

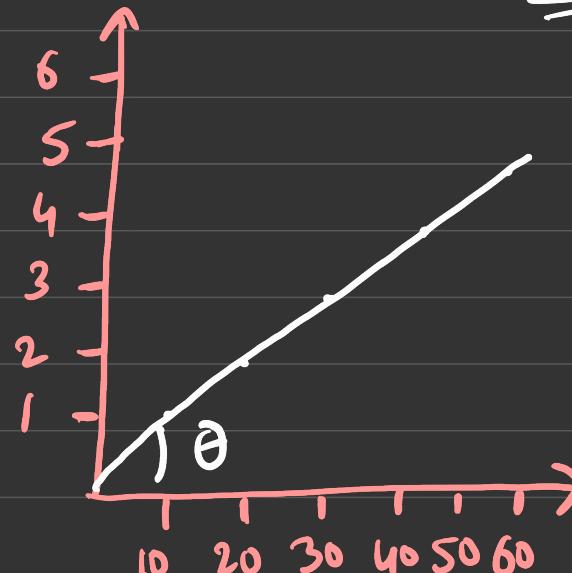
Time complexity \rightarrow TC \neq Time taken
to run a code

W8
with notes

CODE

\hookrightarrow The rate at which time
taken increases with respect
to input size

Mac



How is it calculated?

Big-oh Notation

$O(3N+1)$



$i=1$

while $i \leq N$: $\rightarrow 3N + 1$

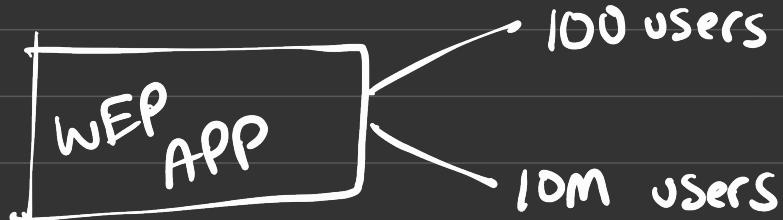
print(i)
 $i = i + 1$

$5 \times 3 + 1$

3 rules :



- 1) Always go for worst-case
- 2) Avoid/Remove constants
- 3) Avoid/Remove lower bounds



2) Ignore constants

$$O(10N^5 + 5N^2 + 15)$$

$$N = 10^5$$

$$O(10 \times 10^{25} + 5 \times 10^{10} + 15)$$

$$1000000$$

3) Avoid lower bounds

$$O(10 \times 10^{25} + 5 \times 10^{10})$$

$$\approx O(10N^5) \approx O(N^5)$$



Big-oh (O)



worst-case
[Upper bound]

Theta (Θ)



Average-case

Omega (Ω)



Best-case
[Lower-bound]

marks = 55 < 60

```
if marks > 90: print("A")
elif marks > 80: print("B")
elif marks > 70: C
elif marks > 60: D
else:
```

F

~~Ex 1~~

```
for i in range(1, n+1):
```

for i in range(1, n+1):

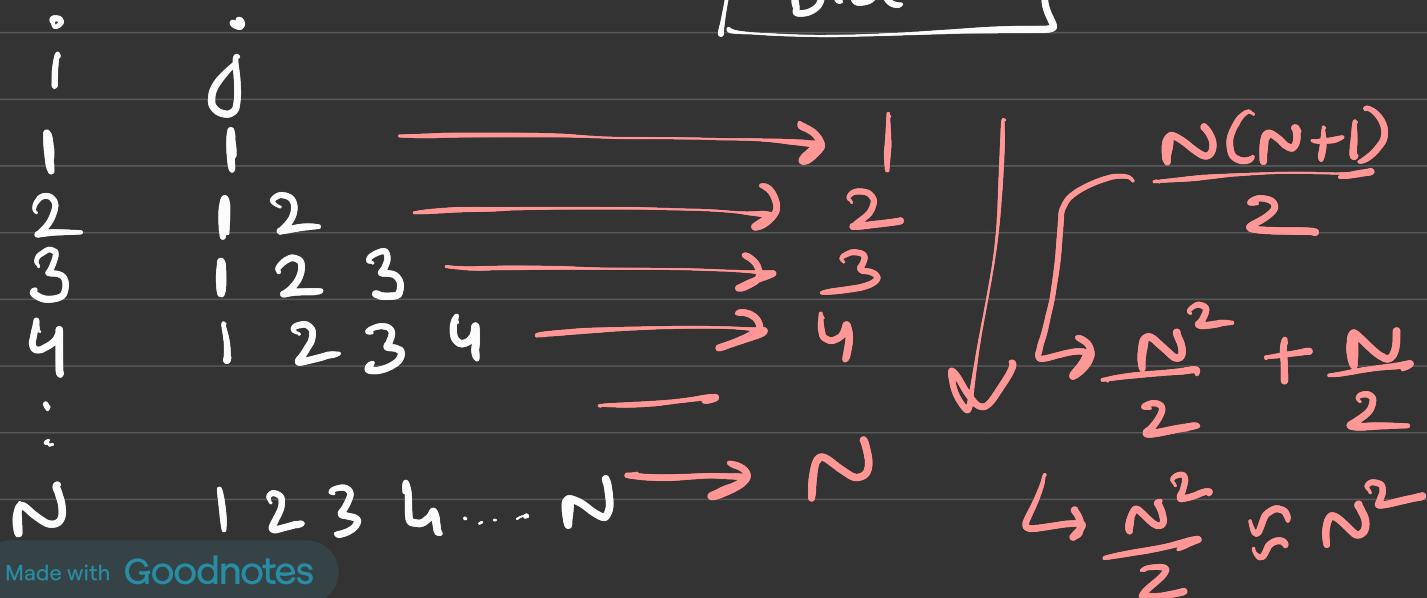
Block

i 0
1 1 2 3 ... N N N N N → N x N
2 1 2 3 ... N
3 1 2 3 ... N
...
N 1 2 3 ... N

Ex 2

for i in range (1, n+1):
 for j in range(1, i+1):

 [Block]



Space Complexity

Big Oh

SC → Auxiliary space → It is the extra space used to solve the algorithm

Input Space

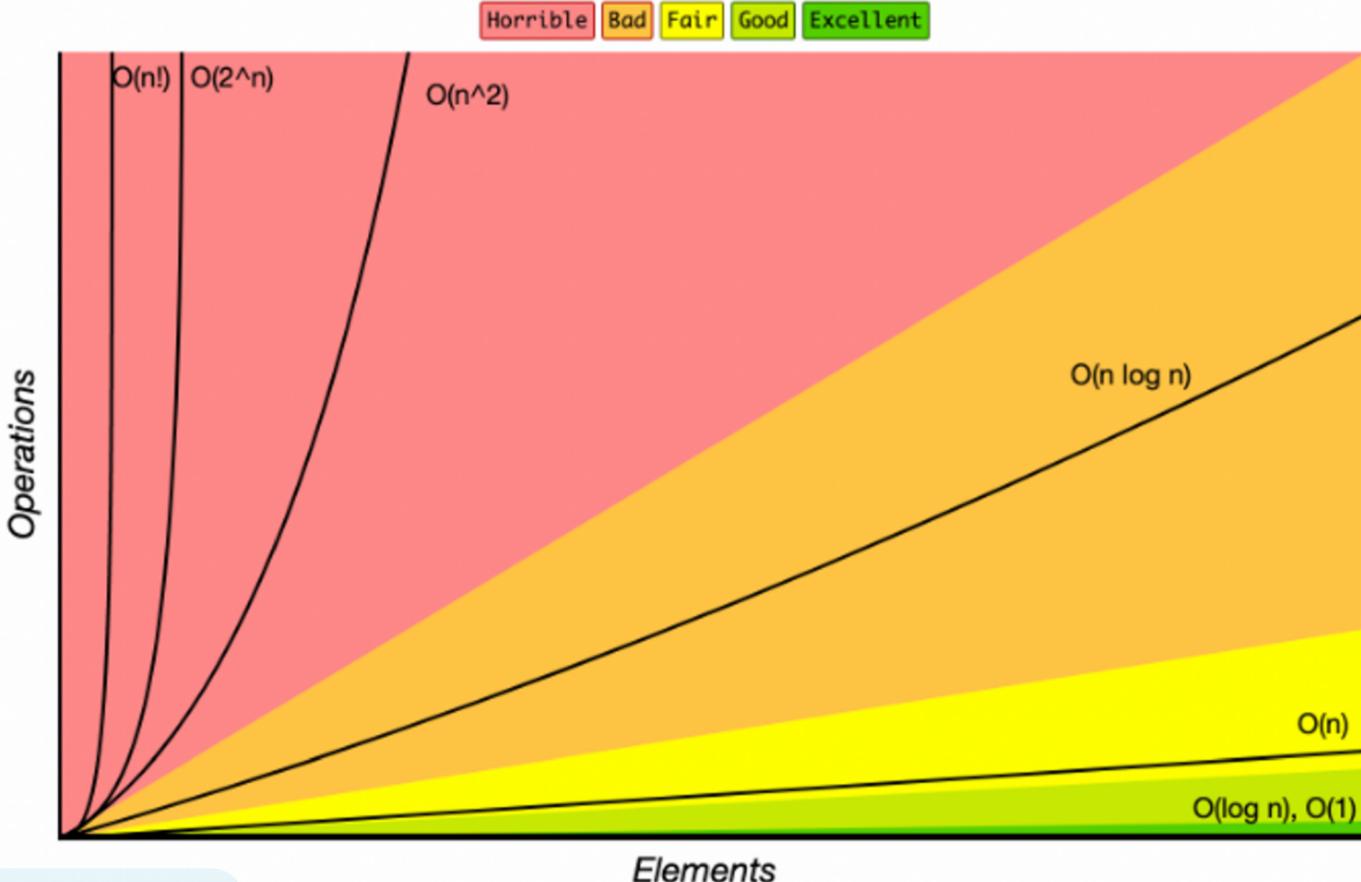
Input + Space

$$a = 5, b = 10, c = 12$$

It is the space to store the input

total = $a+b+c$
print (total)

→ Auxiliary Space



Operation	Average Case	Amortized Worst Case
Copy	$O(n)$	$O(n)$
Append[1]	$O(1)$	$O(1)$
Pop last	$O(1)$	$O(1)$
Pop intermediate[2]	$O(n)$	$O(n)$
Insert	$O(n)$	$O(n)$
Get Item	$O(1)$	$O(1)$
Set Item	$O(1)$	$O(1)$
Delete Item	$O(n)$	$O(n)$
Iteration	$O(n)$	$O(n)$
Get Slice	$O(k)$	$O(k)$
Del Slice	$O(n)$	$O(n)$
Set Slice	$O(k+n)$	$O(k+n)$
Extend[1]	$O(k)$	$O(k)$
Sort	$O(n \log n)$	$O(n \log n)$
Multiply	$O(nk)$	$O(nk)$
x in s	$O(n)$	
min(s), max(s)	$O(n)$	
Get Length	$O(1)$	$O(1)$

arr = {1, 5, 3, 2, 6, 7}

100

[1, 5, 3, 2, 6, 7, 100]
0 1 2 3 4 5 6

a[3]

if 10 in arr: