

Logic and Language: Exercise (Week 5)

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1 LG: continuation semantics

1.1

For each term, we compute the ILL type by first observing the polarity of the sequent and then using the table to interpret complex types.

$$\begin{aligned} \llbracket \text{some} \rrbracket &= \llbracket np/n \rrbracket^\perp && \{ \text{Negative Hypothesis} \} \\ &= (\llbracket np \rrbracket^\perp \otimes \llbracket n \rrbracket)^\perp && \{ \llbracket A/B \rrbracket \text{ with } A \text{ and } B \text{ positive} \} \\ &= (np^\perp \otimes n)^\perp && \{ np \text{ and } n \text{ positive} \} \end{aligned}$$

$$\begin{aligned} \llbracket \text{popular} \rrbracket &= \llbracket n/n \rrbracket^\perp && \{ \text{Negative Hypothesis} \} \\ &= (\llbracket n \rrbracket^\perp \otimes \llbracket n \rrbracket)^\perp && \{ \llbracket A/B \rrbracket \text{ with } A \text{ and } B \text{ positive} \} \\ &= (n^\perp \otimes n)^\perp && \{ n \text{ positive} \} \end{aligned}$$

$$\begin{aligned} \llbracket \text{saint} \rrbracket &= \llbracket n \rrbracket && \{ \text{Positive Hypothesis} \} \\ &= n && \{ n \text{ positive} \} \end{aligned}$$

$$\begin{aligned} \llbracket \text{arrived} \rrbracket &= \llbracket np \backslash s \rrbracket^\perp && \{ \text{Negative Hypothesis} \} \\ &= (\llbracket np \rrbracket \otimes \llbracket s \rrbracket)^\perp && \{ \llbracket B \backslash A \rrbracket \text{ with } B \text{ positive and } A \text{ negative} \} \\ &= (np \otimes s^\perp)^\perp && \{ np \text{ positive, } s \text{ negative} \} \end{aligned}$$

$$\begin{aligned} \alpha &= \llbracket s \rrbracket \\ &= s^\perp && \{ s \text{ negative} \} \end{aligned}$$

$$\begin{aligned} z &= \llbracket np \rrbracket \\ &= np && \{ np \text{ positive} \} \end{aligned}$$

$$\begin{aligned} y &= \llbracket n \rrbracket \\ &= n && \{ n \text{ positive} \} \end{aligned}$$

The ILL types then are:

TERM	TYPE
[some]	$(np^\perp \otimes n)^\perp$
[popular]	$(n^\perp \otimes n)^\perp$
[saint]	n
[arrived]	$(np \otimes s^\perp)^\perp$
α	s^\perp
z	np
y	n

1.2

SOURCE TYPE	CONSTANT	$\lceil \cdot \rceil^\ell$
n/n	popular	$\lambda \langle c, y \rangle . (c \ (\lambda z. \ \wedge \ (y \ z) \ (\text{POPULAR } z)))$

1.3

1. We compute the interpretation below:

[illegible]

2. The adjuged \cdot^ℓ translations are the following:

$$\begin{aligned} [\mathbf{some}]^\ell &= \lambda\langle x, k \rangle. (\exists \lambda z. (\wedge (k \lambda\theta. (\theta z)) (x z))) \\ [\mathbf{popular}]^\ell &= \lambda\langle c, k \rangle. (c (\lambda z. (\wedge (\text{POPULAR } z) (k \lambda\theta. (\theta z))))) \\ [\mathbf{saint}]^\ell &= \lambda c. (c \text{ SAINT}) \\ [\mathbf{arrived}]^\ell &= \lambda\langle k, c \rangle. (k \lambda z. (c (\text{ARRIVED } z))) \end{aligned}$$

We can now corroborate the α -equivalence of the two \cdot^ℓ translations:

$$\begin{aligned}
\llbracket \dagger \rrbracket^\ell &= \lambda a_0. (\lambda \langle k, c \rangle. (\underline{k \lambda z. (c \text{ (ARRIVED } z))}) \\
&\quad \langle \lambda \beta_0. (\lambda \langle x, k \rangle. (\exists \lambda z. (\wedge (k \lambda \theta. (\theta z)) (x z))) \langle \beta_0, \lambda \gamma_0. (\\
&\quad \lambda \langle c, k \rangle. (c (\lambda z. (\wedge (\text{POP } z) (k \lambda \theta. (\theta z)))) \langle \gamma_0, \lambda a_1. (\lambda c. (c \text{ SAINT } a_1)) \rangle \rangle), a_0 \rangle) \\
&\rightarrow_\beta^* \lambda a_0. (\lambda \langle x, k \rangle. (\exists \lambda z. (\wedge (\underline{k \lambda \theta. (\theta z))} (x z))) \langle \lambda z. (a_0 \text{ (ARRIVED } z)), \lambda \gamma_0. (\\
&\quad \lambda \langle c, k \rangle. (c (\lambda z. (\wedge (\text{POPULAR } z) (k \lambda \theta. (\theta z)))) \langle \gamma_0, \lambda a_1. (\lambda c. (c \text{ SAINT } a_1)) \rangle \rangle) \\
&\rightarrow_\beta^* \lambda a_0. (\exists \lambda z. (\wedge (\underline{\lambda \langle c, k \rangle. (c (\lambda z. (\wedge (\text{POPULAR } z) (k \lambda \theta. (\theta z))))}) \\
&\quad \langle \lambda \theta. (\theta z), \lambda a_1. (\lambda c. (c \text{ SAINT } a_1)) \rangle) (a_0 \text{ (ARRIVED } z))) \\
&\rightarrow_\beta^* \lambda a_0. (\exists \lambda z. (\wedge (\underline{\lambda \theta. (\theta z)} (\lambda z. (\wedge (\text{POPULAR } z) (\underline{\lambda c. (c \text{ SAINT } a_1))} (\underline{\lambda \theta. (\theta z))}))) (a_0 \text{ (ARRIVED } z))) \\
&\rightarrow_\beta^* \lambda a_0. (\exists \lambda z. (\wedge (\wedge (\text{POPULAR } z) (\text{SAINT } z)) (\underline{a_0 \text{ (ARRIVED } z)})) \\
&\rightarrow_\alpha \lambda c. (\exists \lambda z. (\wedge (\wedge (\text{POPULAR } z) (\text{SAINT } z)) (c \text{ (ARRIVED } z)))) \\
&\rightarrow_\alpha \lambda c. (\exists \lambda x. (\wedge (\wedge (\text{POPULAR } x) (\text{SAINT } x)) (c \text{ (ARRIVED } x)))) \\
&= \llbracket \dagger \rrbracket^\ell
\end{aligned}$$

2 Pregroups

2.1

$$\begin{aligned}
(4) \quad & \begin{array}{ccc} & (\rightarrow) & (\leftarrow) \\ 1^\ell & \xrightarrow{(1)} 1^\ell \underline{1} \xrightarrow{(2)} 1 & 1 \xrightarrow{(2)} \underline{11}^\ell \xrightarrow{(1)} 1^\ell \end{array}
\end{aligned}$$

$$\begin{aligned}
(5) \quad & \begin{array}{c} (\rightarrow) \\ A^{rl} \xrightarrow{(1)} A^{rl} \underline{1} \xrightarrow{(3)} A^{rl} (\underline{A^r A}) \xrightarrow{(1)} (\underline{A^{rl} A^r}) A \xrightarrow{(2)} 1 A \xrightarrow{(1)} A \\ (\leftarrow) \\ A \xrightarrow{(1)} A \underline{1} \xrightarrow{(2)} A (\underline{A^r A^{rl}}) \xrightarrow{(1)} (\underline{A A^r}) A^{rl} \xrightarrow{(3)} 1 A^{rl} \xrightarrow{(1)} A^{rl} \end{array}
\end{aligned}$$

$$\begin{aligned}
(6) \quad & \begin{array}{c} (\rightarrow) \\ (AB)^\ell \xrightarrow{(1)} (\underline{1AB})^\ell \xrightarrow{(3)} (((\underline{B^\ell A^\ell})^r (\underline{B^\ell A^\ell})) AB)^\ell \xrightarrow{(1^*)} ((B^\ell A^\ell)^r B^\ell (\underline{A^\ell A}) B)^\ell \xrightarrow{(2+1)} \\ ((B^\ell A^\ell)^r \underline{B^\ell B})^\ell \xrightarrow{(2+1)} (B^\ell A^\ell)^{rl} \xrightarrow{(5)} B^\ell A^\ell \\ (\leftarrow) \\ B^\ell A^\ell \xrightarrow{(1)} B^\ell A^\ell \underline{1} \xrightarrow{(2)} B^\ell A^\ell (\underline{AB}) (\underline{AB})^\ell \xrightarrow{(1)} B^\ell (\underline{A^\ell A}) B (AB)^\ell \xrightarrow{(2+1)} (\underline{B^\ell B}) (AB)^\ell \xrightarrow{(2+1)} (AB)^\ell \end{array}
\end{aligned}$$

$$(7) \quad B^\ell \xrightarrow{(1)} B^\ell \underline{1} \xrightarrow{(2)} B^\ell \underline{AA}^\ell \xrightarrow{(A \rightarrow B)} \underline{B^\ell B} A^\ell \xrightarrow{(2+1)} A^\ell$$

2.2

We first calculate the pregroup translation of the given sequent:

$$\begin{array}{ll}
\overline{(p/((q/q)/r))/r} & \\
= \overline{(p/((q/q)/r))} r^\ell & \{\overline{A/B} = \overline{A}(\overline{B})^\ell\} \\
= p \overline{((q/q)/r)}^\ell r^\ell & \{\overline{A/B} = \overline{A}(\overline{B})^\ell\} \\
= p \overline{(q/q)}^\ell r^\ell r^\ell & \{\overline{A/B} = \overline{A}(\overline{B})^\ell\} \\
= p ((q q^\ell) r^\ell)^\ell r^\ell & \{\overline{A/B} = \overline{A}(\overline{B})^\ell\} \\
= p (q (q^\ell r^\ell))^\ell r^\ell & \{\text{rule (1) from 2.1}\} \\
= p (q^\ell r^\ell)^\ell q^\ell r^\ell & \{\text{rule (6) from 2.1}\} \\
= p r^{\ell\ell} q^{\ell\ell} q^\ell r^\ell & \{\text{rule (6) from 2.1}\}
\end{array}$$

Finally, we prove the sequent by drawing a string diagram:

$$\begin{array}{c}
p r^{\ell\ell} q^{\ell\ell} q^\ell r^\ell \\
| \quad \frown
\end{array}$$