D^3 as a 2-MCFL

Orestis Melkonian, Konstantinos Kogkalidis

January 25, 2018

Universiteit Utrecht

1.3 Challenge

[This is not a hand-in exercise. If you can solve it by Dec 5, there will be a present for you!]

Let D^n be the language over an n-symbol alphabet, lexicographically ordered $a_1 < \cdots < a_n$, where words satisfy the following conditions:

1. each word contains an equal number of the n alphabet symbols
2. for every prefix p of a word, the number of a_i in $p \geq$ the number of a_{i+1} $(1 \leq i \leq n-1)$

D^n generalizes the familiar language of balanced brackets, in which case you have an alphabet of size 2, say $\{\mathsf{a}, \mathsf{b}\}$, with 'opening bracket' a preceding 'closing bracket' b in the lexicographic ordering.

The conjecture (Makoto Kanazawa, p.c.) is that for $n \geq 2$, D^n is the language of a non-wellnested $(n-1)$-MCFG.

Give a 2-MCFG for D^3, i.e. words over a 3-letter alphabet $\{\mathsf{a}, \mathsf{b}, \mathsf{c}\}$ (with the usual lexicographic order) satisfying conditions (1) and (2) above. Give the ACG encoding of your MCFG for D^3.

Reference M. Moortgat (2014), A note on multidimensional Dyck languages.

Some examples

Dyck words

- abc
- aabbcc
- abcabcabacbc

Some examples

Dyck words

- abc
- aabbcc
- abcabcabacbc

Non-dyck words

- aabb

Some examples

Dyck words

- abc
- aabbcc
- abcabcabacbc

Non-dyck words

- aabb
- aabbbcc

Some examples

Dyck words

- abc
- aabbcc
- abcabcabacbc

Non-dyck words

- aabb
- aabbbcc
- abcacb

Some examples

Dyck words

- abc
- aabbcc
- abcabcabacbc

Non-dyck words

- aabb
- aabbbcc
- abcacb

$$a\,b\,a\,b\,a\,c\,b\,c\,a\,b\,c\,c$$

First-match policy

Some examples

Dyck words

- abc
- aabbcc
- abcabcabacbc

Non-dyck words

- aabb
- aabbbcc
- abcacb

$a\,b\,a\,b\,a\,c\,b\,c\,a\,b\,c\,c$

First-match policy

Some examples

Dyck words

- abc
- aabbcc
- abcabcabacbc

Non-dyck words

- aabb
- aabbbcc
- abcacb

$$a\,b\,a\,b\,a\,c\,b\,c\,a\,b\,c\,c$$

First-match policy

Some examples

Dyck words

- abc
- aabbcc
- abcabcabacbc

Non-dyck words

- aabb
- aabbbcc
- abcacb

$a\,b\,a\,b\,a\,c\,b\,c\,a\,b\,c\,c$

First-match policy

Some examples

Dyck words

- abc
- aabbcc
- abcabcabacbc

Non-dyck words

- aabb
- aabbbcc
- abcacb

First-match policy

G_0: Grammar of triple insertions

$$\mathrm{S}(xy) \leftarrow \mathrm{W}(x, y). \tag{1}$$

$$\mathrm{W}(\epsilon, xy\mathbf{abc}) \leftarrow \mathrm{W}(x, y). \tag{2}$$

$$\mathrm{W}(\epsilon, x\mathbf{a}y\mathbf{bc}) \leftarrow \mathrm{W}(x, y). \tag{3}$$

$$\ldots\ldots$$

$$\mathrm{W}(\mathbf{ab}x\mathbf{c}y, \epsilon) \leftarrow \mathrm{W}(x, y). \tag{60}$$

$$\mathrm{W}(\mathbf{abc}xy, \epsilon) \leftarrow \mathrm{W}(x, y). \tag{61}$$

$$\mathrm{W}(\epsilon, \mathbf{abc}). \tag{62}$$

$$\mathrm{W}(\mathbf{a}, \mathbf{bc}). \tag{63}$$

$$\mathrm{W}(\mathbf{ab}, \mathbf{c}). \tag{64}$$

$$\mathrm{W}(\mathbf{abc}, \epsilon). \tag{65}$$

G_0: Grammar of triple insertions

$a \quad a \quad b \quad b \quad a \quad c \quad b \quad a \quad c \quad c \quad b \quad c$

