

Q.P
10/04/2024

Assignment No: 1

Q.1 With references to assembler, consider any assembly program and generate Pass-1 and Pass-2 output. Also show contents of Database table involved in it.

⇒

ANITA

ADD

- Assembly program:

START 1000H :

L DAD ALPHA A100H

D STA BETA A101H

HLT

ALPHA RES 1

BETA RES 2

END START

- Pass-1 output (Symbol Table):

Symbol Address

START 1000

ALPHA 100

BETA 101

- Pass-2 output (Machine code):

100 0000 1010

101 0001 1010

102 0010 1010

2.019 Assembly Language

- Database table view of assembly file

In I part discuss how memory addressed

Address Label Opcode Operand

200	START	.hi ni bavani shah
201	LDA	ALPHA
202	STA	BETA
203	HLT	TRAP
204	ALPHA A19JA RES1	1
205	BETA AT38 RES2	2
		TRAP

- In Pass-1, the assembler reads assembly code, assigns addresses to symbols, and creates a symbol table.
- In Pass-2, the assembler translates assembly instruction into machine codes using the symbol table generated in Pass-1.
- The database table shows the address, label, opcode and operand for each instruction or data declaration in the program.

(Label and op) bug two C-2219

0101 0000 0011

0101 1000 0011

0101 0100 0011

- Q. 2. Write a short note on YACC.
- ⇒ Each translation rule input to YACC has a string specification that resembles a production of grammar it has a non-terminal on the LHS and a few alternatives on the RHS.
 - For simplicity, we will refer to a string specification as a production. YACC generates an LALR(1) parser for language L from the productions, which is a bottom-up parser.
 - The parser would operate as follows: For a shift action, it would invoke the scanner to obtain the next token and continue the parse by using token.
 - While performing a reduced action in accordance with a production, it would perform the semantic action associated with that production.
 - The semantic actions associated with productions achieve the building of an intermediate representation or target code as follows:
 - 1) Every non-terminal symbol in the parser has an attribute, say s_1, s_2, \dots, s_n .
 - 2) The semantic action uses the values of these attributes for building the intermediate representation or target code.
 - A parser generator is a program that takes as input a specification of a syntax and produces as output procedure for recognizing that language. Historically, they are also called compile compilers.

- Input File: YACC input file is divided into three parts:
 - **symbol table**: It keeps track of all symbols and their definitions with their respective types.
 - **/* definitions**: It contains declarations for tokens used in the grammar.
 - **/* rules**: It contains the grammar rules in BNF format.
- Auxiliary routines: It contains code for various utility functions.

Definition Part: This part includes the information about the tokens used in syntax definition. The definition part can include C code within the tokens themselves for the definition of the parser and variable declarations, within `#if` and `#endif` blocks in the first column, and `#else`, `#endif`, `#if`, and `#endif` blocks in the last column.

- Rule part: The rule part contains grammar definitions in a modified BNF format. Action parts of code in curly braces {} can be embedded inside, also known as inline actions or local functions.
- Auxiliary Routine Part: The auxiliary routines part is only C code. It includes function definitions for every function needed in the rules part.