

Logic and Language: Exercise (Week 5)

Orestis Melkonian [6176208], Konstantinos Kogkalidis [6230067]

1 LG: continuation semantics

1.1

For each term, we compute the ILL type by first observing the polarity of the sequent and then using the table to interpret complex types.

$$\begin{aligned} & [\mathbf{some}] \\ & = [np \otimes n]^\perp \quad \{ \textit{Negative Hypothesis} \} \\ & = ([np]^\perp \otimes [n])^\perp \quad \{ [A \otimes B] \text{ with } A \text{ and } B \text{ positive} \} \end{aligned}$$

$$\begin{aligned} & [\mathbf{popular}] \\ & = [n \otimes n]^\perp \quad \{ \textit{Negative Hypothesis} \} \\ & = ([n]^\perp \otimes [n])^\perp \quad \{ [A \otimes B] \text{ with } A \text{ and } B \text{ positive} \} \end{aligned}$$

$$\begin{aligned} & [\mathbf{saint}] \\ & = [n] \quad \{ \textit{Positive Hypothesis} \} \end{aligned}$$

$$\begin{aligned} & [\mathbf{arrived}] \\ & = [np \otimes s]^\perp \quad \{ \textit{Negative Hypothesis} \} \\ & = ([np] \otimes [s]^\perp)^\perp \quad \{ [A \otimes B] \text{ with } A \text{ positive and } B \text{ negative} \} \end{aligned}$$

$$\begin{aligned} & \alpha \\ & = [s]^\perp \quad \{ \textit{Negative Hypothesis} \} \end{aligned}$$

$$\begin{aligned} & z \\ & = [np] \quad \{ \textit{Positive Hypothesis} \} \end{aligned}$$

$$\begin{aligned} & y \\ & = [n] \quad \{ \textit{Positive Hypothesis} \} \end{aligned}$$

The ILL types then are:

TERM	TYPE
$\llbracket \text{some} \rrbracket$	$(\llbracket np \rrbracket^\perp \otimes \llbracket n \rrbracket)^\perp$
$\llbracket \text{popular} \rrbracket$	$(\llbracket n \rrbracket^\perp \otimes \llbracket n \rrbracket)^\perp$
$\llbracket \text{saint} \rrbracket$	$\llbracket n \rrbracket$
$\llbracket \text{arrived} \rrbracket$	$(\llbracket np \rrbracket \otimes \llbracket s \rrbracket^\perp)^\perp$
α	$\llbracket s \rrbracket^\perp$
z	$\llbracket np \rrbracket$
y	$\llbracket n \rrbracket$

1.2

SOURCE TYPE	CONSTANT	$\llbracket \cdot \rrbracket^\ell$
n/n	popular	$\lambda \langle c, y \rangle. (c (\lambda z. \wedge (y z) (\text{POPULAR } z)))$

1.3

1. We compute the interpretation below:

$$\llbracket \dagger \rrbracket = \lambda a_0. (\llbracket \text{arrived} \rrbracket \langle \lambda \beta_0. (\llbracket \text{some} \rrbracket \langle \beta_0, \lambda \gamma_0. (\llbracket \text{popular} \rrbracket \langle \gamma_0, \lambda a_1. (\llbracket \text{saint} \rrbracket a_1) \rangle \rangle), a_0 \rangle))$$

2. The adjucted \cdot^ℓ translations are the following:

$$\begin{aligned} \llbracket \text{some} \rrbracket^\ell &= \lambda \langle x, k \rangle. (\exists \lambda z. (\wedge (k \lambda \theta. (\theta z)) (x z))) \\ \llbracket \text{popular} \rrbracket^\ell &= \lambda \langle c, k \rangle. (c (\lambda z. (\wedge (\text{POPULAR } z) (k \lambda \theta. (\theta z)))))) \\ \llbracket \text{saint} \rrbracket^\ell &= \lambda c. (c \text{ SAINT}) \\ \llbracket \text{arrived} \rrbracket^\ell &= \lambda \langle k, c \rangle. (k \lambda z. (c (\text{ARRIVED } z))) \end{aligned}$$

We can now corroborate the α -equivalence of the two \cdot^ℓ translations:

$$\begin{aligned} \llbracket \dagger \rrbracket^\ell &= \lambda a_0. (\lambda \langle k, c \rangle. (k \lambda z. (c (\text{ARRIVED } z)))) \\ &\quad \langle \lambda \beta_0. (\lambda \langle x, k \rangle. (\exists \lambda z. (\wedge (k \lambda \theta. (\theta z)) (x z))) \langle \beta_0, \lambda \gamma_0. (\\ &\quad \lambda \langle c, k \rangle. (c (\lambda z. (\wedge (\text{POP } z) (k \lambda \theta. (\theta z)))))) \langle \gamma_0, \lambda a_1. (\lambda c. (c \text{ SAINT}) a_1) \rangle \rangle), a_0 \rangle) \\ &\rightarrow_\beta^* \lambda a_0. (\lambda \langle x, k \rangle. (\exists \lambda z. (\wedge (k \lambda \theta. (\theta z)) (x z)))) \langle \lambda z. (a_0 (\text{ARRIVED } z)), \lambda \gamma_0. (\\ &\quad \lambda \langle c, k \rangle. (c (\lambda z. (\wedge (\text{POPULAR } z) (k \lambda \theta. (\theta z)))))) \langle \gamma_0, \lambda a_1. (\lambda c. (c \text{ SAINT}) a_1) \rangle \rangle) \\ &\rightarrow_\beta^* \lambda a_0. (\exists \lambda z. (\wedge (\lambda \langle c, k \rangle. (c (\lambda z. (\wedge (\text{POPULAR } z) (k \lambda \theta. (\theta z)))))) \\ &\quad \langle \lambda \theta. (\theta z), \lambda a_1. (\lambda c. (c \text{ SAINT}) a_1) \rangle) (a_0 (\text{ARRIVED } z))) \\ &\rightarrow_\beta^* \lambda a_0. (\exists \lambda z. (\wedge (\lambda \theta. (\theta z)) (\lambda z. (\wedge (\text{POPULAR } z) (\lambda c. (c \text{ SAINT}) (\lambda \theta. (\theta z)))))) (a_0 (\text{ARRIVED } z))) \\ &\rightarrow_\beta^* \lambda a_0. (\exists \lambda z. (\wedge (\wedge (\text{POPULAR } z) (\text{SAINT } z)) (a_0 (\text{ARRIVED } z)))) \\ &\rightarrow_\alpha \lambda c. (\exists \lambda z. (\wedge (\wedge (\text{POPULAR } z) (\text{SAINT } z)) (c (\text{ARRIVED } z)))) \\ &\rightarrow_\alpha \lambda c. (\exists \lambda x. (\wedge (\wedge (\text{POPULAR } x) (\text{SAINT } x)) (c (\text{ARRIVED } x)))) \\ &= \llbracket \dagger \rrbracket^\ell \end{aligned}$$