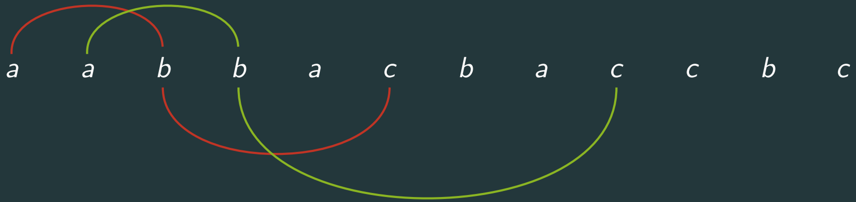
Straddling counter-example

G_0: Grammar of triple insertions

$a \quad a \quad b \quad b \quad a \quad c \quad b \quad a \quad c \quad c \quad b \quad c$

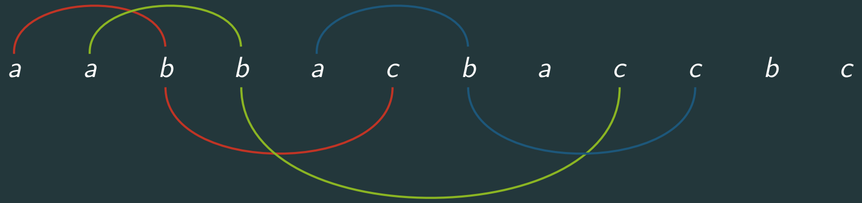
Straddling counter-example

G_0: Grammar of triple insertions



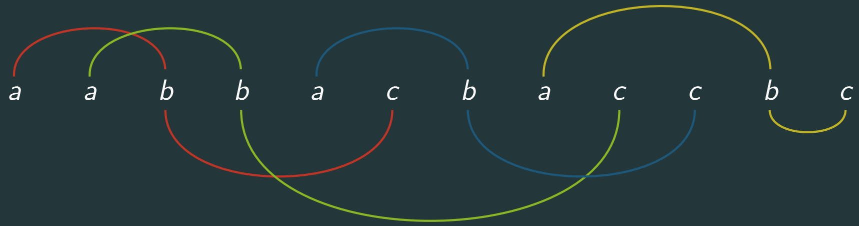
Straddling counter-example

G_0: Grammar of triple insertions



Straddling counter-example

G_0: Grammar of triple insertions



Straddling counter-example

Notation

$$\mathcal{O}_m[\![\textit{conclusion} \leftarrow \textit{premises} \mid \{\textit{partial orderings of inserted elements}\}]\!].$$

Meta-grammars: Introduction

Notation

$$\mathcal{O}_m[\![\mathit{conclusion} \leftarrow \mathit{premises} \mid \{\mathit{partial\ orderings\ of\ inserted\ elements}\}]\!].$$

Meta-grammar G_1

$$\left.\begin{array}{l} \mathcal{O}_2[\![W \leftarrow \epsilon \mid \{a < b < c\}]\!]. \\ \mathcal{O}_2[\![W \leftarrow W_{xy} \mid \{x < y,\ a < b < c\}]\!]. \end{array}\right\} \begin{array}{c} \text{TRIPLE} \\ \text{INSERTION} \end{array}$$

Meta-grammars: Introduction

Notation

$$\mathcal{O}_m[\![\textit{conclusion} \leftarrow \textit{premises} \mid \{\textit{partial orderings of inserted elements}\}]\!].$$

Meta-grammar G_1

$$\left.\begin{array}{l}\mathcal{O}_2[\![W \leftarrow \epsilon \mid \{a < b < c\}]\!].\\ \mathcal{O}_2[\![W \leftarrow W_{xy} \mid \{x < y,\ a < b < c\}]\!].\end{array}\right\} \quad \begin{array}{c}\text{TRIPLE}\\ \text{INSERTION}\end{array}$$

$$+$$

$$\mathcal{O}_2[\![W \leftarrow W_{xy}, W_{zw} \mid \{x < y,\ z < w\}]\!].\ \Big\} \quad \begin{array}{c}\text{INTERLEAVING}\\ \text{WORDS}\end{array}$$

$\mathbf{G}_2$: Adding states

$$\left.\begin{array}{l}\mathcal{O}_2[\![A^+ \leftarrow \epsilon \mid \{a\}]\!].\\ \mathcal{O}_2[\![B^+ \leftarrow \epsilon \mid \{b\}]\!].\\ \mathcal{O}_2[\![C^+ \leftarrow \epsilon \mid \{c\}]\!].\end{array}\right\}\text{BASE CASES}$$

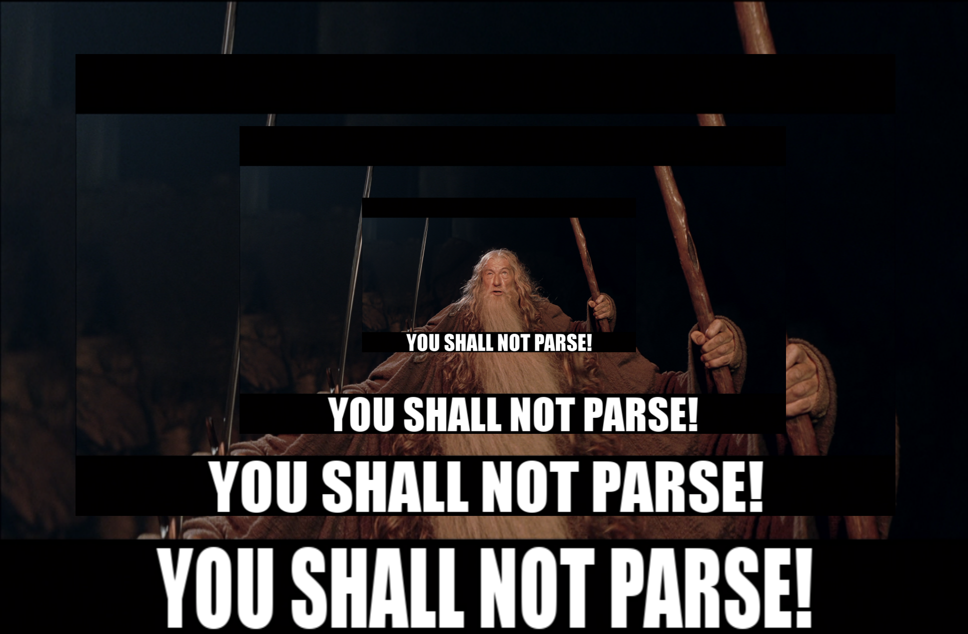
$$\left.\begin{array}{l}\mathcal{O}_2[\![C^- \leftarrow A^+, B^+ \mid \{x < y < z < w\}]\!].\\ \mathcal{O}_2[\![B^- \leftarrow A^+, C^+ \mid \{x < y < z < w\}]\!].\\ \mathcal{O}_2[\![A^- \leftarrow B^+, C^+ \mid \{x < y < z < w\}]\!].\\ \mathcal{O}_2[\![A^+ \leftarrow C^-, B^- \mid \{x < y < z < w\}]\!].\\ \mathcal{O}_2[\![B^+ \leftarrow C^-, A^- \mid \{x < y < z < w\}]\!].\\ \mathcal{O}_2[\![C^+ \leftarrow B^-, A^- \mid \{x < y < z < w\}]\!].\\ \forall\ \mathrm{K} \in \mathcal{S} \setminus \mathrm{W} :\ \mathcal{O}_2[\![\mathrm{K} \leftarrow \mathrm{K}_{xy}, \mathrm{W}_{zw} \mid \{x < y,\ z < w\}]\!].\end{array}\right\}\text{COMBINATIONS}$$

$$\left.\begin{array}{l}\mathcal{O}_2[\![W \leftarrow A^+, A^- \mid \{x < y < z < w\}]\!].\\ \mathcal{O}_2[\![W \leftarrow C^-, C^+ \mid \{x < y < z < w\}]\!].\end{array}\right\}\text{CLOSURES}$$

G_3: G_2 + Universal triple insertion

$$\mathsf{G}_3 = \mathsf{G}_2 + \forall\, \mathrm{K} \in \mathcal{S} \setminus \mathrm{W} : \quad \mathcal{O}_2[\![\mathrm{K} \leftarrow \mathrm{K}_{xy} \mid \{x < y,\ a < b < c\}]\!].$$

Demo


YOU SHALL NOT PARSE!
YOU SHALL NOT PARSE!
YOU SHALL NOT PARSE!
YOU SHALL NOT PARSE!

Refining states

Example

$$\begin{tikzcd}[column sep=small]
& A^+(x, y) \arrow[dl] \arrow[dr] & \\
A^+_{left} & & A^+_{right}
\end{tikzcd}$$

$$\vdots$$

$$\begin{tikzcd}[column sep=small]
& C^-(x, y) \arrow[dl] \arrow[d] \arrow[dr] & \\
C^-_{left} & C^-_{left,\ right} & C^-_{right}
\end{tikzcd}$$

Why?

