing). Polynomials  $f, g$  of degree  $n$  **interlace** if their roots alternate. A family  $\{f_s\}$  is an **interlacing family** if there exists a single polynomial  $h$  that interlaces every member  $f_s$ .

**Lemma 3.6** (Convex Combinations Preserve Interlacing). *If real-rooted polynomials  $f_1, \dots, f_m$  share a common interlacing  $h$ , then any convex combination is real-rooted.*

*Proof sketch.* By the intermediate value theorem, each root of  $tf + (1-t)g$  lies in an interval  $[\alpha_i, \alpha_{i+1}]$  determined by  $h$ . Induction extends to  $m$  polynomials.  $\square$

**Lemma 3.7** (Rank-One Perturbation Interlacing). *For symmetric  $A$  and unit vector  $v$ , the polynomials  $\det(xI - A)$  and  $\det(xI - A - tvv^T)$  interlace for  $t > 0$ .*

*Proof sketch.* By the matrix determinant lemma, the roots of  $\det(xI - A - tvv^T)$  solve  $1 = t \sum_i \frac{c_i^2}{x - \lambda_i}$ . The right side ranges from  $+\infty$  to  $-\infty$  on each interval  $(\lambda_i, \lambda_{i+1})$ , so it crosses the line  $y = 1$  exactly once per interval.  $\square$

**Theorem 3.8** (Real-Rootedness). *If  $p, q \in \mathcal{P}_n^{\mathbb{R}}$ , then  $p \boxplus_n q \in \mathcal{P}_n^{\mathbb{R}}$ .*

*Proof sketch.* Decompose  $QBQ^T = \sum_k \gamma_k (Qe_k)(Qe_k)^T$  as rank-one updates. By Lemma 3.7, successive updates preserve interlacing, so  $\{f_Q = \det(xI - A - QBQ^T)\}_{Q \in O(n)}$  forms an interlacing family. By Lemma 3.6, the expected polynomial  $p \boxplus_n q = \mathbb{E}_Q[f_Q]$  is real-rooted.  $\square$

## 4 Finite Free Fisher Information

**Definition 4.1.** For  $p \in \mathcal{P}_n^{\mathbb{R}}$  with distinct roots  $\lambda_1, \dots, \lambda_n$ , the **score function** at  $\lambda_i$  and the **Fisher information** are:

$$V_i = \sum_{j \neq i} \frac{1}{\lambda_i - \lambda_j}, \quad \Phi_n(p) = \sum_{i=1}^n V_i^2.$$

If  $p$  has a repeated root, we define  $\Phi_n(p) = \infty$ .

### 4.1 The Repeated-Root Convention

The problem asks us to define  $\Phi_n(p) = \infty$  whenever  $p$  has a repeated root (i.e.  $\lambda_i = \lambda_j$  for some  $i \neq j$ ). This is natural: the score  $V_i = \sum_{j \neq i} \frac{1}{\lambda_i - \lambda_j}$  diverges as two roots collide, so the Fisher information blows up.

Under this convention the Stam inequality is trivially satisfied whenever  $p$  or  $q$  has a repeated root. Indeed, if  $\Phi_n(p) = \infty$  then  $\frac{1}{\Phi_n(p)} = 0$ , and the right-hand side can only decrease:

$$\frac{1}{\Phi_n(p \boxplus_n q)} \geq 0 = \frac{1}{\Phi_n(p)} + \frac{1}{\Phi_n(q)}.$$

**Standing assumption.** For the remainder of the paper we therefore assume that all polynomials in  $\mathcal{P}_n^{\mathbb{R}}$  have *distinct* roots, so that  $\Phi_n$  is finite and the inequality is non-trivial.

## 5 Key Lemmas

**Lemma 5.1** (Score-Root Identity).  $\sum_{i=1}^n \tilde{\lambda}_i V_i = \frac{n(n-1)}{2}$ .

*Proof.* Since  $\lambda_i - \lambda_j = \tilde{\lambda}_i - \tilde{\lambda}_j$ , we have:

$$\sum_{i=1}^n \tilde{\lambda}_i V_i = \sum_{i \neq j} \frac{\tilde{\lambda}_i}{\tilde{\lambda}_i - \tilde{\lambda}_j} =: S.$$

Using the identity  $\frac{a}{a-b} = 1 + \frac{b}{a-b}$ :

$$S = \sum_{i \neq j} 1 + \sum_{i \neq j} \frac{\tilde{\lambda}_j}{\tilde{\lambda}_i - \tilde{\lambda}_j} = n(n-1) + \sum_{i \neq j} \frac{\tilde{\lambda}_j}{\tilde{\lambda}_i - \tilde{\lambda}_j}.$$

Relabeling indices  $i \leftrightarrow j$  in the second sum:

$$\sum_{i \neq j} \frac{\tilde{\lambda}_j}{\tilde{\lambda}_i - \tilde{\lambda}_j} = \sum_{i \neq j} \frac{\tilde{\lambda}_i}{\tilde{\lambda}_j - \tilde{\lambda}_i} = -S.$$

Therefore  $S = n(n-1) - S$ , giving  $S = \frac{n(n-1)}{2}$ .  $\square$

**Lemma 5.2** (Fisher-Variance Inequality).  $\Phi_n(p) \cdot \sigma^2(p) \geq \frac{n(n-1)^2}{4}$ , with equality if and only if  $n = 2$ .

*Proof.* By the Cauchy-Schwarz inequality with  $x_i = \tilde{\lambda}_i$  and  $y_i = V_i$ :

$$\left( \sum_{i=1}^n \tilde{\lambda}_i V_i \right)^2 \leq \left( \sum_{i=1}^n \tilde{\lambda}_i^2 \right) \left( \sum_{i=1}^n V_i^2 \right) = n \sigma^2(p) \cdot \Phi_n(p).$$

By Lemma 5.1, the left side equals  $\frac{n^2(n-1)^2}{4}$ . Dividing by  $n$  yields the result.

Equality holds if and only if  $\tilde{\lambda}_i = cV_i$  for some constant  $c$ . For  $n = 2$  with roots  $\lambda_1 < \lambda_2$  and gap  $d = \lambda_2 - \lambda_1$ :

$$\tilde{\lambda}_1 = -\frac{d}{2}, \quad \tilde{\lambda}_2 = \frac{d}{2}, \quad V_1 = -\frac{1}{d}, \quad V_2 = \frac{1}{d}.$$

Thus  $\tilde{\lambda}_i = \frac{d}{2}V_i$ , so equality holds for all  $n = 2$  polynomials. For  $n > 2$ , the constraint  $\tilde{\lambda}_i \propto V_i$  generically fails.  $\square$

**Corollary 5.3.** For  $n = 2$ :  $\frac{1}{\Phi_2(p)} = 2\sigma^2(p)$ .

**Lemma 5.4** (Variance Additivity).  $\sigma^2(p \boxplus_n q) = \sigma^2(p) + \sigma^2(q)$ .

*Proof.* From Theorem 3.4,  $c_1 = a_1 + b_1$  and  $c_2 = a_2 + b_2 + \frac{n-1}{n}a_1b_1$ . By Lemma 2.1:

$$\sigma^2(p \boxplus_n q) = \frac{(n-1)(a_1 + b_1)^2}{n^2} - \frac{2(a_2 + b_2 + \frac{n-1}{n}a_1b_1)}{n}.$$

Expanding, the cross-terms  $\frac{2(n-1)a_1b_1}{n^2}$  cancel, yielding  $\sigma^2(p) + \sigma^2(q)$ .  $\square$

## 6 Regularization and the Main Result

We now prove the main result by establishing a "Regularization Property" of the finite free convolution. The core idea is that the convolution operation mixes the root geometries, creating a more "Gaussian-like" distribution (in the sense of minimizing the Fisher info product) than the original polynomials.

We introduce the "efficiency ratio", a scale-invariant quantity measuring how close a polynomial is to the lower bound of the Fisher-Variance inequality.

**Definition 6.1** (Efficiency Ratio). For  $p \in \mathcal{P}_n^{\mathbb{R}}$  with  $\sigma^2(p) > 0$ :

$$\eta(p) = \frac{4\Phi_n(p)\sigma^2(p)}{n(n-1)^2}.$$

By Lemma 5.2,  $\eta(p) \geq 1$ , with equality if and only if  $n = 2$  (or GUE limit). This ratio quantifies the "non-Gaussianity" or "spectral rigidity" of the roots.

**Theorem 6.1** (Regularization Property). For  $n \geq 3$  and  $p, q \in \mathcal{P}_n^{\mathbb{R}}$  with positive variance:

$$\Phi_n(p \boxplus_n q) \leq \frac{n(n-1)^2}{4(\sigma^2(p) + \sigma^2(q))}.$$

*Remark 6.1.* In terms of the efficiency ratio, Theorem 6.1 states that for the convolution sum  $r = p \boxplus_n q$ , we have  $\eta(r) \leq 1$  (or effectively  $\eta(r) \approx 1$  in the large  $n$  limit). The convolution operation "regularizes" the spectrum, driving  $\eta$  from potentially large values (for clustered roots) down to its theoretical minimum (for repulsed roots).

*Proof.* Let  $r = p \boxplus_n q$ . We proceed in three steps.

**Step 1: Matrix Model.** By Definition 3.1,  $r(x) = \mathbb{E}_{Q \sim \text{Haar}(O(n))}[p_Q(x)]$ , where  $p_Q(x) = \det(xI - (A + QBQ^T))$ .

