

Logic and Language: Exercise (Week 6)

Orestis Melkonian [6176208], Konstantinos Kogkalidis [6230067]

1 Syntax

1.1

First, we define the rules of *rightward extraction* $\widehat{\alpha}_\diamond^r, \widehat{\sigma}_\diamond^r$:

$$\frac{f : A \otimes (B \otimes \diamond C) \rightarrow D}{\widehat{\alpha}_\diamond^r f : (A \otimes B) \otimes \diamond C \rightarrow D} \qquad \frac{f : (A \otimes \diamond C) \otimes B \rightarrow D}{\widehat{\sigma}_\diamond^r f : (A \otimes B) \otimes \diamond C \rightarrow D}$$

We can now proceed with the derivation of

$$n \otimes ((n \setminus n) / (s / \diamond \square np)) \otimes ((np / n) \otimes n) \otimes ((np \setminus s) / np) \rightarrow n$$

as follows:

$$\frac{\overline{np \vdash np} \quad 1_{np} \quad \overline{n \vdash n} \quad 1_n}{\overline{np/n \vdash np/n}} /$$
$$\frac{(np/n) \otimes n \vdash np}{(np/n) \otimes n \vdash np} \triangleright^{-1} \quad \frac{s \vdash s}{s \vdash s} 1_s \quad \frac{\overline{np \vdash np} \quad 1_{np}}{\square np \vdash \square np} \square$$
$$\frac{\overline{np \setminus s \vdash ((np/n) \otimes n) \setminus s}}{\overline{np \setminus s \vdash ((np/n) \otimes n) \setminus s}} \searrow \quad \frac{\overline{\diamond np \vdash np}}{\diamond np \vdash np} \nabla^{-1}$$
$$\frac{\overline{(np \setminus s)/np \vdash (((np/n) \otimes n) \setminus s)/\diamond \square np}}{\overline{(np \setminus s)/np \vdash (((np/n) \otimes n) \setminus s)/\diamond \square np}} /$$
$$\frac{\overline{(np \setminus s)/np \otimes \diamond \square np \vdash ((np/n) \otimes n) \setminus s}}{\overline{(np \setminus s)/np \otimes \diamond \square np \vdash ((np/n) \otimes n) \setminus s}} \triangleright^{-1}$$
$$\frac{\overline{((np/n) \otimes n) \otimes ((np \setminus s)/np \otimes \diamond \square np) \vdash s}}{\overline{((np/n) \otimes n) \otimes ((np \setminus s)/np) \otimes \diamond \square np \vdash s}} \triangleleft^{-1}$$
$$\frac{\overline{(((np/n) \otimes n) \otimes ((np \setminus s)/np)) \otimes \diamond \square np \vdash s}}{\overline{((np/n) \otimes n) \otimes ((np \setminus s)/np) \vdash s/\diamond \square np}} \widehat{\alpha}_{\diamond}^r$$
$$\frac{\overline{n \vdash n} \quad 1_n \quad \overline{n \vdash n} \quad 1_n}{\overline{n \setminus n \vdash n \setminus n}} \searrow \quad \frac{\overline{((np/n) \otimes n) \otimes ((np \setminus s)/np) \vdash s/\diamond \square np}}{\overline{((np/n) \otimes n) \otimes ((np \setminus s)/np) \vdash s/\diamond \square np}} \triangleright$$
$$\frac{\overline{(n \setminus n)/(s/\diamond \square np) \vdash (n \setminus n)/(((np/n) \otimes n) \otimes (np \setminus s)/np)}}{\overline{(n \setminus n)/(s/\diamond \square np) \vdash (n \setminus n)/(((np/n) \otimes n) \otimes (np \setminus s)/np)}} /$$
$$\frac{\overline{(n \setminus n)/(s/\diamond \square np) \otimes (((np/n) \otimes (np \setminus s)/np) \vdash n \setminus n)}}{\overline{n \otimes ((n \setminus n)/(s/\diamond \square np) \otimes (((np/n) \otimes n) \otimes (np \setminus s)/np)) \vdash n}} \triangleleft^{-1}$$

