

Logic and Language: Exercise (Week 3)

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1 LGa

1.1

1. Unfolding the combinator f/g , we get the following derivation:

$$\begin{array}{c}
 \frac{1_{A/B} : A/B \longrightarrow A/B}{\Downarrow^{-1} 1_{A/B} : A/B \otimes B \longrightarrow A} \quad f : A \longrightarrow A' \quad \frac{g : B \longrightarrow B' \quad \frac{1_{A/B'} : A/B \longrightarrow A/B'}{\Downarrow^{-1} 1_{A/B'} : (A/B') \otimes B' \longrightarrow A}}{\Downarrow^{-1} 1_{A/B'} : B' \longrightarrow (A/B') \setminus A}} \\
 \frac{f \circ (\Downarrow^{-1} 1_{A/B}) : A/B \otimes B \longrightarrow A'}{\Downarrow(f \circ (\Downarrow^{-1} 1_{A/B})) : A/B \longrightarrow A'/B} \quad \frac{(\Downarrow \Downarrow^{-1} 1_{A/B'}) \circ g : B \longrightarrow (A/B') \setminus A}{\Downarrow^{-1}((\Downarrow \Downarrow^{-1} 1_{A/B'}) \circ g) : (A/B') \otimes B \longrightarrow A}} \\
 \frac{\Downarrow(f \circ (\Downarrow^{-1} 1_{A/B})) : A/B \longrightarrow A'/B \quad \Downarrow \Downarrow^{-1}((\Downarrow \Downarrow^{-1} 1_{A/B'}) \circ g) : A/B' \longrightarrow A/B}{(\Downarrow(f \circ (\Downarrow^{-1} 1_{A/B}))) \circ (\Downarrow \Downarrow^{-1}((\Downarrow \Downarrow^{-1} 1_{A/B'}) \circ g)) : A/B' \longrightarrow A'/B}
 \end{array}$$

2. Applying the arrow reversal transformation $(\cdot)^\dagger$ to f/g , we get the reverse combinator:

$$\begin{aligned}
 (f/g)^\dagger &= ((\Downarrow(f \circ (\Downarrow^{-1} 1_{A/B}))) \circ (\Downarrow \Downarrow^{-1}((\Downarrow \Downarrow^{-1} 1_{A/B'}) \circ g)))^\dagger \\
 &= (\Downarrow \Downarrow^{-1}((\Downarrow \Downarrow^{-1} 1_{A/B'}) \circ g))^\dagger \circ (\Downarrow(f \circ (\Downarrow^{-1} 1_{A/B})))^\dagger \\
 &= (\Downarrow \Downarrow^{-1}((\Downarrow \Downarrow^{-1} 1_{A/B'}) \circ g))^\dagger \circ (\Downarrow(f \circ (\Downarrow^{-1} 1_{A/B})))^\dagger \\
 &= (\Downarrow \Downarrow^{-1}((\Downarrow \Downarrow^{-1} 1_{A/B'}) \circ g))^\dagger \circ (\Downarrow((\Downarrow^{-1} 1_{A/B})^\dagger \circ f^\dagger)) \\
 &= (\Downarrow \Downarrow^{-1}(g^\dagger \circ (\Downarrow \Downarrow^{-1} 1_{A/B'})^\dagger))^\dagger \circ (\Downarrow((\Downarrow^{-1} 1_{A/B})^\dagger \circ f^\dagger)) \\
 &= (\Downarrow \Downarrow^{-1}(g^\dagger \circ (\Downarrow \Downarrow^{-1} 1_{A/B'})^\dagger))^\dagger \circ (\Downarrow((\Downarrow^{-1} 1_{B \otimes A}) \circ f^\dagger)) \\
 &= (\Downarrow \Downarrow^{-1}(g^\dagger \circ (\Downarrow \Downarrow^{-1} 1_{(A/B')^\dagger}))^\dagger) \circ (\Downarrow((\Downarrow^{-1} 1_{B \otimes A}) \circ f^\dagger)) \\
 &= (\Downarrow \Downarrow^{-1}(g^\dagger \circ (\Downarrow \Downarrow^{-1} 1_{B' \otimes A})))^\dagger \circ (\Downarrow((\Downarrow^{-1} 1_{B \otimes A}) \circ f^\dagger))
 \end{aligned}$$

Unfolding the combinator f/g , we get the following derivation:

$$\begin{array}{c}
\frac{1_{B' \otimes A} : B' \otimes A \longrightarrow B' \otimes A}{\blacktriangleleft^{-1} 1_{B' \otimes A} : A \longrightarrow B' \oplus (B' \otimes A)} \\
\frac{\blacktriangleright \blacktriangleleft^{-1} 1_{B' \otimes A} : A \otimes (B' \otimes A) \longrightarrow B' \quad g^\dagger : B' \longrightarrow B}{g^\dagger \circ (\blacktriangleright \blacktriangleleft^{-1} 1_{B' \otimes A}) : A \otimes (B' \otimes A) \longrightarrow B} \\
\frac{\blacktriangleright^{-1} (g^\dagger \circ (\blacktriangleright \blacktriangleleft^{-1} 1_{B' \otimes A})) : A \longrightarrow B \oplus (B' \otimes A)}{\blacktriangleleft \blacktriangleright^{-1} (g^\dagger \circ (\blacktriangleright \blacktriangleleft^{-1} 1_{B' \otimes A})) : B \otimes A \longrightarrow B' \otimes A} \\
\frac{(\blacktriangleleft \blacktriangleright^{-1} (g^\dagger \circ (\blacktriangleright \blacktriangleleft^{-1} 1_{B' \otimes A}))) \circ (\blacktriangleleft ((\blacktriangleleft^{-1} 1_{B \otimes A}) \circ f^\dagger)) : B \otimes A' \longrightarrow B' \otimes A}{\quad}
\end{array}
\quad
\begin{array}{c}
\frac{1_{B \otimes A} : B \otimes A \longrightarrow B \otimes A}{\blacktriangleleft^{-1} 1_{B \otimes A} : A \longrightarrow B \oplus (B \otimes A)} \\
\frac{f^\dagger : A' \longrightarrow A \quad \blacktriangleleft^{-1} 1_{B \otimes A} : A \longrightarrow B \oplus (B \otimes A)}{(\blacktriangleleft^{-1} 1_{B \otimes A}) \circ f^\dagger : A' \longrightarrow B \oplus (B \otimes A)} \\
\frac{(\blacktriangleleft^{-1} 1_{B \otimes A}) \circ f^\dagger : A' \longrightarrow B \oplus (B \otimes A)}{\blacktriangleleft ((\blacktriangleleft^{-1} 1_{B \otimes A}) \circ f^\dagger) : B \otimes A' \longrightarrow B \otimes A}
\end{array}$$

1.2

1. $(a \oplus b)/c \longrightarrow a/(c \otimes b)$
2. $(b \oplus c) \otimes a \longrightarrow b \oplus (c \otimes a)$