# Operating System Assignment 2

Eran Aflalo 209343722

Omer Luxembourg 205500390

# Part 1 – CPU Scheduling Implementation

In this part, we implemented a scheduling system to work on a given list of tasks.

The file is constructed by rows of task data - <task id>, <priority>, <arrival time>, <burst time>.

**Build(char \*line)** – The functions build convert each line in the file to a struct called task that has the same fields ( task id, priority ...). The struct of a task is as follows:

**Table(char \* file\_name)** - The functions return an array of tasks, we assumed that the MAX length of tasks is equal to 10. (MAX is a macro). The table will be created by a given file name on the system.

**Display(struct task\* array)** - The function prints a table that shows ( like in a practical session )

the tasks in the file.

Schedule(struct task\* tasks\_orig, enum Algorithm alg, int Q) – This function print the scheduling of an array of tasks.

Each algorithm has its function which implements it.

On the following page, is the Schedule main function. The parameter 'Q', is only for Round Robin algorithms.

Note that we have created a new copy of the tasks table, to work on it freely without erasing important information.

### Algorithms:

• **First Come First Serve** – This function is getting the first tasks that had come. This is an iterative function, which means it will only print 1 job at a call.

Note that we used task.id = -1 as a task that has been done.

• **Shortest Job First** – This function will get the shortest job that has arrived until the current CPU time, and then executes it. This is an iterative function, which prints 1 job at a call.

This function will return the current CPU time, for searching for the tasks that have been arrived until that moment.

• **Priority** – This function will get the most prioritized job that has arrived until the current CPU time, and then executes it (higher priority is most urgent). This is an iterative function, which prints 1 job at a call.

Note that this function is very similar to the shortest job first, only that we check the priority field of the task structure, instead of the burst time.

• **Round Robin** – To implementing Round Robin algorithm we created a Queue system to handle incoming tasks.

```
int intArray[MAX];
int front = 0;
int rear = -1;
int itemCount = 0;
 void Qinit() {
   front = 0;
   rear = -1;
bool QisEmpty() {
    return itemCount == 0;
bool QisFull() {
    return itemCount == MAX;
             if(rear == MAX-1) {
    rear = -1;
}
       if(front == MAX) {
```

The queue will save us the index of the wanted task in the tasks table.

After that, we created the following function to handle the scheduling – note that this function is not iterative and prints the whole scheduled events.

```
oid printRoundRobin(struct task* tasks, int curr time, int quantumTime, bool withPrior){
   bool busy = false;
         for (int i = 0; i < MAX; i++) {
         ,
// running job did not finished
if(busy == true && busy_time > 0){
         // running job finished
else if(busy == true && busy_time == 0){
                  Qinsert(busy_index); // inserting Queue

// else: busy task is over - do not add it back to the queue
              busy = false;
         // no running job, and no one is waiting.
if(busy == false && QisEmpty()){
         // no running job but there is one waiting
else if(busy == false) {
              busy = true;
```

While we have tasks to do, we first add the tasks that are in the same arrival time as the CPU time, to the queue.

Next, we have 2 main conditions – busy, which means the CPU is running on a job, or not busy, which means we can do tasks from the queue if there are tasks.

Each task is being executed as the quantum time Q, or less (if it needs less time than Q...).

• **Priority with Round Robin** – To implement it, we only needed to change the way we remove data from the queue. Instead of getting the first item in the queue, we searched for the first item with the highest priority that exists in the queue.

Note that this function is an extended version of the QremoveData function(..) we showed earlier.

### **Results:**

We created the main function to run on simple examples:

```
1, 4, 0, 10
2, 3, 3, 8
3, 3, 2, 9
```

```
processes_SJF.txt

1, 4, 1, 7
2, 3, 3, 5
3, 3, 3, 1
4, 0, 4, 2
5, 0, 5, 8
```

```
processes_prior.txt

1, 2, 0, 4
2, 4, 1, 2
3, 6, 2, 3
4, 10, 3, 5
5, 8, 4, 1
6, 12, 5, 4
7, 9, 6, 6
```

```
processes_roundrobin.txt

1, 0, 0, 4
2, 0, 1, 5
3, 0, 2, 2
4, 0, 3, 1
5, 0, 6, 3
6, 0, 6, 3
```

```
processes_priorroundrobin.txt

1, 0, 0, 4
2, 1, 1, 5
3, 2, 2, 2
4, 0, 3, 1
5, 0, 6, 3
6, 0, 6, 3
```

The terminal output is as we learned and solved by hand:

# Part 2 – XV6 Scheduling implementation

## 2.1: Priority-based Scheduler for XV6

In this part, we replaced the round-robin scheduler for xv6 with a priority-based scheduler. We have changed the *proc.c* and *proc.h* to work with priority – first added the priority field to process structure, then we changed the *scheduler* function in the *proc.c* file to work first on the highest prioritized processes.

```
proc.c
      for(p_find = ptable.proc; p_find < &ptable.proc[NPROC]; p_find++) {
   if (p_find->state != RUNNABLE)
```

As you can see, we used the same pattern as we implemented the priority with round-robin, in the previous section, to find the most prioritized task.

### 2.2: Add a Syscall to Set Priority

As we did in task 1, we added the system call *setpriority* to the XV6 system call and to the user space – by updating *syscall.h*, *user.h*, *usys.S*, *syscall.c*, *sysproc.c* according to the new changes (as done in task 1...).

To test the priority setter, and the scheduler, we created a user space program – *test\_scheduling.c*.

```
#include "types.h"
#include "stat.h"
#include "stat.h"
#include "stat.h"
#include "font1.h"

////passing command line arguments

int main(int argc, char **argv)
{
    int pid = fork();
    for(int j=0; j<1000000; j++){}

    if(pid == 0){
        int newprior = atoi(argv[2]);
        int neyPid = getpid();
        int oldPrior = setpriority(newprior);
        printf(1, "I'm the child - pid %d! change my priority to %d from %d\n",myPid,newprior,oldPrior);
        sleep(*);
        for(int i=0;ic10;i++){
            for(int j=0;j<1000000; j++){}
            printf(1, "wpid = %d run number %d: prior= %d>\n ", myPid,i,newprior);
        }
        exit();
        else if(pid < 0){
            printf(1, "This is fork failed\n");
        exit();
        else if(pid < 0){
            int newprior = atoi(argv[1]);
            int newprior = atoi(argv[1]);
        int myPid = getpid();
        int oldPrior = setpriority(newprior);
        printf(1, "I'm the father - pid %d! change my priority to %d from %d\n",myPid,newprior,oldPrior);
        sleep(%);
        for(int i=0; j<1000000; j++){}
            for(int i=0; j<1000000; j++){}
            printf(1, "Spid = %d run number %d: prior= %d>\n ", myPid,i,newprior);
        }
    }
    wait();
    exit();
}
```

We created a user program that will spend a lot of CPU time on an empty for loop, to see when the CPU is given for each task.

The first try is to give the father lower priority (30) and the child higher priority (20), and the second try is the other way around:

father\_priority = 30, child\_priority=20:

```
sb: size 1000 nblocks 941 ninodes 200 nlog 30 logstart 2 inodestart 32 bmap start
58
init: starting sh
S test schedule 30 20
I'm the father - pid 3! change my priority to 30 from 50
I'm the child - pid 4! change my priority to 20 from 50
<pid = 3 run number 0: prior= 30>
  <pid = 3 run number 1: prior= 30<pid = 4 run number 0: prior= 20>
 <pid = 4 run number 1: prior= 20>
 <pid = 4 run number 2: prior= 20>
 <pid = 4 run number 3: prior= 20>
 <pid = 4 run number 4: prior= 20>
                                          Here the child is stealing the CPU
 <pid = 4 run number 5: prior= 20>
                                          from the father, after it woken up
 <pid = 4 run number 6: prior= 20>
 <pid = 4 run number 7: prior= 20>
                                          from sleep.
 <pid = 4 run number 8: prior= 20>
 <pid = 4 run number 9: prior= 20>
 <pid = 3 run number 2: prior= 30>
 <pid = 3 run number 3: prior= 30>
 <pid = 3 run number 4: prior= 30>
 <pid = 3 run number 5: prior= 30>
 <pid = 3 run number 6: prior= 30>
 <pid = 3 run number 7: prior= 30>
 <pid = 3 run number 8: prior= 30>
 <pid = 3 run number 9: prior= 30>
```

father priority = 20, child priority=30:

```
$ test_schedule 20 30
I'm the father - pid 5! change my priority to 20 from 50
I'm the child - pid 6! change my priority to 30 from 50
<pid = 5 run number 0: prior= 20>
 <pid = 5 run number 1: prior= 20>
<pid = 5 run number 2: prior= 20>
 <pid = 5 run number 3: prior= 20>
 <pid = 5 run number 4: prior= 20>
 <pid = 5 run number 5: prior= 20>
 <pid = 5 run number 6: prior= 20>
 <pid = 5 run number 7: prior= 20>
 <pid = 5 run number 8: prior= 20>
 <pid = 5 run number 9: prior= 20>
 <pid = 6 run number 0: prior= 30>
 <pid = 6 run number 1: prior= 30>
 <pid = 6 run number 2: prior= 30>
 <pid = 6 run number 3: prior= 30>
 <pid = 6 run number 4: prior= 30>
 <pid = 6 run number 5: prior= 30>
 <pid = 6 run number 6: prior= 30>
 <pid = 6 run number 7: prior= 30>
<pid = 6 run number 9: prior= 30>
$
```

As we can see, after changing the priority and going to sleep, the CPU is given to the lower number priority – higher priority process.

We can notice that the default priority is 50 as demanded.

Now we tested the same priority process number to check that they "share" the CPU time:

```
$ test_schedule 1 1
I'm the father - pid 7! change my priority to 1 from 50
I'm the child - pid 8! change my priority to 1 from 50
<pid = 7 run number 0: prior= 1>
<pid = 7 run number 1: prior= 1>
<pid = 7 run number <pid = 8 run number 0: prior= 1>
 <pid = 8 run number 1:2: prior= 1>
 <pid = 7 run numb prior= 1>
<pid = 8 run number 2: prior= 1>
 <pid = 8 run number 3: prior= 1>
 <pier 3: prior= 1>
 <pid = 7 run number 4: prior= 1>
<pid = 7 run number 5: prior= 1>
 <pid = 7 run nud = 8 run number 4: prior= 1>
 <pid = 8 run number 5: prior= 1>
 <pmber 6: prior= 1>
 <pid = 7 run number 7: prior= 1>
 <pid = 7 run numbid = 8 run number 6: prior= 1>
 <pid = 8 run number 7: prior= 1>
 <pid = 8 run number er 8: prior= 1>
 <pid = 7 run number 9: prior= 1>
 8: prior= 1>
<pid = 8 run number 9: prior= 1>
```

As we can see, they share the CPU and interrupting each other.