

Logic and Language: Exercise (Week 5)

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1 LG: continuation semantics

1.1

For each term, we compute the ILL type by first observing the polarity of the sequent and then using the table to interpret complex types.

$$\begin{aligned} \llbracket \text{some} \rrbracket &= \llbracket np/n \rrbracket^\perp && \{ \text{Negative Hypothesis} \} \\ &= (\llbracket np \rrbracket^\perp \otimes \llbracket n \rrbracket)^\perp && \{ \llbracket A/B \rrbracket \text{ with } A \text{ and } B \text{ positive} \} \\ &= (np^\perp \otimes n)^\perp && \{ np \text{ and } n \text{ positive} \} \end{aligned}$$

$$\begin{aligned} \llbracket \text{popular} \rrbracket &= \llbracket n/n \rrbracket^\perp && \{ \text{Negative Hypothesis} \} \\ &= (\llbracket n \rrbracket^\perp \otimes \llbracket n \rrbracket)^\perp && \{ \llbracket A/B \rrbracket \text{ with } A \text{ and } B \text{ positive} \} \\ &= (n^\perp \otimes n)^\perp && \{ n \text{ positive} \} \end{aligned}$$

$$\begin{aligned} \llbracket \text{saint} \rrbracket &= \llbracket n \rrbracket && \{ \text{Positive Hypothesis} \} \\ &= n && \{ n \text{ positive} \} \end{aligned}$$

$$\begin{aligned} \llbracket \text{arrived} \rrbracket &= \llbracket np \backslash s \rrbracket^\perp && \{ \text{Negative Hypothesis} \} \\ &= (\llbracket np \rrbracket \otimes \llbracket s \rrbracket)^\perp && \{ \llbracket B \backslash A \rrbracket \text{ with } B \text{ positive and } A \text{ negative} \} \\ &= (np \otimes s^\perp)^\perp && \{ np \text{ positive, } s \text{ negative} \} \end{aligned}$$

$$\begin{aligned} \alpha &= \llbracket s \rrbracket \\ &= s^\perp && \{ s \text{ negative} \} \end{aligned}$$

$$\begin{aligned} z &= \llbracket np \rrbracket \\ &= np && \{ np \text{ positive} \} \end{aligned}$$

$$\begin{aligned} y &= \llbracket n \rrbracket \\ &= n && \{ n \text{ positive} \} \end{aligned}$$

The ILL types then are:

TERM	TYPE
[some]	$(np^\perp \otimes n)^\perp$
[popular]	$(n^\perp \otimes n)^\perp$
[saint]	n
[arrived]	$(np \otimes s^\perp)^\perp$
α	s^\perp
z	np
y	n

1.2

SOURCE TYPE	CONSTANT	$[\cdot]^\ell$
n/n	popular	$\lambda \langle c, y \rangle . (c \ (\lambda z. \ \wedge (y \ z) \ (\text{POPULAR } z)))$

1.3

1.

$$\begin{array}{c}
\frac{\frac{\frac{\alpha_1}{\boxed{\alpha_1 : n}} \vdash n \quad CoAx}{\boxed{\gamma_0 : n} \vdash n \quad CoAx} \quad \frac{\frac{[saint] \quad \alpha_1 : n \vdash \alpha_1 : n}{n \vdash \boxed{\lambda_{\alpha_1}.([saint] \quad \alpha_1) : n}} \quad \frac{\gamma_0}{\boxed{\gamma_0 : n} \vdash n \quad CoAx} \quad \frac{\frac{\frac{\langle \gamma_0, \lambda_{\alpha_1}.([saint] \quad \alpha_1) \rangle : n/n}{\boxed{\langle \gamma_0, \lambda_{\alpha_1}.([saint] \quad \alpha_1) \rangle : n/n} \vdash \gamma_0 : n \cdot / \cdot n}}{\frac{\beta_0}{\boxed{\beta_0 : np}} \vdash np \quad CoAx \quad \frac{([popular] \langle \gamma_0, \lambda_{\alpha_1}.([saint] \quad \alpha_1) \rangle) : n/n \vdash \gamma_0 : n \cdot / \cdot n}{(n/n) \cdot \otimes \cdot n \vdash \boxed{\lambda_{\gamma_0}.([popular] \langle \gamma_0, \lambda_{\alpha_1}.([saint] \quad \alpha_1) \rangle) : n}} \quad \frac{\frac{\frac{\langle \beta_0, \lambda_{\gamma_0}.([popular] \langle \gamma_0, \lambda_{\alpha_1}.([saint] \quad \alpha_1) \rangle) \rangle) : np/n}{\boxed{\langle \beta_0, \lambda_{\gamma_0}.([popular] \langle \gamma_0, \lambda_{\alpha_1}.([saint] \quad \alpha_1) \rangle) \rangle) : np/n}} \quad \frac{\frac{([some] \langle \beta_0, \lambda_{\gamma_0}.([popular] \langle \gamma_0, \lambda_{\alpha_1}.([saint] \quad \alpha_1) \rangle) \rangle) : np/n \vdash \beta_0 : np \cdot / \cdot ((n/n) \cdot \otimes \cdot n)}}{(np/n) \cdot \otimes \cdot ((n/n) \cdot \otimes \cdot n) \vdash \boxed{\lambda_{\beta_0}.([some] \langle \beta_0, \lambda_{\gamma_0}.([popular] \langle \gamma_0, \lambda_{\alpha_1}.([saint] \quad \alpha_1) \rangle) \rangle) : np}} \quad \frac{\alpha_0}{\boxed{\alpha_0 : s}} \quad \frac{\frac{\frac{([arrived] \langle \beta_0.([some] \langle \beta_0, \lambda_{\gamma_0}.([popular] \langle \gamma_0, \lambda_{\alpha_1}.([saint] \quad \alpha_1) \rangle) \rangle), \alpha_0) : np \setminus s \vdash \alpha_0 : ((np/n) \cdot \otimes \cdot ((n/n) \cdot \otimes \cdot n))}{([arrived] \langle \beta_0.([some] \langle \beta_0, \lambda_{\gamma_0}.([popular] \langle \gamma_0, \lambda_{\alpha_1}.([saint] \quad \alpha_1) \rangle) \rangle), \alpha_0) : np \setminus s \vdash \alpha_0 : ((np/n) \cdot \otimes \cdot ((n/n) \cdot \otimes \cdot n))}}{\frac{(np/n) \cdot \otimes \cdot ((n/n) \cdot \otimes \cdot n) \vdash \boxed{\lambda_{\alpha_0}.([arrived] \langle \lambda_{\beta_0}.[some] \langle \beta_0, \lambda_{\gamma_0}.([popular] \langle \gamma_0, \lambda_{\alpha_1}.([saint] \quad \alpha_1) \rangle) \rangle)}}{
\end{array}$$

2. We compute the interpretation below:

$$[\ddagger] = \lambda a_0.([\text{arrived}] \langle \lambda \beta_0.([\text{some}] \langle \beta_0, \lambda \gamma_0.([\text{popular}] \langle \gamma_0, \lambda a_1.([\text{saint}] a_1) \rangle) \rangle), a_0 \rangle)$$

3. The adjuged \cdot^ℓ translations are the following:

$$\begin{aligned} [\mathbf{some}]^\ell &= \lambda\langle x, k \rangle. (\exists \lambda z. (\wedge (k \lambda \theta. (\theta z)) (x z))) \\ [\mathbf{popular}]^\ell &= \lambda\langle c, k \rangle. (c (\lambda z. (\wedge (\text{POPULAR } z) (k \lambda \theta. (\theta z))))) \\ [\mathbf{saint}]^\ell &= \lambda c. (c \text{ SAINT}) \\ [\mathbf{arrived}]^\ell &= \lambda\langle k, c \rangle. (k \lambda z. (c (\text{ARRIVED } z))) \end{aligned}$$

We can now corroborate the α -equivalence of the two \cdot^ℓ translations:

$$\begin{aligned}
\lceil \dagger \rceil^\ell &= \lambda a_0. (\lambda \langle k, c \rangle. (k \lambda z. (c \text{ (ARRIVED } z)))) \\
&\quad \langle \lambda \beta_0. (\lambda \langle x, k \rangle. (\exists \lambda z. (\wedge (k \lambda \theta. (\theta z)) (x z))) \langle \beta_0, \lambda \gamma_0. (\\
&\quad \lambda \langle c, k \rangle. (c (\lambda z. (\wedge (\text{POP } z) (k \lambda \theta. (\theta z)))) \langle \gamma_0, \lambda a_1. (\lambda c. (c \text{ SAINT } a_1)))) \rangle, a_0 \rangle) \\
&\rightarrow_\beta^* \lambda a_0. (\lambda \langle x, k \rangle. (\exists \lambda z. (\wedge (k \lambda \theta. (\theta z)) (x z))) \langle \lambda z. (a_0 \text{ (ARRIVED } z)), \lambda \gamma_0. (\\
&\quad \lambda \langle c, k \rangle. (c (\lambda z. (\wedge (\text{POPULAR } z) (k \lambda \theta. (\theta z)))) \langle \gamma_0, \lambda a_1. (\lambda c. (c \text{ SAINT } a_1)))) \rangle) \\
&\rightarrow_\beta^* \lambda a_0. (\exists \lambda z. (\wedge (\lambda \langle c, k \rangle. (c (\lambda z. (\wedge (\text{POPULAR } z) (k \lambda \theta. (\theta z)))) \\
&\quad \langle \lambda \theta. (\theta z), \lambda a_1. (\lambda c. (c \text{ SAINT } a_1))) \rangle (a_0 \text{ (ARRIVED } z)))) \\
&\rightarrow_\beta^* \lambda a_0. (\exists \lambda z. (\wedge (\lambda \theta. (\theta z) (\lambda z. (\wedge (\text{POPULAR } z) (\lambda c. (c \text{ SAINT } (\lambda \theta. (\theta z)))))) (a_0 \text{ (ARRIVED } z)))) \\
&\rightarrow_\beta^* \lambda a_0. (\exists \lambda z. (\wedge (\wedge (\text{POPULAR } z) (\text{SAINT } z)) (a_0 \text{ (ARRIVED } z)))) \\
&\rightarrow_\alpha \lambda c. (\exists \lambda z. (\wedge (\wedge (\text{POPULAR } z) (\text{SAINT } z)) (c \text{ (ARRIVED } z)))) \\
&\rightarrow_\alpha \lambda c. (\exists \lambda x. (\wedge (\wedge (\text{POPULAR } x) (\text{SAINT } x)) (c \text{ (ARRIVED } x)))) \\
&= \lceil \dagger \rceil^\ell
\end{aligned}$$

2 Pregroups

2.1

$$\begin{array}{ccc}
& (\rightarrow) & (\leftarrow) \\
(4) & \frac{\frac{}{1^l \rightarrow 1^l 1} (1) \quad \frac{}{1^l 1 \rightarrow 1} (2)}{1^l \rightarrow 1} \rightarrow & \frac{\frac{}{1 \rightarrow 1^l 1} (2) \quad \frac{}{1^l 1 \rightarrow 1^l} (1)}{1 \rightarrow 1^l} \rightarrow
\end{array}$$

(5) (\rightarrow) :

$$\begin{array}{c}
\frac{\frac{}{A^{rl} \rightarrow A^{rl} 1} (1) \quad \frac{}{1 \rightarrow A^r A} (2)}{A^{rl} \rightarrow A^{rl} (A^r A)} \rightarrow \\
\frac{\frac{}{A^{rl} \rightarrow (A^{rl} A^r) A} (1) \quad \frac{}{A^{rl} A^r \rightarrow 1} (2)}{A^{rl} \rightarrow 1 A} \rightarrow \frac{}{1 A \rightarrow A} (1) \rightarrow \\
A^{rl} \rightarrow A
\end{array}$$

(\leftarrow) :

$$\begin{array}{c}
\frac{\frac{}{A \rightarrow A 1} (1) \quad \frac{}{1 \rightarrow A^r A^{rl}} (2)}{A \rightarrow A (A^r A^{rl})} \rightarrow \\
\frac{\frac{}{A \rightarrow (A A^r) A^{rl}} (1) \quad \frac{}{A A^r \rightarrow 1} (3)}{A \rightarrow 1 A^{rl}} \rightarrow \frac{}{1 A^{rl} \rightarrow A^{rl}} (1) \rightarrow \\
A \rightarrow A^{rl}
\end{array}$$

(6) (\rightarrow) :

$$\begin{array}{c}
\frac{\overline{(AB)^l \rightarrow (1(AB))^l} \rightarrow \overline{1 \rightarrow (B^l A^l)^r (B^l A^l)}}{\overline{(AB)^l \rightarrow ((B^l A^l)^r (B^l A^l))(AB))^l} \rightarrow} \quad (3) \\
\frac{\overline{(AB)^l \rightarrow ((B^l A^l)^r ((B^l A^l)(AB)))^l}}{\overline{(AB)^l \rightarrow ((B^l A^l)^r (B^l (A^l (AB))))} \rightarrow} \quad (1) \\
\frac{\overline{(AB)^l \rightarrow ((B^l A^l)^r (B^l (A^l (AB))))}}{\overline{(AB)^l \rightarrow ((B^l A^l)^r (B^l ((A^l A)B)))} \rightarrow} \quad (1) \\
\frac{\overline{(AB)^l \rightarrow ((B^l A^l)^r (B^l (1B)))^l} \rightarrow \overline{1B \rightarrow B} \quad (1)}{\overline{(AB)^l \rightarrow ((B^l A^l)^r (B^l B))^l} \rightarrow \overline{B^l B \rightarrow 1} \quad (2)} \\
\frac{\overline{(AB)^l \rightarrow ((B^l A^l)^r 1)^l} \quad \text{todo} \quad \overline{(B^l A^l)^r 1 \rightarrow (AB)^l}}{\overline{(AB)^l \rightarrow ((B^l A^l)^r)^l} \quad (5)}
\end{array}$$

2.2

We first calculate the pregroup translation of the given sequent:

$$\begin{array}{ll}
\overline{(p/((q/q)/r))/r} & \\
= \overline{(p/((q/q)/r))} r^\ell & \{\overline{A/B} = \overline{A(\overline{B})}^\ell\} \\
= p \overline{((q/q)/r)}^\ell r^\ell & \{\overline{A/B} = \overline{A(\overline{B})}^\ell\} \\
= p \overline{((q/q))}^\ell r^\ell r^\ell & \{\overline{A/B} = \overline{A(\overline{B})}^\ell\} \\
= p ((q q^\ell) r^\ell)^\ell r^\ell & \{\overline{A/B} = \overline{A(\overline{B})}^\ell\} \\
= p (q (q^\ell r^\ell))^\ell r^\ell & \{\text{rule (1) from 2.1}\} \\
= p (q^\ell r^\ell)^\ell q^\ell r^\ell & \{\text{rule (6) from 2.1}\} \\
= p r^{\ell\ell} q^{\ell\ell} q^\ell r^\ell & \{\text{rule (6) from 2.1}\}
\end{array}$$

Finally, we prove the sequent by drawing a string diagram:

$$\begin{array}{c}
p \ r^{\ell\ell} \ q^{\ell\ell} \ q^\ell \ r^\ell \\
| \quad \frown
\end{array}$$