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DSA LAB:5. Theory

Q1. Define linked list. Comment on: The features of linked list over array.

→ Linked list is the linear data structure in which elements are not stored at contiguous memory location.
The linked list node contains two parts head & tail. Head stores information while tail stores address of next node.

Features over array:

- i) It can be accessed in sequential manner.
- ii) In linked list insertion & deletion can be done at any point in list at constant time.
- iii) In linked list we can add any number of elements in list.
- iv) In linked list we can extend list as we needed while in array extension of some declared array is not possible.

Q2. How many pointers are required for implementing a singly linked list?

→ Three types of pointers are required to implement singly linked list:-

i) A head pointer which is required for pointing to start of record.

ii) Info pointer which require to point last of node.

iii) Link pointer which require to point the ~~to~~ next node.

Q3. State and explain different type of linked list. How a linked list node can be represented?

→ following are the different types of linked list:-

i) Simply linked list:

It is simplest type of linked list in which every node contain same data and pointer to next data, it goes only in one direction.

2) Doubly linked list:-

In this type of linked list items can be travelled in both direction.

3) Circular linked list:-

In circular linked list last node contains a pointer to first node of list.

4) Circular doubly linked list:-

In this item can be traversed in both direction as well as last node contains address of 1st node.

Q4. Compare Singly linked list and doubly linked list.

- ① Singly linked list can be traversed only in one direction while Doubly linked list can be traversed in both direction.
- ② Singly linked list contain pointer to next node only while doubly contain pointer to next as well as previous node.

③ Singly consist of two parts i.e. data +
pointer to next node while doubly
consists of three parts i.e. data,
ii) pointer to next node iii) pointer to
next node.

④ Singly uses less memory as it
stores pointer of only one node
where as Doubly uses more memory
as it stores pointer of both previous
and next node.

Q5. Give applications of linked list.

→ Application of linked list:

- i) In implementation of stack & queue.
- ii) In implementation of graphs.
- iii) for representing sparse matrices.
- iv) Manipulation of polynomial by storing coefficient.
- v) for dynamic memory allocation.
- vi) for maintaining directory of names.

Q6. Write down the detailed algorithms for the given problem statement.

→ @ Algorithm for question 1.

Insertion of node at the end of list.

Step - 1 : Start.

Step - 2 : Create a node pointer dynamically
 $\text{node} * \text{ptr} = \text{new Node}();$

Step - 3 : Assign the data to the data part
 $\text{ptr} \rightarrow \text{data} = \text{data};$

Step - 4 : Assign NULL pointer to the pointer of the node.

$\text{ptr} \rightarrow \text{next} = \text{NULL}$

Step - 5 : If head pointer is equal to NULL
assign head pointer to the created node pointer

Step - 6 : Else iterate through the nodes till one of the nodes do not point to NULL

$\text{curr} = \text{head}$.

$\text{while} (\text{curr} \rightarrow \text{next} \neq \text{NULL})$
 $\text{curr} = \text{curr} \rightarrow \text{next}$.

Step - 7 : Assign the pointer to the related node to the node pointer which points to NULL

Step 8 : END.

Deletion of Node from the starting point

Step 1 : Start .

Step 2 : Assign two node pointer one to the head of the linked list and another to NULL
 $\text{Node}^* \text{sub} = \text{Head}$
 $\text{Node}^* \text{last} = \text{ptr}$.

Step 3 : Check if Head is not NULL and the key data is passed is equal to the data which is to be erased .

```
if (sub != NULL & sub->data == data)
{
    Head = sub->next;
    delete sub;
```

Step 4 : Else iterate through the list till the data is not found or the pointer is equal to NULL .

Step 5 : If data not found point "Data not found".

Step 6 : Else point the "last pointer to "sub pointer".

i.e. point the previous node's pointer to the next node's pointer .

```
last->next = sub->next;
delete the current node.
```

```
delete sub;
```

Step 7 : END .

