

## Q.6 Prim's Algorithm For MST.

- i) A spanning tree of a connected undirected graph  $G$ , is a sub-graph of  $G$  which is a tree that connects all vertices together.
- ii) A MST is defined as a spanning tree with weight less than or equal to the weight of every other spanning tree.
- iii) Prim's algorithm is a greedy algorithm that used to form MST for a connected weighted undirected graph.
- iv) In other words, the algorithm builds a tree that includes every vertex and a subset of the edges in such a way that the total weight of all the edges in the tree is minimized.

### • Algorithm :

PRIM ( $G, n$ )

{

cost  $\leftarrow 0$

for  $i \leftarrow 1$  to  $n-1$  do

MST [ $i$ ]  $\leftarrow 0$

MST [ $0$ ]  $\leftarrow 2$

for  $k \leftarrow 1$  to  $n$  do

$\text{min\_dist} \leftarrow \infty$

    for  $i \leftarrow 1$  to  $n-1$  do

        for  $j \leftarrow 1$  to  $n-1$  do

            if  $a[i, j]$  AND ((MST[i] AND MST[j])  
                OR (MST[i] AND MST[j])) then

                if  $a[i, j] < \text{min\_dist}$  then

$\text{min\_dist} \leftarrow a[i, j]$

$u \leftarrow i$

$v \leftarrow j$

        print ( $u, v, i, \text{min\_dist}$ )

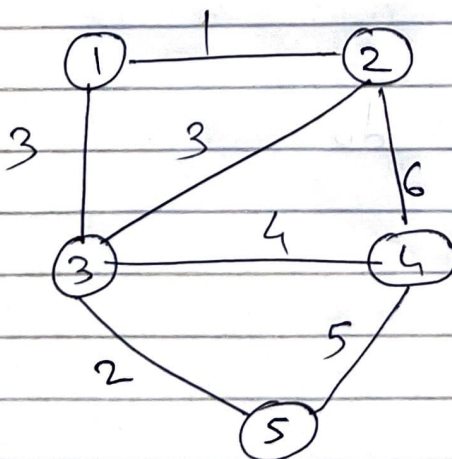
        MST[u]  $\leftarrow$  MST[v]  $\leftarrow$  1

        cost  $\leftarrow$  cost + min\_dist

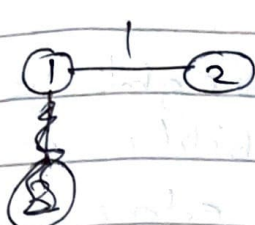
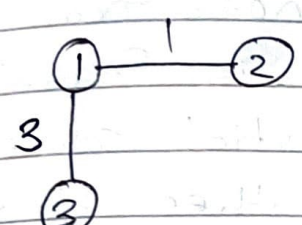
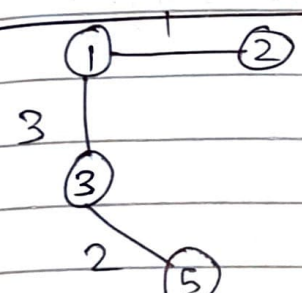
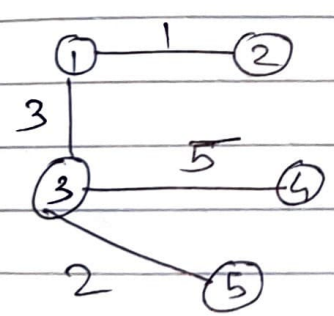
print ("Total cost = ", cost)

• The time complexity is  $O(n^2)$ .

• For e.g.,





Partial Solution	Set of neighbour edges	Cost	Updated $U_v$
1) $U_v = \{1, 2, 3, 4, 5\}$ Initial partial solution ①	$\langle 1, 2 \rangle$ $\langle 1, 3 \rangle$	1 3	$\{2, 3, 4, 5\}$
2) 	$\langle 1, 3 \rangle$ $\langle 2, 3 \rangle$ $\langle 2, 4 \rangle$	3 3 6	$\{3, 4, 5\}$
3) 	$\langle 2, 3 \rangle$ $\langle 2, 4 \rangle$ $\langle 3, 4 \rangle$ $\langle 3, 5 \rangle$	3 6 4 2	$\{4, 5\}$
4) 	$\langle 2, 4 \rangle$ $\langle 3, 4 \rangle$ $\langle 5, 4 \rangle$	6 4 5	$\{4\}$
5) 			

∴ Cost of solution

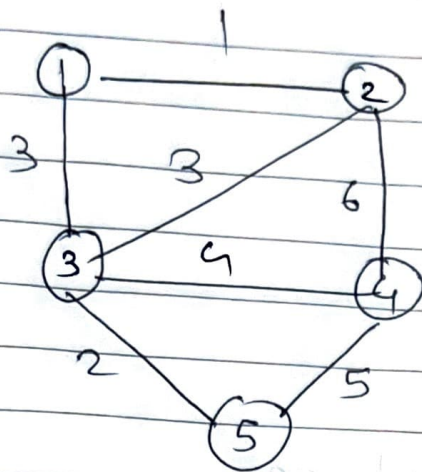
$$\begin{aligned}
 &= w(1, 2) + w(1, 3) + w(3, 4) + w(3, 5) \\
 &= 1 + 3 + 5 + 2 \\
 &= 10
 \end{aligned}$$

### Q.7 Kruskal's Algorithm

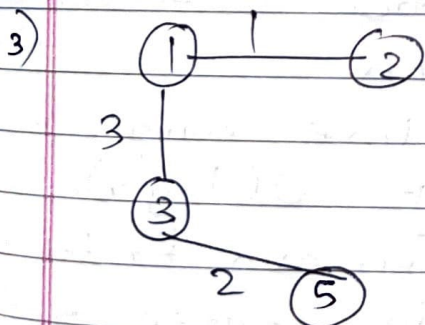
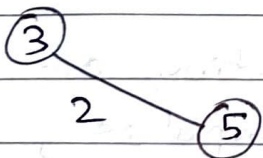
- ⇒ i) Kruskal's algorithm is used to find the MST for a connected weighted graph.
- ii) This algorithm first sorts all the edges in non-decreasing order of their weight.
- iii) Edge with minimum weight is selected and its feasibility is tested.
- iv) IF inclusion of the edge to a partial solution does not form the cycle, then the edge is feasible and added to the partial solution.
- v) IF it is not feasible then skip it and check for the next edge.



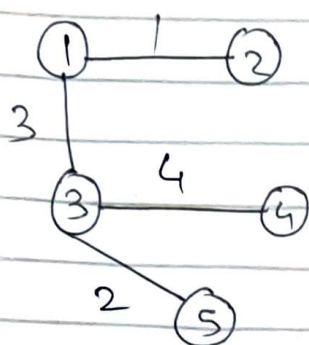
For e.g.,



Edge	$\langle 1, 2 \rangle$	$\langle 3, 5 \rangle$	$\langle 1, 3 \rangle$	$\langle 2, 3 \rangle$	$\langle 3, 4 \rangle$	$\langle 4, 5 \rangle$	$\langle 2, 4 \rangle$
cost	1	2	3	3	4	5	6



$\langle 2, 3 \rangle$  edges creates cycle so remove it



Cost = 10 //

### # Difference between

Prim's Algorithm	Kruskal's Algorithm
i) Prim's algorithm always chooses the next edge which is a neighbour of vertices in partially generated solution.	i) The Kruskal algorithm chooses the edge having a minimum weight.
ii) Prim's algorithm ensures that partial solution is always a tree.	ii) In Kruskal's algorithm, a partial solution can be a forest.
iii) Sorting of edges is not required.	iii) Sorting of edges is compulsory.
iv) In Prim's algorithm graph must be a connected graph.	iv) Kruskal's can be function on disconnected graph.
v) Time complexity of $O(n^2)$	v) Time complexity of $O(\log N)$ .
vi) Can run faster in dense graph	vi) Can run faster in sparse graph.