

Problem 3

step	N'	$D(t),p(t)$	$D(u),p(u)$	$D(v),p(v)$	$D(w),p(w)$	$D(y),p(y)$	$D(z),p(z)$
0	x	∞	∞	3,x	6,x	6,x	8,x
1	xv	7,v	6,v	3,x	6,x	6,x	8,x
2	xvu	7,v	6,v	3,x	6,x	6,x	8,x
3	xvuuv	7,v	6,v	3,x	6,x	6,x	8,x
4	xvuuvy	7,v	6,v	3,x	6,x	6,x	8,x
5	xvuuvyt	7,v	6,v	3,x	6,x	6,x	8,x
6	xvuuvytz	7,v	6,v	3,x	6,x	6,x	8,x

Problem 4

a)

Step	N'	$D(x), p(x)$	$D(u),p(u)$	$D(v),p(v)$	$D(w),p(w)$	$D(y),p(y)$	$D(z),p(z)$
0	t	∞	2,t	4,t	∞	7,t	∞
1	tu	∞	2,t	4,t	5,u	7,t	∞
2	tuv	7,v	2,t	4,t	5,u	7,t	∞
3	tuvw	7,v	2,t	4,t	5,u	7,t	∞
4	tuvwxy	7,v	2,t	4,t	5,u	7,t	15,x
5	tuvwxyz	7,v	2,t	4,t	5,u	7,t	15,x
6	tuvwxyz	7,v	2,t	4,t	5,u	7,t	15,x

Problem 5

		Cost to				
		u	v	x	y	z
From	v	∞	∞	∞	∞	∞
	x	∞	∞	∞	∞	∞
	z	∞	6	2	∞	0

		Cost to				
		u	v	x	y	Z
From	v	1	0	3	∞	6
	x	∞	3	0	3	2
	z	7	5	2	5	0

		Cost to				
		u	v	x	y	z
From	v	1	0	3	3	5
	x	4	3	0	3	2
	z	6	5	2	5	0

		Cost to				
		u	v	x	y	z
From	v	1	0	3	3	5
	x	4	3	0	3	2
	z	6	5	2	5	0

Problem 7

- a) $D_x(w) = 2$, $D_x(y) = 4$, $D_x(u) = 7$
- b) First consider what happens if $c(x,y)$ changes. If $c(x,y)$ becomes larger or smaller (as long as $c(x,y) \geq 1$), the least cost path from x to u will still have cost at least 7. Thus a change in $c(x,y)$ (if $c(x,y) \geq 1$) will not cause x to inform its neighbors of any changes.

If $c(x,y) = \delta < 1$, then the least cost path now passes through y and has cost $\delta + 6$.

Now consider if $c(x,w)$ changes. If $c(x,w) = \varepsilon \leq 1$, then the least-cost path to u continues to pass through w and its cost changes to $5 + \varepsilon$; x will inform its neighbors of this new cost. If $c(x,w) = \delta > 6$, then the least cost path now passes through y and has cost 11; again x will inform its neighbors of this new cost.

- c) Any change in link cost $c(x,y)$ (and as long as $c(x,y) \geq 1$) will not cause x to inform its neighbors of a new minimum-cost path to u.

Problem 15

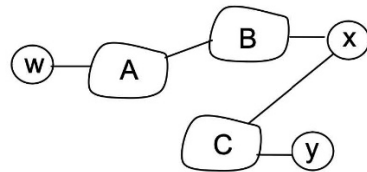
- a) I1 because this interface begins the least cost path from 1d towards the gateway router 1c.
- b) I2. Both routes have equal AS-PATH length but I2 begins the path that has the closest NEXT-HOP router.

c) I1. I1 begins the path that has the shortest AS-PATH.

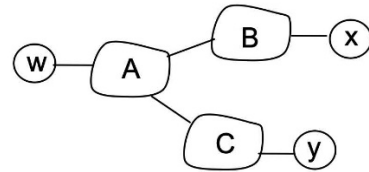
Problem 16

One way for C to force B to hand over all of B's traffic to D on the east coast is for C to only advertise its route to D via its east coast peering point with C.

Problem 17



X's view of the topology



W's view of the topology

In the above solution, X does not know about the AC link since X does not receive an advertised route to w or to y that contain the AC link (i.e., X receives no advertisement containing both AS A and AS C on the path to a destination).