2 Pregroups

2.1

$$(4) \quad \begin{array}{ccc} & (\rightarrow) & (\leftarrow) \\ \frac{\overline{1^l \rightarrow 1^l 1} \quad (1) \quad \overline{1^l 1 \rightarrow 1} \quad (2)}{1^l \rightarrow 1} \rightarrow & & \frac{\overline{1 \rightarrow 1^l 1} \quad (2) \quad \overline{1^l 1 \rightarrow 1^l} \quad (1)}{1 \rightarrow 1^l} \rightarrow \end{array}$$

$$(5) \quad (\rightarrow): \quad \frac{\frac{\overline{A^{rl} \rightarrow A^{rl} 1} \quad (1) \quad \overline{1 \rightarrow A^r A} \quad (2)}{A^{rl} \rightarrow A^{rl}(A^r A)} \rightarrow \quad \frac{\overline{A^{rl} \rightarrow (A^{rl} A^r) A} \quad (1) \quad \overline{A^{rl} A^r \rightarrow 1} \quad (2)}{A^{rl} \rightarrow 1 A} \rightarrow \quad \frac{\overline{1 A \rightarrow A} \quad (1)}{A^{rl} \rightarrow A} \rightarrow$$

$$(\leftarrow): \quad \frac{\frac{\overline{A \rightarrow A 1} \quad (1) \quad \overline{1 \rightarrow A^r A^{rl}} \quad (2)}{A \rightarrow A(A^r A^{rl})} \rightarrow \quad \frac{\overline{A \rightarrow (A A^r) A^{rl}} \quad (1) \quad \overline{A A^r \rightarrow 1} \quad (3)}{A \rightarrow 1 A^{rl}} \rightarrow \quad \frac{\overline{1 A^{rl} \rightarrow A^{rl}} \quad (1)}{A \rightarrow A^{rl}} \rightarrow$$

$$(6) \quad (\rightarrow): \quad \frac{\overline{(AB)^l \rightarrow (1(AB))^l} \rightarrow \quad \overline{1 \rightarrow (B^l A^l)^r (B^l A^l)} \quad (3)}{\overline{(AB)^l \rightarrow (((B^l A^l)^r (B^l A^l))(AB))^l} \rightarrow} \rightarrow$$

$$\frac{\overline{(AB)^l \rightarrow ((B^l A^l)^r ((B^l A^l)(AB)))^l} \quad (1)}{\overline{(AB)^l \rightarrow ((B^l A^l)^r (B^l (A^l (AB))))} \quad (1)} \rightarrow$$

$$\frac{\overline{(AB)^l \rightarrow ((B^l A^l)^r (B^l ((A^l A) B)))} \quad (1) \quad \overline{A^l A \rightarrow 1} \quad (2)}{\overline{(AB)^l \rightarrow ((B^l A^l)^r (B^l (1 B)))^l} \rightarrow \quad \overline{1 B \rightarrow B} \quad (1)} \rightarrow$$

$$\frac{\overline{(AB)^l \rightarrow ((B^l A^l)^r (B^l B))^l} \quad (2) \quad \overline{B^l B \rightarrow 1} \quad (2)}{\overline{(AB)^l \rightarrow ((B^l A^l)^r 1)^l} \rightarrow \quad \overline{(B^l A^l)^r 1 \rightarrow (AB)^l} \quad (5)} \rightarrow$$

2.2

We first calculate the pregroup translation of the given sequent:

$$\begin{aligned}
& \overline{(p/((q/q)/r))/r} \\
&= \overline{(p/((q/q)/r))} r^\ell \\
&= p \overline{((q/q)/r)}^\ell r^\ell \\
&= p \overline{(q/q)}^\ell r^\ell r^\ell \\
&= p ((q q^\ell) r^\ell)^\ell r^\ell \\
&= p (q (q^\ell r^\ell))^\ell r^\ell & \{\text{rule (1) from 2.1}\} \\
&= p (q^\ell r^\ell)^\ell q^\ell r^\ell & \{\text{rule (6) from 2.1}\} \\
&= p r^{\ell^\ell} q^{\ell^\ell} q^\ell r^\ell & \{\text{rule (6) from 2.1}\}
\end{aligned}$$

Finally, we prove the sequent by drawing a string diagram: