CSE 541: Database Systems I

Overview of Database Systems

Database System = Database + DBMS

<u>Database (DB):</u> data organized in some way

- Model real-world applications
- Tables/entities/relations (e.g., students, courses)
- Relationships (e.g., a student is taking CMPT 454)
- A relational DB is a collection of tables

<u>Database Management System (DBMS):</u> software package manages databases

- Facilitate access to databases
- Accesses = queries and updates = reads and writes

Why not use files?

Side by Side: Files vs. DBMS

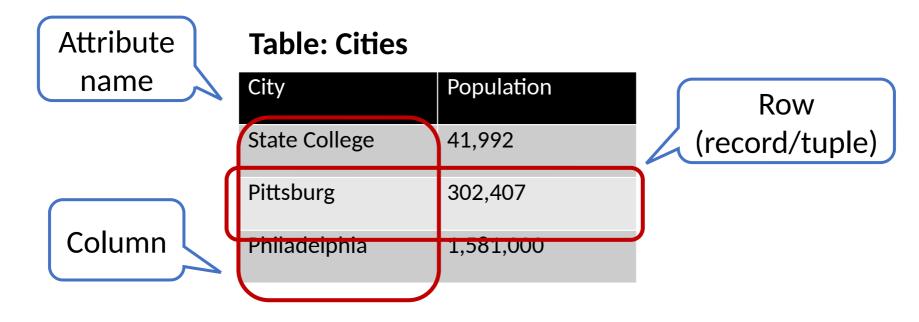
Using files:

- Application handles data access directly
 - Buffering
 - Writing data to disk
 - Optimizing performance
 - Protecting data from inconsistencies (concurrency issues)
 - Security and access control
- Special application code for different queries, and HW/SW platforms
- Custom recovery code

Using DBMS:

- Application calls DBMS for data access
- DBMS handles issues related to data access
 - Data independence
 - Buffering, performance optimization, consistency, security...
- Easier application development
 - Universal interface (e.g., SQL)
 - Same code for different platforms
- No need to worry about recovery

Relational Model



- A table consists of rows (aka records, tuples) and columns
- A row consists of fields, each column has a unique name
- A column includes all fields under a specific attribute
- Tables are instances of <u>schemas</u>

Schema

- DBMS enforces schemas defined by database designer
- Schema is "metadata" that defines what a table looks like:
 - Column definitions: name and data type of each column
 - Constraints: rules that define a valid table, enforced by DBMS

```
CREATE TABLE checklists (
   todo_id INT AUTO_INCREMENT,
   task_id INT,
   todo VARCHAR(255) NOT NULL,
   is_completed TINYINT NOT NULL DEFAULT 0,
   CHECK(task_id > 0),
   PRIMARY KEY (todo_id , task_id),
   FOREIGN KEY (task_id) REFERENCES tasks(task_id)
   ON UPDATE RESTRICT ON DELETE CASCADE);
```

Operations in RDBMS

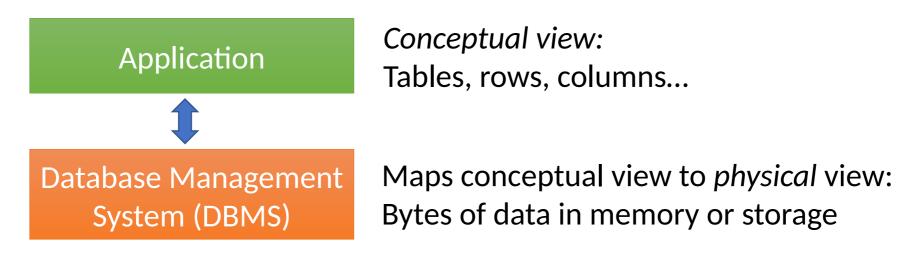
• SQL = DDL + DML.

- DDL: Data Definition Language
 - Define/ALTER Structure of data.
 - Meta info change. (Database, Table, Schema, etc.)
- DML: Data Manipulation Language
 - CRUD on actual data.

```
UPDATE Cities
SET Population = 100000
WHERE City = State
College
SELECT * FROM Cities
WHERE Population <
2000000
```

Data Independence

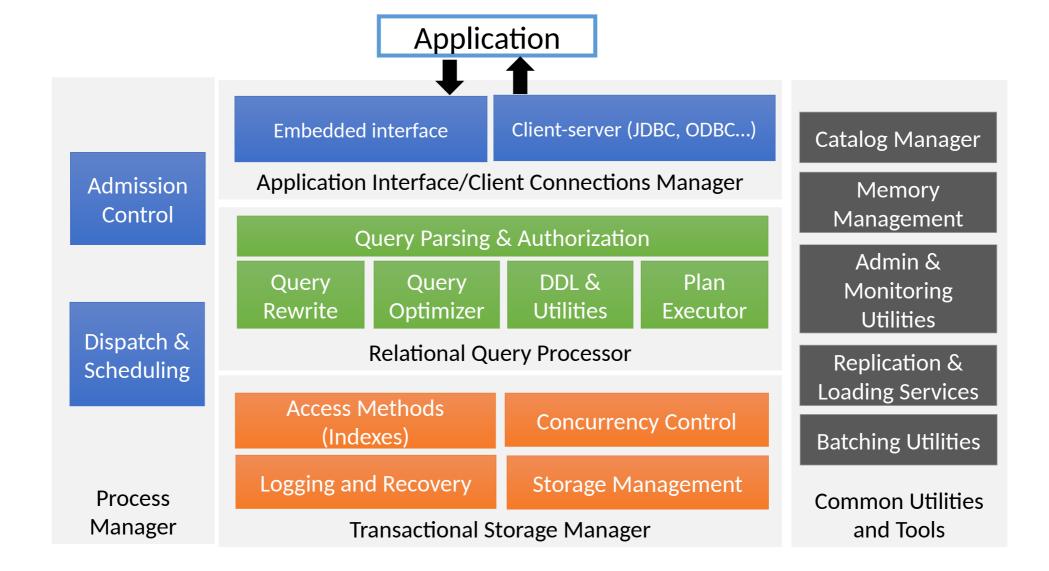
- Application (i.e., user of DBMS) views data under the relational model
- DBMS uses physical, in-memory or on-disk representations of data



- Regardless of implementation details, DBMS maintains the same relational conceptual view for applications
- Logical vs. Physical Data Independence

DBMS Functionality

- Interface to applications
- Manage SQL commands
 - Admission control: when to process the query?
 - Process management: assign thread to run the command
- Parse the command and generate a query plan
- Execute the query plan
 - Storage management and transaction management are involved.
 - Access data records while maintaining constraints
- Return results



Running Example

• Q: Which cities have at least 300k population?

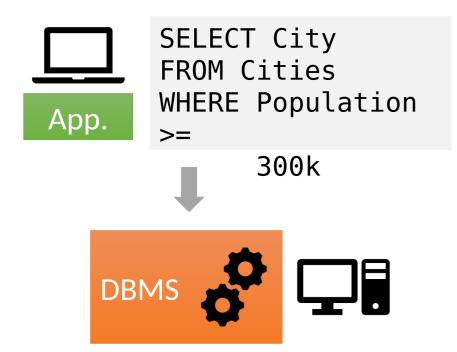
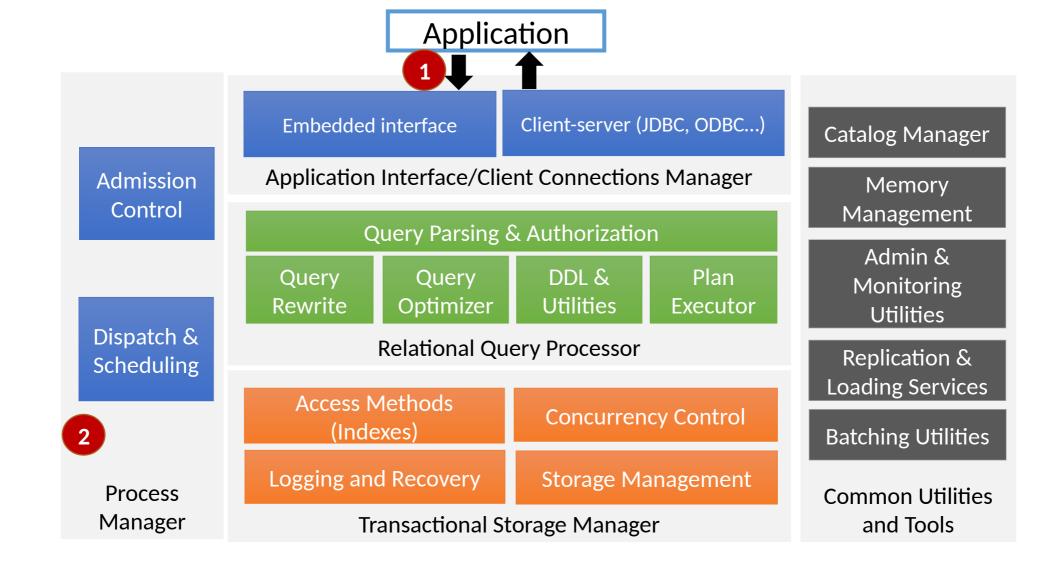


