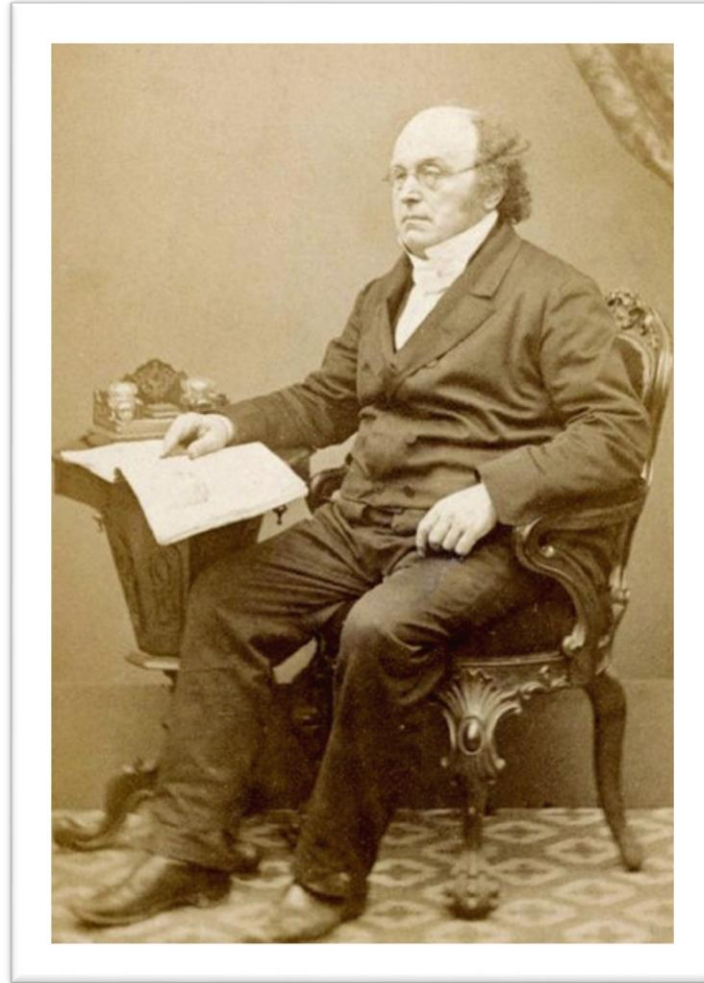


Binary Logic

DeMorgan Theorems

Augustus DeMorgan (1806-1871)



British mathematician born in India

DeMorgan's Theorems

- Are used to simplify complex **negated** binary logical expressions.

Example:

- $Z = \overline{ABC + A + B + C + AB + AC + BC}$

DeMorgan's Theorems

$$\overline{x+y} = \bar{x} \bullet \bar{y}$$

First DeMorgan's Theorem

$$\overline{x \bullet y} = \bar{x} + \bar{y}$$

Second DeMorgan's Theorem

Proof of the DeMorgan's theorems

Proof of the DeMorgan's theorem

- $\overline{x+y} = \overline{x} \bullet \overline{y}$

- $\overline{x \bullet y} = \overline{x} + \overline{y}$

x	y	\overline{x}	\overline{y}	$x y$	$\overline{x y}$	$\overline{x+y}$	$\overline{x} \ \overline{y}$	$x + y$	$\overline{\overline{x} + \overline{y}}$
0	0								
0	1								
1	0								
1	1								

Proof of the DeMorgan's theorems

- $\overline{x+y} = \overline{x} \bullet \overline{y}$

- $\overline{x \bullet y} = \overline{x} + \overline{y}$

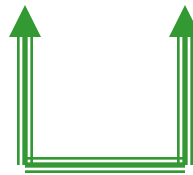
x	y	\overline{x}	\overline{y}	$x y$	$\overline{x y}$	$\overline{x+y}$	$\overline{x} \ \overline{y}$	$x + y$	$\overline{\overline{x} + \overline{y}}$
0	0	1	1	0					
0	1	1	0	0					
1	0	0	1	0					
1	1	0	0	1					

Proof of the first DeMorgan's theorem

- $\overline{x+y} = \overline{x} \bullet \overline{y}$

- $\overline{x \bullet y} = \overline{x} + \overline{y}$

x	y	\overline{x}	\overline{y}	$x y$	$\overline{x y}$	$\overline{x+y}$	$\overline{x} \overline{y}$	$x + y$	$\overline{x + y}$
0	0	1	1	0	1	1			
0	1	1	0	0	1	1			
1	0	0	1	0	1	1			
1	1	0	0	1	0	0			

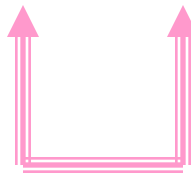


Proof of both DeMorgan's theorems

- $\overline{x+y} = \overline{x} \bullet \overline{y}$

- $\overline{x \bullet y} = \overline{x} + \overline{y}$

x	y	\overline{x}	\overline{y}	$x y$	$\overline{x y}$	$\overline{x+y}$	$\overline{x} \ \overline{y}$	$x + y$	$\overline{x + y}$
0	0	1	1	0	1	1	1	0	1
0	1	1	0	0	1	1	0	1	0
1	0	0	1	0	1	1	0	1	0
1	1	0	0	1	0	0	0	1	0



Perfect Induction

Example: Using DeMorgan's

$$\overline{A \bullet \overline{B} + C}$$

- $\overline{x + y} = \overline{x} \bullet \overline{y}$

- $\overline{x \bullet y} = \overline{x} + \overline{y}$

First application of the theorem

$$\overline{A \bullet B + C} = \overline{A \bullet B} \bullet \overline{C}$$

- $\overline{x + y} = \overline{x} \bullet \overline{y}$

- $\overline{x \bullet y} = \overline{x} + \overline{y}$

First application of the theorem

$$\overline{A \cdot B + C} = \overline{A \cdot B} \cdot \overline{C}$$

- $\overline{x + y} = \overline{x} \cdot \overline{y}$

- $\overline{x \cdot y} = \overline{x} + \overline{y}$

Second application of theorem

$$\begin{aligned}\overline{A \bullet \overline{B} + C} &= \overline{A \bullet \overline{B} \bullet \overline{C}} \\ &= (\overline{A} + \overline{\overline{B}}) \bullet \overline{C}\end{aligned}$$

$$\bullet \overline{x+y} = \overline{x} \bullet \overline{y}$$

$$\bullet \overline{x \bullet y} = \overline{x} + \overline{y}$$

Distribute...

$$\begin{aligned}\overline{A \cdot \overline{B} + C} &= \overline{A \cdot \overline{B}} \cdot \overline{C} \\ &= (\overline{A} + \overline{\overline{B}}) \cdot \overline{C} \\ &= (\overline{A} + B) \cdot \overline{C} \\ &= \overline{A} \cdot \overline{C} + B \cdot \overline{C}\end{aligned}$$

Sum-of-products (SOP) form

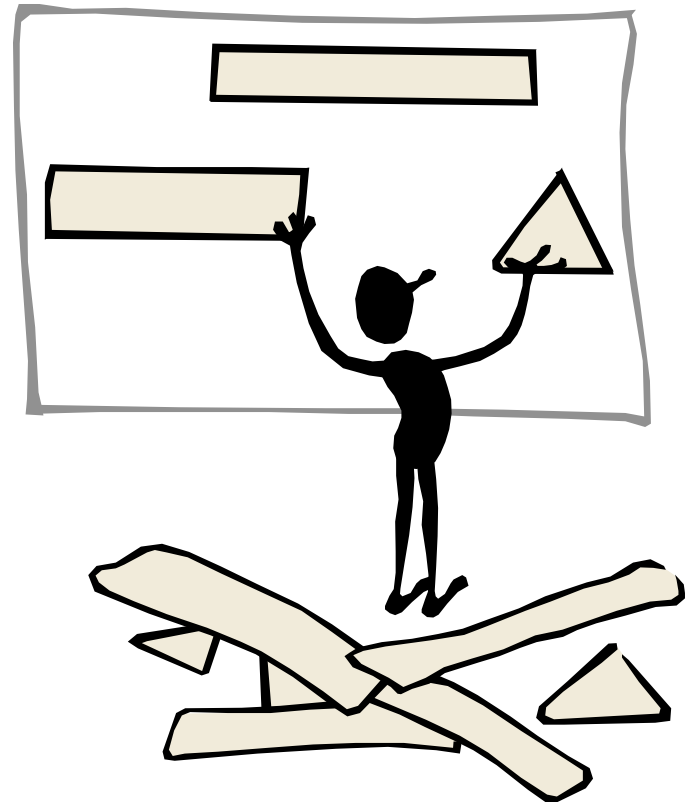
$$\begin{aligned}\overline{A \cdot \overline{B} + C} &= \overline{A \cdot \overline{B} \cdot \overline{C}} \\ &= (\overline{A} + \overline{\overline{B}}) \cdot \overline{C} \\ &= (\overline{A} + B) \cdot \overline{C} \\ &= \overline{A} \cdot \overline{C} + B \cdot \overline{C}\end{aligned}$$



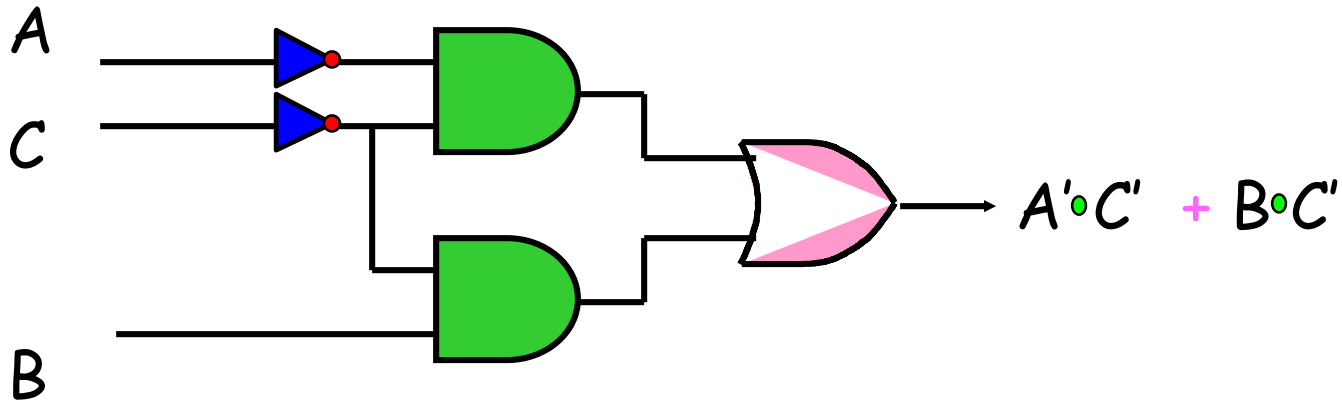
Implement with gates

Sum-of-Products (SOP) form

Ready to see the circuit?