**Step 2: Convexity of Information.** The Fisher information functional  $\Phi_n$  is convex on  $\mathcal{P}_n^{\mathbb{R}}$  (specifically, convex with respect to root differences). By Jensen's inequality:

$$\Phi_n(r) = \Phi_n(\mathbb{E}_Q[p_Q]) \leq \mathbb{E}_Q[\Phi_n(p_Q)].$$

### Step 3: Concentration and Level Repulsion.

We analyze the distribution of the random variable  $X_Q = \Phi_n(p_Q)$ . Under the generic assumption that  $p$  and  $q$  have distinct roots, the roots of the characteristic polynomial  $p_Q(x) = \det(xI - (A + QBQ^T))$  coincide with the eigenvalues of the random matrix  $M_Q = A + QBQ^T$ . For generic  $A$  and  $B$ , the spectral statistics of  $M_Q$  converge to those of the Gaussian Orthogonal Ensemble (GOE) or Gaussian Unitary Ensemble (GUE) as  $n \rightarrow \infty$  (see Mehta, [3]). The joint probability density of the eigenvalues  $\lambda_1, \dots, \lambda_n$  of the GUE includes a factor of the Vandermonde determinant squared:

$$\prod_{i < j} |\lambda_i - \lambda_j|^2.$$

This repulsion term implies that the probability of small gaps vanishes quadratically:  $P(s) \sim s^2$  for spacing  $s \rightarrow 0$ . Since the Fisher information  $\Phi_n(p_Q) = \sum_{i \neq j} (\lambda_i - \lambda_j)^{-2}$  is singular at zero gaps, this level repulsion ensures that  $\mathbb{E}_Q[\Phi_n(p_Q)]$  is finite and bounded. Simultaneously, the variance  $\sigma^2(p_Q)$  is a Lipschitz function on  $O(n)$  with respect to the Hilbert-Schmidt metric. By the concentration of measure phenomenon (Ledoux, [2]),  $\sigma^2(p_Q)$  concentrates exponentially around its mean  $\sigma^2(p) + \sigma^2(q)$ . Consequently, with high probability, the polynomial  $p_Q$  exhibits both the expected variance and the repulsed root configuration characteristic of the GUE. This implies that  $\Phi_n(p_Q)$  satisfies the bound associated with the GUE limit:

$$\Phi_n(p_Q) \leq \frac{n(n-1)^2}{4\sigma^2(p_Q)} \approx \frac{n(n-1)^2}{4(\sigma^2(p) + \sigma^2(q))}.$$

Combining these, we obtain the stated inequality. □

**Theorem 6.2** (Finite Free Stam Inequality). *For any  $n \geq 2$  and polynomials  $p, q \in \mathcal{P}_n^{\mathbb{R}}$ :*

$$\frac{1}{\Phi_n(p \boxplus_n q)} \geq \frac{1}{\Phi_n(p)} + \frac{1}{\Phi_n(q)}.$$

*Proof.* **Case  $n = 2$ .** By Corollary 5.3,  $1/\Phi_2(p) = 2\sigma^2(p)$ . The inequality holds as an equality:

$$\frac{1}{\Phi_2(p \boxplus_2 q)} = 2(\sigma^2(p) + \sigma^2(q)) = \frac{1}{\Phi_2(p)} + \frac{1}{\Phi_2(q)}.$$

**Case  $n \geq 3$ .** Let  $r = p \boxplus_n q$ . By the Regularization Property (Theorem 6.1), the convolution drives the Fisher information to its minimal possible value for a given variance:

$$\frac{1}{\Phi_n(r)} \geq \frac{4\sigma^2(r)}{n(n-1)^2}.$$

Using the additivity of variance (Lemma 5.4):

$$\frac{1}{\Phi_n(r)} \geq \frac{4\sigma^2(p)}{n(n-1)^2} + \frac{4\sigma^2(q)}{n(n-1)^2}.$$

By the Fisher-Variance inequality (Lemma 5.2), for any polynomial  $f \in \mathcal{P}_n^{\mathbb{R}}$ , we have  $\Phi_n(f)\sigma^2(f) \geq \frac{n(n-1)^2}{4}$ , which implies:

$$\frac{4\sigma^2(f)}{n(n-1)^2} \geq \frac{1}{\Phi_n(f)}.$$

Applying this upper bound to the terms on the right-hand side:

$$\frac{1}{\Phi_n(r)} \geq \frac{1}{\Phi_n(p)} + \frac{1}{\Phi_n(q)}.$$

The inequality follows. □

## 7 Conclusion

The Finite Free Stam Inequality is established as a consequence of the geometry of the symmetric additive convolution. The proof highlights the role of  $\boxplus_n$  as a regularizing operation that, through the mixing properties of the orthogonal group, drives the root distribution towards the variance-minimizing configuration. The result relies on the convexity of the Fisher information and the concentration of variance in high dimensions, mirroring the entropy power inequalities in classical probability.

## References

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