Table: Cities

City	Population
State College	41,992
Pittsburg	302,407
Philadelphia	1,581,000

How does the query reach DBMS?

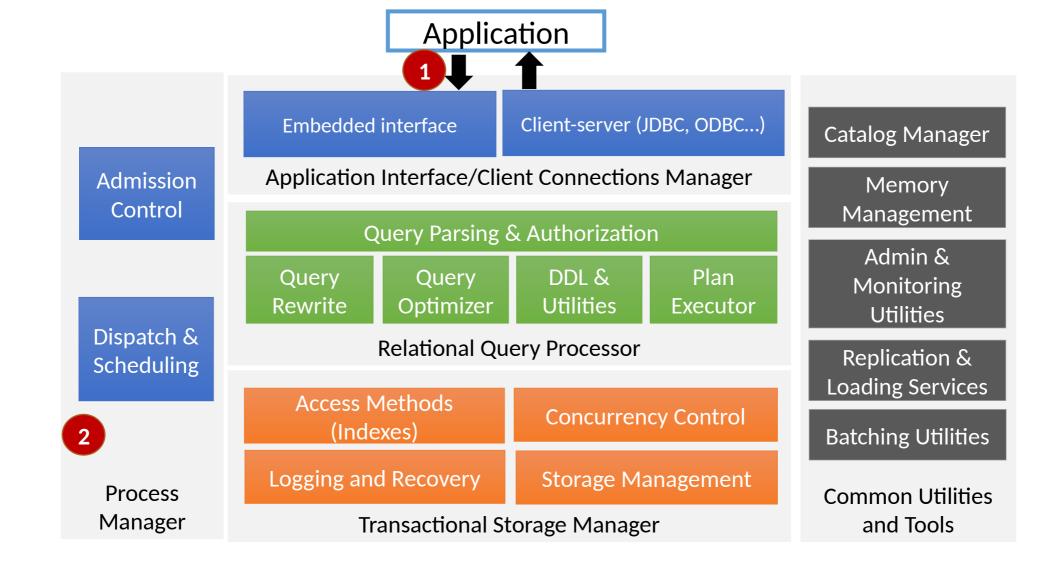
- Depends on architecture:
- 1. Embedded or
- 2. Client-server



Application-DBMS Interface

Two Major Alternatives:

- Embedded
 - Application uses DBMS functionality via direct API function calls
 - Examples: SQLite, RocksDB, etc.
- Tiered client-server
 - Application establishes a connection with the DBMS's client communications manager, via JDBC/ODBC/etc.
 - Connect to some server(s) that connect to the database server
 - Examples: MySQL, Postgres, Oracle, etc.

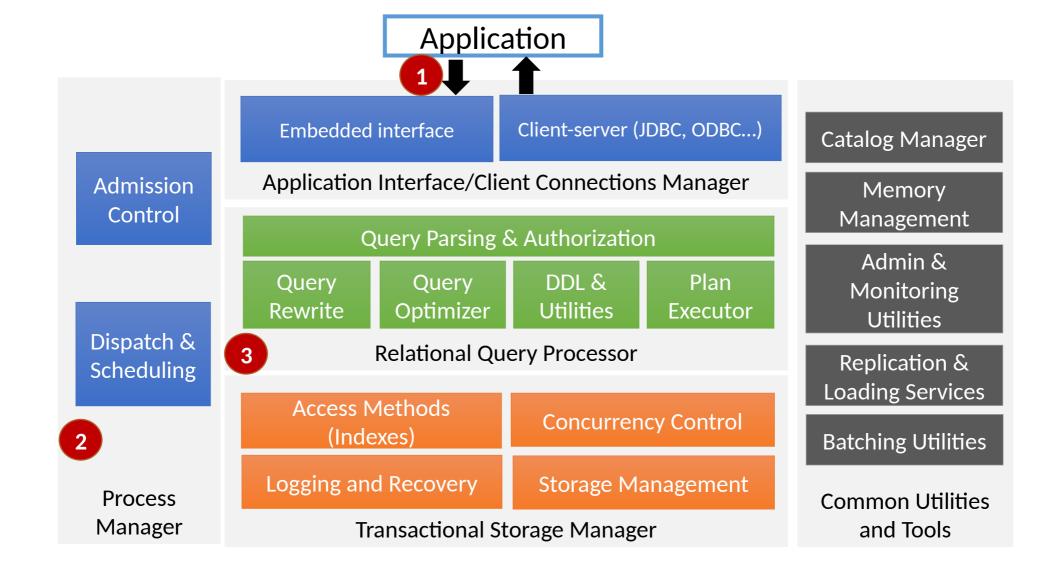


Process Manager

- Queries (usually in the form of SQL commands) arrive through application-DBMS the connection/API calls
- DBMS must assign a thread of computation to handle the commands/transactions
 - I.e., a thread that runs on a CPU core

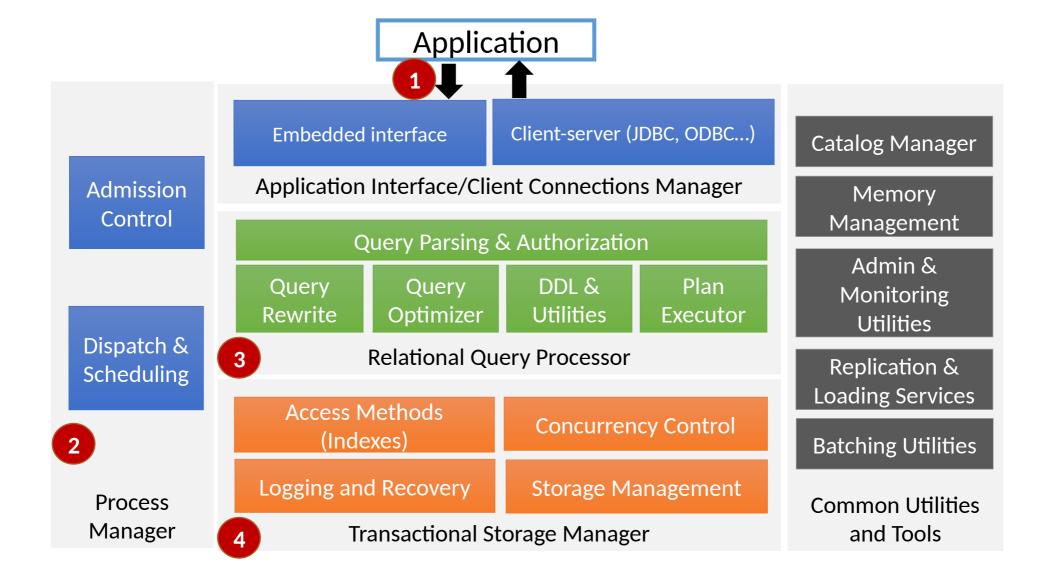
Admission control

- The DBMS and server have a limited amount of resources.
- Must not accept new work unless there is enough resource to handle it to avoid thrashing, allow graceful degradation
 - At peak utilization, throughput remains high, latency increases
- Done in the DBMS or another tier (e.g., app server)



(Relational) Query Processor

- Check the user is authorized to run the query
- Planner: Generate a query plan
- Plan executor executes the query plan
- Optimizer: Choose the best (logical/physical) plan to execute.
- Executor: Consists of a series of general-purpose "operators" that can be used to execute any query, e.g., join, sort, select
- Operators will call into the transactional storage manager for the actual data access

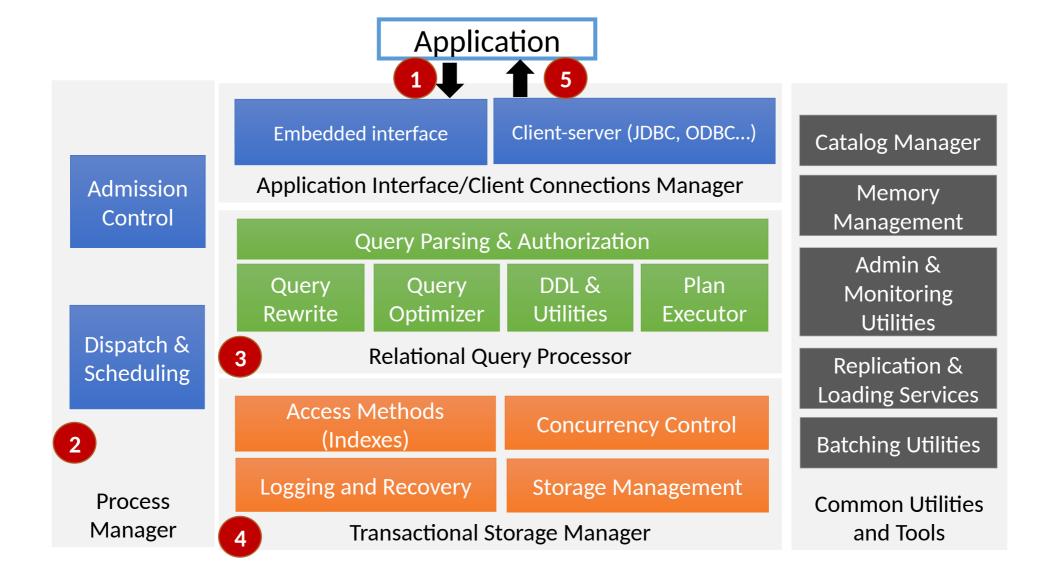


Transactional Storage Manager

- Manage all data accesses and manipulations
- Efficient ways to read and modify data
 - Indexes, tables
 - Buffering and judicious use of storage devices
- Guarantee data integrity
 - Guard against data corruption upon failures and crashes
 - Coordinate concurrent accesses from multiple threads

Key component for OLTP

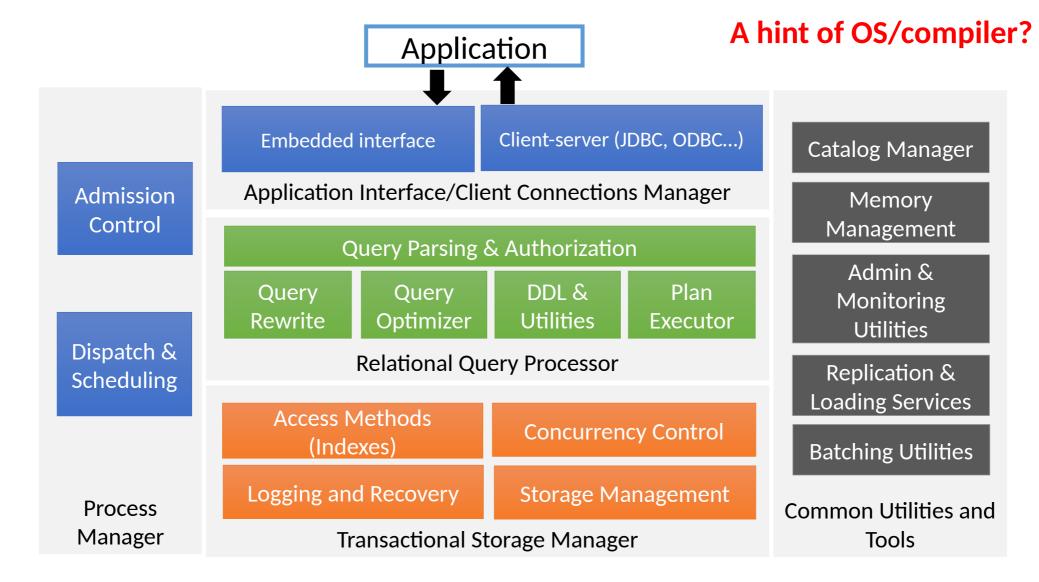
We will spend a fair amount of time on this.



Last Step: Results back to Application

• "Unwinding" the stack of previous activities

- Access method returns control to Query Processor
- Query Processor generates result data records and place them in a buffer for Connection Manager/API Interface
- Connection manager/API Interface ships results back to the caller



Takeaways...

- DBMS maps conceptual and physical views
 - Applications use the relational mode
 - DBMS uses its own physical representation
 - Optimizations can happen in the physical layer without affecting conceptual, application semantics
- Building a DBMS requires a full-stack effort
 - Understanding and utilizing hardware features
 - Sometimes (very rare) propose our own, new hardware
 - Know low-level software primitives well (features, traps, bugs...)
 - Sometimes propose our own, new ones
 - Designing suitable DBMS components