Printing the data in the nodes in reverse order using recursion.

Step 1 : START

Step 2 : Check the current node pointer passed is NULL

Step 3 : If yes, then return

Step 4 : Else call the function by sending the address of the pointer of next node.

Step 5 : Print the data

Step 6 : END

Pseudo Code :

```
void reverse (Node* ptr)
```

```
    IF (ptr == NULL) return
```

```
    else reverse (ptr → next)
```

```
    cout << ptr → data ;
```

```
}
```

* Algorithm of 2nd question:

Inseriton in doubly linked list.

Step 1 :- start

Step 2 :- Assign a new node pointer.

Step 3 :- Assign data passed to the data
variable in pointer.

Step 4 :- Assign the forward pointer to
the Head.

Step 5 :- Assign the backward pointer to
NULL

Step 6 :- Check if head pointer is equal
to NULL.

Step 7 :- If no, point the backward
pointer head to the new node
pointed created.

Step 8 :- Point the head pointer to the
new node created.

Step 9 :- END.

Pseudo Code :

```
Node *ptr = new Node();
```

```
ptr → data = data;
```

```
ptr → next = Head;
```

```
ptr → prev = NULL;
```

```
if (Head != NULL) (head → prev) = ptr  
head = ptr;
```

Sorting of the doubly linked list

Step 1 :- Start

Step 2 - Assign two node pointers to NULL
and declare a int variable

Step 3 - check if the list is empty.

Step 4 - make a do-while loop.

Step 5 - Assign int variable value zero
and assign one pointer to head
flag = 0;

ptr = head

Step 6 :- Iterate ptr till it is not equal to,

Step 7 :- While in loop check the condition
loop if the $\text{ptr} \rightarrow \text{data}$ is greater than
 $\text{ptr} \rightarrow \text{next} \rightarrow \text{data}$

Step 8 :- If the condition is true,

swap ($\text{ptr} \rightarrow \text{data}$; $\text{ptr} \rightarrow \text{next} \rightarrow \text{data}$)
flag = 1.

Step 9 :- Break the loop when the value of
flag is zero.

Step 10 :- END.

Pseudo code :

```
node * ptr = NULL
```

```
node * ptr1 = NULL
```

```
int flag;
```

```
do
```

```
{
```

```
    flag = 0;
```

```
    ptr = head;
```

```

while (ptr->next != ptr+1)
{
    if (ptr->data > ptr->next->data)
    {
        swap (ptr->data, ptr->next->data);
        flag = 1;
    }
    ptr = ptr->next;
}
ptr = ptr->next;
}
while (flag != 1);

```

Addition of two Binary Number

Pseudo code :

```

String add (string s1, string s2)
{
    int len = s1.size();
    string t = " "; int carry = 0;
    for (int i = len - 1; i >= 0; i--)
    {
        if ((s1[i] == '1' + + s2[i] == '1') || (carry == 0))
            t.append ('0');
        carry = 1;
    }
    else if ((s1[i] == '1' + + s2[i] == '1') || (t + carry == 1))
        t.append ('1');
    carry = 1;
}
else if ((s1[i] == '0' + + s2[i] == '1') || (s1[i] == '1' + + s2[i] == '0'))
    t.append ('0');
    carry = 0;
}

```

```
{ t.append('1');  
    carry = 0;  
}
```

```
else if ((s1[i] == '1' + s2[i] == '0') ||  
(s1[i] == '0' + s2[i] == '1') +  
    carry == 1)
```

```
{ t.append('0');  
    carry = 1;  
}
```

```
else if (s1[i] == '0' + s2[i] == '0'  
    carry == 0)
```

```
{ t.append('0');  
    carry = 0;  
}
```

```
else if (s1[i] == '0' + + carry == 1)
```

```
{ t.append('1');  
    carry = 0;  
}
```

```
} if (carry == 1) t.append('1');  
reserve(t.begin(), t.end());  
return t;
```

```
}
```