2 Interpretation

2.1

We start by assigning a temporary variable at each rule application in the proof tree:

$$\begin{array}{c}
\frac{\overline{np \vdash np} \quad \overline{1_{np} \quad n \vdash n}}{f : np/n \vdash np/n} \quad \frac{\overline{1_n}}{\quad} / \\
\frac{g : (np/n) \otimes n \vdash np}{h : np \vdash ((np/n) \otimes n) \setminus s} \quad \frac{\overline{s \vdash s} \quad \overline{1_s} \quad \frac{\overline{np \vdash np} \quad \overline{1_{np}}}{b : \Box np \vdash \Box np} \quad \square}{d : \Diamond np \vdash np} \quad \nabla^{-1} \\
\frac{\quad}{i : (np \setminus s)/np \vdash (((np/n) \otimes n) \setminus s) / \Diamond \Box np} \quad \frac{\quad}{j : (np \setminus s)/np \otimes \Diamond \Box np \vdash ((np/n) \otimes n) \setminus s} \quad \triangleright^{-1} \\
\frac{\quad}{k : ((np/n) \otimes n) \otimes ((np \setminus s)/np \otimes \Diamond \Box np) \vdash s} \quad \frac{\quad}{l : (((np/n) \otimes n) \otimes ((np \setminus s)/np)) \otimes \Diamond \Box np \vdash s} \quad \triangleleft^{-1} \\
\frac{\overline{n \vdash n} \quad \overline{1_n} \quad \overline{n \vdash n} \quad \overline{1_n}}{m : n \setminus n \vdash n \setminus n} \quad \frac{\quad}{n : ((np/n) \otimes n) \otimes ((np \setminus s)/np) \vdash s / \Diamond \Box np} \quad \widehat{\alpha}_{\Diamond}^r \quad \triangleright \\
\frac{\quad}{o : (n \setminus n) / (s / \Diamond \Box np) \vdash (n \setminus n) / (((np/n) \otimes n) \otimes (np \setminus s) / np)} \quad \frac{\quad}{p : (n \setminus n) / (s / \Diamond \Box np) \otimes (((np/n) \otimes (np \setminus s) / np) \vdash n \setminus n)} \quad \triangleright^{-1} \\
\frac{\quad}{q : n \otimes ((n \setminus n) / (s / \Diamond \Box np) \otimes (((np/n) \otimes n) \otimes (np \setminus s) / np)) \vdash n} \quad \triangleleft^{-1}
\end{array}$$

We now work our way top-down through the proof-tree, writing the interpretation of each formula using the rules of 3.1.

$$f : (1_{np} \otimes \eta_n \otimes 1_n) \circ (1_{np \otimes n} \otimes 1_n \otimes 1_n) \circ (1_{np \otimes n} \otimes \epsilon_n) \quad (1)$$

$$g \equiv \triangleright^{-1} f : (f \otimes 1_n) \circ (1_{np} \otimes \epsilon_n) \quad (2)$$

$$h : (1_{np} \otimes \eta_{(np/n) \otimes n} \otimes 1_s) \circ (1_{np} \otimes g \otimes 1_{((np/n) \otimes n) \otimes s}) \circ (\epsilon_{np} \otimes 1_{((np/n) \otimes n) \otimes s}) \quad (3)$$

$$i : (h \otimes \eta_{np} \otimes 1_{np}) \circ (1_{((np/n) \otimes n) \setminus s} \otimes d \otimes 1_{np}) \otimes (1_{((np/n) \otimes n) \setminus s} \otimes \epsilon_s) \quad (4)$$

$$j \equiv \triangleright^{-1} i : (i \otimes 1_{\diamond \square np}) \circ (1_{((np/n) \otimes n) \setminus s} \otimes \epsilon_{\diamond \square np}) \quad (5)$$

$$k \equiv \triangleleft^{-1} i : (1_{(np/n) \otimes n} \otimes i) \circ (\epsilon_{(np/n) \otimes n} \otimes 1_s) \quad (6)$$

$$l \equiv \hat{\alpha}_{\diamond}^r k : \alpha \circ k \quad (7)$$

$$m : (1_n \otimes \eta_n \otimes 1_n) \circ (1_n \otimes 1_n \otimes 1_{n \otimes n}) \circ (\epsilon_n \otimes 1_{n \otimes n}) \quad (8)$$

$$n \equiv \triangleright l : (1_{(np/n) \otimes n} \otimes \eta_{(np \setminus s)/np}) \circ (l \otimes 1_{(np \setminus s)/np}) \quad (9)$$

$$o : (m \otimes \eta_{(((np/n) \otimes n) \otimes (np \setminus s)/np)} \otimes 1_{(s/\diamond \square np)}) \circ (1_{(n \setminus n) \otimes (((np/n) \otimes n) \otimes (np \setminus s)/np)} \otimes n \otimes 1_{s/\diamond \square n}) \quad (10)$$