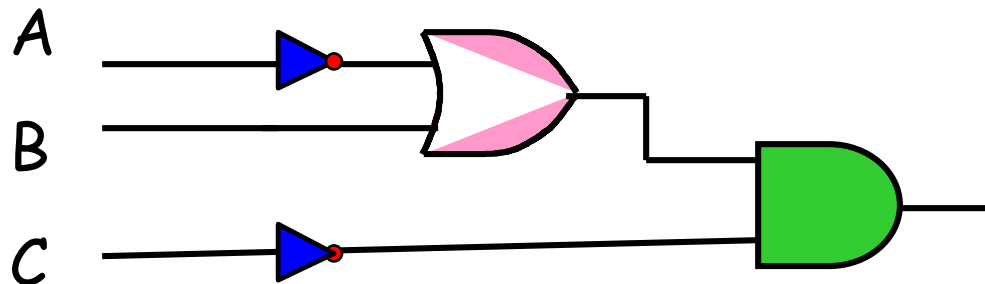


The SOP leads to “two-level-realization”



factor ... “multi-level-realization”

$$A' C' + B C' = C'(A' + B)$$



More gates....

More Gates

- NOR (Not OR)
- NAND (Not AND)

OR

A	B	OR	NOR
0	0	0	
0	1	1	
1	0	1	
1	1	1	

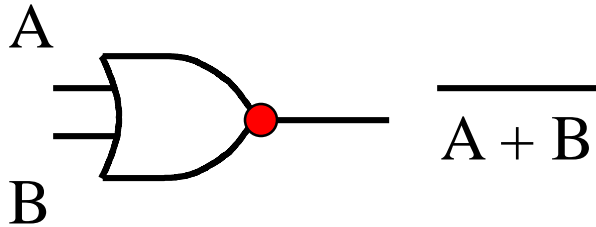


NOR

A	B	OR	NOR
0	0	0	1
0	1	1	0
1	0	1	0
1	1	1	0



NOR



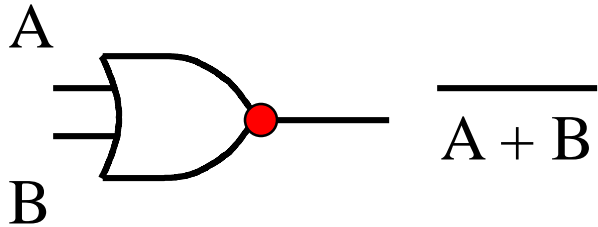
Using DeMorgan's theorem

A	B	OR	NOR
0	0	0	1
0	1	1	0
1	0	1	0
1	1	1	0

$$\bullet \overline{x+y} = \overline{x} \bullet \overline{y}$$

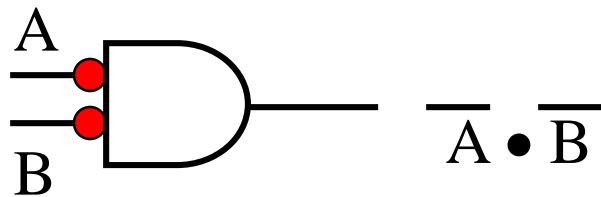
$$\bullet \overline{x \bullet y} = \overline{x} + \overline{y}$$

NOR



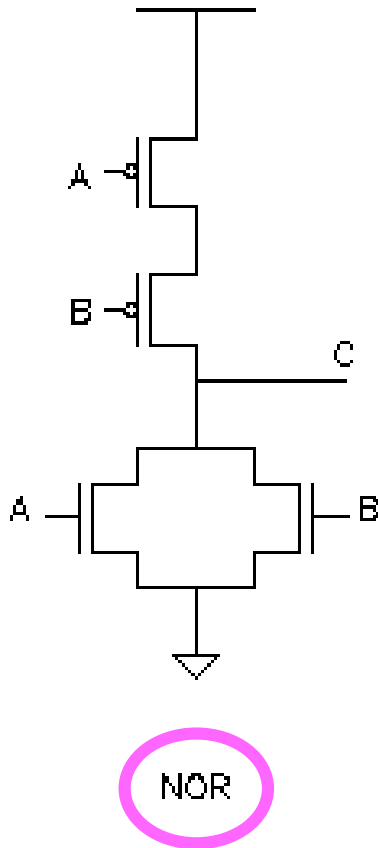
Using DeMorgan's theorem

A	B	OR	NOR
0	0	0	1
0	1	1	0
1	0	1	0
1	1	1	0



$$\overline{A} \bullet \overline{B} = \overline{A + B}$$

NOR: CMOS and gate layout



AND ... Not AND

A	B	AND	NAND
0	0	0	
0	1	0	
1	0	0	
1	1	1	

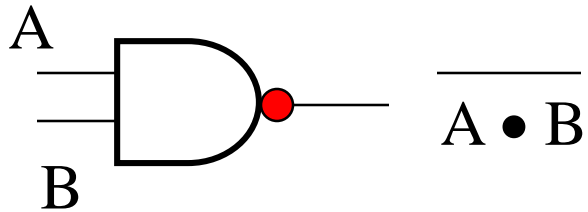


NAND

A	B	AND	NAND
0	0	0	1
0	1	0	1
1	0	0	1
1	1	1	0



NAND



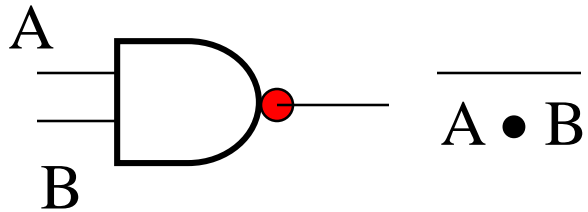
Using DeMorgan's theorem

A	B	AND	NAND
0	0	0	1
0	1	0	1
1	0	0	1
1	1	1	0

$$\bullet \overline{x+y} = \overline{x} \bullet \overline{y}$$

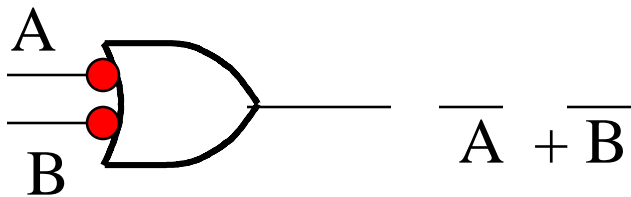
$$\bullet \overline{x \bullet y} = \overline{x} + \overline{y}$$

NAND



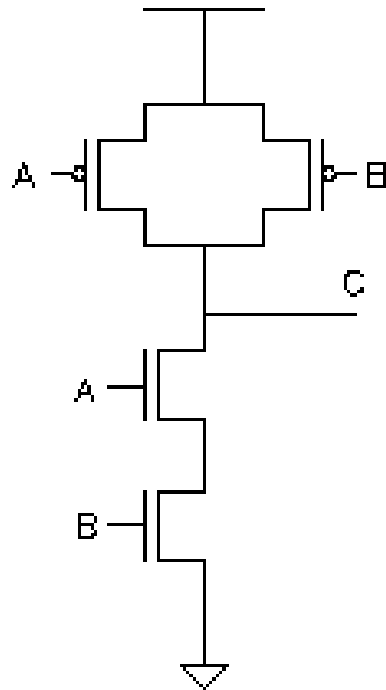
Using DeMorgan's theorem

A	B	AND	NAND
0	0	0	1
0	1	0	1
1	0	0	1
1	1	1	0

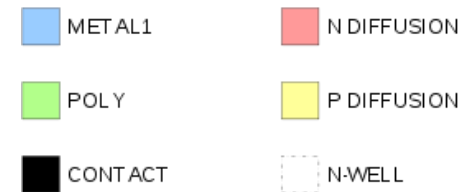
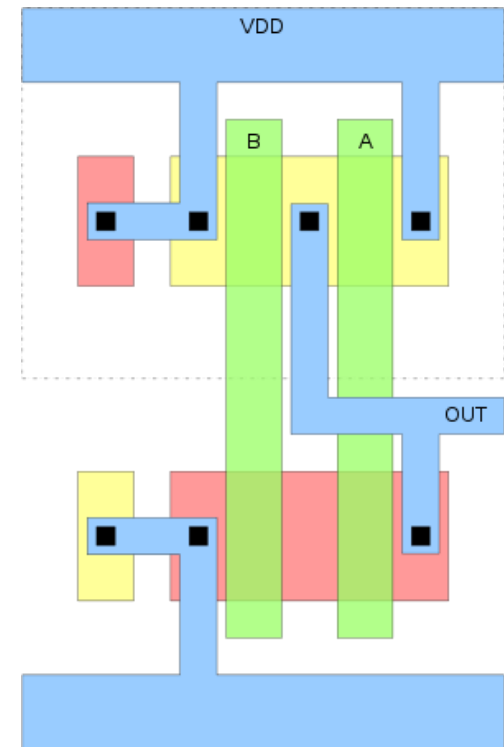


$$\overline{A \bullet B} = \overline{A} + \overline{B}$$

NAND: CMOS and gate layout



NAND

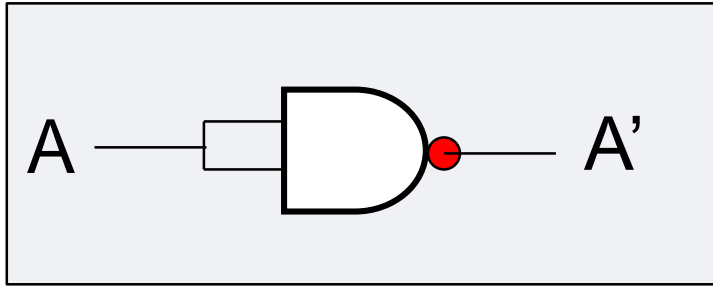


Universality of NAND and NOR Gates

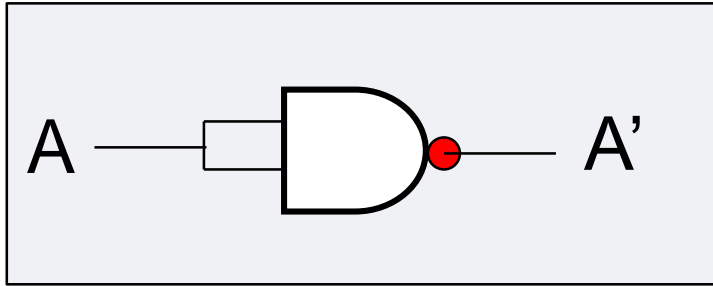
NAND and NOR gates ...

- **Can implement any Boolean expression**
- **Can simulate all three basic gates (AND, OR, NOT)**

Universality of NAND: NOT

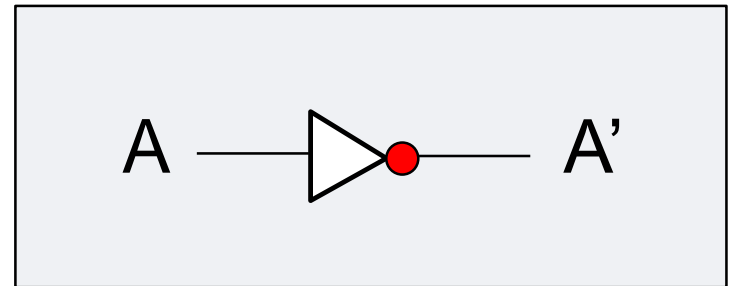
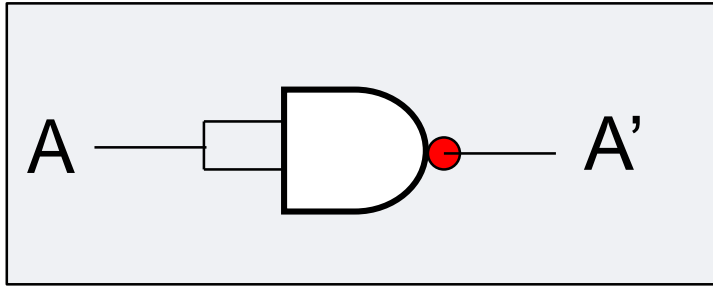


Proof



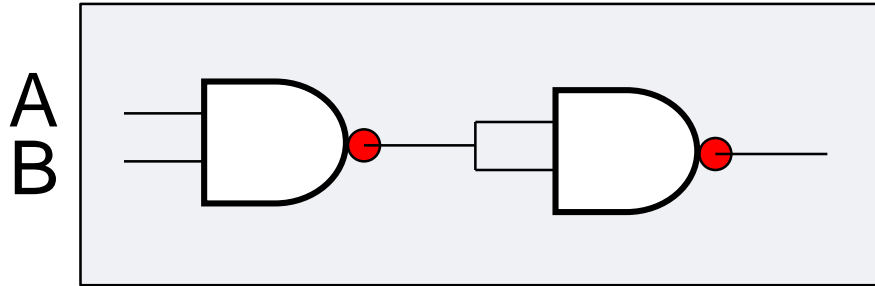
$$(AA)' = A'$$

Proof

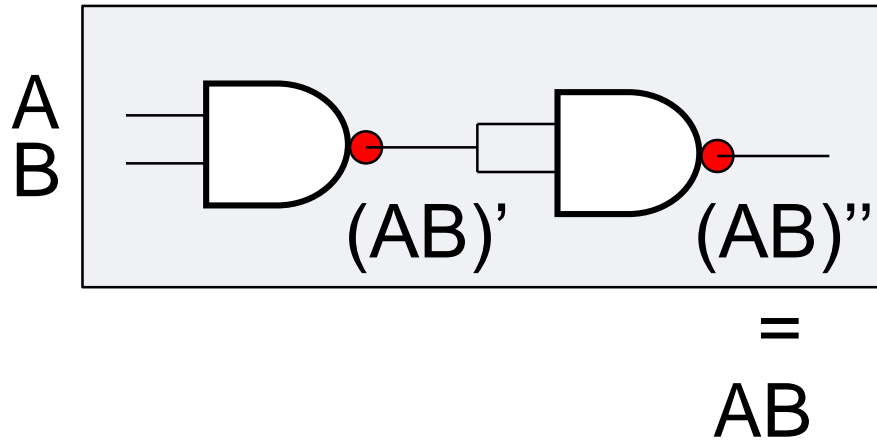


$$(AA)' = A'$$

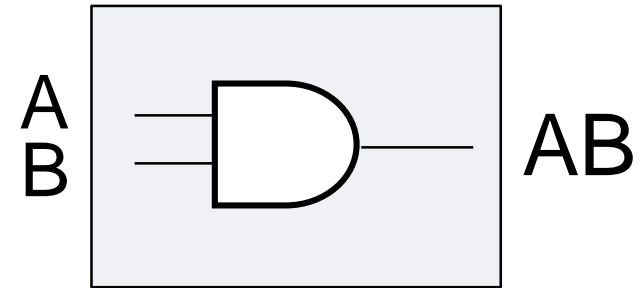
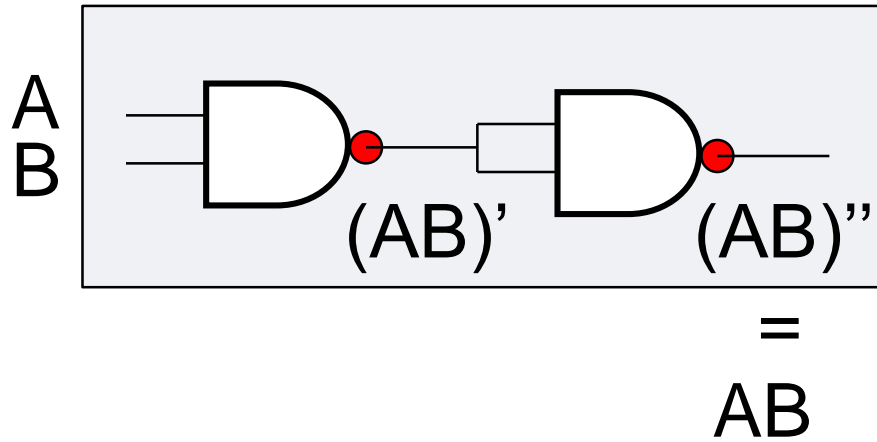
Universality of NAND: AND



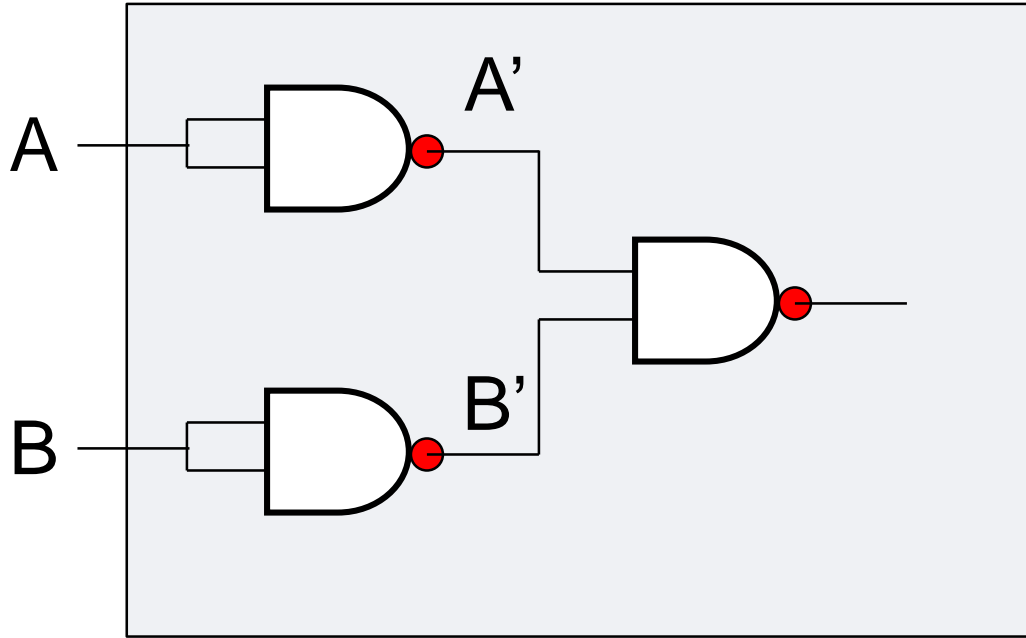
Proof



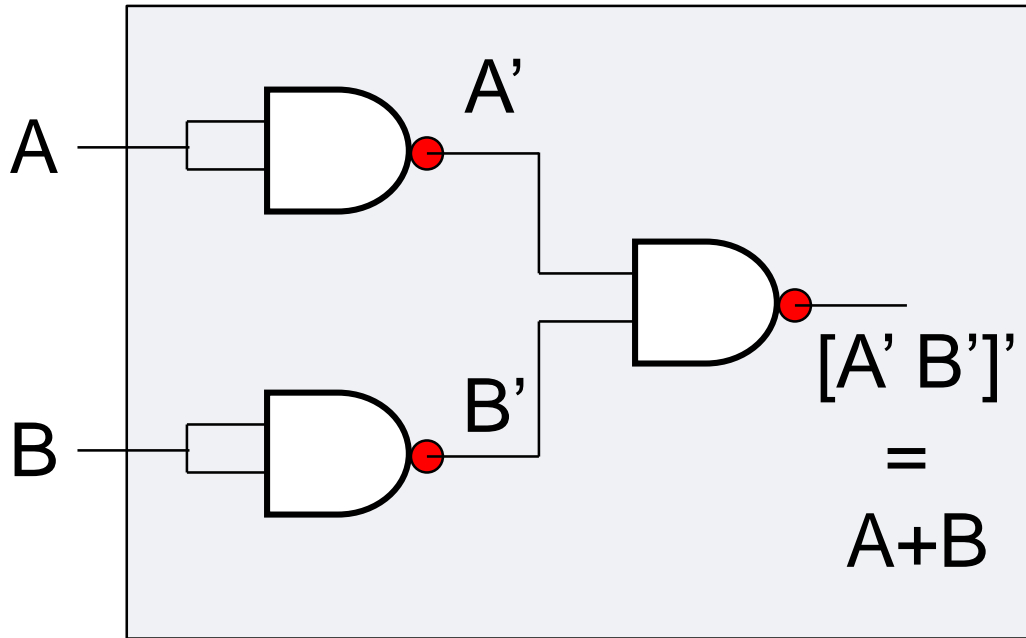
Proof



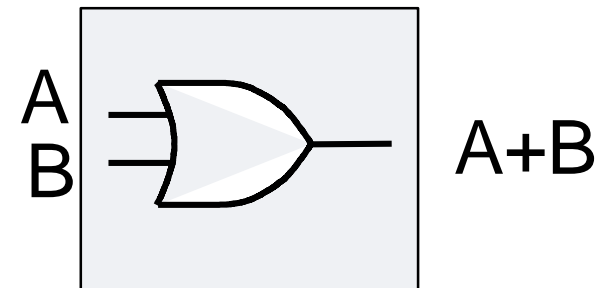
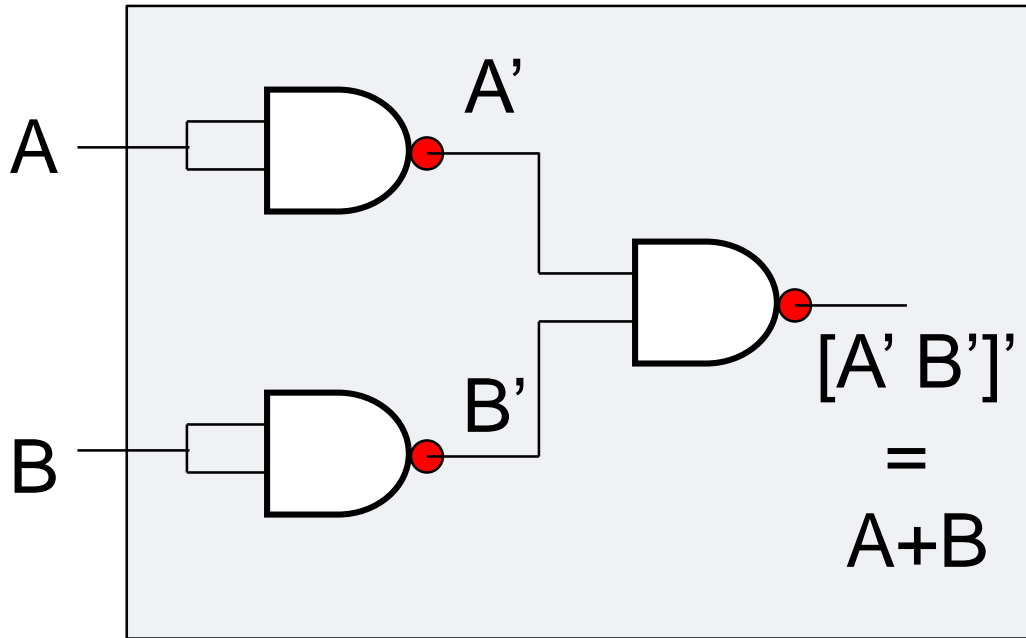
Universality of NAND: OR



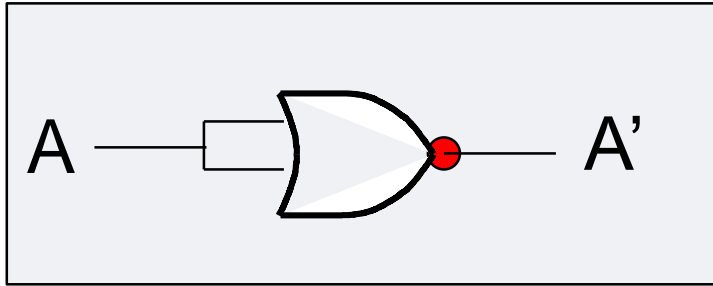
Universality of NAND: OR



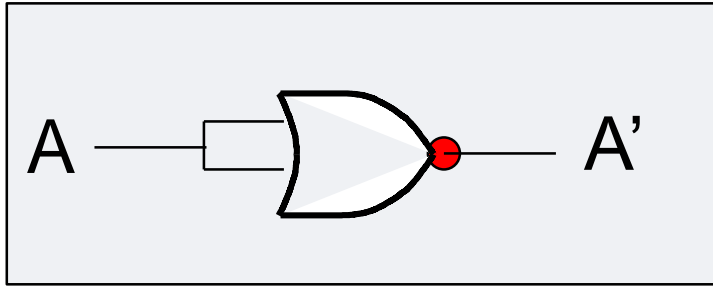
Universality of NAND: OR



Universality of NOR: NOT

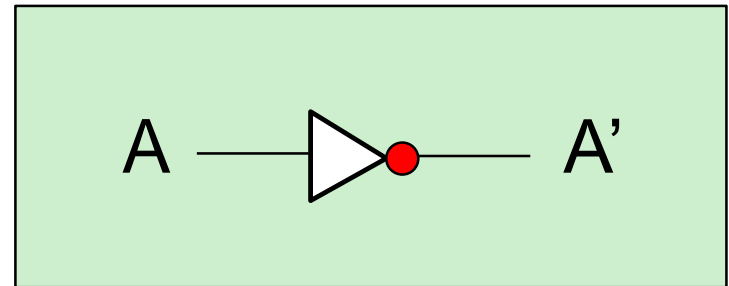
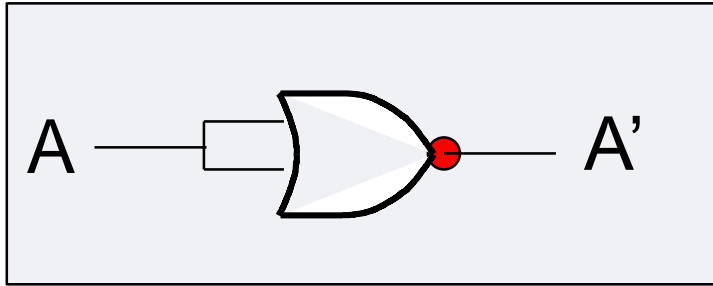


Proof



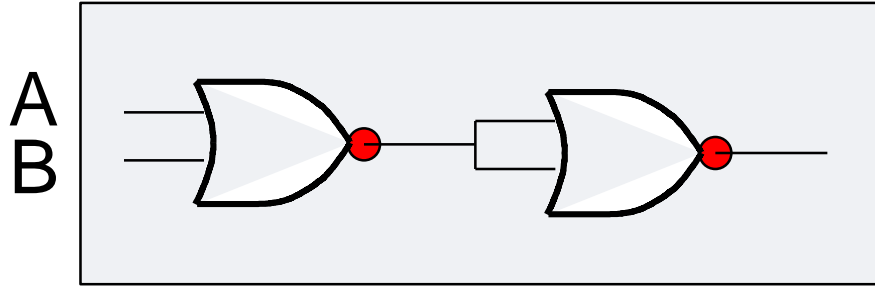
$$(A+A)' = A'$$

Proof

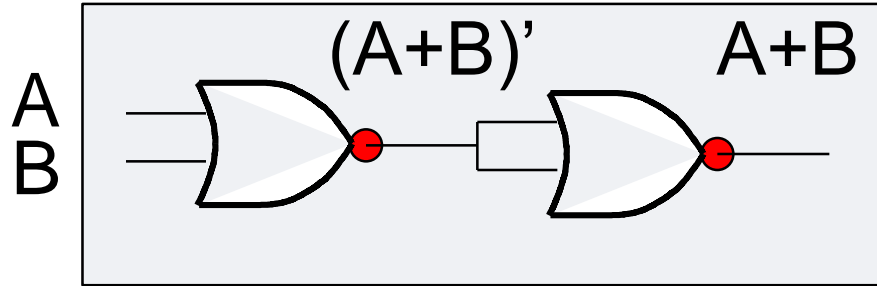


$$(A+A)' = A'$$

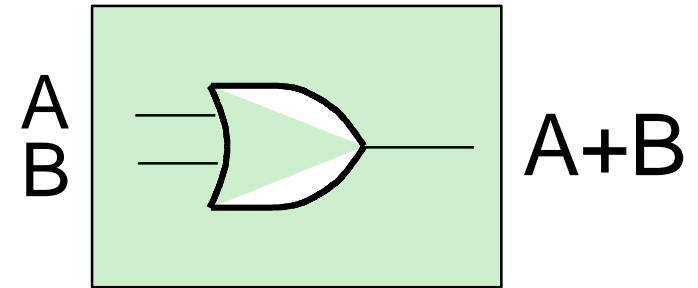
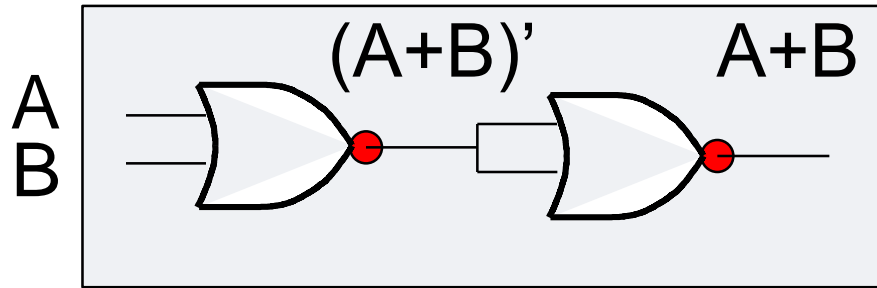
Universality of NOR: OR



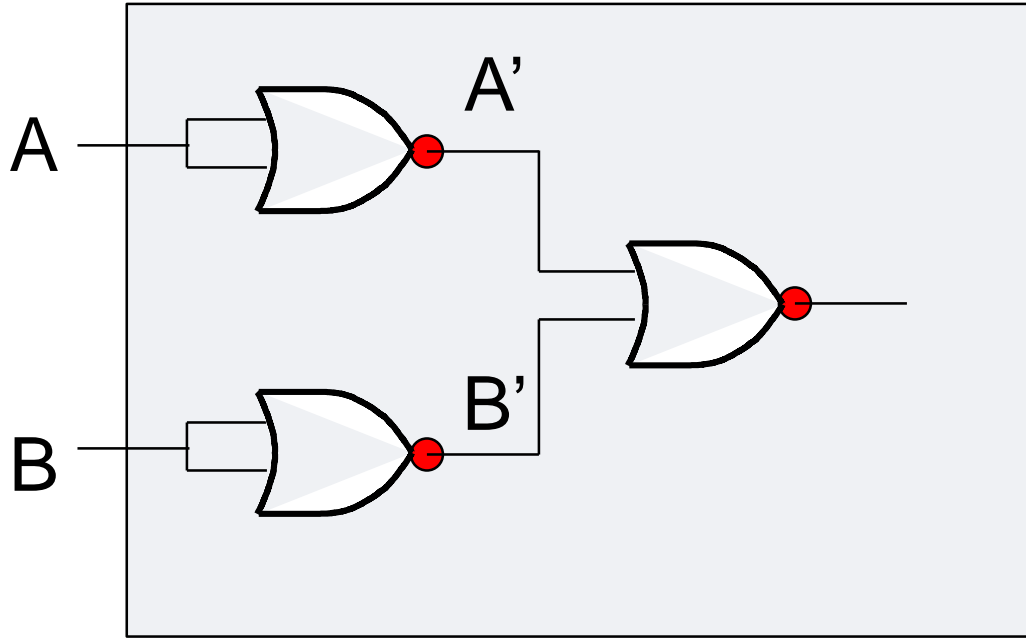
Proof



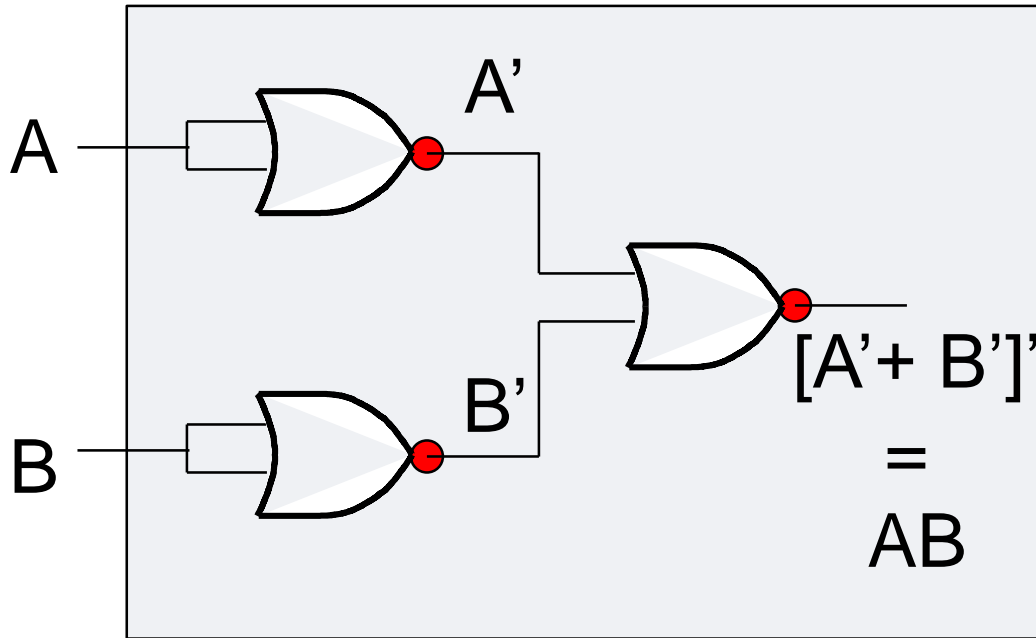
Proof



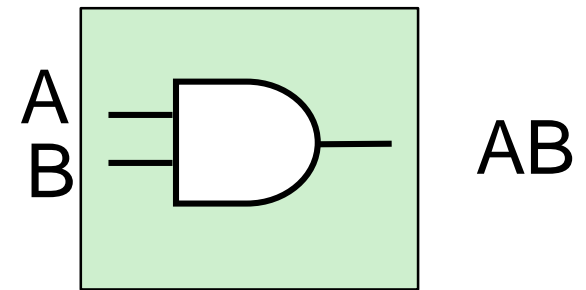
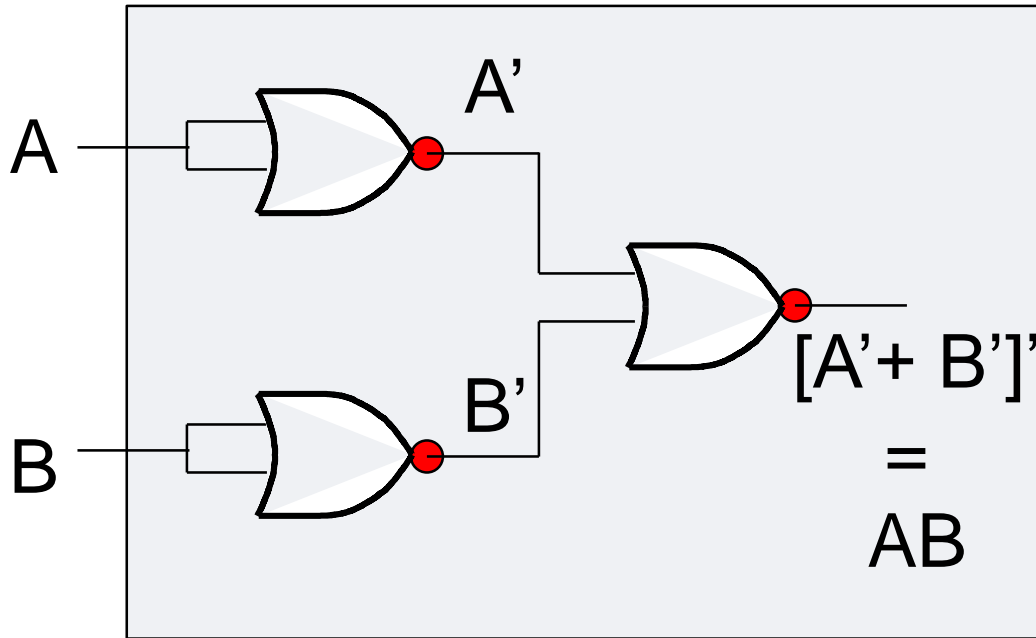
Universality of NOR: AND



Proof



Proof



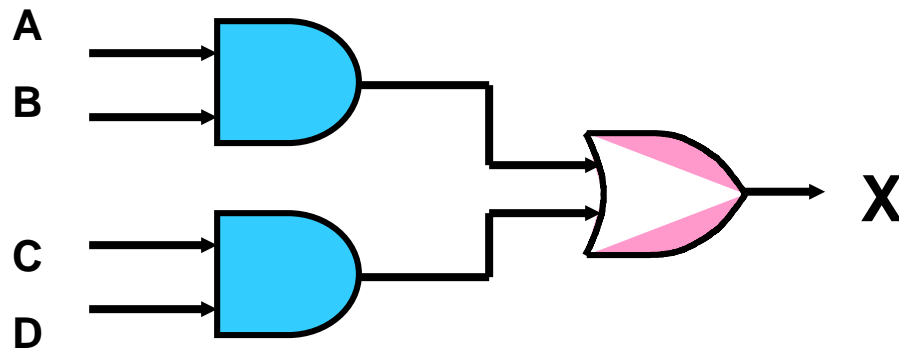
Example

Implement the Boolean function: $X = AB + CD$

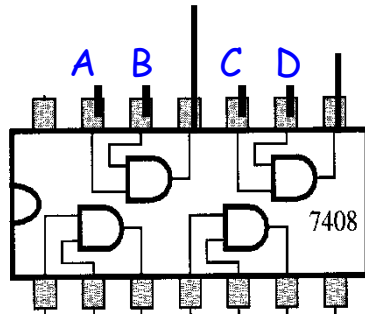
1. Using: AND, OR, NOT gates
2. Using: NAND gates

You have ... 5 minutes ...

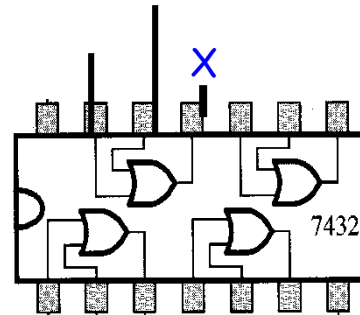
1) $X = AB + CD$; Using AND, OR, NOT gates



$X = AB + CD$; Using Chips (7408-7432)

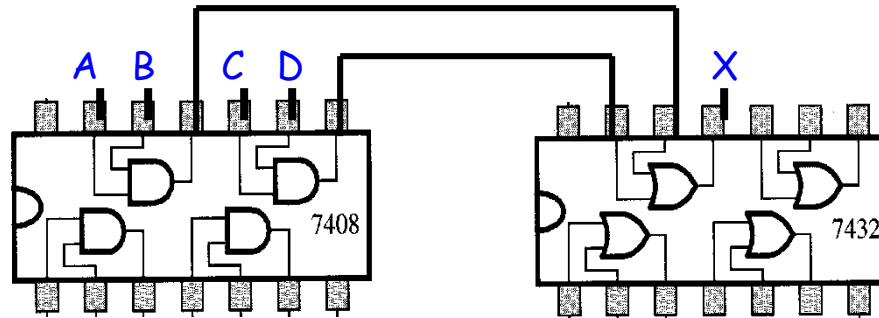


AND gates chip



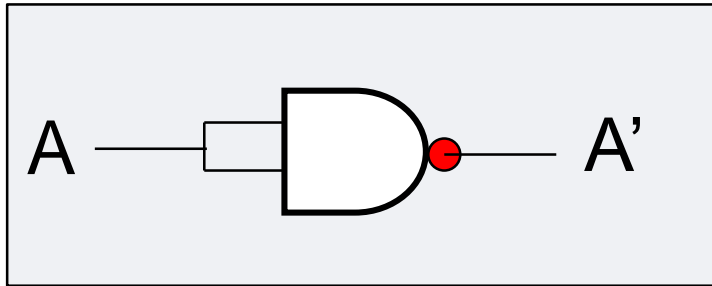
OR gates chip

2 Chips are used

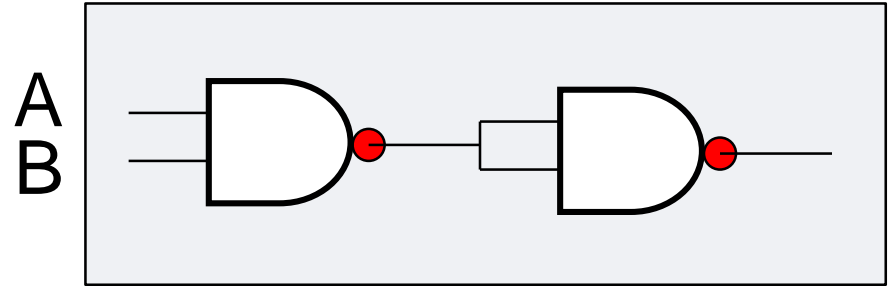


We use 2 different chips

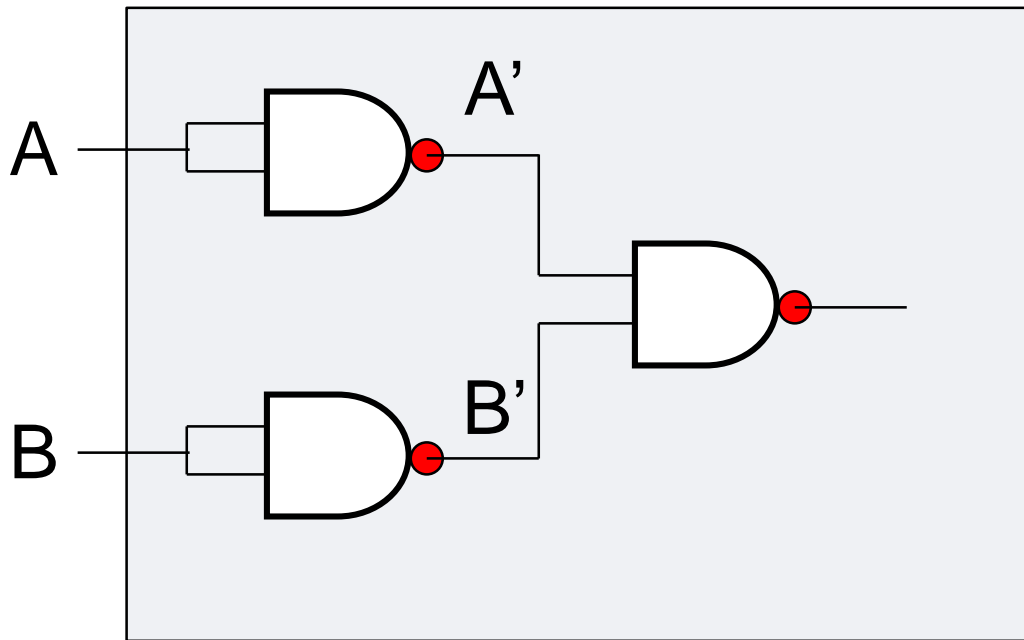
NAND gates



NOT

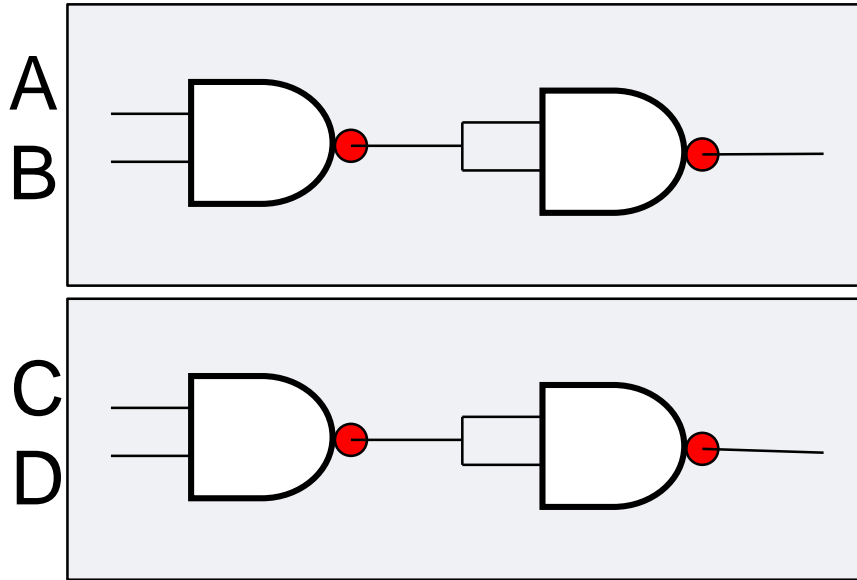


AND

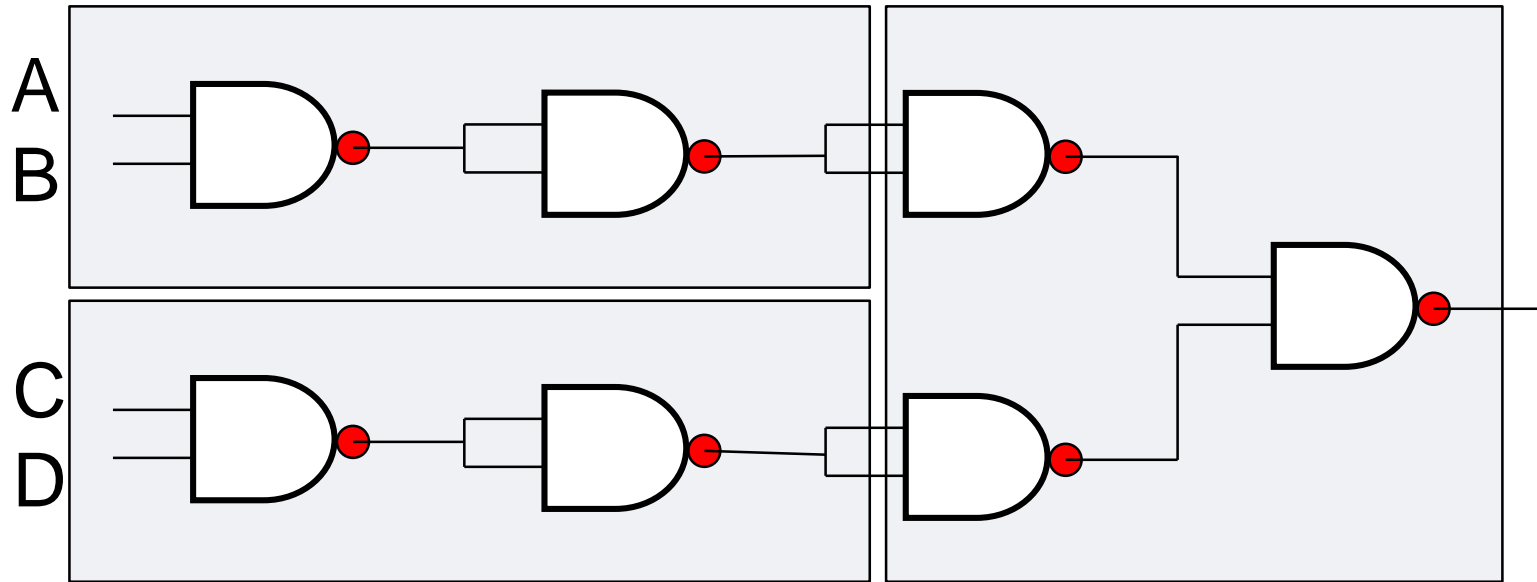


OR

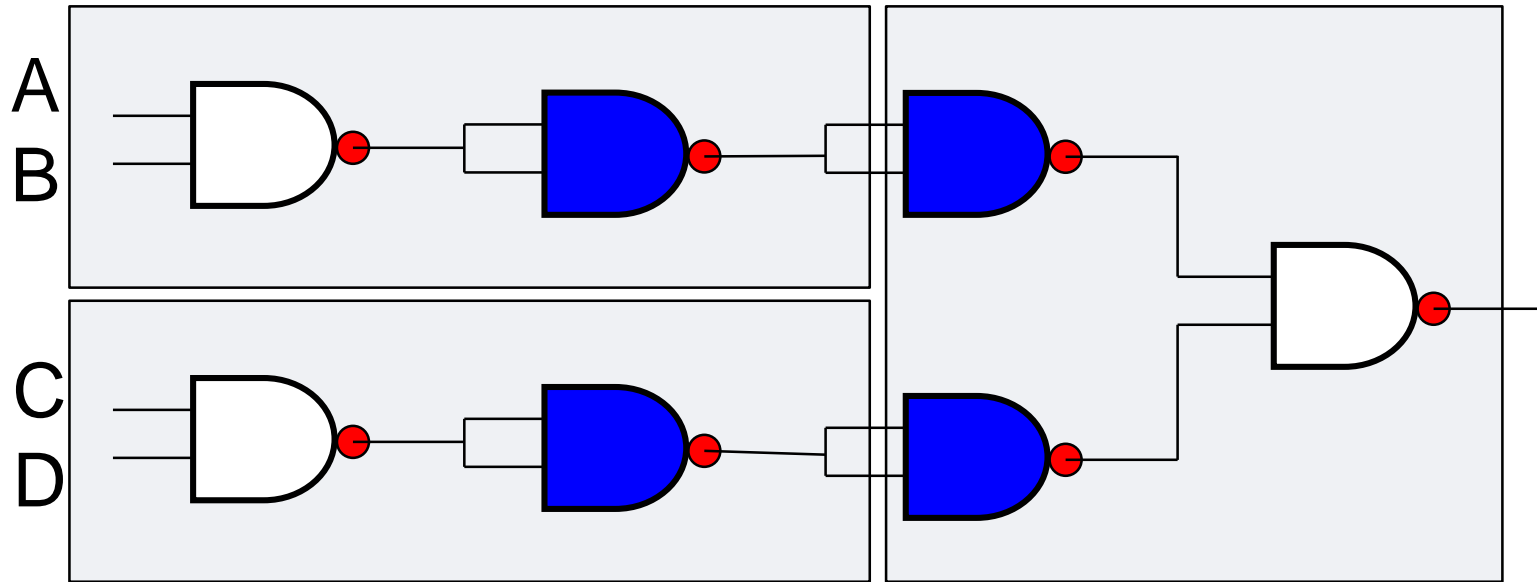
2) $X = AB + CD$; Using NAND gates



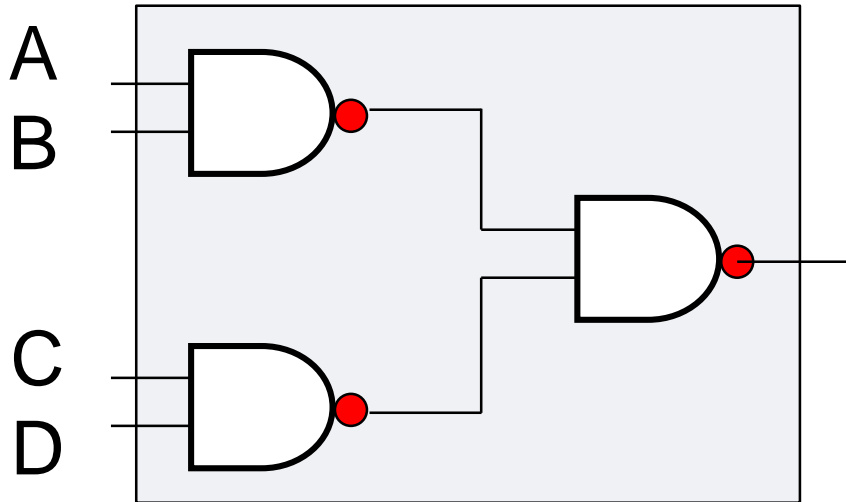
$X = AB + CD$; Using NAND gates



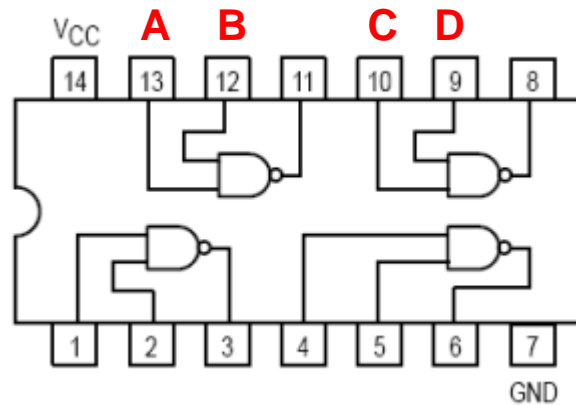
Two inverts cancel each other



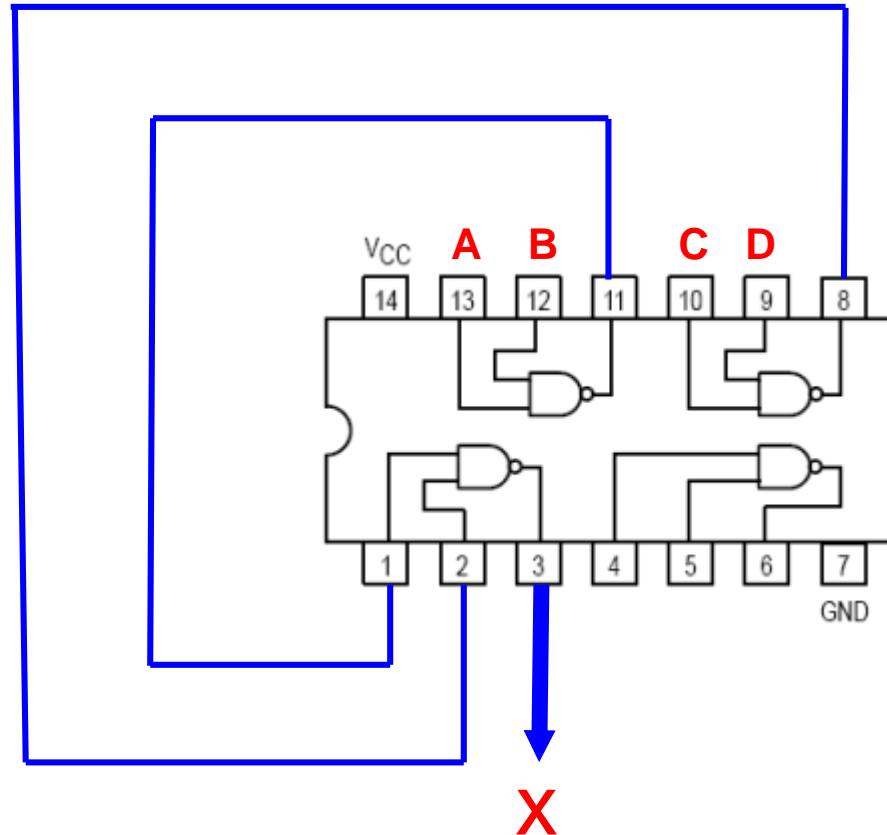
$X = AB + CD$; Using NAND gates



$X = AB + CD$; Using NAND Chip (74LS00)



$X = AB + CD$; Using NAND Chip (74LS00)



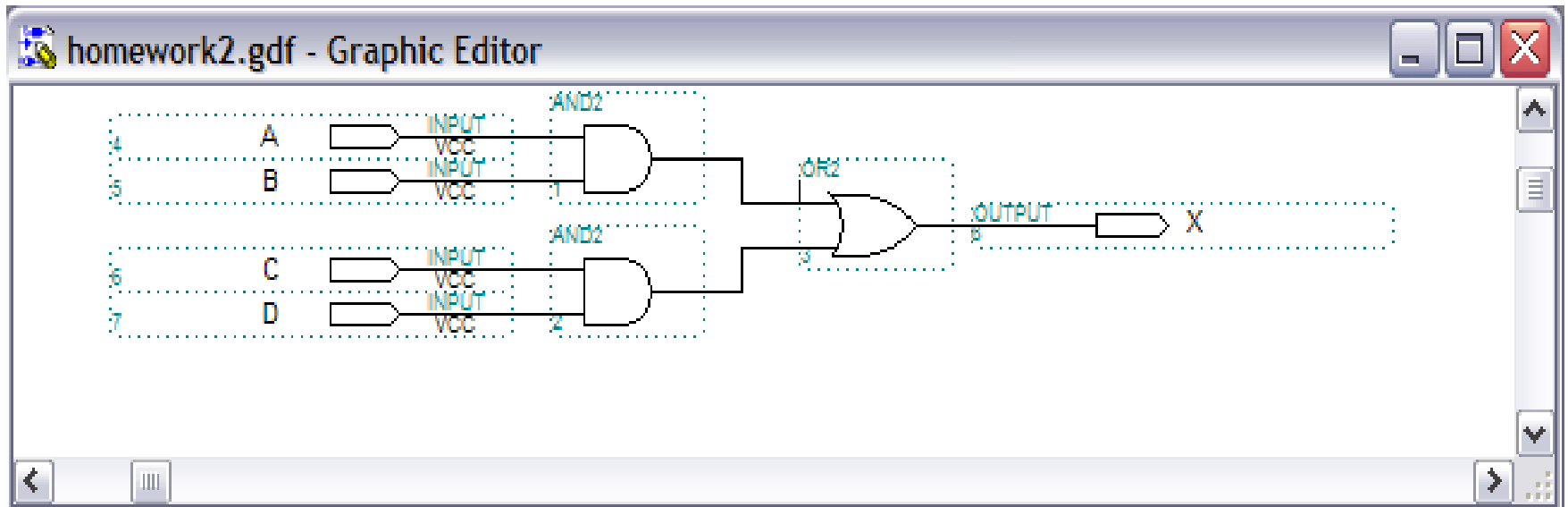
We only use 1 chip

NAND technology allows us to use less chips

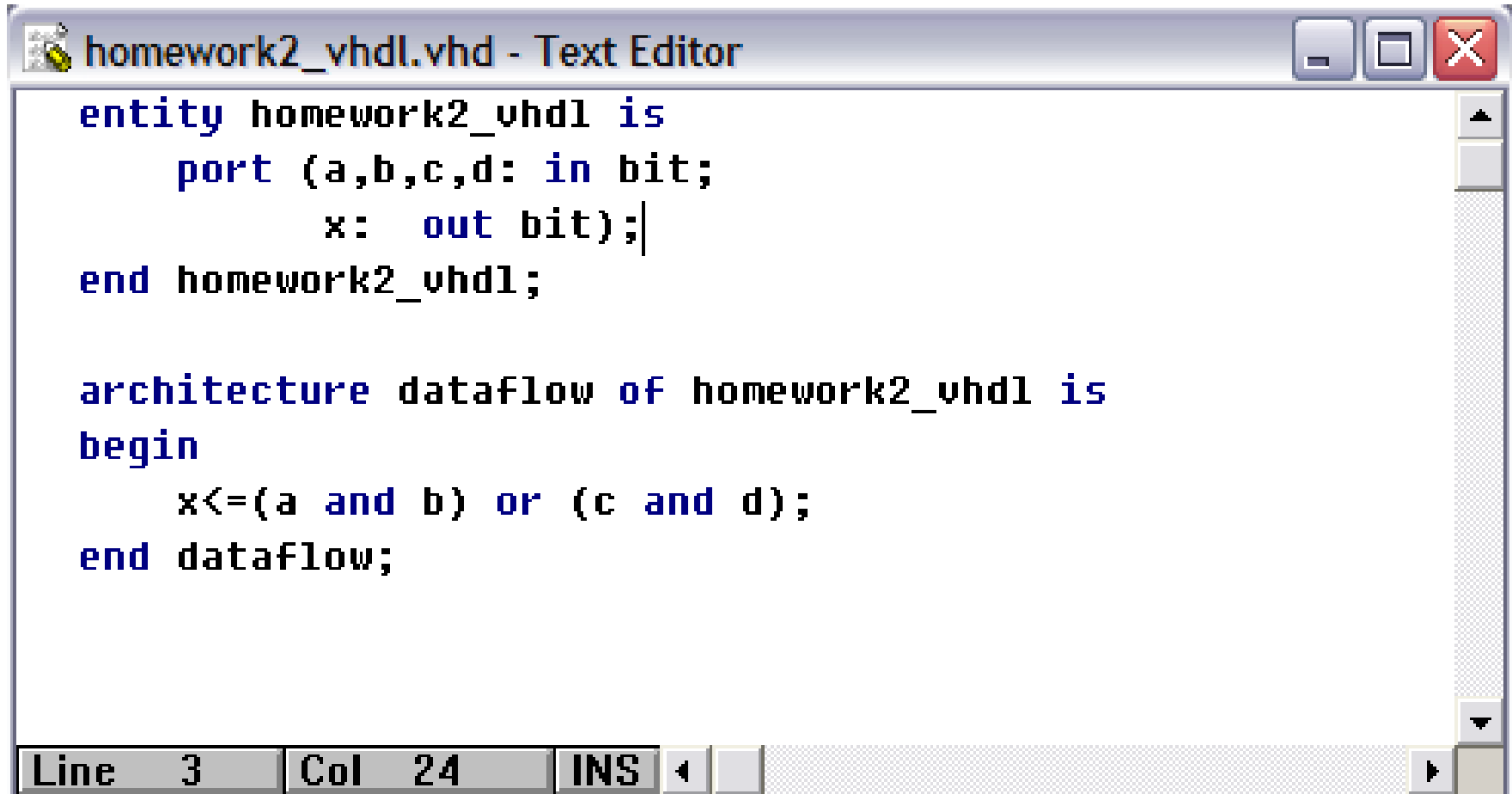
Why today's technology uses NAND gates?

- Reduces the integrated circuit complexity since NAND gates can be implemented with less transistors than the basic AND/OR gates
- Increases integrated circuit's speed
- Minimizes:
 - Production chip cost
 - Packing density of the chip.

$X = AB + CD$; Using VHDL



VHDL Code: $X = AB + CD$

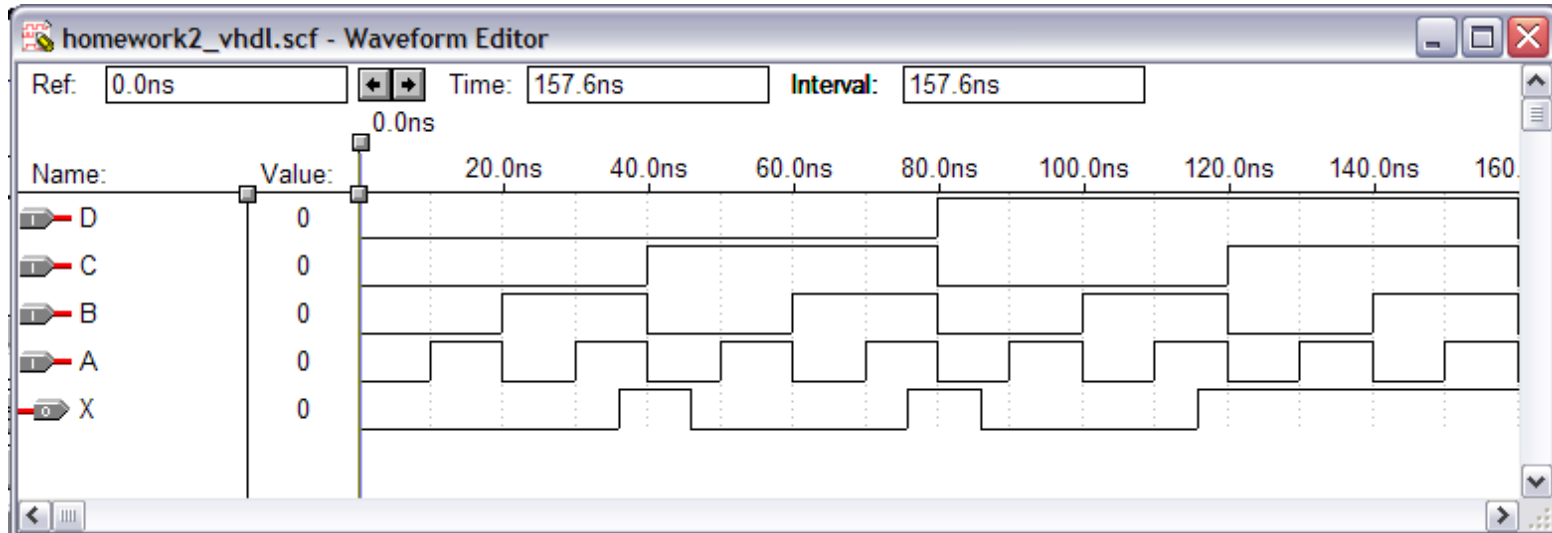


```
entity homework2_vhdl is
    port (a,b,c,d: in bit;
          x: out bit);
end homework2_vhdl;

architecture dataflow of homework2_vhdl is
begin
    x<=(a and b) or (c and d);
end dataflow;
```

Line 3 Col 24 INS

Waveform



... Two more gates

✓ XOR

✓ XNOR



OR gate

A	B	OR gate
0	0	0
0	1	1
1	0	1
1	1	1

XOR (eXclusiveOR) gate

A	B	OR gate	XOR gate
0	0	0	0
0	1	1	1
1	0	1	1
1	1	1	0

XOR (eXclusiveOR) gate



A	B	OR gate	XOR gate
0	0	0	0
0	1	1	1
1	0	1	1
1	1	1	0

$$A \text{ XOR } B = A \oplus B$$
$$= \overline{A} B + A \overline{B}$$

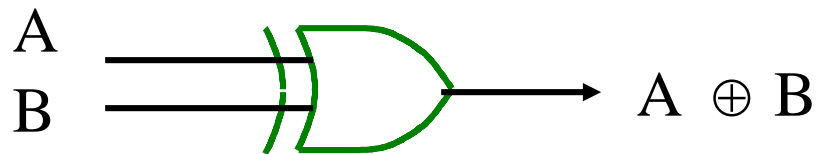
It produces a high output whenever the two inputs are at opposite levels

XOR (eXclusiveOR) gate

A	B	OR gate	XOR gate
0	0	0	0
0	1	1	1
1	0	1	1
1	1	1	0

$$A \text{ XOR } B = A \oplus B$$
$$= \overline{A} B + A \overline{B}$$

It produces a high output whenever the two inputs are at opposite levels



Another gate ... **XNOR**

$$\overline{A \oplus B} = ?$$

XNOR (eXclusiveNOR) gate

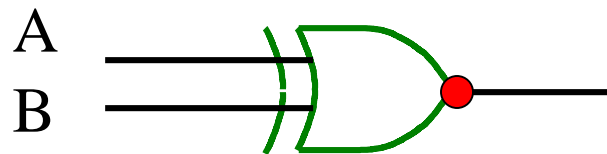
A	B	XOR gate	XNOR gate
0	0	0	1
0	1	1	0
1	0	1	0
1	1	0	1

XNOR (eXclusiveNOR) gate



A	B	XOR gate	XNOR gate
0	0	0	1
0	1	1	0
1	0	1	0
1	1	0	1

$$\begin{aligned} A \text{ XNOR } B &= \overline{A \oplus B} \\ &= \overline{A} \overline{B} + A B \end{aligned}$$



$$\overline{A \oplus B} = A \odot B$$

Total we have $2^4 = 16$ gates ...

AB
00
01
10
11

Age Group	Percentage
18-24	15%
25-34	20%
35-44	25%
45-54	20%
55-64	15%
65-74	10%
75-84	5%
85+	5%

Total we have $2^4 = 16$ gates ...

AB	0	AND		A		B	XOR	OR	NOR	XNOR	NOTB		NOTA		NAND	1
00	0	0	0	0	0	0	0	0	1	1	1	1	1	1	1	1
01	0	0	0	0	1	1	1	1	0	0	0	0	1	1	1	1
10	0	0	1	1	0	0	1	1	0	0	1	1	0	0	1	1
11	0	1	0	1	0	1	0	1	0	1	0	1	0	1	0	1

Only 7 gates are useful

AND, XOR, OR, NOR, XNOR, NOT, NAND