

1.
 - 1.1. Suppose at the current cycle, $r[2] = 12$, $r[3] = 45$, $r[4] = 9$, and $r[5] = 22$, and the following commands are written in the `run()` function:
`sprn->r[4] = spro->r[2] + spro->r[3];`
`sprn->r[5] = spro->r[3] - spro->r[4];`
After execution of 2 cycles, what will be the values of $r2$, $r3$, $r4$, and $r5$?

In the end of the first cycle the values of the registers will remain the same.
In the beginning of the second cycle the values will be: , $r[2] = 12$, $r[3] = 45$, $r[4] = 57$, and $r[5] = 22$. At the end of this cycle the values will remain the same.
At the beginning of the 3rd cycle the values will be: , $r[2] = 12$, $r[3] = 45$, $r[4] = 57$, and $r[5] = -12$.
 - 1.2. The code is wrong because `spro` should not be modified during a clock cycle (like Q port of the FF). It should only get the `sprn` value in the next cycle.
2.
 - 2.1. The execution of the first instruction will take 7 clock cycles (Including the IDLE state). The other $n-1$ instructions will take 6 clock cycles each.
So execution of n instructions will take $7+6*(n-1) = 6n+1$ clock cycles.
 - 2.2. No, because reading from the memory can be done only when the state machine is in "EXEC0" state and the state machine does not support remaining in this state for more than 1 clock cycle.
 - 2.3. The disadvantage of this microarchitecture is that each instruction execution takes at least 6 clock cycles in comparison to pipelined microarchitecture.
The advantage of this microarchitecture is that it is more simple to implement this microarchitecture than a pipelined microarchitecture.
3. The `sp.c` file is attached.
4. The outputs for multiplication test are under the **outputs** directory,
5. The outputs for Fibonacci test are under the **outputs** directory,
6. The `sp.c` file with `dma` is located in the **with_dma** directory, the outputs of the `dma` test are located under **outputs** directory and `bin` file under **bin** directory. Build with **make dma**.
7. ASM code:

```
asm_cmd(ADD, 3, 0, 1,100); // 0: R3 = 100
asm_cmd(ADD, 2, 0, 1,0); // 1: R2 = 0
asm_cmd(ADD, 6, 0, 1,2); // 2: R6 = 2
asm_cmd(LD, 4, 0, 3,0); // 3: R4 = MEM[R3]
asm_cmd(ADD, 5, 0,1,1000); // 4: R5 = 1000
asm_cmd(CMB, 0, 5, 4,10); // 5: copy the content (size 2) from address 1000 to
address that was stored in MEM[100]
asm_cmd(POL, 6, 0, 0,0); // 6: poll the DMA status into R6
asm_cmd(ST, 0, 6, 3,0); // 7: MEM[R3] = R6
asm_cmd(JNE, 0, 2, 6,6); // 9: if R6 != 0 jump to pc 6
asm_cmd(HLT, 0, 0, 0,0); // 10: HALT
```