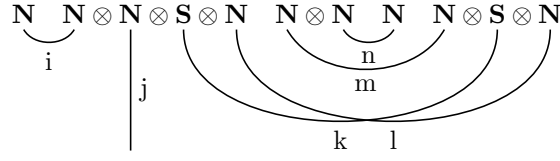
Recursively unwrapping the above interpretations, we obtain the final interpretation:

2.2

By working our way from the leaves of the proof tree, we get the following generalized Kronecker delta:

$$\begin{array}{l} \text{island}_i \otimes \text{that}_{j,k,l,m} \otimes \text{the}_{n,o} \otimes \text{hurricane}_p \otimes \text{destroyed}_{q,r,s} \xrightarrow{\delta_{j,t,r,s,q,p}^{i,k,l,m,n,o}} \mathbf{v}_r^{obj} \in \mathbf{N} \\ \mathbf{v}_r^{obj} = \text{island}_i \otimes \text{that}_{i,j,k,l} \otimes \text{the}_{m,n} \otimes \text{hurricane}_n \otimes \text{destroyed}_{m,k,l} \quad (\text{relabelled}) \end{array}$$

We give the matching diagram in the figure below:



2.3

In order to calculate the semantic value for the relative clause body 'the hurricane destroyed', we first apply **the**_{MN} to **hurricane**_N. The operation yields the noun-phrase **the hurricane**_M, represented by a row-vector equal to that of **hurricane**. The verb **destroyed**_{MKL} is then applied to the resulting vector, thereupon we obtain the final result **the hurricane destroyed**_{KL}. Concretely, **the hurricane destroyed**_{KL} = **destroyed**_{MKL}(**the**_{MN}**hurricane**_N) is a 2 by 3 matrix, the elements of which are:

$$\begin{pmatrix} 12 & -19 & 3 \\ 5 & 10 & 1 \end{pmatrix}$$

given by:

$$\mathbf{the_hurricane_destroyed}(k, l) = \sum_{m \in M} \mathbf{hurricane}(m) \times \mathbf{destroyed}(m, k, l), \forall k \in K, l \in L$$

The corresponding Python code is given below:

2.4

The interpreted type for the relative pronoun is:

$$[(n \setminus n) / (s / \Diamond \Box np)] = [n \setminus n] \otimes [s / \Diamond \Box np] = [n] \otimes [n] \otimes [s] \otimes [\Diamond \Box np] = N \otimes N \otimes S \otimes N$$

We can now give the following Frobenius recipe for **that**:

$$I \cong I \otimes I \xrightarrow{\eta_N \otimes \eta_N} N \otimes N \otimes N \otimes N \cong N \otimes N \otimes N \otimes I \otimes N \xrightarrow{1_N \otimes \mu_N \otimes \zeta_S \otimes 1_N} N \otimes N \otimes S \otimes N$$

In order to obtain the final interpretation, we do the following (dictated from the above recipe):

1. Reduce the rank of the transitive verb by summing over the S component, thus obtaining the following matrix:

$$\mathbf{collapsed_destroyed} = \begin{pmatrix} \begin{pmatrix} 9 & 2 & 3 \end{pmatrix} \\ \begin{pmatrix} 1 & -5 & 2 \end{pmatrix} \\ \begin{pmatrix} -1 & -8 & 1 \end{pmatrix} \end{pmatrix}$$

2. Apply **collapsed_destroyed** to **the_hurricane** in subject position:

$$\mathbf{the_hurricane_destroyed} = \begin{pmatrix} 17 & -9 & 4 \end{pmatrix}$$

3. Multiply **the_hurricane_destroyed** element-wise with **island**:

$$\mathbf{island_that_the_hurricane_destroyed} = \begin{pmatrix} -85 & -36 & 0 \end{pmatrix}$$

The corresponding Python code is given below: