

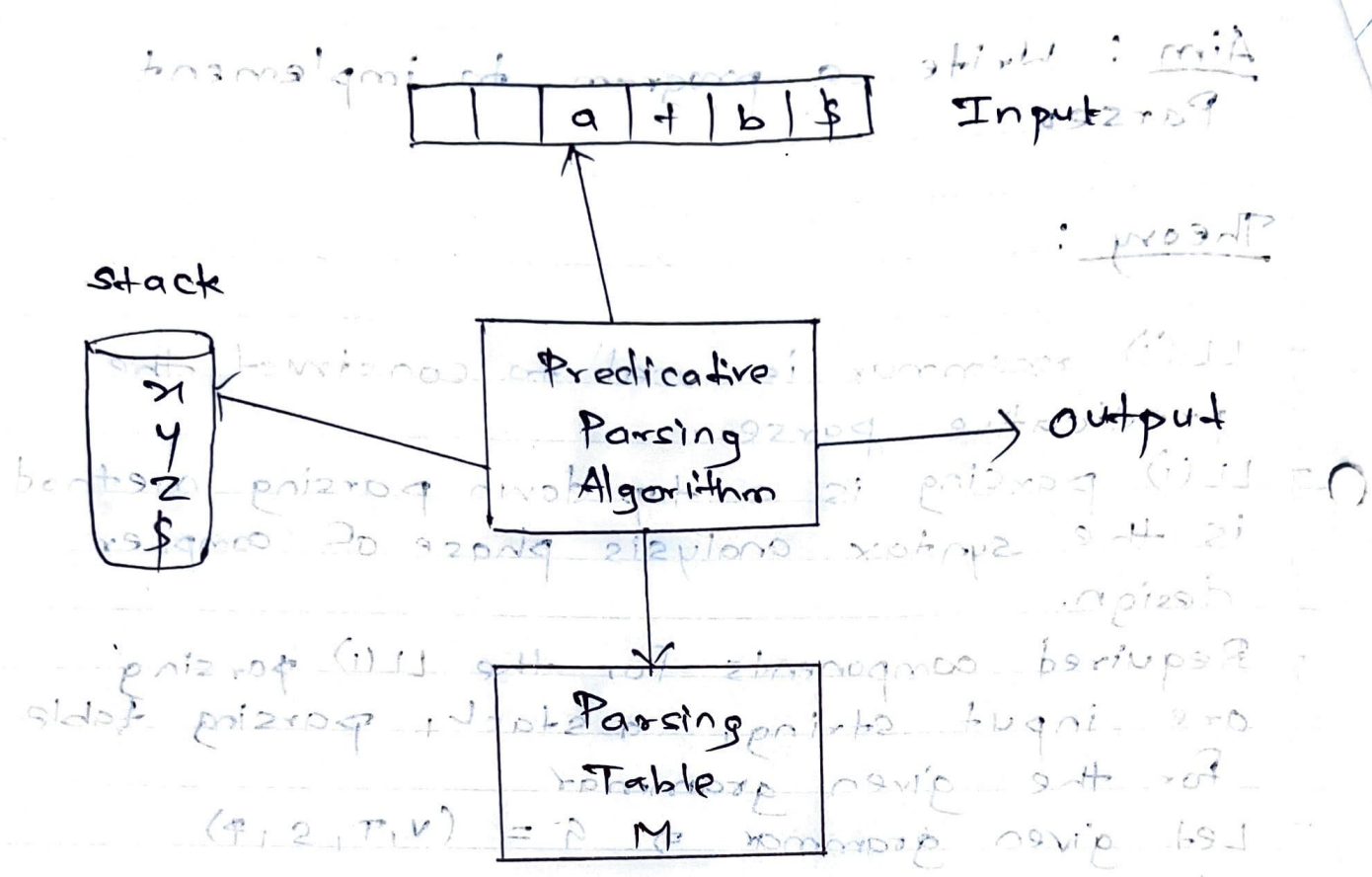
Experiment No. 4

Aim : Write a program to implement Parser.

Theory :

- LL(1) grammar is used to construct the predictive parser.
- LL(1) parsing is a top down parsing method is the syntax analysis phase of compiler design.
- Required components for the LL(1) parsing are input string, a stack, parsing table for the given grammar.
- Let given grammar $G = (V, T, S, P)$ where,
 - V = set of variables
 - T = terminal symbols set
 - S = start symbol
 - P = production set.
- Non-recursive predictive parser is also known as LL(1) parser.
- The LL(1) parser stands for :
 - "1" : left to right scan of input
 - "L" : Leftmost derivation
 - An LL parser is called an LL(K) parser if uses.
 - "K" : It uses K tokens of look ahead when parsing sentence.

Model of Non-Recursive Predicative Parser :



$V = \text{set of variables}$
 $T = \text{terminal symbols}$
 $Z = \text{start symbol}$
 $P = \text{production set}$

The Non-recursive predicative parser is also known as LL(1) parser. The LL(1) parser stands for: "L" is left to right scan of input, "L" is leftmost derivation, and "1" is An LL parser is called an LL(1) parser. It uses k tokens of look ahead when parsing sentence.

Algorithm :

Input : Grammar G

Output : Parsing Table M

Step 1 : For each terminal 'a' in $FIRST(\alpha)$
ADD $A \rightarrow \alpha$ to $M[A, a]$

Step 2 :

CASE 1

\Rightarrow IF ϵ is in $FIRST(\alpha)$ then
for each terminal b in $FOLLOW(A)$
ADD $A \rightarrow \alpha$ to $M[A, b]$

Case 2 :

\Rightarrow IF ϵ is in $FIRST(\alpha)$ and $\$$ in $FOLLOW(A)$
then for each terminal b in $FOLLOW(A)$
Add $A \rightarrow \alpha$ to $M[A, b]$

Step 3 : Make each undefined entry of M be an error.

- Conclusion : In this we have implemented LL(1) parser algorithm successfully.

28/03/2024

For e.g.,
 $E \mapsto TE'$, $E' \mapsto TE' \mid \varepsilon$, $T \mapsto FT'$,

$$T' \rightarrow * F T' \mid \varepsilon, F \rightarrow (E) \mid \text{id}$$
$$F_1(E) \text{ lid}$$

Recommend : 4000

code (x) P29P7 / ci 21 3 7E 6

$[d, A]M \neq \lambda \cdot A \cdot dA$

(*) ω_1 is a limit ordinal $\Rightarrow \omega_1$ is a limit ordinal

0

$$[d, A]M \quad \text{at} \quad x \in A \quad b.b.A$$

M 3d no record

11/17/2022 Multiphase (MPP)

Code:

```

from collections import OrderedDict

def isterminal(char):
    if (char.isupper() or char == "`"):
        return False
    else:
        return True

def insert(grammar, lhs, rhs):
    if (lhs in grammar and rhs not in grammar[lhs] and grammar[lhs] !=
    "null"):
        grammar[lhs].append(rhs)
    elif (lhs not in grammar or grammar[lhs] == "null"):
        grammar[lhs] = [rhs]
    return grammar

def first(lhs, grammar, grammar_first):
    rhs = grammar[lhs]
    for i in rhs:
        k = 0
        flag = 0
        current = []
        confirm = 0
        flog = 0
        if (lhs in grammar and "`" in grammar_first[lhs]):
            flog = 1
        while (1):
            check = []
            if (k >= len(i)):
                if (len(current) == 0 or flag == 1 or confirm == k or flog ==
1):
                    grammar_first = insert(grammar_first, lhs, "`")
                    break
            if (i[k].isupper()):
                if (grammar_first[i[k]] == "null"):
                    grammar_first = first(i[k], grammar, grammar_first)
                for j in grammar_first[i[k]]:
                    grammar_first = insert(grammar_first, lhs, j)
                    check.append(j)
            else:
                grammar_first = insert(grammar_first, lhs, i[k])

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        check.append(i[k])
    if (i[k] == "`"):
        flag = 1
    current.extend(check)
    if ("`" not in check):
        if (flog == 1):
            grammar_first = insert(grammar_first, lhs, "`")
            break
        else:
            confirm += 1
            k += 1
            grammar_first[lhs].remove("`")
    return (grammar_first)

def rec_follow(k, next_i, grammar_follow, i, grammar, start,
grammar_first,
            lhs):
    if (len(k) == next_i):
        if (grammar_follow[i] == "null"):
            grammar_follow = follow(i, grammar, grammar_follow, start)
            for q in grammar_follow[i]:
                grammar_follow = insert(grammar_follow, lhs, q)
        else:
            if (k[next_i].isupper()):
                for q in grammar_first[k[next_i]]:
                    if (q == "`"):
                        grammar_follow = rec_follow(k, next_i + 1, grammar_follow, i,
                                                    grammar, start, grammar_first,
lhs)
                    else:
                        grammar_follow = insert(grammar_follow, lhs, q)
            else:
                grammar_follow = insert(grammar_follow, lhs, k[next_i])

    return (grammar_follow)

def follow(lhs, grammar, grammar_follow, start):
    for i in grammar:
        j = grammar[i]
        for k in j:
            if (lhs in k):
                next_i = k.index(lhs) + 1

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        grammar_follow = rec_follow(k, next_i, grammar_follow, i,
grammar,
                                start, grammar_first, lhs)

    if (lhs == start):
        grammar_follow = insert(grammar_follow, lhs, "$")
    return (grammar_follow)

def show_dict(dictionary):
    for key in dictionary.keys():
        print(key + " : ", end="")
        for item in dictionary[key]:
            if (item == "`"):
                print("Epsilon, ", end="")
            else:
                print(item + ", ", end="")
        print("\b\b")

def get_rule(non_terminal, terminal, grammar, grammar_first):
    for rhs in grammar[non_terminal]:
        #print(rhs)
        for rule in rhs:
            if (rule == terminal):
                string = non_terminal + "~" + rhs
                return string

            elif (rule.isupper() and terminal in grammar_first[rule]):
                string = non_terminal + "~" + rhs
                return string

def generate_parse_table(terminals, non_terminals, grammar,
grammar_first,
                        grammar_follow):
    parse_table = [[""] * len(terminals) for i in
range(len(non_terminals))]

    for non_terminal in non_terminals:
        for terminal in terminals:
            if terminal in grammar_first[non_terminal]:
                rule = get_rule(non_terminal, terminal, grammar, grammar_first)
            elif ("`" in grammar_first[non_terminal]
                and terminal in grammar_follow[non_terminal]):
                rule = non_terminal + "~`"

```

```

        elif (terminal in grammar_follow[non_terminal]):
            rule = "Sync"

        else:
            rule = ""

        parse_table[non_terminals.index(non_terminal)][terminals.index(
            terminal)] = rule

    return (parse_table)

def display_parse_table(parse_table, terminal, non_terminal):
    print("\t\t\t\t", end="")
    for terminal in terminals:
        print(terminal + "\t\t", end="")
    print("\n\n")

    for non_terminal in non_terminals:
        print("\t\t" + non_terminal + "\t\t", end="")
        for terminal in terminals:
            print(parse_table[non_terminals.index(non_terminal)][terminals.index(
                terminal)] + "\t\t",
                  end="")
        print("\n")

def parse(expr, parse_table, terminals, non_terminals):
    stack = ["$"]
    stack.insert(0, non_terminals[0])

    print("\t\t\t\tMatched\t\t\t\tStack\t\t\t\tInput\t\t\t\tAction\n")
    print("\t\t\t\t-\t\t\t\t", end="")
    for i in stack:
        print(i, end="")
    print("\t\t\t\t", end="")
    print(expr + "\t\t\t\t", end="")
    print("-")

    matched = "-"
    while (True):
        action = "-"

```



```

    if (stack[0] == expr[0] and stack[0] == "$"):
        break

    elif (stack[0] == expr[0]):
        if (matched == "-"):
            matched = expr[0]
        else:
            matched = matched + expr[0]
            action = "Matched " + expr[0]
            expr = expr[1:]
            stack.pop(0)

    else:
        action =
parse_table[non_terminals.index(stack[0])][terminals.index(
    expr[0])]
        stack.pop(0)
        i = 0
        for item in action[2:]:
            if (item != "`"):
                stack.insert(i, item)
                i += 1

    print("\t\t\t" + matched + "\t\t\t", end="")
    for i in stack:
        print(i, end="")
    print("\t\t\t", end="")
    print(expr + "\t\t\t", end="")
    print(action)

grammar = OrderedDict()
grammar_first = OrderedDict()
grammar_follow = OrderedDict()

f = open('grammar.txt')
for i in f:
    i = i.replace("\n", "")
    lhs = ""
    rhs = ""
    flag = 1
    for j in i:
        if (j == "~"):

```

```

        flag = (flag + 1) % 2
        continue
    if (flag == 1):
        lhs += j
    else:
        rhs += j
    grammar = insert(grammar, lhs, rhs)
    grammar_first[lhs] = "null"
    grammar_follow[lhs] = "null"

print("Grammar\n")
show_dict(grammar)

for lhs in grammar:
    if (grammar_first[lhs] == "null"):
        grammar_first = first(lhs, grammar, grammar_first)

print("\n")
print("First\n")
show_dict(grammar_first)

start = list(grammar.keys())[0]
for lhs in grammar:
    if (grammar_follow[lhs] == "null"):
        grammar_follow = follow(lhs, grammar, grammar_follow, start)

print("\n")
print("Follow\n")
show_dict(grammar_follow)

non_terminals = list(grammar.keys())
terminals = []

for i in grammar:
    for rule in grammar[i]:
        for char in rule:

            if (isterminal(char) and char not in terminals):
                terminals.append(char)

terminals.append("$")

print("Parse Table\n\n")

```

```

parse_table = generate_parse_table(terminals, non_terminals, grammar,
                                    grammar_first, grammar_follow)
display_parse_table(parse_table, terminals, non_terminals)

expr = "i+i*i$"

print("\n\n")
print("Parsing Expression\n\n")
parse(expr, parse_table, terminals, non_terminals)

```

Output:

Grammar

```

E : TL,
L : +TL, Epsilon,
T : FK,
K : *FK, Epsilon,
F : i, (E),

```

First

```

E : i, (,
L : +, Epsilon,
T : i, (,
K : *, Epsilon,
F : i, (,

```

Follow

```

E : ), $,
L : ), $,
T : +, ), $,
K : +, ), $,
F : *, +, ), $,

```

Parse Table

	+	*	i	()	\$	
E			E~TL	E~TL		Sync	Sync
L	L~+TL					L~`	L~`
T	Sync		T~FK	T~FK		Sync	Sync
K	K~`	K~*FK				K~`	K~`
F	Sync	Sync	F~i	F~(E)		Sync	Sync

Parsing Expression

Matched	Stack	Input	Action
-	E\$	i+i*i\$	-
-	TL\$	i+i*i\$	E~TL
-	FKL\$	i+i*i\$	T~FK
-	iKL\$	i+i*i\$	F~i
i	KL\$	+i*i\$	Matched i
i	L\$	+i*i\$	K~`
i	+TL\$	+i*i\$	L~+TL
i+	TL\$	i*i\$	Matched +
i+	FKL\$	i*i\$	T~FK
i+	iKL\$	i*i\$	F~i
i+i	KL\$	*i\$	Matched i
i+i	*FKL\$	*i\$	K~*FK
i+i*	FKL\$	i\$	Matched *
i+i*	iKL\$	i\$	F~i
i+i*i	KL\$	\$	Matched i
i+i*i	L\$	\$	K~`
i+i*i	\$	\$	L~`