Refining states

Example

$$A^+(x,y)$$

$$\swarrow \qquad \searrow$$

$$A^+_{left} \qquad\qquad A^+_{right}$$

$$\vdots$$

$$C^-(x,y)$$

$$\swarrow \qquad \downarrow \qquad \searrow$$

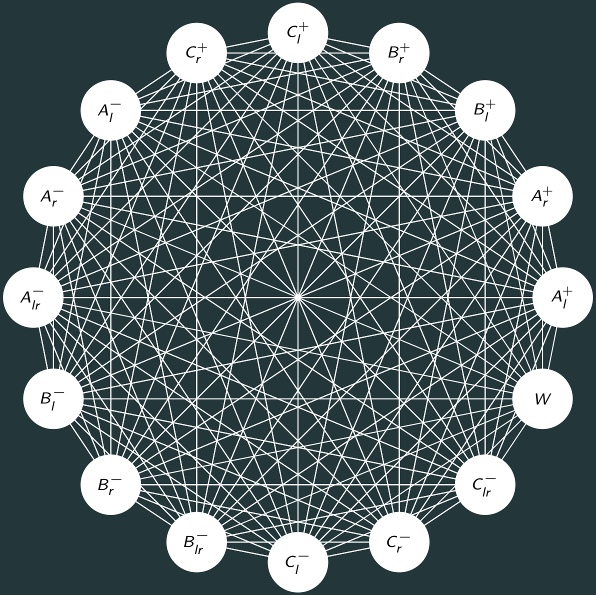
$$C^-_{left} \qquad C^-_{left,\ right} \qquad C^-_{right}$$

Why?

New orders in interactions

$$C^-(xz,\ wy) \leftarrow A^+_{left}(x,\ y),\ B^+_{left}(z,\ w).$$

Refining states: Interactions



G_4: Automatic Rule Inference

State descriptors $\mathcal{D}$

$$
\begin{aligned}
\mathrm{W} &\mapsto (\epsilon, \epsilon) & \mathrm{A}_r^- &\mapsto (\epsilon, bc) \\
\mathrm{A}_l^+ &\mapsto (a, \epsilon) & \mathrm{A}_{lr}^- &\mapsto (b, c) \\
\mathrm{A}_r^+ &\mapsto (\epsilon, a) & \mathrm{B}_l^- &\mapsto (ac, \epsilon) \\
\mathrm{B}_l^+ &\mapsto (b, \epsilon) & \mathrm{B}_r^- &\mapsto (\epsilon, ac) \\
\mathrm{B}_r^+ &\mapsto (\epsilon, b) & \mathrm{B}_{lr}^- &\mapsto (a, c) \\
\mathrm{C}_l^+ &\mapsto (c, \epsilon) & \mathrm{C}_l^- &\mapsto (ab, \epsilon) \\
\mathrm{C}_r^+ &\mapsto (\epsilon, c) & \mathrm{C}_r^- &\mapsto (\epsilon, ab) \\
\mathrm{A}_l^- &\mapsto (bc, \epsilon) & \mathrm{C}_{lr}^- &\mapsto (a, b)
\end{aligned}
$$

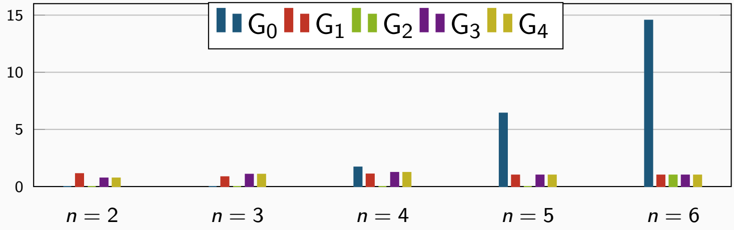
Automatic Rule Inference: Example

$$\textsc{Case of } \mathrm{A}^{-}_{lr}(x_b,\ y_c) + \mathrm{B}^{-}_{lr}(x_a,\ y_c)$$

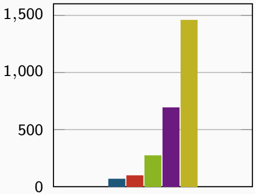
$$\begin{tikzcd}
\text{PERMUTATION} & \text{DESCRIPTOR} & \text{ELIMINATED} \\
\vdots & \vdots & \vdots \\
 & & (c,\ \epsilon) \mapsfrom \mathrm{C}^{+}_{l} \\
(zxw,\ y) \arrow[r] & (abc,\ c) \arrow[ur] \arrow[dr] & \\
 & & (c,\ \epsilon) \mapsfrom \mathrm{C}^{+}_{r} \\
\vdots & \vdots & \vdots \\
(xzw,\ y) \arrow[r] & (bac,\ c) \arrow[r] & (bac,\ c) \notin \mathcal{D} \\
\vdots & \vdots & \vdots
\end{tikzcd}$$

Results

Counter-examples (%)



Rule size



Computation Time (sec)