D^3 AS A 2-MCFL

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August 8, 2018

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2-MCFG

Generalization of the CFG over tuples of strings

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N-dimensional Dyck Language D^N

Defined over an ordered alphabet of N symbols: $\{\alpha_1 < \cdots < \alpha_N\}$ s.t. words satisfy two conditions:

1. Equal number of occurrences of all alphabet symbols
2. Any prefix of a word must contain at least as many α_i as $\alpha_{i+1} \quad \forall\, i \ \leq N-1$

D^3 - Some examples

Dyck words

- abc
- aabbcc
- abcabcabacbc

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$$a\,b\,a\,b\,a\,c\,b\,c\,a\,b\,c\,c$$

First-match policy

D^3 - Some examples

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$$a b a b a c b c a b c c$$

First-match policy

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First-match policy

Motivation

Natural Languages

Free word order respecting linear order constraints

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Programming Languages

Static Analysis of non-standard control flows (e.g. *yield*)

G_0: Grammar of triple insertions

$$S(xy) \leftarrow W(x, y). \tag{1}$$

$$W(\epsilon, xy\mathbf{abc}) \leftarrow W(x, y). \tag{2}$$

$$W(\epsilon, x\mathbf{a}y\mathbf{bc}) \leftarrow W(x, y). \tag{3}$$

$$\ldots$$

$$W(\mathbf{ab}x\mathbf{c}y, \epsilon) \leftarrow W(x, y). \tag{60}$$

$$W(\mathbf{abc}xy, \epsilon) \leftarrow W(x, y). \tag{61}$$

$$W(\epsilon, \mathbf{abc}). \tag{62}$$

$$W(\mathbf{a}, \mathbf{bc}). \tag{63}$$

$$W(\mathbf{ab}, \mathbf{c}). \tag{64}$$

$$W(\mathbf{abc}, \epsilon). \tag{65}$$

G_0: Grammar of triple insertions

$a \quad a \quad b \quad b \quad a \quad c \quad b \quad a \quad c \quad c \quad b \quad c$

Straddling counter-example

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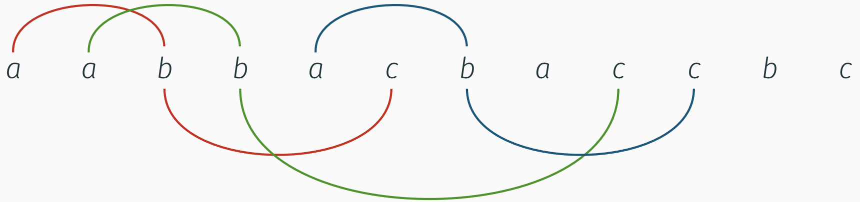

Straddling counter-example

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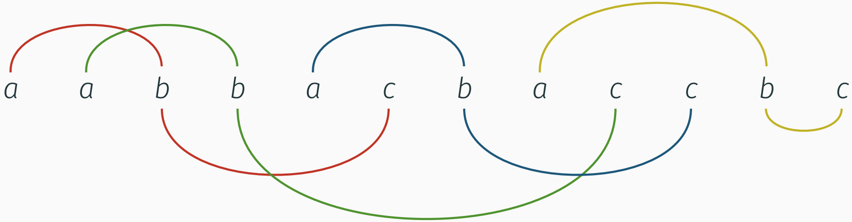

Straddling counter-example

Notation

$$\mathcal{O}_m[\![\text{conclusion} \leftarrow \text{premises} \mid \text{partial orders}]\!].$$

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Meta-grammar G_1

$$\left.\begin{array}{l} \mathcal{O}_2[\![W \leftarrow \epsilon \mid \{a < b < c\}]\!]. \\ \mathcal{O}_2[\![W \leftarrow W_{xy} \mid \{x < y,\ a < b < c\}]\!]. \end{array}\right\} \quad \begin{array}{c}\text{TRIPLE}\\ \text{INSERTION}\end{array}$$

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$+$

$$\mathcal{O}_2[\![W \leftarrow W_{xy}, W_{zw} \mid \{x < y,\ z < w\}]\!]. \quad \Big\} \quad \text{INTERLEAVING WORDS}$$

