CS 451/551 Quiz 6 Annotated Solution

Date: May 4, 2016 Max Points: 15

Important Reminder As per the course Academic Honesty Statement, cheating of any kind will minimally result in receiving an F letter grade for the entire course.

- 1. Which one of the following IPC mechanisms is a good choice when efficiency is a primary concern?
 - a) Anonymous pipes.
 - b) FIFO's.
 - c) Shared memory.
 - d) Message queues.
 - e) Files.

Answer: (c)

All of the above facilities require kernel involvement during set up, but shared memory does not require any kernel involvement during the actual IPC. All the alternatives require a user-space to system-space transition during the IPC. Hence shared memory is a leading contender for an efficient IPC mechanism.

- 2. Which one of the following statements about threads within a single process is necessarily **false**?
 - a) It is always possible for any thread to cancel any other arbitrary thread.
 - b) If control returns from main(), then all threads are terminated.
 - c) If a thread is detached, then it's exit status will not be available to other threads.
 - d) Code which is MT-safe can be executed correctly by multiple concurrent threads.
 - e) Each thread has its own stack.

Answer: (a)

A thread can cancel another thread only if the subject thread permits cancellation. The other statements are true because:

- Control returning from main() terminates the entire process and consequently all contained threads.
- If a thread is detached, then it is not possible for other threads to join with it and access it's termination status.

- The definition of MT-safe is that the code can be run correctly by concurrent threads.
- Each thread has a separate locus of control, it needs to have its own stack.
- 3. Which one of the following statements about the correspondence between pthread mutexes and condition variables is necessarily **true**?
 - a) There is a 1:1 correspondence between each condition variables and its controlling mutex.
 - b) Each condition variable can have one-or-more controlling mutexes.
 - c) Each mutex can control one-or-more condition variables.
 - d) Each condition variable can have one-or-more controlling mutexes and each mutex can control one-or-more condition variables.
 - e) A condition variable can be used without any corresponding mutex.

Answer: (c)

A condition variable must wait and be signalled under the same mutex, but a mutex can control multiple conditions (a queue can be full or empty).

- 4. Which one of the following constitutes a correct use of condition variables to check a shared condition() within a mult-threaded context?
 - a) pthread_mutex_t mutex = PTHREAD_MUTEX_INIITIALIZER;
 pthread_cond_t cond = PTHREAD_COND_INITIALIZER;
 if (!condition()) pthread_cond_wait(&cond, &mutex);
 //process condition
 - b) pthread_mutex_t mutex = PTHREAD_MUTEX_INIITIALIZER;
 pthread_cond_t cond = PTHREAD_COND_INITIALIZER;
 while (!condition()) pthread_cond_wait(&cond, &mutex);
 //process condition
 - c) pthread_mutex_t mutex = PTHREAD_MUTEX_INIITIALIZER;
 pthread_cond_t cond = PTHREAD_COND_INITIALIZER;
 pthread_mutex_lock(&mutex);
 while (!condition()) pthread_cond_wait(&cond, &mutex);
 //process condition
 pthread_mutex_unlock(&mutex);
 - d) pthread_mutex_t mutex = PTHREAD_MUTEX_INIITIALIZER;
 pthread_cond_t cond = PTHREAD_COND_INITIALIZER;
 pthread_mutex_lock(&mutex);
 if (!condition()) pthread_cond_wait(&cond, &mutex);
 //process condition
 pthread_mutex_unlock(&mutex);

```
e) pthread_mutex_t mutex = PTHREAD_MUTEX_INIITIALIZER;
   pthread_cond_t cond = PTHREAD_COND_INITIALIZER;
   pthread_cond_wait(&cond, &mutex);
   //process condition
```

Answer: (c)

A condition variable should be waited for only when its controlling mutex is locked; this precludes (a), (b) and (e). It is necessary the condition should be checked again after the wait; this preclused (d). This leaves (c).

5. Which one of the following sequence of calls for creating and using a streaming **server** socket is valid?

```
a) socket(); bind(); listen(); accept().
b) socket(); bind(); listen(); connect(); accept().
c) socket(); bind(); connect(); accept().
d) socket(); connect(); listen(); accept().
e) socket(); connect().
```

Answer: (a)

connect () is used only for client sockets.

6. Which one of the following sequence of calls for creating and using a streaming **client** socket is valid?

```
a) socket(); bind(); listen(); accept().
b) socket(); bind(); listen(); connect(); accept().
c) socket(); bind(); connect(); accept().
d) socket(); connect(); listen(); accept().
e) socket(); connect().
```

Answer: (e)

The correct sequence of calls is socket(); connect().bind(),listen() and accept() are only used for server sockets.

7. Assuming no errors and that printf() always runs atomically, what will be printed by the following program?

```
pthread_mutex_t lock = PTHREAD_MUTEX_INITIALIZER;
void *fn(void *p) {
   int *ip = p;
   pthread_mutex_lock(&lock);
   *ip += 5;
   pthread_mutex_unlock(&lock);
   return NULL;
}

int main() {
   int v = 5;
   pthread_t t;
   pthread_create(&t, NULL, fn, &v);
   fn(&v);
   printf("%d\n", v);
   return 0;
}
```

- a) 10 followed by 15.
- b) Only 15 will always be printed.
- c) One of 5, 10 or 15 will always be printed.
- d) Only 10 will always be printed.
- e) One of 10 or 15 will always be printed.

Answer: (e).

[During the exam it was announced that if alternatives overlapped, the more accurate alternative should be chosen.]

The main() thread always calls fn(); hence v will always be incremented at least once to 10. However there is a race condition between execution of the auxiliary thread and the execution of the printf() by the main thread (it is also possible that the auxiliary thread cannot even increment v before the main thread exits, terminating the entire process including the auxiliary thread). Hence the auxiliary thread may or may not have a chance to increment v before the printf(). Hence the output will be one of 10 or '15.