Sorting ⇒ Arranging elements in order based on a property.

Eg: 1) Words in a Dictionary
2) Books in a library.

Integers.

[-5, 10, 2, 3, 0]

[-5, 0, 2, 3, 10] → Increasing
[10, 3, 2, 0, -5] → Decreasing.

[-5, 10, 0, -5, 0]

[-5, -5, 0, 0, 10] → Non decreasing
[10, 0, 0, -5, -5] → Non increasing

Ascending
Descending.

Characters → 1 Byte → 8 bits → $2^8 = 256$ → [0, 255]

a = 97	A = 65
b = 98	B = 66
⋮	⋮
z = 122	Z = 90

ASCII values

'a' < 'c'

'A' < 'a'

Strings — Alphabetic order / Lexographical

1) a l l
2) a p e
3) a p p l e

Custom Comparisom

1) Arrange the given strings in ascending order of length.

[apple, bat, dog, carrot, it]
 5 3 3 6 2

[it, bat, dog, apple, carrot]

2) Arrange integers in descending order of their absolute value.

[-10, 6, 2, -5, 0, 8]
 10 6 2 5 0 8

|-2| = 2

x → x if +ve
  → -x if -ve

[-10, 8, 6, -5, 2, 0]

Sorting functions

JAVA ⇒ Collections.sort(al) → Array List
Arrays.sort(arr) → array.

C++ ⇒ sort(a.begin(), a.end())

Python ⇒ A.sort()
A = sorted(A)

JS ⇒ A.sort()

T.C. = O(NlogN)
S.C. = O(1)
(for intermediate only)

Stable Sort

⇒ Relative order of equal elements should not change.

Q. Given an integer array of size N. Find the min cost of removing all elements from the array.

Cost of removing an element = Sum of all elements before removing this element.

A = [2, 4, 1]

Case 1

Ele	Cost	
2	2+4+1 = 7	[4, 1]
4	4+1 = 5	[1]
1	1 = 1	[]
	13	

Case 2

Ele	Cost	
4	2+4+1 = 7	[2, 1]
1	2+1 = 3	[2]
2	2 = 2	
	12	

Case 3

Ele	Cost	
4	2 + 4 + 1 = 7	[2, 1]
2	2 + 1 = 3	[1]
1	1 = 1	[]
	11	

Observations

↳ Remove elements in decending order.

[a, b, c, d]

Ele	Cost	
a → Largest	a + b + c + d	[b, c, d]
b	b + c + d	[c, d]
c	c + d	[d]
d → Smallest	d	[]

Descending. a + 2b + 3c + 4d

Goal: Cost of removing elements in descending order.

$$A = \underset{0}{[} \underset{1}{10}, \underset{2}{8}, \underset{3}{2}, \underset{4}{0}, \underset{5}{-1} \underset{}{]}$$

A[i] $\Rightarrow$ (i+1) times.

Code

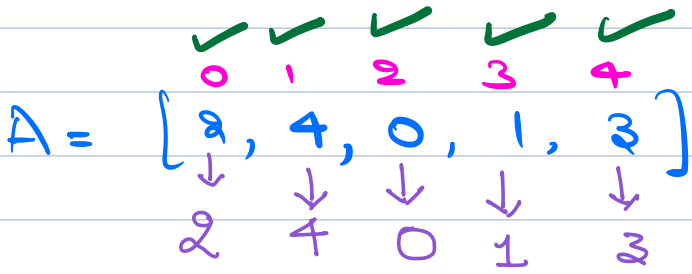
```
ans = 0
for (i = 0; i < N; i++) {
    ans += A[i] * (i+1);
}
return ans;
```

$$T.C. = O(N \log N + N)$$
$$= O(N \log N)$$

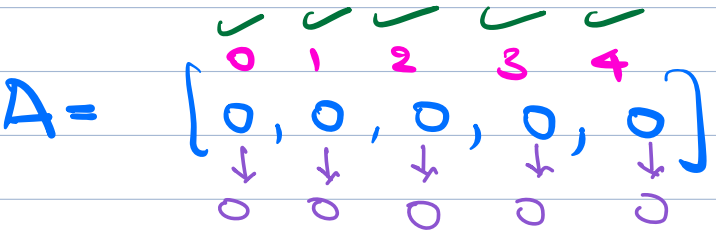
Given an integer array of size N. Count the no. of noble integers present.

Distinct elements.

Noble integers $\Rightarrow$ Count of elements $< A[i]$ = A[i]



Ans = 5



Ans = 5.

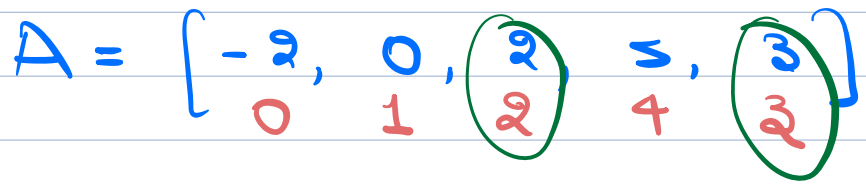
Soln

1) Brute force

⇒ $\forall i$, iterate & count no. of elements < A[i]

$O(N)$

T.C. = $O(N^2)$



2) Optimise

Count of no.s less than any element.

⇒ Sorted.

$$A = [\underset{0}{-2}, \underset{1}{0}, \underset{2}{2}, \underset{3}{3}, \underset{4}{5}]$$

$0 \quad (i-1)$ $\quad i$

$< \quad =$

$\# = i$

Code

```
count = 0;
A.sort();
for (i=0; i<N; i++) {
    if (i == A[i])
        count++;
}
return count;
```

T.C. = O(N log N)

S.C. = O(1)

If array has duplicate elements.

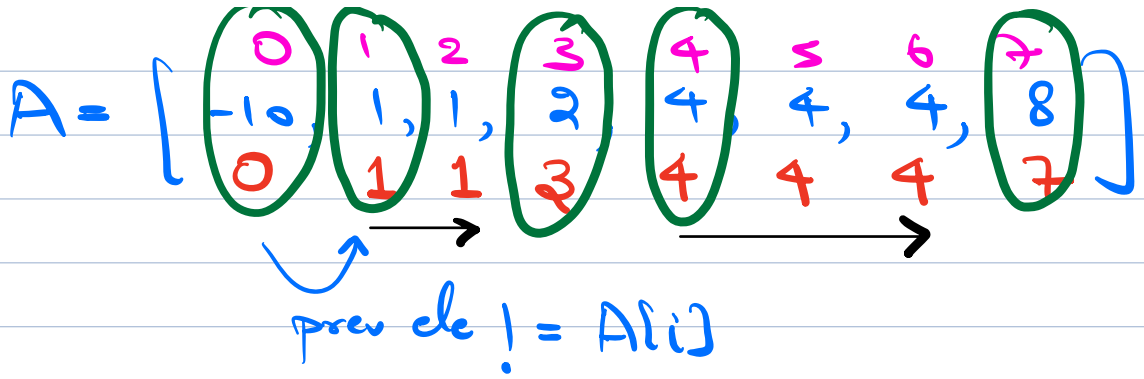
$A = [\underset{0}{-10}, \underset{1}{1}, \underset{2}{3}, \underset{3}{1}, \underset{4}{100}]$

$A = [-10, 1, 1, 3, 100]$

$\quad\; 0 \quad 1 \quad 1 \quad 3 \quad 4$

$A = [\underset{0}{0}, \underset{1}{0}, \underset{2}{0}, \underset{3}{0}, \underset{4}{0}]$

$\quad\; 0 \quad 0 \quad 0 \quad 0 \quad 0$



Code

```
A.sort();
ans = 0;
count = 0;          if (A[0] == 0) { ans++; }

for (i = 1; i < N; i++) {
        if (A[i] != A[i-1]) {
            count = i;
        }
        if (count == A[i]) {
            ans++;
        }
}
return ans;
```

T.C. = $O(N \log N)$

S.C. = $O(1)$

Given an integer array of size N (all +ve)

Sort the array (ASC) on the basis of count of factors.

A[i] <<< N

	0	1	2	3	4	5	6
A =	7	13	9	12	36	16	1
	2	2	3	6	9	5	1

A = [1, 7, 13, 9, 16, 12, 36]

Custom Comparator

```
int compare (x, y) {
```

return → -ve (x should be on left of y)

→ 0 (x = y ⇒ Keep x, y in same order)

→ +ve (y should be on left of x)

}

CC for factor Question.

```
A.sort (custom);

int custom (x, y) {                    //O(1)
    int countx = factor(x);
    int county = factor(y);            > √A[i]
    if (countx < county) {
        // x to come on left of y.
        return -1;
    }
    else if (countx == county) {
        return 0;
    }
    else {
        return 1;
    }
}
```

T.C. = $O(N \log N)$