G_2: Adding states

$$\left.\begin{array}{l}\mathcal{O}_2[\![A^+ \leftarrow \epsilon \mid \{a\}]\!].\\ \mathcal{O}_2[\![B^+ \leftarrow \epsilon \mid \{b\}]\!].\\ \mathcal{O}_2[\![C^+ \leftarrow \epsilon \mid \{c\}]\!].\end{array}\right\} \text{BASE CASES}$$

$$\left.\begin{array}{l}\mathcal{O}_2[\![C^- \leftarrow A^+, B^+ \mid \{x < y < z < w\}]\!].\\ \mathcal{O}_2[\![B^- \leftarrow A^+, C^+ \mid \{x < y < z < w\}]\!].\\ \mathcal{O}_2[\![A^- \leftarrow B^+, C^+ \mid \{x < y < z < w\}]\!].\\ \mathcal{O}_2[\![A^+ \leftarrow C^-, B^- \mid \{x < y < z < w\}]\!].\\ \mathcal{O}_2[\![B^+ \leftarrow C^-, A^- \mid \{x < y < z < w\}]\!].\\ \mathcal{O}_2[\![C^+ \leftarrow B^-, A^- \mid \{x < y < z < w\}]\!].\\ \forall\, \mathrm{K} \in \mathcal{S} \setminus \mathrm{W} :\ \mathcal{O}_2[\![\mathrm{K} \leftarrow \mathrm{K}_{xy}, \mathrm{W}_{zw} \mid \{x < y,\ z < w\}]\!].\end{array}\right\} \text{COMBINATIONS}$$

$$\left.\begin{array}{l}\mathcal{O}_2[\![W \leftarrow A^+, A^- \mid \{x < y < z < w\}]\!].\\ \mathcal{O}_2[\![W \leftarrow C^-, C^+ \mid \{x < y < z < w\}]\!].\end{array}\right\} \text{CLOSURES}$$

G_3: G_2 + Universal triple insertion

$$\mathsf{G}_3 = \mathsf{G}_2 + \forall\, \mathsf{K} \in \mathcal{S} \setminus \mathsf{W} : \quad \mathcal{O}_2[\![\mathsf{K} \leftarrow \mathsf{K}_{xy} \mid \{x < y,\ a < b < c\}]\!].$$

YOU SHALL NOT PARSE!

Example

$$\begin{tikzcd}[column sep=small] & A^{+}(x, y) \arrow[dl] \arrow[dr] & \\ A^{+}_{\mathit{left}} & & A^{+}_{\mathit{right}} \end{tikzcd}$$

$$\vdots$$

$$\begin{tikzcd}[column sep=small] & C^{-}(x, y) \arrow[dl] \arrow[d] \arrow[dr] & \\ C^{-}_{\mathit{left}} & C^{-}_{\mathit{left},\ \mathit{right}} & C^{-}_{\mathit{right}} \end{tikzcd}$$

Refining states

Example

$$
\begin{tikzcd}[column sep=small]
 & A^{+}(x, y) \arrow[dl] \arrow[dr] & \\
A^{+}_{left} & & A^{+}_{right}
\end{tikzcd}
$$

$$\vdots$$

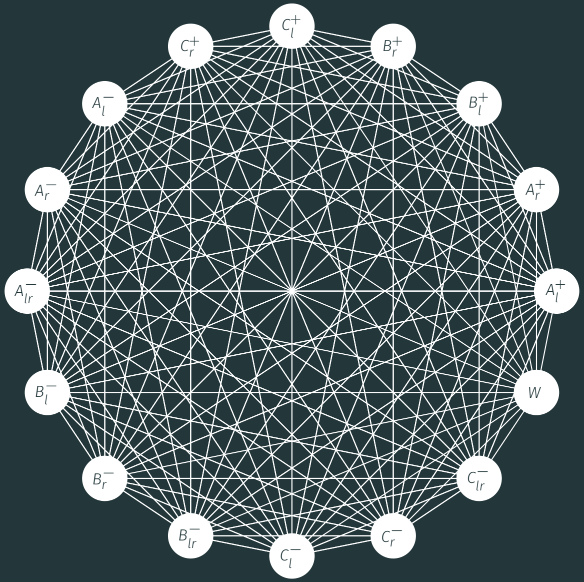
$$
\begin{tikzcd}[column sep=small]
 & C^{-}(x, y) \arrow[dl] \arrow[d] \arrow[dr] & \\
C^{-}_{left} & C^{-}_{left,\ right} & C^{-}_{right}
\end{tikzcd}
$$

Why?

New orders in interactions

$$C^{-}(xz,\ wy) \leftarrow A^{+}_{left}(x,\ y),\ B^{+}(z,\ w).$$

Refining states: Interactions



G_4: Automatic Rule Inference

State descriptors $\mathcal{D}$

$$\mathrm{W} \mapsto (\epsilon, \epsilon)$$
$$\mathrm{A}_l^+ \mapsto (a, \epsilon)$$
$$\mathrm{A}_r^+ \mapsto (\epsilon, a)$$
$$\mathrm{B}_l^+ \mapsto (b, \epsilon)$$
$$\mathrm{B}_r^+ \mapsto (\epsilon, b)$$
$$\mathrm{C}_l^+ \mapsto (c, \epsilon)$$
$$\mathrm{C}_r^+ \mapsto (\epsilon, c)$$
$$\mathrm{A}_l^- \mapsto (bc, \epsilon)$$
$$\mathrm{A}_r^- \mapsto (\epsilon, bc)$$
$$\mathrm{A}_{lr}^- \mapsto (b, c)$$
$$\mathrm{B}_l^- \mapsto (ac, \epsilon)$$
$$\mathrm{B}_r^- \mapsto (\epsilon, ac)$$
$$\mathrm{B}_{lr}^- \mapsto (a, c)$$
$$\mathrm{C}_l^- \mapsto (ab, \epsilon)$$
$$\mathrm{C}_r^- \mapsto (\epsilon, ab)$$
$$\mathrm{C}_{lr}^- \mapsto (a, b)$$

Automatic Rule Inference: Example

Case of $A^{-}_{lr}(x_b, y_c) + B^{-}_{lr}(z_a, w_c)$

permutation

$$\vdots$$

$$(zxw, y)$$

$$\vdots$$

$$(xzw, y)$$

$$\vdots$$

Automatic Rule Inference: Example

Case of $A^{-}_{lr}(x_b,\ y_c) + B^{-}_{lr}(z_a,\ w_c)$

Permutation		Descriptor
$\vdots$		$\vdots$
$(zxw,\ y)$	$\longrightarrow$	$(abc,\ c)$
$\vdots$		$\vdots$
$(xzw,\ y)$	$\longrightarrow$	$(bac,\ c)$
$\vdots$		$\vdots$

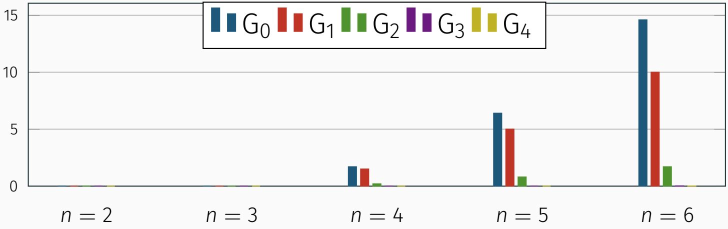
Automatic Rule Inference: Example

Case of $\mathsf{A}^{-}_{lr}(x_b,\ y_c) + \mathsf{B}^{-}_{lr}(z_a,\ w_c)$

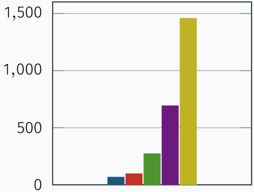
$$
\begin{array}{ccccc}
\underline{\text{PERMUTATION}} & & \underline{\text{DESCRIPTOR}} & & \underline{\text{ELIMINATED}} \\
\vdots & & \vdots & & \vdots \\
 & & & & (c,\ \epsilon) \hookleftarrow \mathsf{C}^{+}_{l} \\
(zxw,\ y) & \longrightarrow & (abc,\ c) & \rightrightarrows & \\
 & & & & (\epsilon,\ c) \hookleftarrow \mathsf{C}^{+}_{r} \\
\vdots & & \vdots & & \vdots \\
(xzw,\ y) & \longrightarrow & (bac,\ c) & \longrightarrow & (bac,\ c) \notin \mathcal{D} \\
\vdots & & \vdots & & \vdots
\end{array}
$$

Results

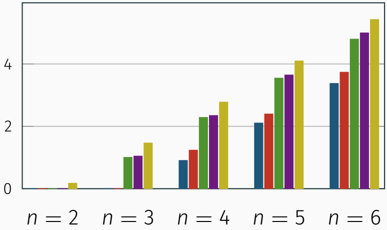
Counter-examples (%)



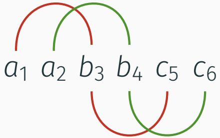
Rule size



Computation Time (in log(sec))



Correspondences: Promotion on Young Tableaux



1	2
3	4
5	6

·	1
2	3
4	5

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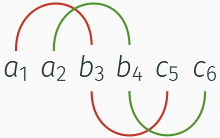

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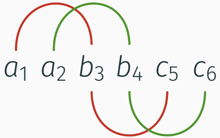

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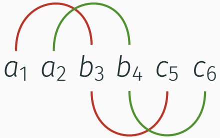

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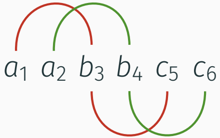

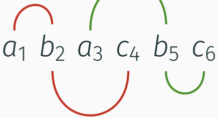

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Correspondences: Spider Webs

Growth rules